Thread (I) Kernel Threads

Example of Multi-Process Collaboration

- Finding the lines containing key word "Linux" in all .txt files in the current directory
- Linux command line:
 grep "Linux" *.txt
- Using multi-process collaboration to do the job?

Multi-Process

Stack

Heap

Data

Code (Text)

PCB

Register Values (including PC)

Information about memory, files,

. . .

Stack

Heap

Data

Code (Text)

PCB

Register Values (including PC)

Information about memory, files, Stack

Heap

Data

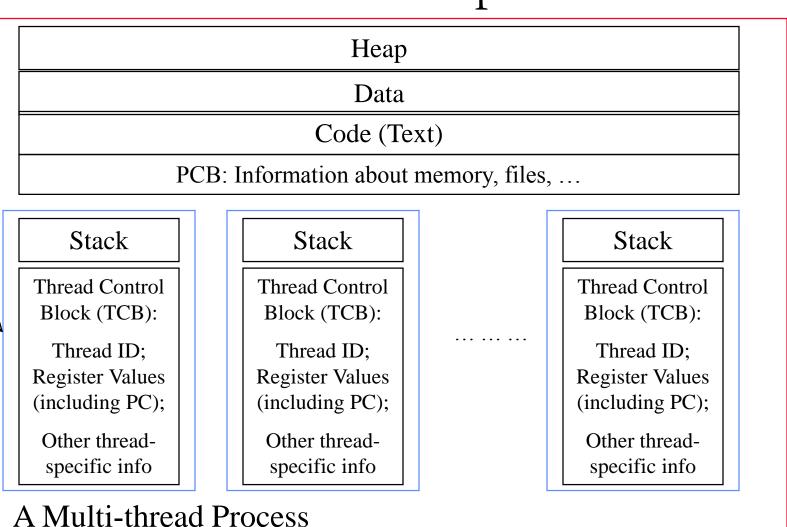
Code (Text)

PCB

Register Values (including PC)

Information about memory, files,

A Process with Multiple Threads

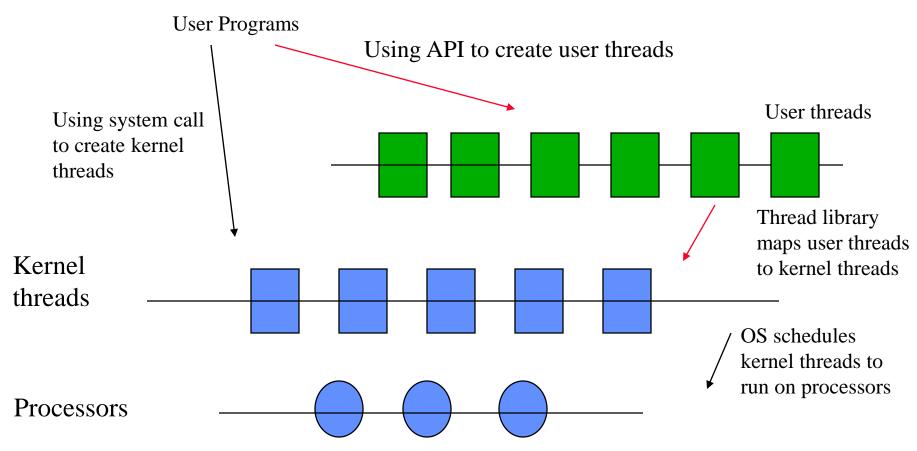


Thread

Benefits

- Responsiveness (similar to multi-process)
 - Allowing a program to continue running even if part of it is blocked
 - Allowing a program to perform multiple tasks concurrently
- Resource sharing
 - Threads of the same process share the memory and resources
- Economy
 - Allocating memory and resources for process creation is costly.
- Scalability (similar to multi-process)
 - A single-thread process can only run on one processor regardless how many are available; multi-threaded process can increase parallelism on a multi-processor machine.

