



RC Days 2023

Advanced topics in HPC:

Numerical accuracy, lies and statistics

Dr. Frank T Willmore
Boise State University Research Computing
FrankWillmore@boisestate.edu

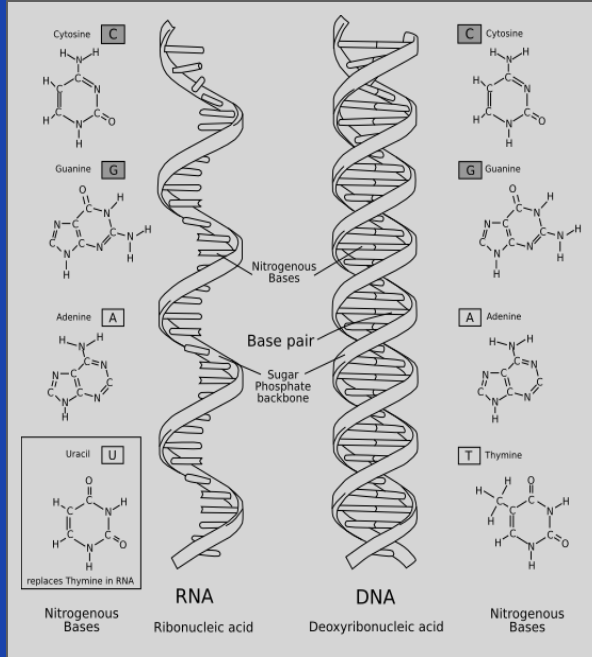


BOISE STATE UNIVERSITY

What is a number?

- Real
- Imaginary
- Complex
- Integer
- Countable
- Random
- Finite/Infinite
- Positive/Negative
- Rational/Irrational/Transcendental
- Binary/Decimal/Octal/Hexadecimal
- Exact/Approximate
- Floating point

What is a number? Storage and representation



https://en.wikipedia.org/wiki/Data_storage



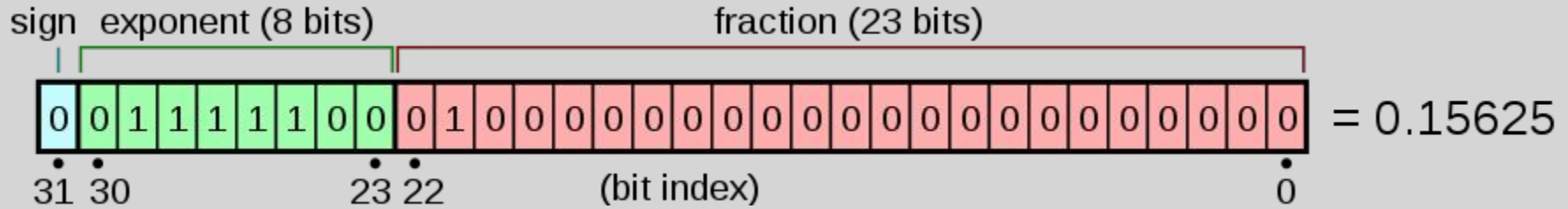
0x0	0	0	1	1	0	0	0	0
0x1	0	0	0	0	0	0	0	0
0x2	0	0	0	0	0	1	0	0
0x3	0	0	0	0	0	0	0	0
0x4	0	0	1	0	0	0	0	0
0x5	0	0	0	0	0	0	0	0
0x6	0	0	0	0	0	0	0	0
0x7	0	0	0	0	0	0	0	0
0x8	0	0	0	1	0	0	0	0
0x9	0	0	0	0	0	0	0	0
0xA	0	0	1	0	0	0	0	0
0xB	0	0	0	1	0	0	0	0
0xC	0	1	0	0	0	0	0	0
0xD	0	0	0	0	1	0	0	0
0xE	0	0	0	1	0	1	0	0
0xF	0	1	0	0	0	1	0	0

