SGPC Programmer's Manual

Steven Vroom

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About This Manual

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Overview

- 1.1 SGPC Architecture Overview
- 1.2 Registers

User-accessible Registers

Internal Registers

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- 1.4 Instruction Set
- 1.5 Interrupt Model
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Register Set

This chapter describes the registers seperated in four groups based on accessability. however, the internal registers are omitted from this chapter since these are implementation specific.

2.1 Foreground Registers

The foreground registers are the registers all regular instructions can read from and write to. There are eight 8-bit and eight 16-bit foreground registers. These registers are preserved in interrupts.

Table 2.1: List of Foreground Registers

ID	Mnonic	Descriptive Name	Length in bits
0x0	al	The lower byte of ax	8
0x1	ah	The higher byte of ax	8
0x2	bl	The lower byte of bx	8
0x3	bh	The higher byte of bx	8
0x4	cl	The lower byte of cx	8
0x5	ch	The higher byte of cx	8
0x6	dl	The lower byte of dx	8
0x7	$\mathrm{d}\mathrm{h}$	The higher byte of dx	8
0x8	ax	The first GPR	16
0x9	bx	The second GPR	16
0xA	cx	The third GPR	16
0xB	dx	The fourth GPR	16
0xC	ex	The fifth GPR	16
0xD	${ m tm}$	Temporary data	16
0xE	sp	Stack pointer	16
0xF	pc	Program counter	16

General-Purpose Registers (GPRs)

These registers are meant for computing storage. The first four 16-bit registers are all splitted into two 8-bit registers. So software can directly access the upper and lower byte of these 16-bit registers.

Temporary Data Register

This registers is meant to facilitate call procedures. So it won't be preserved in a function call. However this nothing more than a suggestion to the user, software can use this register as a regular GPR.

Stack Pointer Register (SP)

This register is meant to keep track of the end of the stack. However this nothing more than a suggestion to the user, software can use this register as a regular GPR.

Program Counter Register (PC)

This register holds the address of the next instruction to run. Writing to this registers means jumping to other code.

2.2 Background Registers

Background registers can only accessed with the instructions BSTR and BLD.

ID	Mnonic	Descriptive Name	Length in bits
0x0	n/a	Reserved	8
0x1	n/a	Reserved	8
0x2	n/a	Reserved	8
0x3	n/a	Reserved	8
0x4	n/a	Reserved	8
0x5	n/a	Reserved	8
0x6	n/a	Reserved	8
0x7	n/a	Reserved	8
0x8	bpc	Program counter backup	16
0x9	bsf	Segments and flags backup	16
0xA	ipc	Interrupt program counter	16
0xB	is	Interrupt segments	16
0xC	n/a	Reserved	16
0xD	n/a	Reserved	16
0xE	n/a	Reserved	16
0xF	n/a	Reserved	16

Table 2.2: List of Background Registers

Backup Registers

The backup registers are used to backup the state of the CPU (see section 7.2). The foreground registers, segment registers and flags register all have their own backup register. However, most backup register aren't background registers. Only the segment registers, program counter register and flags register have directly accessable backup registers. Note that both segment registers and the flag register share one 16-bit backup register.

CS: The code segment, This holds the index of the segment that is used to fetch instructions from the memory (see ??).

DS: The data segment, This holds the index of the segment that is used to read data from and write data to the memory (see ??).

Z, **S**, **P**, **O**, **C** and **I**: The flags, these hold data about the state of the processor or control how the processor behaves (see section 2.3)

Figure 2.1: Segments and Flags Backup Register's Layout

16	5 1	5 1	4 1	3 1	2 1	1 1	0	5	0
	Ι	С	О	Р	S	Z	DS (5)	CS (5)	

Interrupt Registers

The Interrupt registers hold the new state the CPU should jump to when an interrupt is triggered (see chapter 7). Only the segement registers and the program counter have an interrupt register.

CS: The code segment, This holds the index of the segment that is used to fetch instructions from the memory (see ??).

DS: The data segment, This holds the index of the segment that is used to read data from and write data to the memory (see ??).

Figure 2.2: Segments Interrupt Register's Layout

16	10	5	0
RESERVEI	D(6)	S (5)	CS (5)

2.3 Indirect Registers

Indirect registers are registers that software can't directly read nor write to with any instruction. All the backup registers that aren't background registers fall under this catogory. All the indirect registers can only be written to and read from via the state preservation system (see section 7.2).

Flags Register

The flags register holds multiple flags (see Table 2.3). There are two types of flags: status flags and control flags. Status flags give software extra information about the last computation made, while control flags control how the cpu behaves. There is only one control flag in the [insert name here] processor, the interrupts enabled flag. If this flag is set to zero interrupt requests will be ignored (see section 7.1).

Zero flag: If the result of the last computation is equal to zero this flag will be set, otherwise it will be cleared.

Sign flag: If the most significant bit of the result of the last computation is set this flag will be set, otherwise it will be cleared.

Parity flag: If the least significant bit of the result of the last computation is clear this flag will be set, otherwise it will be cleared.

