A brute-force solution to problem "tsp"

As noted in the "Five example optimization problems" handout, problem tsp is intrinsically a discrete optimization problem. In fact it is a *combinatorial* and graph-theoretic optimization problem. Such problems are a topic in MATH 663 Graph Theory.

For that reason, and also because tsp fits neither with the continuous optimization problems of the rest of the course, nor the techniques in the textbook, I will not cover it further. Beyond the brute-force solution tsp.m on the next page, that is! *There is no claim that this code represents a good approach*, but it does solve the problem. Please see a course or book on combinatorial optimization or graph theory for more information on this kind of problem.

The Matlab/Octave code tsp.m starts with a generally-useful and basic idea, namely that one may represent an edge-weighted graph by a *matrix*. First, the vertices (cities) A, F, J, N, S, W have indices 1, 2, 3, 4, 5, 6, respectively, in this data structure. Then the matrix M = edgew has entry $m_{ij} = -1$ if there is no flight between city (vertex) i and city j; otherwise m_{ij} is the cost of the flight from city i to city j.

If the seven-city itinerary $p = i_1 i_2 i_3 i_4 i_5 i_6 i_7$ is feasible, in the sense that the corresponding flights exist, then the cost of p is the sum of the six flights,

$$C(p) = a_{i_1 i_2} + a_{i_2 i_3} + a_{i_3 i_4} + a_{i_4 i_5} + a_{i_5 i_6} + a_{i_6 i_7} = \sum_{k=1}^{6} a_{i_k i_{k+1}}.$$

This cost is computed by a for loop:

Note that if the itinerary is not feasible we "break" out of the for loop and move on to the next case.

The above explains how to find the optimal itinerary, given a way of generating all the feasible ones. How do we do that? The brute force approach here is to pre-generate all itineraries by using MATLAB/OCTAVE's perms function, which generates all possible permutations of a given vector.² In this case the possible itineraries start and end in city

¹A *graph* is a collection of vertices with *edges* which can be defined as pairs of vertices. It is one of the basic objects of (discrete) mathematics.

²This aspect of the solution avoids having to generate a better data structure for holding a graph.

S, so we generate the permutations of the other cities ([1 2 3 4 6]) and then prepend and append S = 5. Then, in a big for loop, we test whether the itinerary is feasible.

Note that there is a small anonymous function printcase, taking three arguments, whose only purpose is to print out an itinerary. We print the ones that are feasible and then, at the end, we print the first of the optimal itineraries.

```
tsp.m
% TSP Setup and solve a traveling salesperson problem
label = ['A' 'F' 'J' 'N' 'S' 'W'];
edgew = [-1\ 100\ 100\ 150\ 250\ 150;
        100 -1 150 200 300 250; % F
        100 150 -1 -1 200 200; % J
        150 200 -1 -1 -1;
        250 300 200 -1 -1 -1; % S
        150 250 200 -1 -1 -1]; % W
% generate permutations
notS = [1 2 3 4 6]; % not including Seattle
N = factorial(5); % this many permutations of notS
p = [5*ones(N,1) perms(notS) 5*ones(N,1)]; % start and end in Seattle
% print one case on a single line
printcase = @(a, p, s) \dots
    fprintf('%s %c%c%c%c%c%c%c : %d\n',
           a, label(p(1)), label(p(2)), label(p(3)), ...
           label(p(4)), label(p(5)), label(p(6)), label(p(7)), s);
% evaluate cost of each (feasible) permutation
mincost = 1.0e100; % bigger than any
for j = 1:N
                   % C = cost of path (or -1)
   C = 0;
   for k = 1:6
       w = edgew(p(j,k),p(j,k+1));
       if w < 0
           C = -1;
           break
       else
           C = C + w;
       end
    end
    if (C > 0)
                   % report if feasible; update optimal
       printcase('feasible ',p(j,:),C)
       if C < mincost
           mincost = C;
           minp = p(j,:);
        end
   end
end
printcase('optimal ',minp,mincost)
```