What is a number? Storage and representation

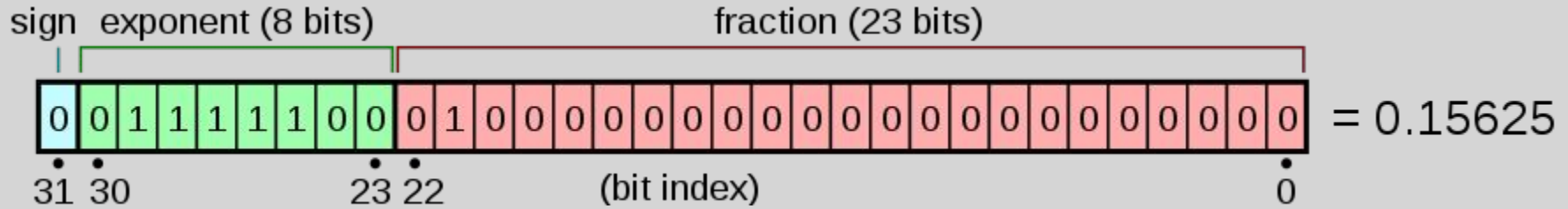


0x0	0	0	1	1	0	0	0	0
0x1	0	0	0	0	0	0	0	0
0x2	0	0	0	0	0	1	0	0
	0	0	0	0	0	0	0	0
	0	0	1	0	0	0	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
0x8	0	0	0	1	0	0	0	0
0x9	0	0	0	0	0	0	0	0
0xA	0	0	1	0	0	0	0	0
0xB	0	0	0	1	0	0	0	0
0xC	0	1	0	0	0	0	0	0
0xD	0	0	0	0	1	0	0	0
0xE	0	0	0	1	0	1	0	0
0xF	0	1	0	0	0	1	0	0

What is a number? Machine representation of a floating point number

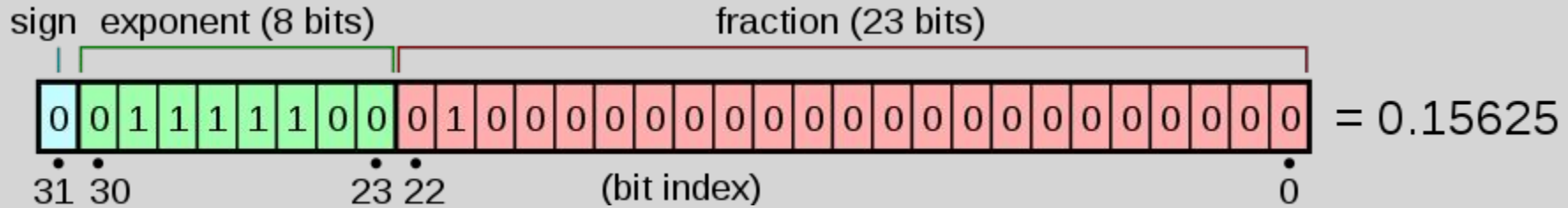


What is a number? Machine representation of a floating point number



-1^{sign}
 $= 1$

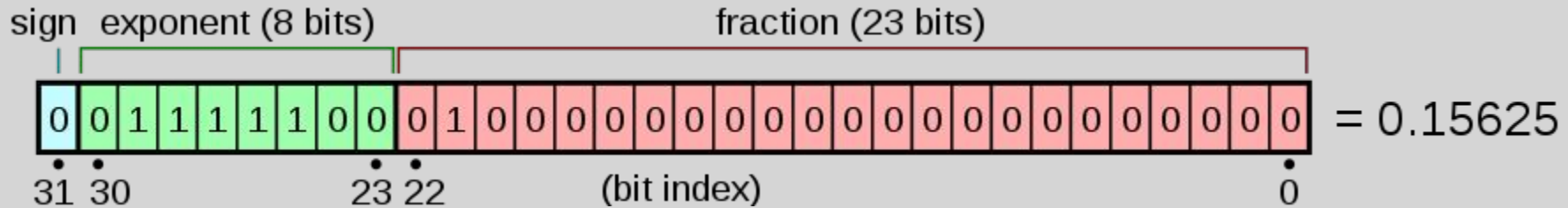
What is a number? Machine representation of a floating point number



$$-1^{\text{sign}} = 1$$

The exponent is 'offset binary' format with a shift of -127, i.e. the exponent is $(0 \cdot 128) + (1 \cdot 64) + (1 \cdot 32) + (1 \cdot 16) + (1 \cdot 8) + (1 \cdot 4) + (0 \cdot 2) + (0 \cdot 1) - 127$, or **-3**. This means a factor of 2^{-3} .