Overflow flag: If the signed two's-complement result of the last computation is too large to fit in operand A (see ??) this flag will be set, otherwise it will be cleared. (Not all computational instructions change this flag)

Carry flag: If the unsigned result of the last computation is too large to fit in operand A (see ??) this flag will be set, otherwise it will be cleared. (Not all computational instructions change this flag)

Descriptive NameType of flagLength in bitsZero flagStatus flag1Sign flagStatus flag1Parity flagStatus flag1

Status flag

Status flag

Control flag

1

Table 2.3: List of Flags

Segment Registers

Overflow flag

Interrupts enabled flag

Carry flag

Mnonic

Ζ

S

Р

O

 \mathbf{C}

The segment registers (see Table 2.4) hold the index of the segments currently used (see chapter 6).

Table 2.4: List of Segment Registers

Mnonic	Descriptive Name	Length in bits
CS	Code segment index	5
DS	Data segment index	5

2.4 Output Registers

The output registers (see Table 2.5) hold information that is sent to the pheripherals (see chapter 8)

Table 2.5: List of Output Registers

Mnonic	Descriptive Name	Length in bits
AO	The first output register	16
ВО	The second output register	16
CO	The third output register	16
DO	The fourth output register	16
EO	The fifth output register	16
FO	The sixth output register	16
GO	The seventh output register	16
НО	The eigth output register	16

Operand Conventions

This chapter describes the possible operands and how to use them. The operands are separated into normal operands and special operands.

3.1 Operands in Instructions

All instructions have two operands, however, some instructions don't use them or only use one. The two operands have a set role. Operand A (the first operand after the instruction) is the output and the secondary input. Operand B (the second operand after the instruction) is the primary input. In other words the processor generally follows this behaviour: read B, optionally read A, calculate and finally write to A.

3.2 Normal Operands

Normal operands simply access the foreground registers. This type of operands is ensured to take the least amount of read and write time, so it is advised to use this operand type for computations.

Exception: ISTR & ILD

The ISTR and ILD alter the behaviour of the normal A and B operands respectively. Instead of the foreground registers the operands access the background registers. There are only four background registers. If ISTR or ILD try to access any of the reserved background register the behaviour of that instruction is undefined.

3.3 Special Operands

Special operands make use of the immidiate of an instruction. If both operands of an instruction are special operands then they share the $\underline{\text{same}}$ immidiate. The A and B operand have different special operands.

Special Operand A

Special operand A is used to access the memory. The operands adds a register to the immidiate and uses the answer of that calculation as memory address. It's also possible to use the immidiate as address directly without the addition. It's possible to read and write both 8-bits and 16-bit.

Special Operand B

Special operand B can both be used access the memory or directly. The operands either adds a register to the immidiate and uses the answer of that calculation as memory address, or just uses the answer of the calculation directly as the value of the operand. It's also possible to use the immidiate as address or value directly without the addition.

3.4 List of Operands

Table 3.1: List of Operands

ID	Mnonic	Description	${f Length} \ {f in \ bits}$		
		Normal Operands			
0x00	al	The lower byte of ax	8		
0x01	ah	The higher byte of ax	8		
0x02	bl	The lower byte of bx	8		
0x03	$_{ m bh}$	The higher byte of bx	8		
0x04	cl	The lower byte of cx	8		
0x05	ch	The higher byte of cx	8		
0x06	dl	The lower byte of dx	8		
0x07	$\mathrm{d}\mathrm{h}$	The higher byte of dx	8		
0x08	ax	The first GPR	16		
0x09	bx	The second GPR	16		
0x0A	cx	The third GPR	16		
0x0B	dx	The fourth GPR	16		
0x0C	ex	The fifth GPR	16		
0x0D	${ m tm}$	Temporary data	16		
0x0E	$_{ m sp}$	Stack pointer	16		
0x0F	pc	Program counter	16		
Ex	Exceptions: Operand A of ISTR & Operand B of ILD				
0x00	n/a	Reserved	8		
0x01	n/a	Reserved	8		
0x02	n/a	Reserved	8		
0x03	n/a	Reserved	8		
0x04	n/a	Reserved	8		
0x05	n/a	Reserved	8		
0x06	n/a	Reserved	8		
0x07	n/a	Reserved	8		
0x08	bpc	Program counter backup	16		
0x09	bsf	Segments and flags backup	16		
0x0A	ipc	Interrupt program counter	16		