Threads in a Computer System



Kernel Threads

- Directly created/managed/scheduled by the OS kernel
- Virtually all contemporary OSes support kernel threads

Linux Kernel Threads

- Linux refers to them as tasks rather than threads
- Thread creation is done through clone() system call

```
#define _GNU_SOURCE
```

#include <sched.h>

int clone(int (*fn)(void *), void *child_stack, int flags, void *arg);

flag	meaning
CLONE_FS	File-system information is shared.
CLONE_VM	The same memory space is shared.
CLONE_SIGHAND	Signal handlers are shared.
CLONE_FILES	The set of open files is shared.

```
#include _GNU_SOURCE
#include <stdio.h>
#include <sched.h>
#include <unistd.h>
#include <stdlib.h>
#include <fcntl.h>

int variable, fd;
int do_something(); //function to be executed by the new task (thread)
```

```
int main() {
void *child_stack;
char tempch;
variable = 9;
fd = open("test.file", O_RDONLY);
child_stack = (void *)malloc(16384); child_stack += 16383;
printf("The variable was %d\n", variable);
clone(do_something, child_stack, CLONE_VM | CLONE_FILES, NULL); /*A*/
sleep(1);
printf("The variable is now %d\n", variable);
if (read(fd, &tempch, 1) \leq 1) {
      perror("File Read Error"); exit(1);}
      printf("Parent could read from the file\n");}
```

```
int do_something() {
  char tempch;
  variable = 42;
 if (read(fd, &tempch, 1) \leq 1) {
    perror("File Read Error");
   _exit(1);
  printf("Child could read from the file\n");
  close(fd);
 _exit(0);
```

1. What is the output of the program?

```
2. What if Line A is changed to: clone(do_something, child_stack, CLONE_VM, NULL); or clone(do_something, child_stack, CLONE_FILES, NULL); or
```

Clone(do_something, child_stack, 0, NULL);

Thread (II) User Threads & User/Kernel Mapping

User Threads

☑ Created by a thread library and scheduling is managed by the library itself in user space (the existence of user threads is unknown to the OS)

User Thread Libraries

- Pthreads (POSIX Threads)
 - A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization
 - API specifies behavior of the thread library, implementation is up to development of the library
 - Common in UNIX operating systems (Solaris, Linux, Mac OS X)
- Java Threads
 - Managed by the JVM
 - May be created by:
 - Extending Thread class
 - Implementing the Runnable interface
- Win32 Threads
 - Similar to pthread, differing in function names

PThreads # include <pthread.h>

- * Creates a new thread
- * 1st argument: the address of the created thread
- * 2nd argument: the address of the structure containing the attributes of the created thread
 - * For default attribute values, make this argument as **NULL**
- * 3rd argument: the pointer to the function that is the code for the created thread
- * 4th argument: the pointer to the argument of the start_routine function. If multiple arguments are needed, define a data structure that contains all arguments.

PThreads # include <pthread.h>

- * void pthread_exit (void *status)
 - * Terminate execution of the calling thread
 - * 1st argument: Points to an optional termination status. If no termination status is desired, its value should be NULL.
- pthread_t pthread_self (void)
 - * returns the *pthread* handle (ID) of the calling thread

PThreads # include <pthread.h>

- * int pthread_join (pthread_t tid, void **status)
 - ★ "Joining" is one way to accomplish synchronization between threads.
 - * Blocks the caller until the specified thread terminates
 - * 1st argument: the *id* of the thread to be waited
 - * 2nd argument: the address of the variable to receive the thread's exit status
 - **★** Usually set to **NULL** or **int***

Example: Hello World!

This simple example code creates 5 threads with the pthread_create() routine. Each thread prints a "Hello World!" message, and then terminates with a call to pthread_exit().

Example: Hello World! (con't)

```
int main (int argc, char *argv[])
     pthread_t threads[NUM_THREADS];
     int rc, t;
     for (t=0;t < NUM_THREADS;t++) {
        printf("Creating thread %d\n", t);
        rc = pthread_create(&threads[t], NULL, PrintHello, (void *) t);
        if (rc){
          printf("ERROR; return code from pthread_create() is %d\n", rc);
          exit(-1);
     for(t=0; t<NUM_THREADS;t++) pthread_join(threads[t],NULL);
     pthread_exit(NULL);
                                                                    20
```

Example: Compilation

Suppose the program file is named: HelloWorld.c

\$gcc –o HelloWorld HelloWorld.c -lpthread

Java Threads

```
Defining a class X that implements the Runnable
  interface to contain the implementation of a thread
  public interface Runnable
      public abstract void run();
Creating a Thread object to wrap the class X
Manipulate the Thread object by calling its methods
   亙 run( ) method
   join() method
   51 . . . . . .
```

Example: Summation

```
class Summation implements Runnable
  private int upper;
  public int sum;
  public Summation(int upper, int sum) {
       this.upper = upper;
       this.sum = sum;
  public void run() {
       int sum = 0;
       for (int i=0; i\le upper; i++) sum +=i;
       System.out.println("Sum of 1 through " + upper +" is "+sum);
                                                                23
```

