Simple ACC Assembler

Assembly language is machine dependent, depending on a specific computer architecture.

Assembly language can be translated into lower level machine language, or interpreted.

- Software interpreter
- Virtual computer/machine

This virtual computer is a simple computer with

- ALU, RAM, just one register called ACCumulator, simple IO, and a built-in stack
- Word size is 2 bytes, and the addressability is thus 64k.

Program Format

- Fach line
 - o independent and self-contained
 - may be blank
 - a complete instruction
 - a complete storage directive
 - all delimiters are WS
- Instruction
 - o for an accumulator machine

(left argument and result are in an implicit accumulator ACC register, except for COPY), with the following format

[Label:] XXX arguments

- XXX is the reserved name
- arguments as needed separated by spaces
- Storage directive
 - XXX val
 - XXX is storage name
 - val is the initial value
 - all storage are signed 2 byte integers
- Names
 - o Instruction names are reserved in upper case
 - Variables start will letter and continue with letters and digits up to 8
- Numbers
 - The computer uses standard data 2's complement data representation so thus data range is -32k to +32k.

Instruction Set (# arguments, meaning)

BR (1, jump to arg) 0 BRNEG (1, jump to arg if ACC < 0) 0 o BRZNEG (1, jump to arg if ACC <=0) **BRPOS** (1, jump to arg if ACC >0) (1, jump to arg if ACC >=0) **BRZPOS** BRZERO (1, jump to arg if ACC == 0)COPY (2, arg1 = arg2)0 ADD (1, ACC = ACC + arg0 o SUB (1, ACC = ACC - arg)DIV (1, ACC = ACC / arg)o MULT (1, ACC = ACC * arg)(1, arg=input integer) READ 0 o WRITE (1, put arg to output as integer) o STOP (0, stop program) o STORE (1, arg = ACC)o LOAD (1, ACC=arg) (0, nothing) o NOOP

Immediate Values

- ADD, DIV, MULT, WRITE, LOAD, SUB
 - o can take either variable or immediate value as the argument
 - o immediate value is positive integer or negative 2-byte integer

Stack

PUSH (0, tos++)
POP (0, tos--)
STACKW (1,stack[tos-arg]=ACC)
STACKR (1,ACC=stack[tos-arg])

PUSH/POP

are only means to reserve/delete automatic storage.

STACKW/STACKR n

- these are stack access (random access) instructions.
- n must be a non-negative number
- the access is to nth element down from TOS

NOTE: TOS points to the topmost element on the stack

Semantics

Execution begins with the first line and continues until STOP is reached

Invocation

- > virtMach // read from stdin
- > virtMach file.asm // read from file.asm

Location

/accounts/classes/janikowc/cs4280/asmInterpreter/virtMach

• readable executable

Examples

Simple flows

Read a number and print number+1

```
READ X
LOAD X
ADD 1
STORE X
WRITE X
STOP
X 0
```

Read 3 numbers, add them and display the total, using 3 variables

```
READ X
READ Y
READ Z
LOAD X
ADD Y
ADD Z
STORE X
WRITE X
STOP
X 0
Y 0
Z 0
```

Conditional Flow

- Simple relational conditions are processed by subtracting sides and implementing proper jump
- This is 'if' without else

High Level If (arg1 RO arg2) Stat1 Stat 2 // if condition true then Stat1 and move to Stat2, else skip Stat1 and move to Stat2

```
VM Assembly

Evaluate arg2
STORE results2
Evaluate arg1
SUB result2 // arg1 – arg2
BR??? Out // BR on false
Stat1
Out: Stat2
```

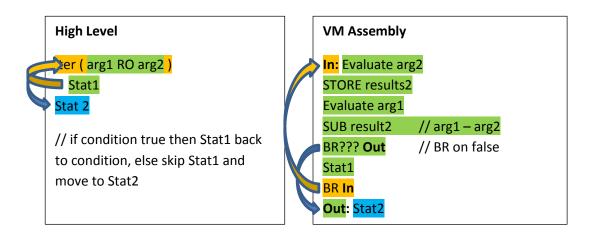
Read a number and print it if >= 1

```
\begin{array}{l} \text{program} \\ \text{var } x = 0 \; . \\ \text{start} \\ \text{read } x \; . \\ \text{iff } (x >> 1) \; \& \; \text{assume this is } >= \& \\ \text{print } x \; . \\ \text{stop} \end{array}
```

```
READ x
LOAD x
SUB 1
BRNEG Out
WRITE x
Out: STOP
x 0
```

Iteration

- Iteration is similar to conditional except that upon executing the 'true' statement there is unconditional jump back to evaluate the condition again.
- This will continue (iterate) until the condition is false



Read input number and print number then number-1 ... down to 1

```
program
var x = 0 .
start
  read x .
  iter ( x >> 1 )
  start
  print x .
  let x = x - 1 .
  stop
stop
```

```
READ x
In: LOAD x
SUB 1
STORE x
BRNEG Out
WRITE x
LOAD x
SUB 1
BR In
Out: STOP
x 0
```