Table 3.1: List of Operands

		Table 5.1. List of Operands	
ID	Mnonic	Description	Length
O _{rr} OD	:-	Interment game ente	in bits
0x0B	is	Interrupt segments	16
0x0C	n/a	Reserved	16
0x0D	n/a	Reserved	16
0x0E	n/a	Reserved	16
0x0F	n/a	Reserved	16
0.00		exception: Operand A of OUT	10
0x00	AO	The first output register	16
0x01	BO	The second output register	16
0x02	CO	The third output register	16
0x03	DO	The fourth output register	16
0x04	EO	The fifth output register	16
0x05	FO	The sixth output register	16
0x06	GO	The seventh output register	16
0x07	НО	The eigth output register	16
0x08	n/a	Reserved	16
0x09	n/a	Reserved	16
0x0A	n/a	Reserved	16
0x0B	n/a	Reserved	16
0x0C	n/a	Reserved	16
0x0D	n/a	Reserved	16
0x0E	n/a	Reserved	16
0x0F	n/a	Reserved	16
	•	Special Operands A	
0x10	[ax+#]	The byte at: ax plus the immidiate	8
0x11	[bx+#]	The byte at: bx plus the immidiate	8
0x12	[cx+#]	The byte at: cx plus the immidiate	8
0x13	[dx+#]	The byte at: dx plus the immidiate	8
0x14	[fx+#]	The byte at: fx plus the immidiate	8
0x15	[pc+#]	The byte at: pc plus the immidiate	8
0x16	[sp+#]	The byte at: sp plus the immidiate	8
0x17	[#]	The byte at: the immidiate	8
0x18	@[ax+#]	The word at: ax plus the immidiate	16
0x19	@[bx+#]	The word at: bx plus the immidiate	16
0x1A	@[cx+#]	The word at: cx plus the immidiate	16
0x1B	@[dx+#]	The word at: dx plus the immidiate	16
0x1C	@[fx+#]	The word at: fx plus the immidiate	16
0x1D	@[pc+#]	The word at: pc plus the immidiate	16
0x1E	@[sp+#]	The word at: sp plus the immidiate	16
0x1E	@[#]	The word at: the immidiate	16
$\frac{Special Operands B}{Special Operands B}$			
0x10	[ax+#]	Memory from: ax plus the immidiate	same as A
0x10	[bx+#]	Memory from: bx plus the immidiate	same as A
0x11	[cx+#]	Memory from: cx plus the immidiate	same as A
0x12 $0x13$	$[\mathrm{dx}+\#]$	Memory from: ex plus the immidiate Memory from: dx plus the immidiate	same as A
0x13 $0x14$	[fx+#]	Memory from: dx plus the immidiate Memory from: fx plus the immidiate	same as A
0x14 $0x15$			same as A
OXIO	[pc+#]	Memory from: pc plus the immidiate	same as A

Table 3.1: List of Operands

ID	Mnonic	Description	Length in bits
0x16	[sp+#]	Memory from: sp plus the immidiate	same as A
0x17	[#]	Memory from: the immidiate	same as A
0x18	ax+#	ax plus the immidiate	16
0x19	bx+#	bx plus the immidiate	16
0x1A	cx+#	cx plus the immidiate	16
0x1B	dx+#	dx plus the immidiate	16
0x1C	fx+#	fx plus the immidiate	16
0x1D	pc+#	pc plus the immidiate	16
0x1E	$\mathrm{sp}+\#$	sp plus the immidiate	16
0x1F	#	the immidiate	16

Addressing

- 4.1 Bit and Byte Ordering
- 4.2 Aligned and Misaligned Memory Access

Instruction Set Summary

This chapter gives a summary on the instruction set of the ABCD processor. It will describe the types of instructions and the binary layout of the instructions.

5.1 Instruction Types

The instructions can be separated in five types. Two of these types are not meant to be used.

Move Instructions

The move instructions copy data from it's source operand to it's target operand. The simplest move instruction is MOV and does just that. The rest of the move instructions are conditional move instructions, these only copy data when certain conditions are met. These conditions all have to do with previous calculations and use the flags register.

 $\mathbf{MOVZ}\,:\, \mathbf{Move}$ if the zero flag is set.

 \mathbf{MOVNZ} : Move if the zero flag is clear.

MOVS: Move if the sign flag is set.

 \mathbf{MOVNS} : Move if the sign flag is clear.

MOVP: Move if the parity flag is set.

MOVNP: Move if the parity flag is clear.

MOVO: Move if the overflow flag is set.

 $\mathbf{MOVNO}\,:\, \mathbf{Move}$ if the overflow flag is clear.

MOVC: Move if the carry flag is set.

MOVNC: Move if the carry flag is clear.

MOV= : Move if the previous comparison was equal.

MOV!= : Move if the previous comparison was unequal.

MOVU< : Move if the previous comparison of unsigned integers was smaller.

MOVU<=: Move if the previous comparison of unsigned integers was smaller or equal.

MOVU>=: Move if the previous comparison of unsigned integers was greater or equal.

MOVU> : Move if the previous comparison of unsigned integers was greater.

MOVS< : Move if the previous comparison of signed integers was smaller.

MOVS<=: Move if the previous comparison of signed integers was smaller or equal.

MOVS>=: Move if the previous comparison of signed integers was greater or equal.

MOVS>: Move if the previous comparison of signed integers was greater.

Addition and Subtraction Instructions

There preform binary addition and subtraction for both signed and unsigned integers. In all the addition instructions you can two's compliment negate one of the operands. A substraction instruction is an addition instruction with the B operand negated. There are three groups of addition/subtraction instructions:

ADD / **SUB** : Add/subtract the two operands.

ADD1 / SUB1 : Add/subtract the two operands and add/subtract 1.

ADDC / **SUBB** : Add/subtract the two operands and, if the carry flag is set, add/subtract 1.

Logical Instructions

The logical instructions preform simple bitwise logical functions. In all logical Instructions you can one's compliment negate (NOT) one or both of the operands.

OR : Perform bitwise logical OR.

AND: Perform bitwise logical AND.

XOR : Perform bitwise logical XOR.

Bit Shifting Instructions

The bit shifting instructions shift or rotate the bits of its operand by one bit.

SHL: Shifts all bits one bit left and clears the new least significant bit.

SHR: Shifts all bits one bit right and clears the new most significant bit.

SHL1: Shifts all bits one bit left and sets the new least significant bit.

