Lecture 2 Linear Data Structures

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Linear Data Structures

- Data structures support the following operations on a set S of items

 - Remove an item from S
 - Search for an item in S
- What is a linear data structure?
 - Data items are stored sequentially in a linear data structure
 - E.g., stack, queue, linked list

search key 5

st
set

Stack

(Today)

Queue

(Today)

Linked List

(Today)

Tree

(*Lecture 3*)

Heap

(Lecture 4)

.







<u>Operation</u>	Stack S	Queue Q	Linked List <i>L</i>	
Insert an item	push(e)	enqueue(e)	insert(x)	
Remove an item	pop() ◆	dequeue()	remove(x)	
Search for a key	N/A	N/A	$\operatorname{search}(k)$	fixed deletion order

- We are going to study the **properties** and the **operations** of stack, queue, and linked list
- Operations with O(1) running time
 - \diamond All operations in stack and queue, insert(x), remove(x) in linked list
- \diamond Operations with O(n) running time
 - \Rightarrow search(k) in linked list

n denotes the number of data items stored

Outline



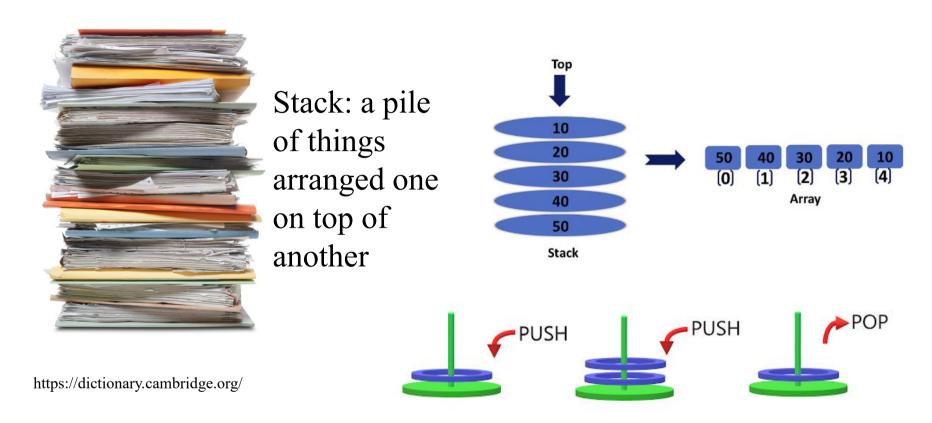
Stacks

Queues

Linked lists

Variants of linked lists

Stack



Simple Representation of a Stack with PUSH and POP operations

Built-in Stack in Java

https://docs.oracle.com/en/java/javase/21/docs/api/java.base/java/util/Stack.html

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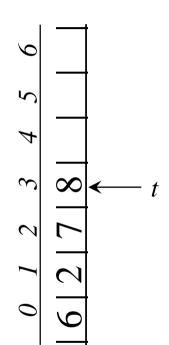
Stack



- The last-in first-out (LIFO) property
 - Only the last inserted object can be removed
 - E.g., insert 6, insert 2, insert 7, insert 8, remove → return 8



- Syntax parsing (e.g., for XML, HTML)
- Searching in puzzle problems
- Memory management (for procedure call) during program execution

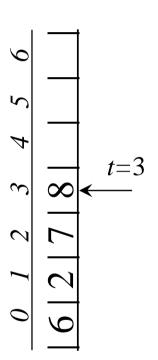


Stack: Array Implementation

- The last-in first-out (LIFO) property
- Instance variables
 - ⋄ S: an array of items
 - \diamond S.length: the length of the array S
 - $\diamond t$: the position of the top item
- Operations
 - \bullet **push**(e): insert item e (at the top of the stack)

 - size(): return the number of items stored
 - isEmpty(): check whether the stack is empty
 - top(): return the top item without removing it





Stack: Auxiliary Operations

- isEmpty(): check whether the stack is empty
- size(): return the number of items stored
- top(): return the top item without removing it

size ()