Example: Summation

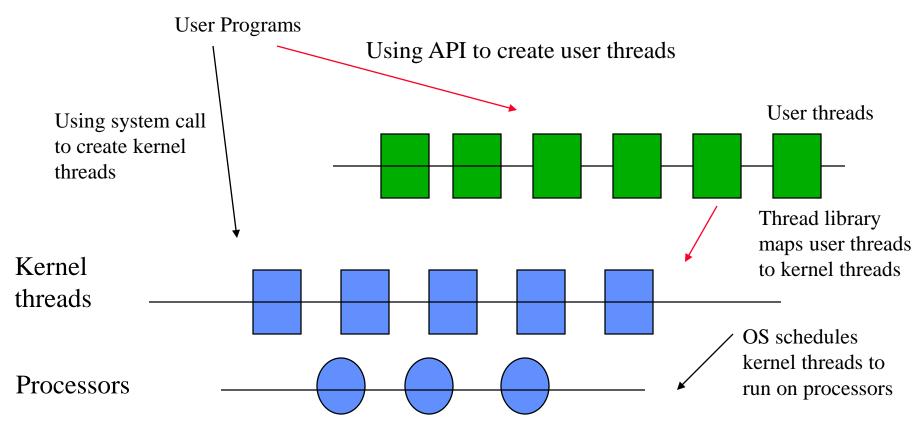
```
public class Driver{
  public static void main(String[] args){
       if (args.length > 0)
        if (Integer.parseInt(args[0])<0)
          System.err.println(args[0]+"must be >=0");
         else{
          int upper = Integer.parseInt(args[0]);
          Summation sumObj = new Summation(upper,0);
           Thread thrd=new Thread(sumObj);
           thrd.start();
          try{thrd.join();}catch(InterruptedException ie){ }
```

Example: Summation

\$javac Summation.java \$javac Driver.java \$java Driver 100 Sum of 1 through 100 is 5050

Threads in a Computer System

■ A thread library provides the programmer with an API for creating and managing threads.



Kernel Threads vs. User Threads

- Kernel threads:
 - Directly created/managed by the OS kernel
- User threads
 - created by a thread library and scheduling is managed by the library itself (the existence of user threads is unknown to the OS)
 - low-cost in thread creation; portable
 - may not utilize multi-processors efficiently (depends on mapping strategy between user and kernel threads)
 - When executed, must be mapped to kernel threads

Mapping User Threads to Kernel Threads

- Many to One
- One to One
- Many to Many

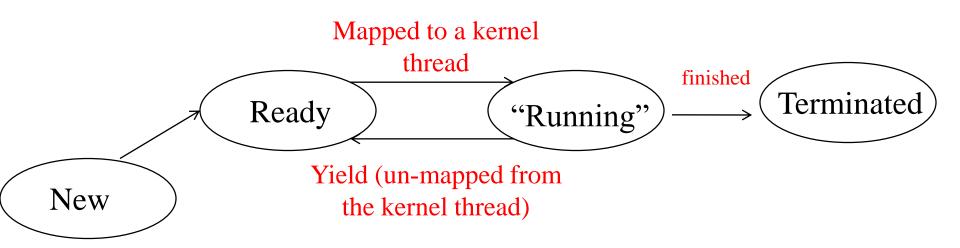
user thread kernel thread

Many-to-One

- Many user-level threads within the same process are mapped to single kernel thread
- User thread library schedules user threads to the kernel thread
- The multiple threads "time-share" the single kernel thread
 - A kernel thread is created for a process
 - Upon the thread yields, switch to another thread in the same process (kernel is unaware of it)
 - Upon the thread blocks, all threads in the process block

Thread Scheduling in User Space

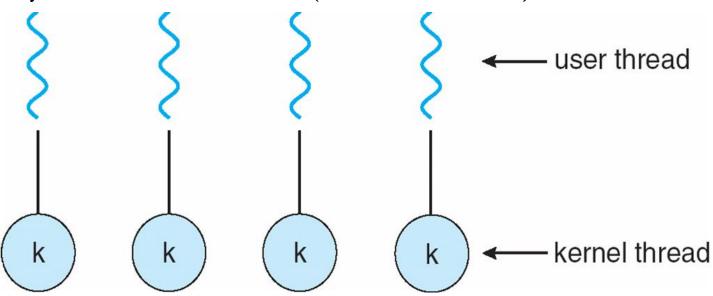
User-space scheduling: managing the (dynamic) mapping from user threads to kernel threads.



