

Boise State University Research Computing: Who we are



Elizabeth Leake, Director Research Computing Services

Directs day-to-day operations, personnel, policies, programs and infrastructure; assists with proposal development and external relations.

Steve Cutchin, Faculty Director RCS, and Associate Professor

Faculty liaison; architect for campus and regional research CI; PoC for WestNet and NSF West Big Data Innovation Hub; Scientific Viz expertise.

Jason Watt, Senior Cyberinfrastructure Engineer

Provides expertise in parallel software development; developing solutions for big data transfer, storage and retrieval; and community engagement.

Kyle Shannon, Senior Scientific Software Engineer

Assists researchers with programmingoptimization, debugging, and creating software for specific applications; provides training, and much more.



James Nelson, HPC Engineer II

Designs and supports a wide range of compute and data storage solutions; serves as an ACCESS Campus Champion; Geosciences domain expertise.

Frank Willmore, HPC Engineer IV

Primary sysadmin for the Falcon HPC system at Idaho National Lab; serves as an ACCESS Campus Champion; Chemical Engineering domain expertise.

Jenny Fothergill, HPC Engineer II

Master Software Carpentry trainer; provides support for R2 and Borah users; ACCESS Campus Champion; Chemistry and Scientific Viz expertise.

Eli Woodard, RCS Student Intern

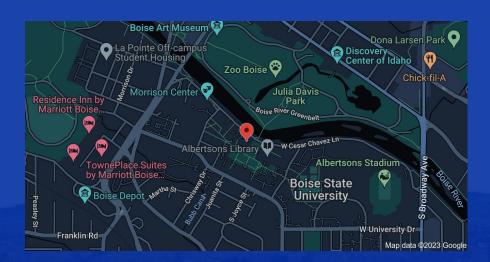
Oversees account provisioning on R2 and Borah HPC, assists with systems support and documentation, and as a student ACCESS Campus Champion.





Boise State University Research Computing: Where we are

- researchcomputing@boisestate.edu
- (208) 426-3904
- Riverfront Hall, Suite 319, 1987 W Cesar Chavez Ln, Boise, ID 83725
- https://boisestate.edu/rcs







Boise State University Research Computing: What we do

- Provider of Computing Resources (HPC clusters and more)
- Research Software Installation
- Data Management and data publication
- Help with computational components of proposal development
- Liaison for ACCESS (NSF sponsored) resources
- Training and one-to-one consultation
- Unsure? Just ask us!

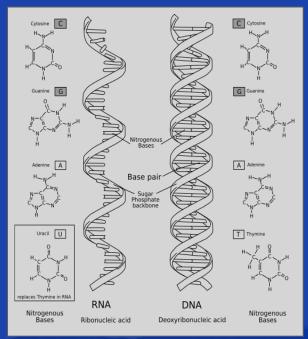


What is a number?

- Real
- Imaginary
- Complex
- Integer
- Countable
- Random
- Finite/Infinite
- Positive/Negative
- Rational/Irrational/Transcendental
- Binary/Decimal/Octal/Hexadecimal
- Exact/Approximate
- Floating point



What is a number? Storage and representation







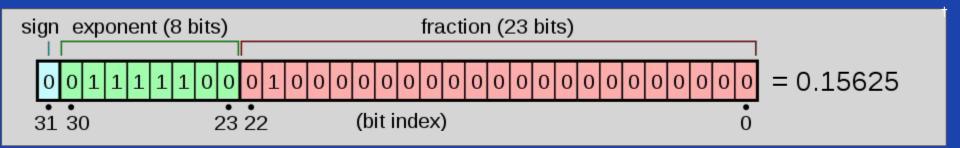
	0x0	o	0	1	1	o	o	o	0
	0x1	0	0	0	0	0	0	0	0
	0 x 2	0	0	0	0	0	1	0	0
	0x3	0	0	0	0	0	0	0	0
	0x4	0	0	1	0	0	0	0	0
	0 x 5	0	0	0	0	0	0	0	0
	0×6	0	0	0	0	0	0	0	0
	0 x 7	0	0	0	0	0	0	0	0
	0 x 8	0	0	0	1	0	0	0	0
	0x9	0	0	0	0	0	0	0	0
	0xA	0	0	1	0	0	0	0	0
	0 x B	0	0	0	1	0	0	0	0
	0xC	0	1	0	0	0	0	0	0
	0xD	0	0	0	0	1	0	0	0
	0xE	0	0	0	1	0	1	0	0
	0xF	0	1	0	0	0	1	0	0