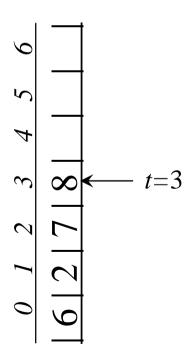
1. return t+1

isEmpty()

1. return (*t*<0)

<u>top()</u>

- 1. if isEmpty()
- 2. throw an exception
- 3. return S[t]

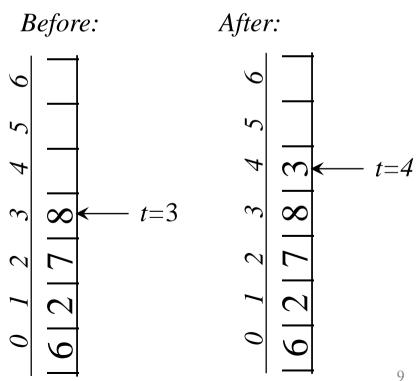


Stack: push(e)

- \bullet **push**(e): insert e to the top of the stack
 - Example: push(3)

push(e)

- 1. if size() = S.length
- throw an exception
- 3. $t \leftarrow t + 1$
- 4. $S[t] \leftarrow e$



Stack: pop()

pop():

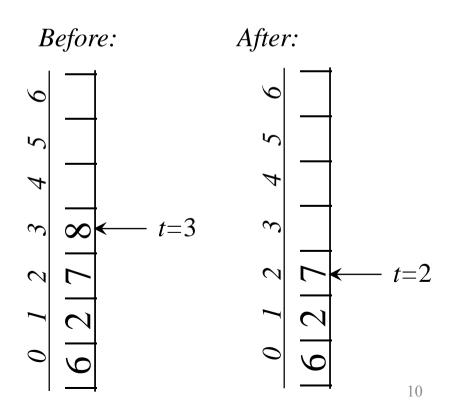
extract the top item

Example:

pop() returns 8

pop()

- 1. if isEmpty()
- 2. throw an exception
- 3. $temp \leftarrow S[t]$
- $4. S[t] \leftarrow \text{null}$
- 5. $t \leftarrow t 1$
- 6. return temp



```
public class ArrayStack<E> {
  private E[ ] S;
  private int t = -1;
  public ArrayStack(int max size)
  public int size() {
  public boolean isEmpty()
  public E top() throws Exception {
  public void push(E e) throws Exception {
  public E pop() throws Exception {
```

Stack: Java Code

- The generic type <E>
 - ⋄ <E> can be provided at run time
 - E.g., create a stack of integers by
 new ArrayStack<int>(10);

Class ArrayStack<E>: defines the type of ArrayStack<E> object

- Instance variables
- Constructor for creating an object
- Methods

- Remark: This version is simpler than the code in the textbook
 - We use one class and one type of Exception only

Examples of using a stack

Reverse an array by using a stack

⋄ Test case: [6,2,7,8] \rightarrow [8,7,2,6]

Match parentheses by using a stack

 \diamond Test case #1: [](()) \rightarrow true

 \diamond Test case #2: (. \rightarrow false

Source codes

- Download at https://bcs.wiley.com/he-bcs/Books?action=resource&bcsId=8635&itemId=1118771338
 &resourceId=35121
- ReverseWithStack.java, MatchDelimiters.java

Outline

Stacks



Queues

Linked lists

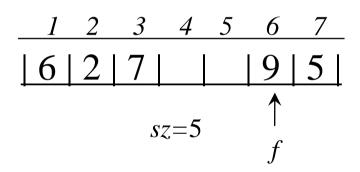
Variants of linked lists

Queue

- The first-in first-out (FIFO) property
 - Only the first inserted object can be removed
 - E.g., insert 9, insert 5, insert 6, insert 2, insert 7 remove → return 9



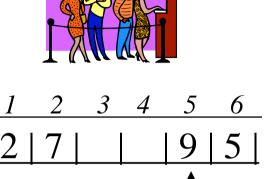
- Applications
 - CPU scheduling
 - Disk scheduling
 - Playlist in media player