What is a number? Machine representation of a floating point number

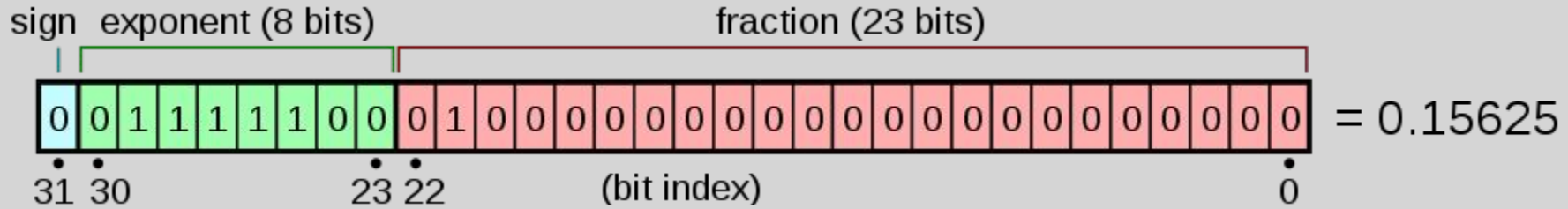


$$-1^{\text{sign}} = 1$$

The exponent is 'offset binary' format with a shift of -127, i.e. the exponent is $(0 \cdot 128) + (1 \cdot 64) + (1 \cdot 32) + (1 \cdot 16) + (1 \cdot 8) + (1 \cdot 4) + (0 \cdot 2) + (0 \cdot 1) - 127$, or **-3**. This means a factor of **2^{-3}** .

Just like base 10 scientific notation, the exponent is shifted so that the first non-zero digit in the fraction is to the left of the decimal point. In binary, this value ($2^0=1$) is implicit and not recorded. Since only the 2^{-2} bit is set above, the fraction value here is $2^0 + 2^{-2} = 1 + \frac{1}{4} = \mathbf{1.25}$.

What is a number? Machine representation of a floating point number



$$1 \times 2^{-3} \times 1.25 = 0.15625$$

So what does it all matter? How does this affect my research?

- Machine precision is inherently limited
- This means rounding errors
- Rounding errors are a source of experimental error (in addition to modeling shortfalls) for computational science and can affect results
- Good research accounts for measurement error



So what does it all matter? How does this affect my research?

- Machine precision is inherently limited
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What can I do?

- Recognize the limitations of machine precision
- Make responsible and efficient choices regarding precision
 - AI/ML applications often use very low precision
 - GROMACS uses single precision by default
- Document choices
- Do better research

EXAMPLE: SUMMATION AND ROUNDING

$$1 + x \rightarrow 1 + x; x > \delta$$

$$1 + x \rightarrow 1; x \leq \delta$$

$$y + x \rightarrow \max(x, y); x/y > \delta$$

$$y + x \rightarrow x + y; x/y \leq \delta$$

EXERCISE: summing floats

```
#include <stdio.h>

int main()
{
    float ttf = 16777216;    // 2^24
    float delta = 1.0f/ttf;  // 0.000000059604645
    float sum = 0.0f;

    // check inputs
    printf("Starting with sum=%1.18f, delta=%1.18f:\n\n", sum, delta);

    // should add to one
    for (int i=0; i<ttf; i++) sum += delta;
    printf("sum=%1.18f\n", sum);

    // what happens if we keep adding?
    for (int i=0; i<ttf; i++) sum += delta;
    printf("sum=%1.18f\n", sum);
}
```


EXERCISE: summing floats

```
id@borah-login$ dev-session-bsu
...
id@cpulxx$ module load borah-base openmpi/4.1.3/gcc/12.1.0
id@cpulxx$ git clone https://github.com/bsurc/AdvancedHPC
id@cpulxx$ cd AdvancedHPC/machine_numbers_and_MPI
id@cpulxx$ vi add_float.c
id@cpulxx$ make add_float
id@cpulxx$ ./add_float
Starting with sum=0.000000000000000000, delta=0.000000059604644775:

sum=1.000000000000000000
sum=1.000000000000000000
```

EXERCISE: summing floats

- What happens if you increase the size of ttf?
 - Edit add_float.c and recompile
 - What is the result?
- What happens if you decrease the size of ttf?
 - Edit add_float.c and attempt to recompile
 - What does the compiler say?
 - What does it mean?

