

## 1. Language Overview

- **Statically typed, stack-only** programming language (PL).
- Compiles to a custom **register-based assembly** (12-bit and 6-bit variants).
- **Single stack frame** per thread; no heap, no recursion, no VLAs, no `alloca`.
- All values live on the stack; pointers/references to stack slots are allowed but cannot escape their frame.
- **Functions** take one argument and return one value (both can be compound types: tuples or vectors).
- **Key syntax patterns:** comma-separated `key:value` lists enclosed in block markers.
- **RHS expressions in** `SET ... WITH` **blocks are evaluated in parallel**, unless they contain impure operations (which is a compile-time error).
- **Operators** can be applied to tuples or vectors if compatible:
  - Vectors: same length and compatible base types
  - Tuples: same fields with compatible types

### 1.1. Core PL Syntax

```
VARIABLES
  name1: type1
, name2: type2
, ...
END
```

```
TYPEDEF
  alias1: existing_type1
, alias2: existing_type2
END
```

```
SET
  lval1: rval1
, lval2: rval2
WITH
  symbol1: rval3
, symbol2: rval4
END
```

```
FUNCTION_HEADERS
  name1: in_type1:out_type1
```

```
, name2: in_type2:out_type2
END
```

```
FUNCTION_BODIES
  name1: BEGIN
    ...statements...
  END
, name2: instruction2
END
```

```
EXPORT
  out_name: in_name
END
```

```
IMPORT
  FROM module_name
  in_name: out_name
END
```

```
SUM_TYPE
  OptionA: TypeA
, OptionB: TypeB
END
```

```
MATCH value WITH
  OptionA:x: stmtA
, OptionB:y: stmtB
END
```

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## 2. Assembly Language (12-bit variant)

### 2.1. Instruction Format

- **12 bits total:**
- **4 MSBs:** Command family
- **8 LSBs:** Payload (arguments, type codes, etc.)

### 2.2. Registers

- **Three core registers:**

- `left` — primary pointer/address
- `right` — secondary pointer/address
- `param1` — auxiliary (e.g., passing syscall parameters)
- **Optional fourth register:** `result`
- **Zeroing behavior:** On jumps, calls, syscalls, and when selecting a register.

## 2.3. Register Input Commands

- **Push immediate:** pushes a byte into the currently selected register.
- 12-bit version: 8-bit immediates
- 6-bit version: 4-bit immediates

## 2.4. ALU Operations and Casting

- ALU is **stateful**: holds a current data type (u8, i8, u16, i16, u32, i32, f32).
- **Cast/Copy Command** takes two 4-bit type codes (`source`, `destination`):
- **Case 1:** `source == destination` and `left == right` → default operation only.
- **Case 2:** `source == destination`, `left != right` → copy `sizeof(type)` bytes from `right` to `left`, then default.
- **Case 3:** `source != destination` → cast from `source` to `destination`, store at `left`, then default.
- **Default operation** (performed after any copy/cast):

```
*(destination*)left = *(source*)right;
left += sizeof(destination);
right += sizeof(source);
type_of_ALU = destination;
```

- **Vectorized operators** also use `left` / `right` cursor advancement, enabling efficient tuple/vector processing.

## 2.5. Memory Model

- **Stack layout per function:**
  - Address **0**: return value
  - Next addresses: function parameter, then local variables
- **Address encoding:**
  - LSB = 0 → stack-local address (offset from frame start)
  - LSB = 1 → module-level static/global variable

## 2.6. Indirection and Pointers

- **Indirect load/store**: operate on the memory address stored *at* the address in a register.
- No heap; all indirection is within the stack or static data.

## 2.7. Control Flow

- **Relative jumps**: offset from function's first instruction.
- Jump to **0** → restart function.
- Jumps reset registers (zero them).

## 2.8. Function Calls

- **Syscall-based call**:
    - Place **magic number** in a designated register.
    - Provide return-value address in another register.
    - Other regs (param1) for argument info if needed.
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# 3. Executable Format and Linking

## 3.1. Global Tables

1. `global_to_local`: maps each **global\_id** (implicit index) → `(module_id, local_id)`
2. `module_metadata`: rows keyed by **module\_id** (implicit) with columns:
3. `first_extern_function`: index of first external function
4. `module_location`: file offset
5. `tables_length`, etc.

## 3.2. Per-Module Tables

1. `function_ends`: end offsets of each function (implicit `function_id` ordering). → can derive start offset from previous end.
2. `extern_to_global`: for each external index (implicit), maps → `global_id`.

## 3.3. Function Magic Number Resolution

```
if (magic_number < first_extern_function) {
    module_id = current_module_id;
    function_id = magic_number;
} else {
    extern_id = magic_number - first_extern_function;
    global_id = extern_to_global[extern_id];
    module_id = global_to_local[global_id].module;
    function_id = global_to_local[global_id].function;
```

```
}  
// Use (module_id, function_id) to look up jump offset
```

- **System module** ID is 0xFFFF for syscalls and runtime services.
- Enables **compact** function references (small magic numbers) with **fast native lookups**.

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*Next steps:* - Define **function-module magic number encoding** details. - Describe **executable file layout**: header, table sections, code segments, static data. - Outline **compiler** phases: parsing, type-checking, IR lowering, assembly emission, linking.