Queue: Array Implementation

- The first-in first-out (FIFO) property
- Instance variables
 - $\diamond Q$: an array of items
 - \diamond Q.length: length of array Q

 - \diamond sz: the number of items stored



$$sz=5$$
 f

- Operations
 - \bullet enqueue(e): insert item e (at the rear of the queue)
 - dequeue(): extract the item from the front
 - size(): return the number of items stored
 - isEmpty(): check whether the queue is empty
 - top(): return the front item without removing it

Queue: Auxiliary Operations

- Use Q as a circular array by using the modulo operation
 - Modulo means the reminder of integer division

```
• The rear position is r = (f+sz) \mod Q. length
```

 \bullet E.g., in this example, the items are stored at 5, 6, 0, 1, 2

size ()

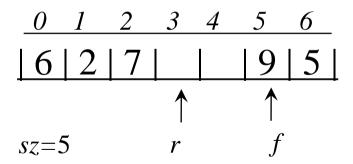
isEmpty ()

1. return sz

1. return (sz = 0)

front ()

- 1. if isEmpty ()
- throw an exception
- 3. return Q[f]

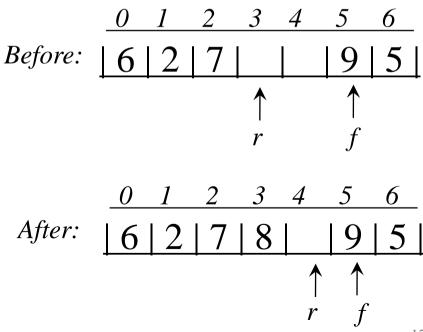


Queue: enqueue(e)

- \bullet enqueue(e): insert item e (at the rear of the queue)
 - Example: enqueue(8)

enqueue (e)

- 1. if size() = Q.length 1
- 2. throw an exception
- 3. else
- 4. $r \leftarrow (f + sz) \mod Q.length$
- 5. $Q[r] \leftarrow e$
- 6. $sz \leftarrow sz + 1$

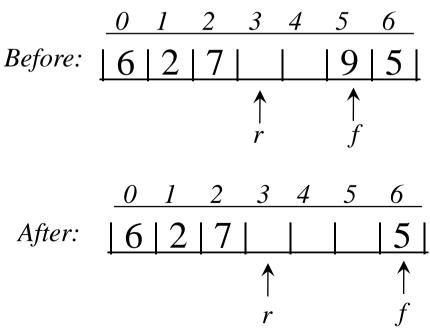


Queue: dequeue()

- dequeue(): extract the item from the front
 - Example: dequeue() returns 9

dequeue ()

- 1. if isEmpty()
- 2. throw an exception
- 3. else
- 4. $temp \leftarrow Q[f]$
- 5. $Q[f] \leftarrow \text{null}$
- 6. $f \leftarrow (f+1) \mod Q.length$
- 7. $sz \leftarrow sz 1$
- 8. return temp



```
public class ArrayQueue<E> {
  private E[ ] O;
  private int f=0;
  private int sz=0;
  public ArrayQueue (int max size)
  public int size() {
  public boolean isEmpty()
  public E front() throws Exception {
  public void enqueue (E e) throws Exception {
  public E dequeue() throws Exception {
```

Queue: Java Code

- Class ArrayQueue<E>: defines the type of ArrayQueue<E> object
 - Instance variables
 - Constructor for creating an object
 - Methods
- *Remark*: This version is simpler than the code in the textbook
 - We use one class and one type of Exception only

Outline

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Queues

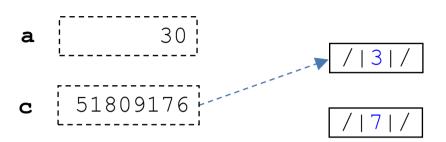


Linked lists

Variants of linked lists

Java: primitive type vs. reference type

```
int a = 30;
Node c = x_node;
a = 50;
c = y_node;
```