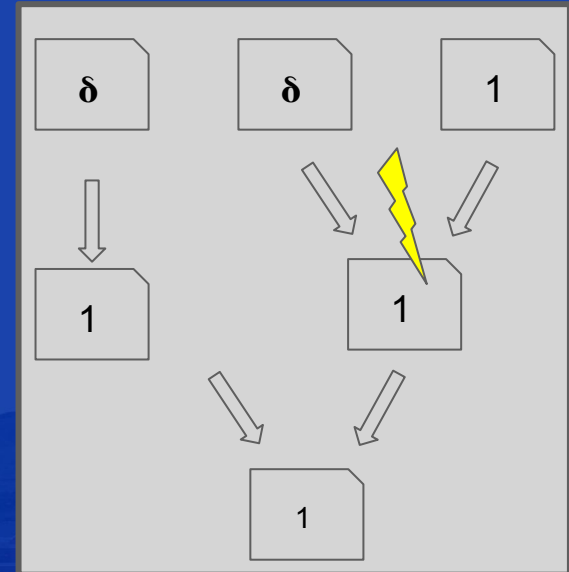
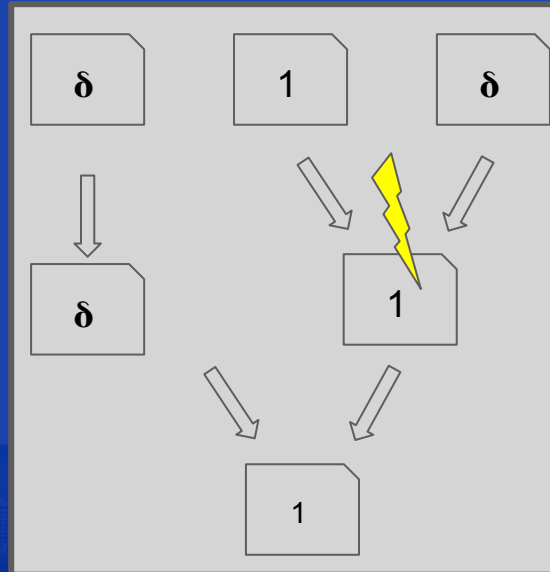
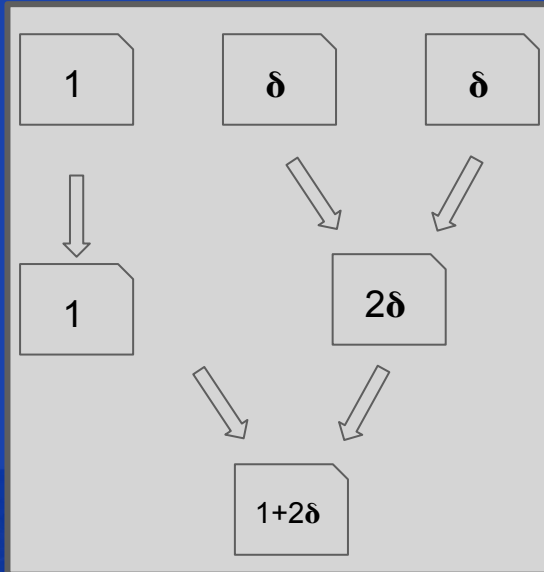
Example of MPI Reduction

- MPI parallel processes run and complete in an indeterminate order
- Because of rounding, results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run



Example of MPI Reduction

- MPI parallel processes run and complete in an indeterminate order
- Because of rounding(⚡), results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run



EXERCISE: MPI reduction

```
#include <stdio.h>
#include <mpi.h>

int main(int argc, char* argv[])
{
    MPI_Init(NULL, NULL);

    float delta = 0.000000059604645;
    int mpi_rank, mpi_size;
    MPI_Comm_rank(MPI_COMM_WORLD, &mpi_rank);
    MPI_Comm_size(MPI_COMM_WORLD, &mpi_size);

    for (int i=0; i<mpi_size; i++)
    {
        float global_sum, local_val;
        if (mpi_rank == i) local_val = 1.0f;
        else local_val = delta;
        MPI_Reduce(&local_val, &global_sum, 1, MPI_FLOAT, MPI_SUM, 0, MPI_COMM_WORLD);
        if (!mpi_rank) printf("big val to rank %d, gives global sum = %1.10f\n", i, global_sum);
    }

    MPI_Finalize( );
}
```


EXERCISE: MPI reduction

```
id@cpulxx$ vi reduce.c
id@cpulxx$ make reduce
id@cpulxx$ mpirun -n 3 ./reduce
rank 1/3 running.
rank 2/3 running.
rank 0/3 running.
big val to rank 0, gives global sum = 1.0000001192
big val to rank 1, gives global sum = 1.0000000000
big val to rank 2, gives global sum = 1.0000000000
```

EXERCISE: MPI reduction

- What happens if you run it again?
 - Do the processes still return in the same order?
 - What is the result?
- What happens if you increase the number of processes, e.g. to 4?
- What happens if you load a different MPI (e.g. mpich) and build/run?
 - Do OpenMPI and MPICH give the same results?
 - Why or why not?