SHR1: Shifts all bits one bit right and sets the new most significant bit.

ROL: Shifts all bits one bit left and writes the old most significant bit to the new least significant bit.

ROR : Shifts all bits one bit right and writes the old least significant bit to the new most significant bit

RCL: Shifts all bits one bit left and clears new the least significant bit.

RCR: Shifts all bits one bit right and clears new the most significant bit.

Control Instructions

Control Instruction change specific parts of the processor state. Unlike the move and computational instructions these instructions don't necessarily read from operand B and write/read from A .

WRI: Sets or clears the I flag.

ISTR : Writes to an background register.

ILD : Writes from an background register.

OUT: Writes to an output register.

IN : Request input from a pheriperal.

BACK: Backup the most registers.

FRET: Fully restore the last backup and marks and marks any running interrupt done.

PRET: Restore the program counter and the stack pointer from the last backup and marks any running interrupt done.

 \mathbf{FJMP} : Restore the program counter and the stack pointer from the last backup.

HLT: Halts the processor untill a interrupt is triggered.

NOP: Does nothing but delay the processor a bit.

Reserved Instructions

Reserved instructions are instructions with undefined behaviour, meaning they can do anything to the state of the CPU or the RAM. Reserved instructions should never be used in code. However, reserved instructions will never message pheripers, so it's impossible to change non-volatile memory or damage any components. Thus any problems that reserved instructions invoke directly will not persist after restarting the system.

Artifact Instructions

Artifact instructions are instructions with defined behaviour, but aren't considererd a permanent part of the ABCD processor family. These instructions tend to only exist because it would take more parts to remove them than it takes to leave them in.

5.2 Instruction Format

All instructions follow one single format. The format has four parts:

- ${\bf OP}$: The opcode, this describes what operation will be excecuted by this instruction (see Appendix A).
- **A** : Operand A, This describes what the instruction will use as output and secondary input (see Table 3.1).
- ${\bf B}$: Operand B, This describes what the instruction will use primary input (see Table 3.1).
- # : The immediate, This is number is used by some operands to add to it's data or offset it's address.

32 # (16) # (16) OP (6) 11 6 0

Figure 5.1: Instruction Layout

Memory Management

6.1 Segments

Base

Limit

- 6.2 Code Segment (CS)
- 6.3 Data Segment (DS)
- 6.4 Segment Switching
- 6.5 Changing Segments

Interrupts

7.1 Interrupt Enabling/Disabling

Interrupt Enabled Flag

Internal Interrupt Mask

Programmable Interrupt Controller (PIC)

7.2 State Preservation

Backup

Full Restore

Partial Restore

7.3 Interrupt Service Routine (ISR)

Interrupt Far Jump

Return to Context

Switch Context

I/O Conventions

- 8.1 Reading Input
- 8.2 Writing Output

Problem with Interrupts

Instruction Set

$9.1 \quad MOVZ / MOV =$

Move if zero / Move if equal

Usages

```
0x00: MOVZ a b
0x00: MOV= a b (Alternative notation)
```

Description

Writes B to A if the zero flag is set. This instruction can also be used to only write B to A if, in the last comparison, B was equal to A. The alternative Memnonic MOV= was provided for this use. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if } Z \textbf{ then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if } b_{previous} = a_{previous} \textbf{ then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \textbf{Where } b_{previous} \textbf{ and } a_{previous} \textbf{ are the} \\ \\ \textbf{end if} \\ \end{array}
```

Flags Affected

$9.2 \quad MOVNZ / MOV! =$

Move if not zero / Move if not equal

Usages

0x01: MOVNZ a b

0x01: MOV!= a b (Alternative notation)

Description

Writes B to A if the zero flag is clear. This instruction can also be used to only write B to A if, in the last comparison, B was unequal to A. The alternative Memnonic MOV= was provided for this use. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
if \neg Z then a \leftarrow b end if
```

alternative pseudocode:

```
if b_{previous} \neq a_{previous} then b_{previous} and a_{previous} are the operands of a previous calculation a \leftarrow b end if
```

Flags Affected

9.3 MOVS

$\underline{\text{Mov}}$ e if $\underline{\text{sign}}$

Usages

 $0\mathrm{x}02\mathrm{:}\ \mathrm{MOVS}$ a b

Description

Writes B to A if the sign flag is set.

Operation

 $\begin{array}{c} \textbf{if} \ S \ \textbf{then} \\ a \leftarrow b \\ \textbf{end} \ \textbf{if} \end{array}$

Flags Affected

9.4 MOVNS

$\underline{Mov}e \ if \ \underline{n}ot \ \underline{s}ign$

Usages

 $0\mathrm{x}03\mathrm{:}\ \mathrm{MOVNS}$ a b

Description

Writes B to A if the sign flag is clear.

Operation

```
\label{eq:stress} \begin{array}{c} \mathbf{if} \ \neg S \ \mathbf{then} \\ a \leftarrow b \\ \mathbf{end} \ \mathbf{if} \end{array}
```

Flags Affected

9.5 MOVP

$\underline{\mathbf{Mov}}\mathbf{e}$ if $\underline{\mathbf{p}}\mathbf{arity}$

Usages

0x04: MOVP a b

Description

Writes B to A if the parity flag is set.

