Assignment 4

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1. Basketball

A. Probability Curry makes his 14^{th} shot

Let X_i represent Curry's i^{th} shot. Then

$$X_i = \begin{cases} 1 & \text{if he makes it} \\ 0 & \text{if he misses} \end{cases}$$

Now,

$$P_{X_{i+1}}(1) = \begin{cases} .6 & \text{if } X_i = 1\\ .3 & \text{if } X_i = 0 \end{cases}$$

Thus, the only shot that matters when predicting the $i+1^{st}$ shot is the outcome of the i^{th} shot. Therefore, when predicting the outcome of the 14^{th} shot given the outcomes of the 2^{nd} and 12^{th} , the outcome of the 12^{th} shot provides all necessary information- knowing the outcome of the second shot gives no additional information. (More formally, let $m, n \in \{1, 2, ... j - 1\}$ and m > n. Then $P(X_j | X_m, X_n) = P(X_j | X_m)$.) Now, since $X_{12} = 1$,

$$P_{X_{13}|X_{12}=1}(x) = \begin{cases} .6 & \text{for } x = 1\\ .4 & \text{for } x = 0 \end{cases}$$

Moreover,

$$P_{X_{14}|X_{13}=1}(1) = .6$$

 $P_{X_{14}|X_{13}=0}(1) = .3$

Therefore,

$$P_{X_{14}}(1) = P_{X_{14}|X_{13}=0}(1) \cdot P_{X_{13}|X_{12}=1}(0) + P_{X_{14}|X_{13}=1}(1) \cdot P_{X_{13}|X_{12}=1}(1)$$

= .3 \cdot .4 + .6 \cdot .6 = .12 + .36 = .48

B. Probability Curry made his 1^{st} shot

First, note
$$P(X_1 = 1 | X_3 = 0, X_9 = 0) = P(X_1 = 1 | X_3 = 0)$$
. Then $P(X_1 = 1 | X_3 = 0)$

$$= \frac{P(X_3 = 0|X_1 = 1)P(X_1 = 1)}{P(X_3 = 0)}$$

$$= \frac{(P(X_3 = 0|X_2 = 0)P(X_2 = 0|X_1 = 1) + P(X_3 = 0|X_2 = 1)P(X_2 = 1|X_1 = 1))P(X_1 = 1)}{P(X_3 = 0)}$$

And

$$P(X_3 = 0) = P(X_3 = 0|X_2 = 0)P(X_2 = 0|X_1 = 0)P(X_1 = 0)$$

$$+ P(X_3 = 0|X_2 = 1)P(X_2 = 1|X_1 = 0)P(X_1 = 0)$$

$$+ P(X_3 = 0|X_2 = 0)P(X_2 = 0|X_1 = 1)P(X_1 = 1)$$

$$+ P(X_3 = 0|X_2 = 1)P(X_2 = 1|X_0 = 1)P(X_1 = 1)$$

$$= .7 \cdot .7 \cdot .6 + .4 \cdot .3 \cdot .6 + .7 \cdot .4 \cdot .4 + .4 \cdot .6 \cdot .4$$

$$= .294 + .072 + .112 + .096$$

$$= .574$$

Thus,
$$P(X_1 = 1 | X_3 = 0)$$

$$= \frac{(P(X_3 = 0|X_2 = 0)P(X_2 = 0|X_1 = 1) + P(X_3 = 0|X_2 = 1)P(X_2 = 1|X_1 = 1))P(X_1 = 1)}{P(X_3 = 0)}$$

$$= \frac{(.7 \cdot .4 + .4 \cdot .6) \cdot .4}{.574}$$

$$= \frac{.52 \cdot .4}{.574} = .36237$$

C. EXPECTED NUMBER OF CONSECUTIVE BASKETS

Let Y be the number of shots made in a row following the first shot. Then

$$E(Y) = E(E(Y|X_1))$$

$$= E(Y|X_1 = 0) \cdot P(X_1 = 0) + E(Y|X_1 = 1) \cdot P(X_1 = 1)$$

$$= E(Y|X_1 = 0) \cdot .6 + E(Y|X_1 = 1) \cdot .4$$

Next, note

$$E(Y|X_1 = 0) = \sum_{y=0}^{\infty} y \cdot P(Y = y|X_1 = 0)$$

and

$$P_{Y|X_1=0}(y) = \begin{cases} .7 & \text{for } y = 0\\ .3 \cdot .6^{y-1} \cdot .4 & \text{for } y \ge 1 \end{cases}$$