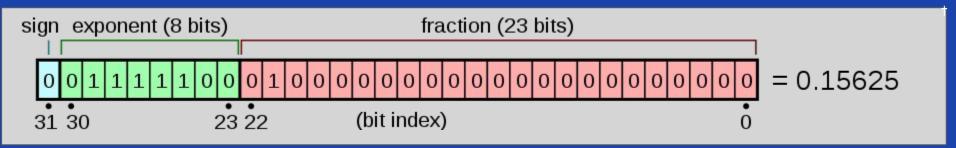
What is a number? Storage and representation





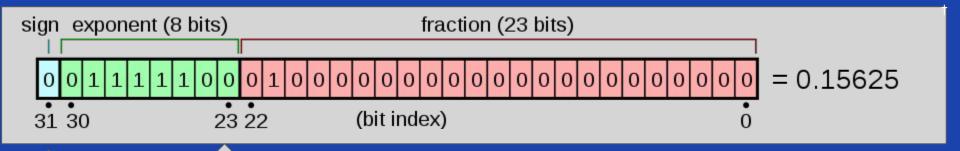








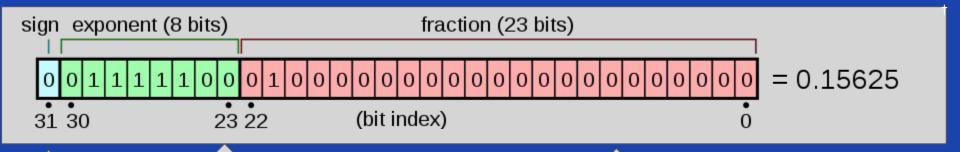






The exponent is 'offset binary' format with a shift of -127, i.e. the exponent is (0*128) + (1*64) + (1*32) + (1*16) + (1*8) + (1*4) + (0*2) + (0*1) - 127, or -3. This means a factor of 2^{-3} .

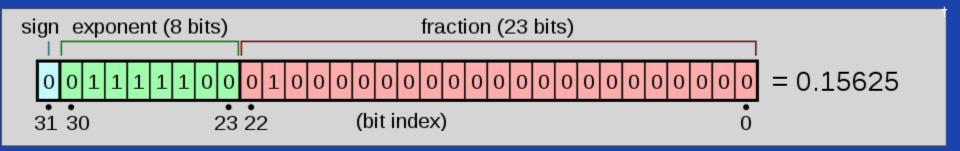




-1^{sign} = **1** The exponent is 'offset binary' format with a shift of -127, i.e. the exponent is (0*128) + (1*64) + (1*32) + (1*16) + (1*8) + (1*4) + (0*2) + (0*1) - 127, or -3. This means a factor of 2^{-3} .

Just like base 10 scientific notation, the exponent is shifted so that the first non-zero digit in the fraction is to the left of the decimal point. In binary, this value $(2^{-0}=1)$ is implicit and not recorded. Since only the 2^{-2} bit is set above, the fraction value here is $2^{-0} + 2^{-2} = 1 + \frac{1}{4} = 1.25$.





$$1 \times 2^{-3} \times 1.25 = 0.15625$$



So what does it all matter? How does this affect my research?

- Machine precision is inherently limited
- This means rounding errors
- Rounding errors are a source of experimental error (in addition to modeling shortfalls) for computational science and can affect results
- Good research accounts for measurement error



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- Good research accounts for measurement error

What can I do?