Operation

 $\begin{array}{c} \textbf{if} \ P \ \textbf{then} \\ a \leftarrow b \\ \textbf{end} \ \textbf{if} \end{array}$

Flags Affected

9.6 MOVNP

$\underline{Mov}e \ if \ \underline{n}ot \ \underline{p}arity$

Usages

 $0\mathrm{x}05\mathrm{:}\ \mathrm{MOVNP}$ a b

Description

Writes B to A if the parity flag is clear.

Operation

```
\label{eq:continuous} \begin{array}{c} \mathbf{if} \ \neg P \ \mathbf{then} \\ a \leftarrow b \\ \mathbf{end} \ \mathbf{if} \end{array}
```

Flags Affected

9.7 MOVO

$\underline{\text{Mov}}$ e if $\underline{\text{o}}$ verflow

Usages

 $0\mathrm{x}06\mathrm{:}\ \mathrm{MOVO}$ a b

Description

Writes B to A if the overflow flag is set.

Operation

 $\begin{array}{c} \textbf{if} \ O \ \textbf{then} \\ a \leftarrow b \\ \textbf{end} \ \textbf{if} \end{array}$

Flags Affected

9.8 MOVNO

$\underline{\text{Mov}}$ e if $\underline{\text{no}}$ overflow

Usages

 $0\mathrm{x}07\mathrm{:}\ \mathrm{MOVNO}$ a b

Description

Writes B to A if the overflow flag is clear.

Operation

```
\label{eq:continuous} \begin{array}{c} \mathbf{if} \ \neg O \ \mathbf{then} \\ a \leftarrow b \\ \mathbf{end} \ \mathbf{if} \end{array}
```

Flags Affected

9.9 MOVC / MOVU>

Move if carry / Move if unsigned greater

Usages

```
0x08: MOVC a b
0x08: MOVU> a b (Alternative notation)
```

Description

Writes B to A if the carry flag is set. This instruction can also be used to only write B to A if, in the last comparison of unsigned numbers, B was greater than A. The alternative Memnonic MOVU> was provided for this use. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if } C \textbf{ then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if } b_{previous} > a_{previous} \textbf{ then} \\ \\ \textbf{operands of a previous calculation} \\ \\ a \leftarrow b \\ \textbf{end if} \end{array} \Rightarrow \textbf{Where } b_{previous} \textbf{ and } a_{previous} \textbf{ are the} \\ \\ \textbf{end if} \end{array}
```

Flags Affected

$9.10 \quad MOVNC / MOVU <=$

Move if no carry / Move if unsigned smaller or equal

Usages

```
0x09: MOVNC a b
0x09: MOVU<= a b (Alternative notation)
```

Description

Writes B to A if the carry flag is clear. This instruction can also be used to only write B to A if, in the last comparison of unsigned numbers, B was smaller than or equal to A. The alternative Memnonic MOVU<= was provided for this use. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if} \ \neg C \ \textbf{then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if} \ b_{previous} \leq a_{previous} \ \textbf{then} \\ \\ \textbf{operands of a previous calculation} \\ a \leftarrow b \\ \\ \textbf{end if} \end{array} \Rightarrow \textbf{Where} \ b_{previous} \ \textbf{and} \ a_{previous} \ \textbf{are the} \\ \\ \textbf{operands of a previous calculation} \\ \\ a \leftarrow b \\ \\ \textbf{end if} \end{array}
```

Flags Affected

9.11 MOVU>=

Move if unsigned greater or equal

Usages

0x0A: MOVU >= a b

Description

Writes B to A if, in the last comparison of unsigned numbers, B was greater than or equal to A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if } C \vee Z \textbf{ then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if } b_{previous} \geq a_{previous} \textbf{ then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \textbf{Where } b_{previous} \textbf{ and } a_{previous} \textbf{ are the} \\ \textbf{operands of a previous calculation} \\ \\ a \leftarrow b \\ \textbf{end if} \end{array}
```

Flags Affected

9.12 MOVU<

Move if unsigned smaller

Usages

0x0B: MOVU < a b

Description

Writes B to A if, in the last comparison of unsigned numbers, B was smaller than A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if} \ \neg(C \lor Z) \ \textbf{then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if} \ b_{previous} < a_{previous} \ \textbf{then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \text{Where } b_{previous} \ \text{and } a_{previous} \ \text{are the} \\ \\ \textbf{end if} \\ \end{array}
```

Flags Affected

9.13 MOVS>

Move if signed greater

Usages

0x0C: MOVS > a b

Description

Writes B to A if, in the last comparison of signed numbers, B was greater than A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if} \ \neg Z \wedge (O \oplus S) \ \textbf{then} \\ a \leftarrow b \\ \textbf{end if} \\ \text{alternative pseudocode:} \\ \textbf{if} \ b_{previous} > a_{previous} \ \textbf{then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \end{array} \Rightarrow \text{Where } b_{previous} \ \text{and } a_{previous} \ \text{are the} \\ \text{operands of a previous calculation} \\ \end{array}
```

Flags Affected

$9.14 \quad MOVS \le =$

Move if signed smaller or equal

Usages

 $0x0D: MOVS \le a b$

Description

Writes B to A if, in the last comparison of signed numbers, B was smaller than or equal to A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if } Z \vee \neg (O \oplus S) \textbf{ then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if } b_{previous} \leq a_{previous} \textbf{ then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \textbf{Where } b_{previous} \textbf{ and } a_{previous} \textbf{ are the} \\ \\ \textbf{end if} \\ \end{array}
```