Central Limit Theorem:

The average values of sets of samples of a randomly distributed variable will tend to be distributed normally:

$\mathcal{R} \in [0,1)$; for a uniformly distributed random variable \mathcal{R}

$R_j \triangleq (1/N) \sum_i \mathcal{R}_i$; with average values R_j for sets of N samples of \mathcal{R}

$p(R_j) \propto \exp\{-(R_j - \langle R_j \rangle)^2/2\}$; those average values R_j are distributed normally

Central Limit Theorem example:

$\mathcal{R} \in [0,1)$; for a uniformly distributed random variable \mathcal{R}

$R_j \triangleq (1/N) \sum_i \mathcal{R}_i$; with average values for a sets of N samples of \mathcal{R}

$p(R_j) \propto \exp\{-(R_j - \langle R_j \rangle)^2/2\}$; those average values are distributed normally

The experiment:

- Using an MPI code with different random number seeds scattered to different MPI processes, sample values of R_j are generated and gathered.
- MPI parallel processes then run and complete, although in an indeterminate order
- Because of rounding, results of a reduction depend on the order in which operations are performed
- MPI operations can give results that vary between runs and depend on size of run

EXERCISE: central limit theorem

```
#include <stdio.h>
#include <mpi.h>
#include <math.h>
#include <stdlib.h>

int main(int argc, char* argv[])
{
    MPI_Init(NULL, NULL);

    int elements_per_proc = 8;
    int initial_seed = 123456;

    int mpi_rank, mpi_size;
    MPI_Comm_rank(MPI_COMM_WORLD, &mpi_rank);
    MPI_Comm_size(MPI_COMM_WORLD, &mpi_size);

    if (!mpi_rank) printf("running %d elements per task with %d tasks.\n", elements_per_proc, mpi_size);

    // create seeds
    int scatter_seeds[mpi_size];
    scatter_seeds[0] = initial_seed;
    if (!mpi_rank) for (int i=1; i<mpi_size; i++) scatter_seeds[i] = scatter_seeds[0] + i;

    // scatter the seeds
    int seed;
    MPI_Scatter(scatter_seeds, 1, MPI_INT, &seed, 1, MPI_INT, 0, MPI_COMM_WORLD);
```


EXERCISE: central limit theorem

```
// seed the RNG
srand(seed);

float sum = 0;
for (int i=0; i<elements_per_proc; i++)
{
    float sample = (float)rand() / (float)RAND_MAX;
    sum += sample;
}
sum /= elements_per_proc;

// gather and bin the results:
float results[mpi_size];
MPI_Gather(&sum, 1, MPI_FLOAT, results, 1, MPI_FLOAT, 0, MPI_COMM_WORLD);
```

EXERCISE: central limit theorem

```
int bins[10];
for (int i=0; i<10; i++) bins[i] = 0;

if (!mpi_rank) for (int i=0; i<mpi_size; i++)
{
    printf("binning %f\n", results[i]);
    bins[(int)(results[i]*10)]++;
}

if (!mpi_rank) for (int i=0; i<10; i++)
{
    for (int j=0; j<bins[i]; j++) printf("X");
    printf("\n");
}

MPI_Finalize();
}
```

EXERCISE: central limit theorem

```
id@cpulxx$ vi central_limit.c
id@cpulxx$ make central_limit
id@cpulxx$ mpirun -n 48 ./central_limit
running 8 elements per task with 48 tasks.
binning 0.488702
binning 0.511513
. . .

XX
XX
XXXXXXXXXXXXXXXXXX
XXXXXXXXXXXXXXXXXXXXXXXXXXXX
XXXXXXX
XX
X
```

EXERCISE: central limit theorem

- What are the likely sources of error and uncertainty?
- What happens if you run it again for a different seed value?
- What happens if you increase the number of MPI processes?
- What happens if you use a different number of samples for each R_j ?
- What happens if we have very small values of \mathcal{R} ?