Primitive-type variable

- \diamond Types: byte, short, int, long, float, double, char, boolean
- A variable stores a value
 - E.g., a stores the integer value 30

Reference-type variable

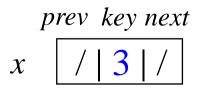
- ♦ Types that are not primitive (e.g., Object, String, Node)
- A variable stores an address of a data object
 - E.g., c stores the address of a Node object
 - It can also store the **null** value, meaning invalid reference
- ♦ Use the dot notation to access an instance variable (e.g., c.key)

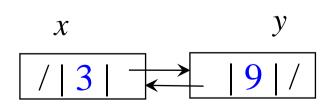
[Doubly] Linked List: Structure

- In a [doubly] linked list, each node x stores:
 - x.key: key value of the node
 - ⋄ x.prev: reference to the previous node
 - * *x.next*: reference to the next node



- A reference can be set to:
 - The memory address of another node, OR
 - null value (shown as '/' in our figures)





[Doubly] Linked List: Structure



- \diamond Linked list L is a chain of nodes (see the previous page)
 - head is the first node in the linked list
 - ♦ First node: prev = null; Last node: next = null
 - Consecutive nodes x & y should satisfy:

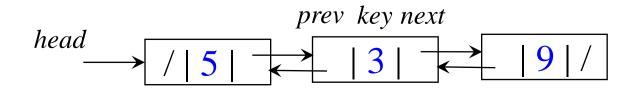
$$x.next = y$$
 and $y.prev = x$

Operations

 \bullet search(k): search a node with key value k O(n) time

 \bullet insert (x): insert a node into L O(1) time

 \bullet remove(x): remove an (existing) node from L O(1) time



[Doubly] Linked List Structure: Java Code

Remark: To make the code simple, we just store an integer key in a node and don't use get/set methods

```
class Node {
  int key;
  Node prev;
  Node next;
}
```

A linked list just stores the head

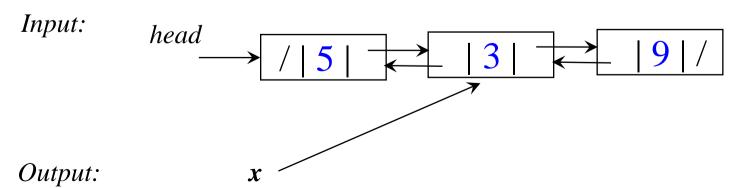
```
class LinkedList {
  Node head;
  LinkedList() {
      head=null;
 Node search(int k) {
  void insert(Node x) {
  void remove(Node x) {
```

[Doubly] Linked List: search(k)

- \bullet search(k): search a node with key value k
 - Example: search(3)

search (k)

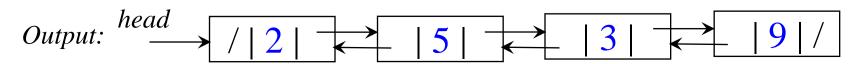
- $1. x \leftarrow head$
- 2. while $x \neq \text{null}$ and $x.key \neq k$
- $3. \qquad x \leftarrow x.next$
- 4. return *x*



[Doubly] Linked List: insert(x)

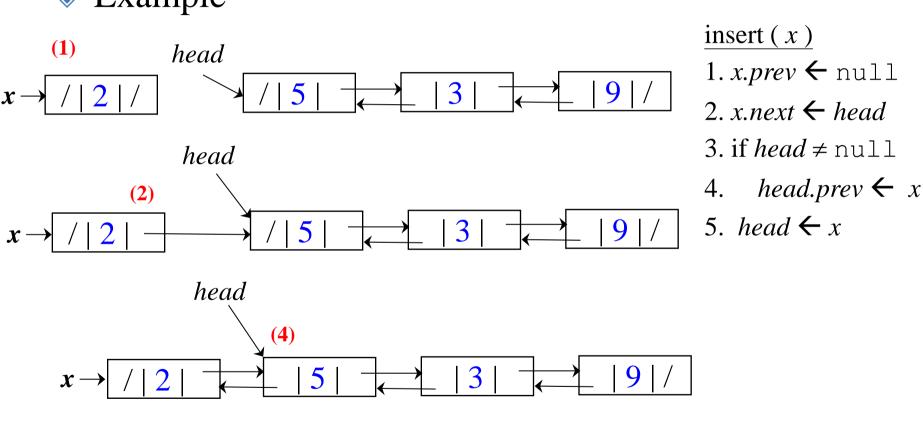
insert(x): insert a node into a linked list