Flags Affected

9.15 MOVS<

Move if signed smaller

Usages

0x0E: MOVS < a b

Description

Writes B to A if, in the last comparison of signed numbers, B was smaller than A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if } O \oplus S \textbf{ then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if } b_{previous} < a_{previous} \textbf{ then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \textbf{Where } b_{previous} \textbf{ and } a_{previous} \textbf{ are the} \\ \textbf{operands of a previous calculation} \\ \\ a \leftarrow b \\ \textbf{end if} \\ \end{array}
```

Flags Affected

9.16 MOVS>=

Move if signed greater or equal

Usages

0x0F: MOVS >= a b

Description

Writes B to A if, in the last comparison of signed numbers, B was greater than or equal to A. The comparison use of this instruction can only be expected to work if the last calculation was a compare or subtract instruction, and the used status flags haven't changed in the meanwhile.

Operation

```
\begin{array}{l} \textbf{if} \ \neg O \oplus S \ \textbf{then} \\ a \leftarrow b \\ \textbf{end if} \\ \\ \textbf{alternative pseudocode:} \\ \textbf{if} \ b_{previous} \geq a_{previous} \ \textbf{then} \\ \text{operands of a previous calculation} \\ a \leftarrow b \\ \textbf{end if} \\ \end{array} \Rightarrow \text{Where } b_{previous} \ \text{and } a_{previous} \ \text{are the} \\ \\ \textbf{end if} \\ \end{array}
```

Flags Affected

9.17 MOV

$\underline{\mathbf{Mov}}\mathbf{e}$

Usages

0x10: MOV a b

Description

Writes B to A unconditionally.

Operation

 $a \leftarrow b$

Flags Affected

9.18 WRI

$\underline{\mathbf{Wr}}$ ite $\underline{\mathbf{i}}$ nterupts enabled

Usages

0x12: WRI $_$ b

Description

Writes the least significant bit of B to the interrupts enabled flag. This allows you to enable and disable interrupts in a single instruction.

Operation

 $I \leftarrow b \ \& \ 0001_{16}$

 \triangleright Everything but the least significant bit is omitted

Flags Affected

Ι

9.19 ISTR

$\underline{\mathbf{I}}\mathbf{n}\mathbf{t}\mathbf{e}\mathbf{r}\mathbf{n}\mathbf{a}\mathbf{l}\ \underline{\mathbf{s}}\mathbf{t}\mathbf{o}\underline{\mathbf{r}}\mathbf{e}$

Usages

 $0\mathrm{x}14\mathrm{:}\ \mathrm{ISTR}$ a b

Description

Writes B to internal register A.

Operation

 $a \leftarrow b$

 \triangleright Where a is an internal register

Flags Affected

9.20 ILD

$\underline{\mathbf{I}}\mathbf{n}\mathbf{t}\mathbf{e}\mathbf{r}\mathbf{n}\mathbf{a}\mathbf{l}\ \underline{\mathbf{l}}\mathbf{o}\mathbf{a}\underline{\mathbf{d}}$

Usages

0x15: ILD a b

Description

Writes internal register B to A.

Operation

 $a \leftarrow b$

 ${\,\vartriangleright\,}$ Where b is an internal register

Flags Affected

9.21 OUT

\underline{Out} put

Usages

0x16: OUT a b

Description

Writes B to output register A.

Operation

 $a \leftarrow b$

 \triangleright Where a is an output register

Flags Affected

9.22 IN

$\underline{\mathbf{Input}}$

Usages

0x17: IN a b

Description

Request input from pheriperal B and write the returned input to A.

Operation

 $a \leftarrow \text{RequestInputFrom}(b)$

Flags Affected

9.23 BACK

Backup

Usages

0x18: BACK $_$ $_$

Description

Backups all the registers that have a backup registers

Operation

```
bax \leftarrow ax
```

 $bbx \leftarrow bx$

 $bcx \leftarrow cx$

 $bdx \leftarrow dx$

 $bex \leftarrow ex$

 $btm \leftarrow tm$

 $bsp \leftarrow sp$

 $bpc \leftarrow pc$

 $bsf \leftarrow sf$

 \triangleright Backup the segment and flag registers

Flags Affected

9.24 FRET

$\underline{\mathbf{F}}\mathbf{ull}\ \underline{\mathbf{rest}}\mathbf{ore}$

Usages

0x19: FRET $_$ $_$

Description

Restores some of the registers and marks that the there is no interrupt being handled

Operation

```
\begin{array}{lll} ax \leftarrow bax \\ bx \leftarrow bbx \\ cx \leftarrow bcx \\ dx \leftarrow bdx \\ ex \leftarrow bex \\ tm \leftarrow btm \\ sp \leftarrow bsp \\ pc \leftarrow bpc \\ sf \leftarrow bsf \\ & \rhd \text{Restore the segment and flag registers} \\ & \texttt{MARKInterruptEnded}(\ ) \\ & \rhd \text{Unlock the internal lock that prevents} \end{array}
```

Flags Affected

Z, S, P, O, C and I

9.25 PRET

Partial restore

Usages

```
0x1A: PRET \_ \_
```

Description

Restores some of the registers and marks that the there is no interrupt being handled.

Operation