Takeaways:

- Greater precision isn't always 'better'
- Knowing when you can use a smaller type can often mean better performance
- Making good choices here depends on understanding the limits of machine numbers
- Operations can give results that vary between runs and depend on size of run
- 'Randomness' is not always an implicit source of error, it depends on how it's used.
- There is a new floating point standard in play (posits) which looks to gain traction
- Limits of machine precision are inherent, no matter the standard



Wait, there's more!



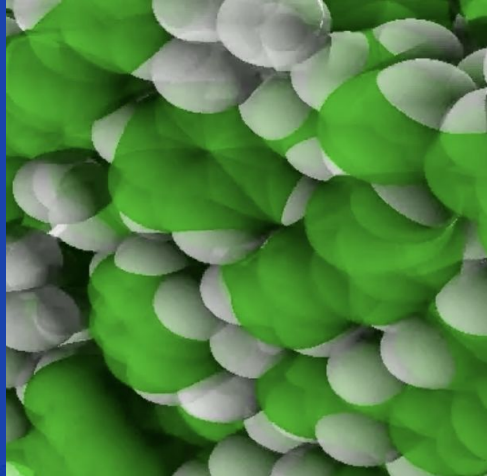
BONUS TOPIC! CUDA/GPU application

- Q: When is it a good idea to use a GPU?
- A: For anything that looks (or can look) like a rendering problem.



BONUS TOPIC! CUDA/GPU application

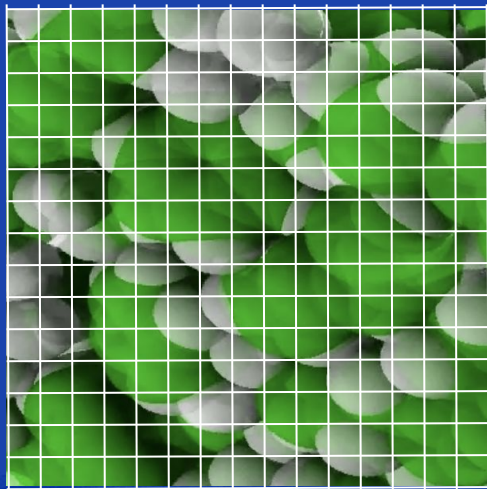
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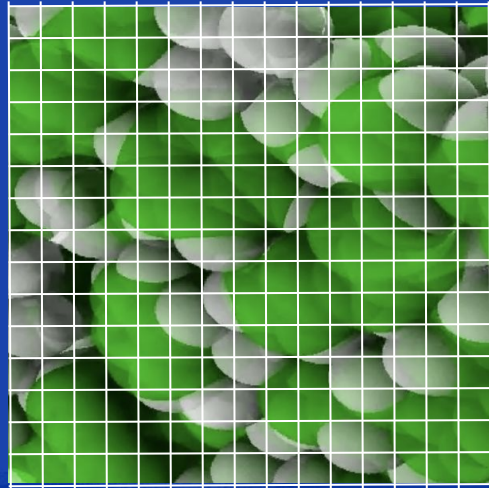
Domain decomposition



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Domain decomposition



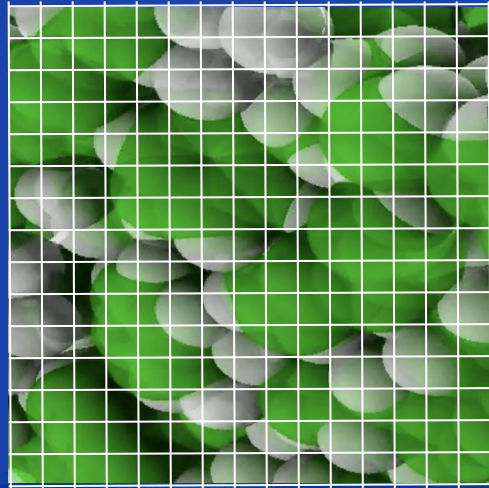
Chemical potential:

$$\mu_i = \left(\frac{\partial U}{\partial N_i} \right)_{S,V,N_{j \neq i}} \quad \mu_i = -k_B T \ln \left(\frac{\mathbf{B}_i}{\rho_i \lambda^3} \right)$$