Lifetime of a user thread

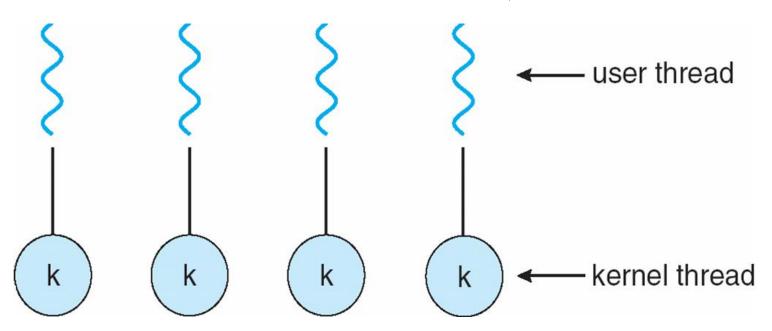
One-to-One

- Each user-level thread maps to a kernel thread
- User thread library simply provides a portable interface for thread creation/management, of which the implementation rely on the kernel thread (kernel of the OS)

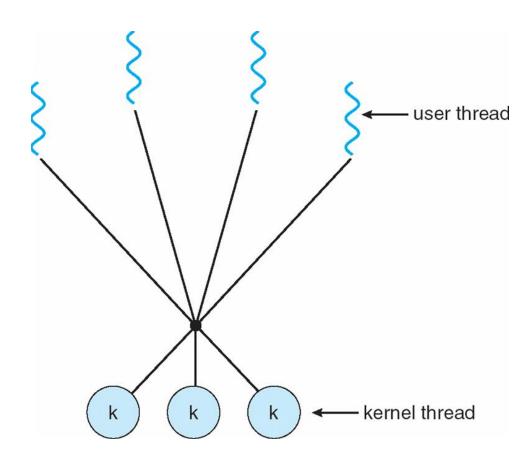


One-to-One

Upon a thread currently running on a CPU yields or blocks,
CPU can be switched to another thread (kernel is aware of this)

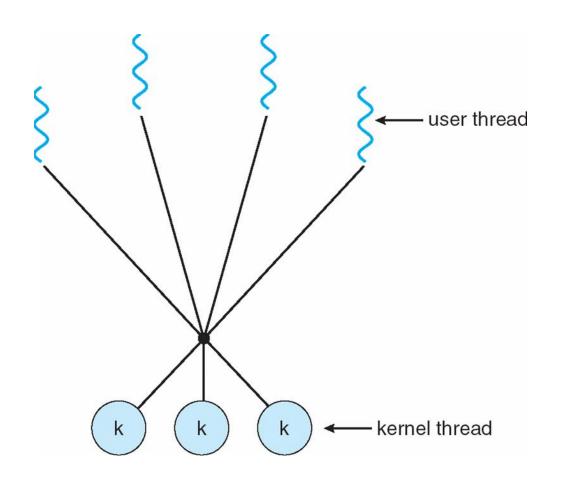


Many-to-Many Model



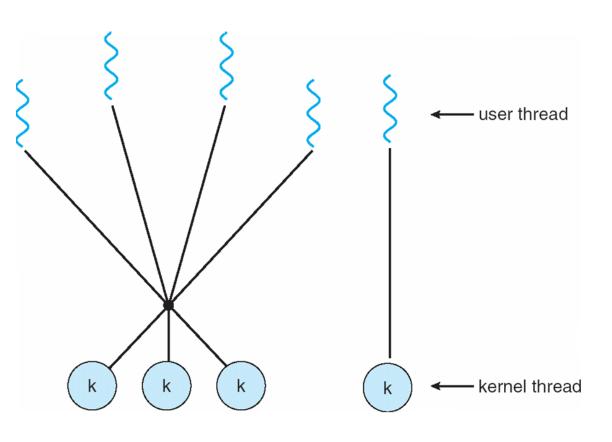
- Allows many user level threads to be mapped to many kernel threads
- User thread library schedules user threads to kernel threads

Many-to-Many Model



- Upon a currentlyrunning user thread
 yields, another thread
 may be mapped to
 the yielding user
 thread's kernel thread
 and runs
- Upon a currentlyrunning user thread blocks, other thread may run on the CPU

Two-level Model



- A sub-type of M:M
- It allows a user thread (e.g., with high priority) to be **bound** to a kernel thread