Fundamental Algorithms CSCI-GA.1170-001/Summer 2016

Homework 10

Problem 1 (CLRS 32.1-2). (1 point) Suppose that all characters in the pattern P are different. Show how to accelerate NAIVE-STRING-MATCHER to run in time O(n) on an n-character text T.

Problem 2 (CLRS 32.1-4). (2 points) Suppose we allow the pattern P to contain occurrences of a *gap character* \Diamond that can match an arbitrary string of characters (even one of zero length). For example, the pattern $ab\Diamond ba\Diamond c$ occurs in the text cabccbacbacab as $c\underline{ab}cc\underline{ba}cba\underline{c}ab$ and as cabccbacbacab.

Note that the gap character may occur an arbitrary number of times in the pattern but not at all in the text. Give a polynomial-time algorithm to determine whether such a pattern P occurs in a given text T, and analyze the running time of your algorithm.

Problem 3 (CLRS 32.2-1). (1 point) Working modulo q = 11, how many spurious hits does the Rabin-Karp matcher encounter in the text T = 3141592653589793 when looking for the pattern P = 26?

Problem 4 (CLRS 32.3-1). (1 point) Construct the string-matching automaton for the pattern P = aabab and illustrate its operation on the text string T = aaababaabaabaaba.

Problem 5 (CLRS 32.3-3). (1 point) We call a pattern P nonoverlappable if $P_k \supset P_q$ implies k = 0 or k = q. Describe the state-transition diagram of the string-matching automaton for a nonoverlappable pattern.

Problem 7 (CLRS 32.4-7). (1 point) Give a linear-time algorithm to determine whether a text T is a cyclic rotation of another string T'. For example, arc and car are cyclic rotations of each other.

Problem 8. (3 points) The *longest palindromic substring* is a maximum-length contiguous substring of a given string that is a palindrome. For example, the longest palindromic substring of ultramarine is ramar.

Give an efficient algorithm to determine the longest palindromic substring of a given string. Explain the algorithm and illustrate its operation on the string evenness.