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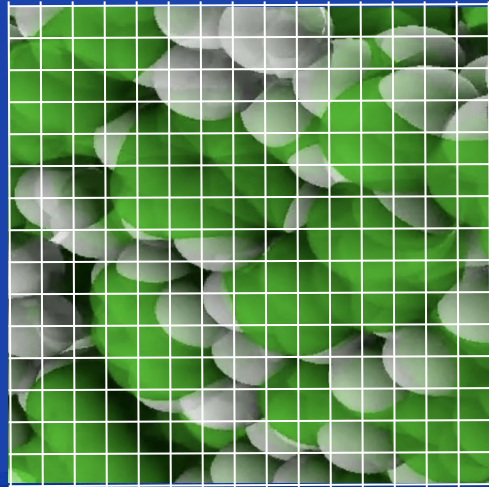
Wisdom insertion parameter:

$$\mathbf{B}_i = \frac{\rho_i}{a_i} = \left\langle \exp \left(-\frac{\psi_i}{k_B T} \right) \right\rangle$$

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$$\mu_i = \left(\frac{\partial U}{\partial N_i} \right)_{S,V,N_{j \neq i}} \cdot \mu_i = -k_B T \ln \left(\frac{\mathbf{B}_i}{\rho_i \lambda^3} \right)$$

Wisdom insertion parameter:

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Free volume index:

$$\phi_i(x,y,z) = e^{-\beta \psi_i}$$

EXERCISE: free volume *via* GPU

```
username@borah-login$ gpu-session
username@gpuxxx$ module load borah-misc vacuumms
username@gpuxxx$ cd ~/AdvancedHPC/GPU_fun
username@gpuxxx$ head PS.gfg
33.637309    36.915309    35.467309    3.520530    0.073457
36.121309    38.629309    34.794309    2.373410    0.028294
34.127309    39.183309    34.275309    3.581180    0.066301
33.785309    38.464309    35.581309    3.581180    0.066301
39.593309    33.823309    36.362309    2.373410    0.028294
33.439309    38.942309    33.477309    2.373410    0.028294
35.546309    38.832309    33.826309    3.581180    0.066301
35.729309    40.915309    33.360309    2.373410    0.028294
34.054309    40.296309    34.378309    2.373410    0.028294
34.658309    38.544309    36.220309    2.373410    0.028294
username@gpuxxx$ ./rungfg2fvi.sh &
reading configuration
calculating resolution = 256 for 612 potential
using sigma = 0.000000 and epsilon = 1.000000
username@gpuxxx$
```

EXERCISE: free volume *via* GPU

```
username@gpuxxx$ nvidia-smi
Thu Mar 23 11:17:06 2023

+-----+
| NVIDIA-SMI 510.39.01      Driver Version: 510.39.01      CUDA Version: 11.6      |
+-----+
| GPU  Name                Persistence-M| Bus-Id        Disp.A | Volatile Uncorr. ECC |
| Fan  Temp  Perf    Pwr:Usage/Cap|      Memory-Usage | GPU-Util  Compute M. |
|                                       |                    |    MIG M.     |
+=====+
|   0   Tesla V100-PCIE...    On      | 00000000:3B:00.0 Off |                    |    Off           |
| N/A   47C    P0     215W / 250W | 1365MiB / 16384MiB |   100%      Default |
|                                       |                    |    N/A           |
+-----+
|   1   Tesla V100-PCIE...    On      | 00000000:D8:00.0 Off |                    |    Off           |
| N/A   30C    P0     23W / 250W |   4MiB / 16384MiB |    0%      Default |
|                                       |                    |    N/A           |
+-----+

+-----+
| Processes:                                     |
|  GPU   GI    CI          PID    Type   Process name                  GPU Memory |
|        ID    ID                                   Usage                    |
+=====+
|    0   N/A  N/A       99430     C      gfg2fvi                      1361MiB |
+-----+

username@gpuxxx$
```


Wrap-up

Interesting GPU fact:

- The first GPUs didn't perform compliant arithmetic operations
- IEEE Compliant operations require:
 - padding 32 bit float to a 40 bit type
 - Performing arithmetic on 40 bit type
 - Rounding the result back to a 32 bit type
- Since they were only shading pixels, a little round-off error didn't matter much

Additional resources:

- https://en.wikipedia.org/wiki/IEEE_754
- https://posithub.org/docs/posit_standard-2.pdf
- <https://www.mpi-forum.org/docs/>