```
\begin{array}{ll} pc \leftarrow bpc \\ sf \leftarrow bsf \\ \text{MarkInterruptEnded()} \end{array} \Rightarrow \text{Restore the segment and flag registers} \\ \text{bullock the internal lock that prevents} \\ \text{interrupts} \end{array}
```

Flags Affected

Z, S, P, O, C and I

9.26 FJMP

$\underline{\mathbf{F}}\mathbf{ar}\ \underline{\mathbf{j}}\mathbf{u}\underline{\mathbf{mp}}$

Usages

0x1B: FJMP $_$ $_$

Description

Restores some of the registers. $\,$

Operation

$$\begin{array}{c} pc \leftarrow bpc \\ sf \leftarrow bsf \end{array}$$

 \triangleright Restore the segment and flag registers

Flags Affected

Z, S, P, O, C and I

9.27 HLT

$\underline{\mathbf{H}}\mathbf{a}\underline{\mathbf{l}}\mathbf{t}$

Usages

0x1C: HLT $_$ $_$

Description

Halts the CPU until an external interrupt is triggered and resumes.

Operation

HaltCPU()

Flags Affected

9.28 NOP

$\underline{\mathbf{N}}\mathbf{o}$ operations

Usages

0x1D: NOP $_$ $_$

Description

Does nothing but delay the cpu a minimal amount of cycles.

Operation

if false then

 \triangleright NOP is implemented as a do never conditional

end if

Flags Affected

9.29 OR

Bitwise or

Usages

```
0x20: OR a b
0x21: OR !a b
0x22: OR a !b
0x23: OR !a !b
0x23: NAND a b (Alternative notation)
```

Description

preforms bitwise logical or operation on A and B and writes the answer to A. The answer is also evaluated to update the status flags. Both operands can be One's Complement Negated (NOT'ed) before the operation in the same instruction. This is marked by the '!' before the operand.

Operation

```
\begin{aligned} y &\leftarrow (!)a \mid (!)b \\ a &\leftarrow y \\ Z, S, P &\leftarrow \text{EvaluateAnswer}(y) \end{aligned}
```

Flags Affected

```
Z, S and P
```

9.30 AND

Bitwise and

Usages

```
0x24: AND a b
0x25: AND !a b
0x26: AND a !b
0x27: AND !a !b
0x27: NOR a b (Alternative notation)
```

Description

preforms bitwise logical and operation on A and B and writes the answer to A. The answer is also evaluated to update the status flags. Both operands can be One's Complement Negated (NOT'ed) before the operation in the same instruction. This is marked by the '!' before the operand.

Operation

```
\begin{aligned} y &\leftarrow (!)a \ \& \ (!)b \\ a &\leftarrow y \\ Z, S, P &\leftarrow \text{EvaluateAnswer}(y) \end{aligned}
```

Flags Affected

```
Z, S and P
```

9.31 XOR

Bitwise exclusive or

Usages

0x28: XOR a b 0x29: XOR !a b 0x29: XAND a b (Alternative notation) 0x2A: XOR a !b (Artifact instruction)

0x2B: XOR !a !b (Artifact instruction)

Description

preforms bitwise logical exclusive or operation on A and B and writes the answer to A. The answer is also evaluated to update the status flags. Both operands can be One's Complement Negated (NOT'ed) before the operation in the same instruction. This is marked by the '!' before the operand. However because of the logical properties of exclusive or, "XOR !a !b" and "XOR a !b" behave exactly the same as "XOR a b" and "XOR !a b" respectively. Thus instruction "XOR !a !b" and "XOR a !b" are deemed to be artifact instructions and shouldn't be used.

Operation

```
\begin{aligned} y &\leftarrow (!)a \ ^{\wedge} (!)b \\ a &\leftarrow y \\ Z, S, P &\leftarrow \text{EvaluateAnswer}(y) \end{aligned}
```

Flags Affected

Z, S and P

9.32 ADD / SUB

Addition / Subtraction

Usages

0x28: ADD a b 0x29: ADD -a b 0x2A: ADD a -b

0x2A: SUB a b (Alternative notation)

Description

adds A to B and writes the answer to A. The addition is checked on overflow and carry and the answer is also evaluated to update the status flags. Both operands can be Two's Complement Negated before the operation in the same instruction. This is marked by the '-' before the operand.

Operation

```
\begin{aligned} y &\leftarrow (-)a + (-)b \\ a &\leftarrow y \\ O, C &\leftarrow \text{CheckAddition}() \\ Z, S, P &\leftarrow \text{EvaluateAnswer}(y) \end{aligned}
```

Flags Affected

Z, S, P, O and C

9.33 ADD1 / SUB1

Addition plus $\underline{1}$ / Subtraction minus $\underline{1}$

Usages

```
0x28: ADD1 a b (Artifact instruction)
0x29: ADD1 -a b (Artifact instruction)
0x2A: ADD1 a -b (Artifact instruction)
0x2A: SUB1 a b (Artifact instruction) (Alternative notation)
```

Description

adds A and 1 to B and writes the answer to A. The addition is checked on overflow and carry and the answer is also evaluated to update the status flags. Both operands can be Two's Complement Negated before the operation in the same instruction. This is marked by the '-' before the operand. If one of the operands is negated the extra 1 will be negated as well

Operation