- Recognize the limitations of machine precision
- Make responsible and efficient choices regarding precision
 - Al/ML applications often use very low precision
 - GROMACS uses single precision by default
- Document choices
- Do better research



EXAMPLE: SUMMATION AND ROUNDING

1 + x
$$\rightarrow$$
 1 + x; x > δ
1 + x \rightarrow 1; x <= δ

$$y + x \rightarrow max(x,y); x/y > \delta$$

 $y + x \rightarrow x + y; x/y <= \delta$



EXERCISE: summing floats

```
#include <stdio.h>
int main()
  float ttf = 16777216; // 2^24
  float delta = 1.0f/ttf; // 0.000000059604645
  float sum = 0.0f;
  printf("Starting with sum=%1.18f, delta=%1.18f:\n\n", sum, delta);
  for (int i=0; i<ttf; i++) sum += delta;</pre>
  printf("sum=%1.18f\n", sum);
  for (int i=0; i<ttf; i++) sum += delta;</pre>
  printf("sum=%1.18f\n", sum);
```



EXERCISE: summing floats

```
id@borah-login$ dev-session-bsu
id@cpu1xx$ module load borah-base openmpi/4.1.3/gcc/12.1.0
id@cpu1xx$ git clone https://github.com/bsurc/AdvancedHPC
id@cpu1xx$ cd AdvancedHPC/machine numbers and MPI
id@cpulxx$ vi add float.c
id@cpu1xx$ make add float
id@cpu1xx$ ./add float
Starting with sum=0.000000000000000000, delta=0.000000059604644775:
```



EXERCISE: summing floats

- What happens if you increase the size of ttf?
 - Edit add_float.c and recompile
 - O What is the result?
- What happens if you decrease the size of ttf?
 - Edit add_float.c and attempt to recompile
 - What does the compiler say?
 - What does it mean?



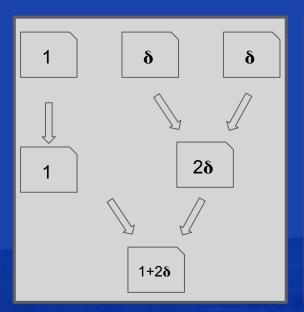
Example of MPI Reduction

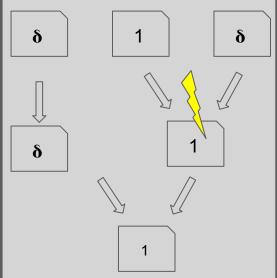
- MPI parallel processes run and complete in an indeterminate order
- Because of rounding, results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run

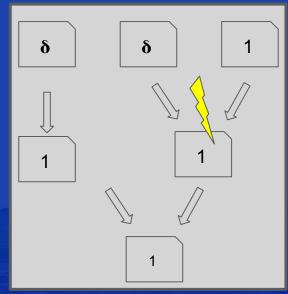


Example of MPI Reduction

- MPI parallel processes run and complete in an indeterminate order
- Because of rounding(), results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run









EXERCISE: MPI reduction

```
int main(int argc, char* argv[])
 MPI Init(NULL, NULL);
  float delta = 0.000000059604645;
  int mpi rank, mpi size;
  MPI Comm rank (MPI COMM WORLD, &mpi rank);
  MPI Comm size (MPI COMM WORLD, &mpi size);
  for (int i=0; i<mpi size; i++)</pre>
    float global sum, local val;
    if (mpi rank == i) local val = 1.0f;
    else local val = delta;
    MPI Reduce(&local val, &global sum, 1, MPI FLOAT, MPI SUM, 0, MPI COMM WORLD);
    if (!mpi rank) printf("big val to rank %d, gives global sum = %1.10f\n", i, global sum);
  MPI Finalize();
```



EXERCISE: MPI reduction

```
id@cpulxx$ vi reduce.c
id@cpulxx$ make reduce
id@cpulxx$ mpirun -n 3 ./reduce
rank 1/3 running.
rank 2/3 running.
rank 0/3 running.
big val to rank 0, gives global sum = 1.0000001192
big val to rank 1, gives global sum = 1.0000000000
big val to rank 2, gives global sum = 1.00000000000
```



EXERCISE: MPI reduction

- What happens if you run it again?
 - O Do the processes still return in the same order?
 - O What is the result?
- What happens if you increase the number of processes, e.g. to 4?
- What happens if you load a different MPI (e.g. mpich) and build/run?
 - Do OpenMPI and MPICH give the same results?
 - Why or why not?