insert (x)1. $x.prev \leftarrow null$ 2. $x.next \leftarrow head$ 3. if $head \neq null$ 4. $head.prev \leftarrow x$ 5. head $\leftarrow x$ Input: head



[Doubly] Linked List: insert(x)

Example



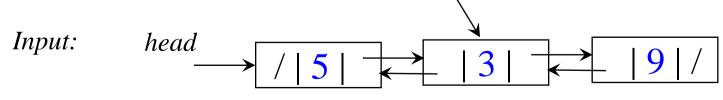
(5) head $x \rightarrow /|2| \rightarrow |5| \rightarrow |3| \rightarrow |9|/$

[Doubly] Linked List: remove(x)

- remove(x): remove an (existing) item from a linked list
 - x must be an existing node in L
 - If we wish to remove a node with the key k, we must search it first

remove (x)

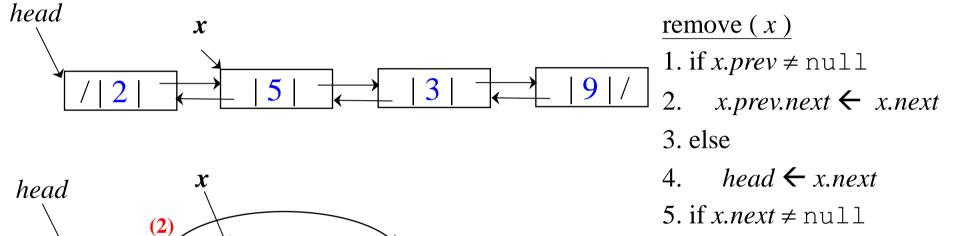
- 1. if $x.prev \neq null$
- 2. $x.prev.next \leftarrow x.next$
- 3. else
- 4. $head \leftarrow x.next$
- 5. if $x.next \neq null$
- 6. $x.next.prev \leftarrow x.prev$

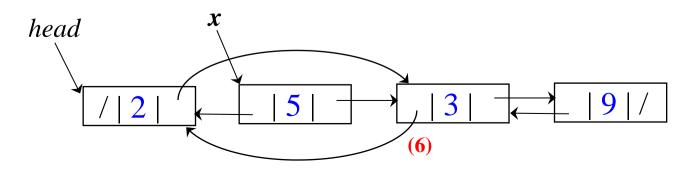




[Doubly] Linked List: remove(x)

Case I: x is not the head of the linked list

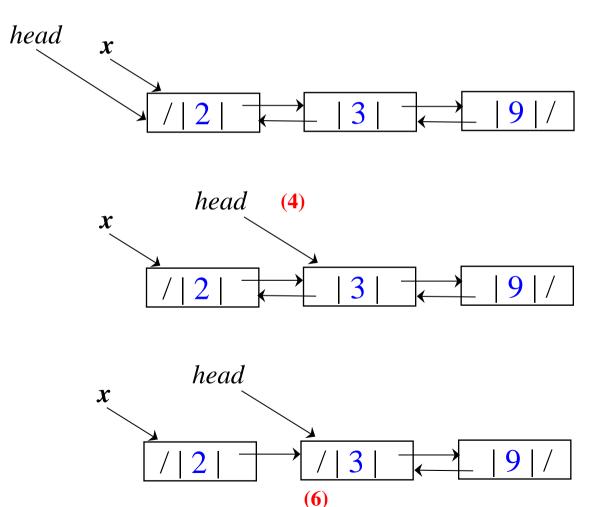




 $x.next.prev \leftarrow x.prev$

[Doubly] Linked List: remove(x)