```
\begin{aligned} y &\leftarrow (-)a + (-)b + (-)1 \\ a &\leftarrow y \\ O, C &\leftarrow \text{CheckAddition}() \\ Z, S, P &\leftarrow \text{EvaluateAnswer}(y) \end{aligned}
```

Flags Affected

Z, S, P, O and C

9.34 ADDC / SUBB

<u>Add</u>ition with <u>carry</u> / <u>Sub</u>traction with <u>b</u>orrow

Usages

0x28: ADDC a b 0x29: ADDC -a b 0x2A: ADDC a -b

0x2A: SUBB a b (Alternative notation)

Description

adds A to B and writes the answer to A, however, if the carry flag is set an extra 1 will be added to the sum. The addition is checked on overflow and carry and the answer is also evaluated to update the status flags. Both operands can be Two's Complement Negated before the operation in the same instruction. This is marked by the '-' before the operand. If one of the operands is negated the extra 1 will be negated as well.

Operation

```
y \leftarrow (-)a + (-)b + (-)C
a \leftarrow y
O, C \leftarrow \text{CheckAddition}()
Z, S, P \leftarrow \text{EvaluateAnswer}(y)
```

Flags Affected

Z, S, P, O and C

9.35

Usages

0x: a b

Description

Operation

 $a \leftarrow b$

Flags Affected

Appendix A

Instruction Set Listings

Table A.1: List of instructions sorted by opcode

Decimal	Opcode Hex	Binary	Memnonic	Operand A	Operand B
0	0x00	000000	MOVZ & MOV=	W?	R?
1	0x01	000001	MOVNZ & MOV!=	W?	R?
2	0x02	000010	MOVS	W?	R?
3	0x03	000011	MOVNS	W?	R?
4	0x04	000100	MOVP	W?	R?
5	0x05	000101	MOVNP	W?	R?
6	0x06	000110	MOVO	W?	R?
7	0x07	000111	MOVNO	W?	R?
8	0x08	001000	MOVC & MOVU>	W?	R?
9	0x09	001001	$MOVNC \& MOVU \le$	W?	R?
10	0x0A	001010	MOVU>=	W?	R?
11	0x0B	001011	MOVU <	W?	R?
12	0x0C	001100	MOVS>	W?	R?
13	0x0D	001101	$MOVS \le $	W?	R?
14	0x0E	001110	MOVS <	W?	R?
15	0x0F	001111	MOVS > =	W?	R?
16	0x10	010000	MOV	W	R
17	0x11	010001	n/a	n/a	n/a
18	0x12	010010	WRI	_	R
19	0x13	010011	n/a	n/a	n/a
20	0x14	010100	STRB	O	\mathbf{R}
21	0x15	010101	LDB	W	O
22	0x16	010110	OUT	O	R
23	0x17	010111	IN	W	O
24	0x18	011000	BACK	_	_
25	0x19	011001	FRET	_	_
26	0x1A	011010	PRET	_	_
27	0x1B	011011	FJMP	_	_
28	0x1C	011100	HLT	_	_
29	0x1D	011101	NOP	_	_
30	0x1E	011110	CMP	R	R

Table A.1: List of instructions sorted by opcode

Decimal	Opcode Hex	Binary	Memnonic		Operand B
31	0x1F	011111	TEST	R	R
32	0x20	100000	OR	R&W	R
33	0x21	100001	OR	!R&W	R
34	0x22	100010	OR	R&W	!R
35	0x23	100011	OR	!R&W	!R
36	0x24	100100	AND	R&W	R
37	0x25	100101	AND	!R&W	R
38	0x26	100110	AND	R&W	!R
39	0x27	100111	AND	!R&W	!R
40	0x28	101000	XOR	R&W	R
41	0x29	101001	XOR	!R&W	R
42	0x2A	101010	XOR	R&W	!R
43	0x2B	101011	XOR	!R&W	!R
44	0x2C	101100	ADD	R&W	R
45	0x2D	101101	ADD	!R&W	R
46	0x2E	101110	ADD & SUB	R&W	!R
47	0x2F	101111	n/a	n/a	n/a
48	0x30	110000	ADD1	R&W	R
49	0x31	110001	ADD1	!R&W	R
50	0x32	110010	ADD1 & SUB1	R&W	!R
51	0x33	110011	n/a	n/a	n/a
52	0x34	110100	ADDC	R&W	R
53	0x35	110101	ADDC	!R&W	R
54	0x36	110110	ADDC & SUBB	R&W	!R
55	0x37	110111	n/a	n/a	n/a
56	0x38	111000	SHL	W	R
57	0x39	111001	SHL1	W	R
58	0x3A	111010	RCL	W	R
59	0x3B	111011	ROL	W	R
60	0x3C	111100	SHR	W	R
61	0x3D	111101	SHR1	W	R
62	0x3E	111110	RCR	W	R
63	0x3F	111111	ROR	W	R

Appendix B

Simplified Mnemonics

Appendix C

Common Procedures

Appendix D

Standard Peripherals

- D.1 Programmable Interrupt Controller (PIC)
- D.2 Keyboard
- D.3 Programmable Interrupt Timer (PIT)
- D.4 Sound Card
- D.5 Graphical Card
- D.6 Memory Control Hub (MCH)
- D.7 Segements and Out Of Bounds Exception