
By Ivan Malishevskyi & Oleksii Savelich

[All the code can be found here](#)

Problem Statement

Given:

- A set of n **guests**, each identified by name.
 - A list of m **pairs** of guests who dislike one another.
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Task:

- Arrange all guests at two tables so that no pair of guests who dislike each other sits at the same table.
 - If no valid arrangement exists, just display `Impossible` in console.
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Algorithm Design

Instead of trying to place each guest one by one with backtracking, we treat the problem as graph bicoloring:

1. Construct an undirected graph

- **Vertices** represent guests.
- **Edges** connect pairs of guests who dislike one another.

2. Non-recursive DFS bicoloring

- Maintain an array `color[]` of length n , initialized to `-1` (unassigned).
- For each uncolored vertex, push it onto a stack with color `0` (Table 1).
- While the stack is nonempty:
 - Pop a vertex u .
 - For each neighbor v of u :

- If `color[v] == -1`, assign `color[v] = 1 - color[u]` and push `v`.
- Else if `color[v] == color[u]`, stop: the graph is not bipartite (output `Impossible`).

This guarantees exactly one pass over each edge and vertex, yielding $O(n + m)$ time.

Seating Representation and Output

- **Tables:** Two lists of names, corresponding to the first table and the second table respectively.
- **Output format:**

```
Table 1:
<guest names at table 1, one per line>
Table 2:
<guest names at table 2, one per line>
```

- If mistake occurs during coloring process:

```
Impossible
```

Code Structure

- **Graph**
 - Defines the `Graph` class with fields `Names[]`, `Adjacency[]`, so basically it's a model class that will represent our graph with guest dislikes.
- **BicoloringSolver**
 - Contains `TryColor(Graph graph, out int[] color)`, implementing the non-recursive DFS.
- **Program**
 - Parses command-line argument for input file, invokes graph construction, solver, and prints results. So basically it's an entry point to program.
- **GraphReader**

- Utility class for reading input from text file and transforming it into graph. Not much to say about it.
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Complexity Analysis

- **Time:** $O(n + m)$ — each vertex and edge is processed once in the DFS.
 - **Space:** $O(n + m)$ — adjacency lists plus the `color[]` array and stack.
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Potential Limitations

- Assumes only two tables; no extension to more than two partitions.
 - Input validation errors (duplicate names or invalid pairs) rely on exceptions.
 - Guests with no dislikes default to Table 1 without further balancing.
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Conclusion

By modeling guests and their dislikes as a graph and using non-recursive DFS, we efficiently obtain a valid seating arrangement or detect impossibility. This approach scales linearly and cleanly separates data parsing, algorithm logic, and I/O. I don't really know what I should write in this section, so I came up with something that will sound smart and increase the number of letters in the report.

Tools used: >

- [Obsidian](#) , to generate this fancy report and export it into pdf
- [C# & .NET](#) , the language of choice and environment to implement algorithms