Case II: x is the head of the linked list



remove (x)

- 1. if $x.prev \neq null$
- 2. $x.prev.next \leftarrow x.next$
- 3. else
- 4. $head \leftarrow x.next$
- 5. if $x.next \neq null$
- 6. $x.next.prev \leftarrow x.prev$

Outline

Stacks

Queues

Linked lists



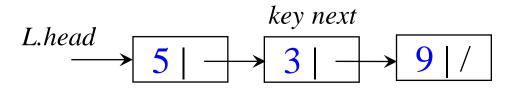
Variants of linked lists

Variants: Singly Linked List

- In a doubly linked list, each node stores key, prev, next
 - This is the type of linked list used in previous slides
 - Require more space ☺
 - Simple algorithm for deleting a node ©
- In a singly linked list, each node stores key, next only

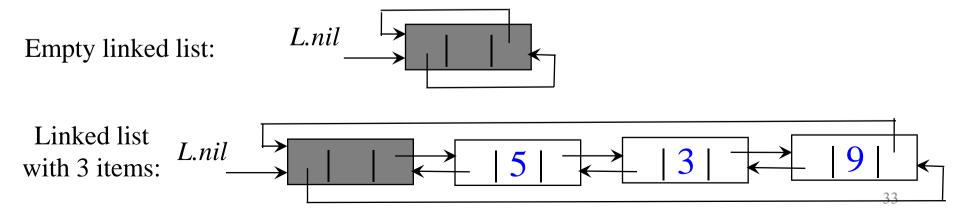
 - More complicated algorithm for deletion <a>©
 [Question] How do we remove the node "3" in this linked list?

Singly linked list:



Variants: Linked List with Dummy Node

- How to simplify the operations on linked lists?
- Use a circular linked list, with a dummy node
 - Replace the null value by a dummy item L.nil
 - ♦ The "next" of *L.nil* is the first node (i.e., *L.head* is *L.nil.next*)
 - ⋄ The "prev" of *L.nil* is the last node
- Simpler algorithms can be used on such linked list
 - ♦ E.g., L.nil.next is always not null (no need to check null)



Variants: Array as Linked List

- We can also implement a linked list by using an array
 - Use an array A[0..n-1] of items; each A[i] is an item
 - ⋄ A[i].prev, A[i].next store array positions instead of reference values (use −1 to represent null)
- head_pos: the head position of the linked list (of used items)
- free_pos: the head position of the free list (of unused items)
 - When we create an item, just extract it from the free list
- Consider the following example:

```
\bullet head_pos = 0 \rightarrow positions of used items: 0, 3, 5
```

$$*$$
 free_pos = 1 \rightarrow positions of unused items: 1, 2, 4, 6

$$\frac{0}{1} \frac{2}{3} \frac{3}{4} \frac{5}{6}$$

$$key: |5| |3| |9|$$

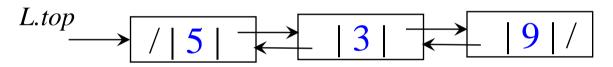
$$prev: |-1|-1| 1 |0| 2 |3| 4$$

$$free_pos = 1$$

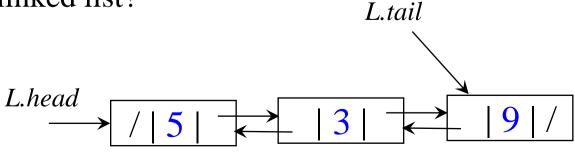
next:

Questions to you

We have to implement a stack by a linked list? How about the stack operations (push, pop) on this linked list?



We How to implement a queue by a linked list?
How about the queue operations (enqueue, dequeue)
on this linked list?



Summary

- Stack: LIFO property, push, pop
- Queue: FIFO property, enqueue, dequeue
- Linked list and its operations

Please read Chapters 3 and 6 in the book "Data Structures and Algorithms in Java"