Central Limit Theorem:

The average values of sets of samples of a randomly distributed variable will tend to be distributed normally:

 $\Re \in [0,1)$; for a uniformly distributed random variable $\Re R_j \triangleq (1/N) \Sigma_i \Re_j$; with average values R_j for sets of N samples of $\Re P(R_j) \propto \exp\{-(R_j - \langle R_j \rangle)^2/2\}$; those average values R_j are distributed normally



Central Limit Theorem example:

 $\Re \in [0,1)$; for a uniformly distributed random variable \Re $R_j \triangleq (1/N) \Sigma_i \Re_i$; with average values for a sets of N samples of $\Re p(R_i) \propto \exp\{-(R_i - \langle R_i \rangle)^2/2\}$; those average values are distributed normally

The experiment:

- Using an MPI code with different random number seeds scattered to different MPI processes, sample values of R_i are generated and gathered.
- MPI parallel processes then run and complete, although in an indeterminate order
- Because of rounding, results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run



```
#include <stdio.h>
#include <mpi.h>
#include <stdlib.h>
int main(int argc, char* argv[])
 MPI Init(NULL, NULL);
  int elements per proc = 8;
  int initial seed = 123456;
  int mpi rank, mpi size;
  MPI Comm rank (MPI COMM WORLD, &mpi rank);
  MPI Comm size (MPI COMM WORLD, &mpi size);
  if (!mpi rank) printf("running %d elements per task with %d tasks.\n", elements per proc, mpi size);
  int scatter seeds[mpi size];
  scatter seeds[0] = initial seed;
  if (!mpi rank) for (int i=1; i<mpi size; i++) scatter seeds[i] = scatter seeds[0] + i;
  int seed;
  MPI Scatter(scatter seeds, 1, MPI INT, &seed, 1, MPI INT, 0, MPI COMM WORLD);
```



```
// seed the RNG
srand(seed);

float sum = 0;
for (int i=0; i<elements_per_proc; i++)
{
    float sample = (float)rand() / (float)RAND_MAX;
    sum += sample;
}
sum /= elements_per_proc;

// gather and bin the results:
float results[mpi_size];
MPI_Gather(&sum, 1, MPI_FLOAT, results, 1, MPI_FLOAT, 0, MPI_COMM_WORLD);</pre>
```



```
int bins[10];
for (int i=0; i<10; i++) bins[i] = 0;
if (!mpi rank) for (int i=0; i<mpi size; i++)</pre>
  printf("binning %f\n", results[i]);
  bins[(int) (results[i]*10)]++;
if (!mpi rank) for (int i=0; i<10; i++)</pre>
  for (int j=0; j<bins[i]; j++) printf("X");</pre>
  printf("\n");
MPI Finalize();
```



```
id@cpu1xx$ vi central limit.c
id@cpu1xx$ make central limit
id@cpu1xx$ mpirun -n 48 ./central limit
running 8 elements per task with 48 tasks.
binning 0.488702
binning 0.511513
XX
XX
XXXXXXXXXXXXX
XXXXXXX
XX
X
```



- What are the likely sources of error and uncertainty?
- What happens if you run it again for a different seed value?
- What happens if you increase the number of MPI processes?
- What happens if you use a different number of samples for each R_i?
- What happens if we have very small values of

 π?



Takeaways:

- Greater precision isn't always 'better',
- Knowing when you can use a smaller type can often mean better performance
- Making good choices here depends on understanding the limits of machine numbers
- Operations can give results that vary between runs and depend on size of run
- 'Randomness' is not always an implicit source of error, it depends on how it's used.
- There is a new floating point standard in play (posits) which looks to gain traction
- Limits of machine precision are inherent, no matter the standard



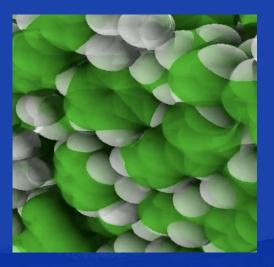
Wait, there's more!



- Q: When is it a good idea to use a GPU?
- A: For anything that looks (or can look) like a rendering problem.



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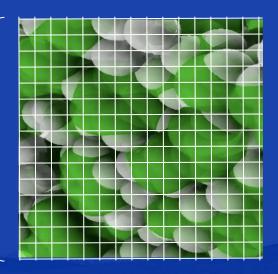
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Domain decomposition



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Domain decomposition



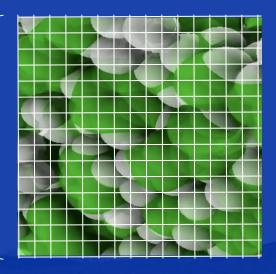
Chemical potential:

$$\mu_i = \left(rac{\partial U}{\partial N_i}
ight)_{S,V,N_{j
eq i}}. \;\; \mu_i = -k_BT\ln\!\left(rac{{f B}_i}{
ho_i\lambda^3}
ight)$$



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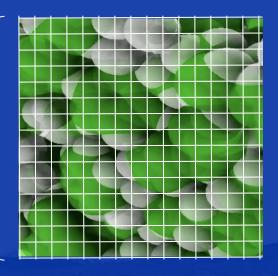
Wisdom insertion parameter:

$$\mathbf{B}_i = rac{
ho_i}{a_i} = \left\langle \exp\!\left(-rac{\psi_i}{k_B T}
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ight)
ight
angle$$

Free volume index:

$$\phi_i(x,y,z) = e^{-\beta \gamma_i}$$



EXERCISE: free volume *via* GPU

```
username@borah-login$ qpu-session
username@gpuxxx$ module load borah-misc vacuumms
username@gpuxxx$ cd ~/AdvancedHPC/GPU fun
username@gpuxxx$ head PS.gfg
33.637309
           36.915309
                       35.467309
                                   3.520530
                                               0.073457
36.121309 38.629309
                       34.794309
                                   2.373410
                                               0.028294
34.127309 39.183309
                       34.275309
                                   3.581180
                                               0.066301
33.785309
           38.464309
                       35.581309
                                   3.581180
                                               0.066301
39.593309
           33.823309
                       36.362309
                                   2.373410
                                               0.028294
33.439309
           38.942309
                       33.477309
                                   2.373410
                                               0.028294
35.546309
           38.832309
                       33.826309
                                   3.581180
                                               0.066301
35.729309
           40.915309
                       33.360309
                                   2.373410
                                               0.028294
34.054309
           40.296309
                       34.378309
                                   2.373410
                                               0.028294
34.658309
           38.544309
                       36.220309
                                   2.373410
                                               0.028294
username@gpuxxx$ ./rungfg2fvi.sh &
reading configuration
calculating resolution = 256 for 612 potential
using sigma = 0.000000 and epsilon = 1.000000
username@gpuxxx$
```



EXERCISE: free volume *via* GPU

```
username@gpuxxx$ nvidia-smi
Thu Mar 23 11:17:06 2023
 NVIDIA-SMI 510.39.01
                    Driver Version: 510.39.01 CUDA Version: 11.6
 GPU Name
               Persistence-M| Bus-Id
                                      Disp.A | Volatile Uncorr. ECC |
 Fan Temp Perf Pwr:Usage/Cap| Memory-Usage | GPU-Util Compute M. |
                                                         MIG M. |
0 Tesla V100-PCIE... On | 00000000:3B:00.0 Off |
                                                            Off
 N/A 47C
            PO 215W / 250W | 1365MiB / 16384MiB | 100%
                                                        Default |
                                                            N/A I
   1 Tesla V100-PCIE... On | 00000000:D8:00.0 Off |
                                                            Off
                               4MiB / 16384MiB |
 N/A 30C
                23W / 250W |
                                                  0 %
                                                        Default |
                                                            N/A I
 Processes:
  GPU GI
          CI
                   PID Type
                              Process name
                                                      GPU Memory
    0 N/A N/A
                 99430
                          C gfg2fvi
username@gpuxxx$
```



Wrap-up

Interesting GPU fact:

- The first GPUs didn't perform compliant arithmetic operations
- IEEE Compliant operations require:
 - padding 32 bit float to a 40 bit type
 - Performing arithmetic on 40 bit type
 - Rounding the result back to a 32 bit type
- Since they were only shading pixels, a little round-off error didn't matter much

Additional resources:

- https://en.wikipedia.org/wiki/IEEE_754
- https://posithub.org/docs/posit_standard-2.pdf
- https://www.mpi-forum.org/docs/

Contact us!

- https://www.boisestate.edu/rcs/
- researchcomputing@boisestate.edu