Therefore,

$$E(Y|X_1 = 0) = \sum_{y=0}^{\infty} y \cdot P(Y = y|X_1 = 0)$$

$$= 0 \cdot .7 + \sum_{y=1}^{\infty} (y \cdot .3 \cdot .6^{y-1} \cdot .4)$$

$$= .3 \cdot .4 \cdot \sum_{y=1}^{\infty} (y \cdot .6^{y-1})$$

$$= .3 \cdot .4 \cdot \frac{1}{.6} \cdot \sum_{y=1}^{\infty} (y \cdot .6^y)$$

$$= .3 \cdot .4 \cdot \frac{1}{.6} \cdot \frac{.6}{(1 - .6)^{.2}}$$

$$= .75$$

Similarly, note

$$E(Y|X_1 = 1) = \sum_{y=0}^{\infty} y \cdot P(Y = y|X_1 = 1)$$

and

$$P_{Y|X_1=1}(y) = .6^y \cdot .4 \text{ for } y \ge 0$$

Therefore

$$E(Y|X_1 = 1) = \sum_{y=0}^{\infty} y \cdot P(Y = y|X_1 = 1)$$

$$= \sum_{y=0}^{\infty} (y \cdot .6^y \cdot .4)$$

$$= .4 \cdot \sum_{y=0}^{\infty} (y \cdot .6^y)$$

$$= .4 \cdot \frac{.6}{(1 - .6)^{.2}}$$

$$= 1.5$$

Finally,

$$E(Y) = E(Y|X_1 = 0) \cdot .6 + E(Y|X_1 = 1) \cdot .4 = .75 \cdot .6 + 1.5 \cdot .4 = 1.05$$

Therefore, we expect Curry to make 1.05 shots in a row regardless of whether he makes the first shot or not.

2. Model

A. FINDING P_N AND P_C

First, we find P_N :

$$P_N(n) = \sum_{k=0}^n \sum_{c \in \{1/4, 4/5\}} P_{N,C,K}(n, c, k)$$

$$= \sum_{k=0}^n \left[\frac{1}{30} \binom{n}{k} (1/4)^k (3/4)^{n-k} + \frac{1}{60} \binom{n}{k} (4/5)^k (1/5)^{n-k} \right]$$

$$= \frac{1}{30} \sum_{k=0}^n \left[\binom{n}{k} (1/4)^k (3/4)^{n-k} \right] + \frac{1}{60} \sum_{k=0}^n \left[\binom{n}{k} (4/5)^k (1/5)^{n-k} \right]$$

$$= \frac{1}{30} \cdot 1 + \frac{1}{60} \cdot 1 = 1/20$$

Next, we find P_C :

$$P_C(c) = \begin{cases} \sum_{n=1}^{20} \left[\sum_{k=0}^n \left(\frac{1}{30} \binom{n}{k} (1/4)^k (3/4)^{n-k} \right) \right] & \text{for } c = 1/4 \\ \sum_{n=1}^{20} \left[\sum_{k=0}^n \left(\frac{1}{60} \binom{n}{k} (4/5)^k (1/5)^{n-k} \right) \right] & \text{for } c = 4/5 \end{cases}$$
$$= \begin{cases} \sum_{n=1}^{20} \frac{1}{30} = \frac{20}{30} = \frac{2}{3} & \text{for } c = 1/4 \\ \sum_{n=1}^{20} \frac{1}{60} = \frac{20}{60} = \frac{1}{3} & \text{for } c = 4/5 \end{cases}$$

Finally, we will show $P_{N,C}(n,c) = P_N(n) \cdot P_C(c)$ (and, as such, C and N are independent):

$$P_{N,C}(n,c) = \sum_{k=0}^{n} P_{N,C,K}(n,c,k)$$

$$= \begin{cases} \sum_{k=0}^{n} \frac{1}{30} \binom{n}{k} (1/4)^{k} (3/4)^{n-k} & \text{for } c = 1/4 \\ \sum_{k=0}^{n} \frac{1}{60} \binom{n}{k} (4/5)^{k} (1/5)^{n-k} & \text{for } c = 4/5 \end{cases}$$

$$= \begin{cases} \frac{1}{30} & \text{for } c = 1/4 \\ \frac{1}{60} & \text{for } c = 4/5 \end{cases}$$

But then note:

$$P_{N,C}(n,c) = \begin{cases} \frac{1}{30} = \frac{1}{20} \cdot \frac{2}{3} = P_N(n) \cdot P_C(c) \text{ for } c = 1/4, n \in \{1, 2, \dots 20\} \\ \frac{1}{60} = \frac{1}{20} \cdot \frac{1}{3} = P_N(n) \cdot P_C(c) \text{ for } c = 4/5, n \in \{1, 2, \dots 20\} \end{cases}$$

Thus, C and N are independent

B. REAL-LIFE EXAMPLE

This model could correspond to the following scenario:

- 1. C: Pick one of two biased coins. Specifically, pick:
 - Coin 1 with probability 2/3. This coin flips heads with probability 1/4.
 - Coin 2 with probability 1/3. This coin flips heads with probability 4/5.
- 2. N: Spin a fair (i.e. random uniform) spinner, numbered 1 to 20, to determine how many times to flip said coin.
- 3. K: Count the number of heads in your N flips.

This scenario would be well modeled by the joint pmf given in the problem.

C. CONDITION PMF $P_{N|K,C}$

First, note

$$P(N = n | C = c, K = k) = \frac{P(N = n, C = c, K = k)}{P(C = c, K = k)}$$

Thus, to proceed we first determine:

$$P_{C,K}(c,k) = \sum_{n=0}^{20} P_{N,C,K}(n,c,k)$$

$$= \begin{cases} \sum_{n=1}^{20} 1/30 \binom{n}{k} (1/4)^k (3/4)^{n-k} & \text{for } c = 1/4 \\ \sum_{n=1}^{20} 1/60 \binom{n}{k} (4/5)^k (1/5)^{n-k} & \text{for } c = 4/5 \end{cases}$$

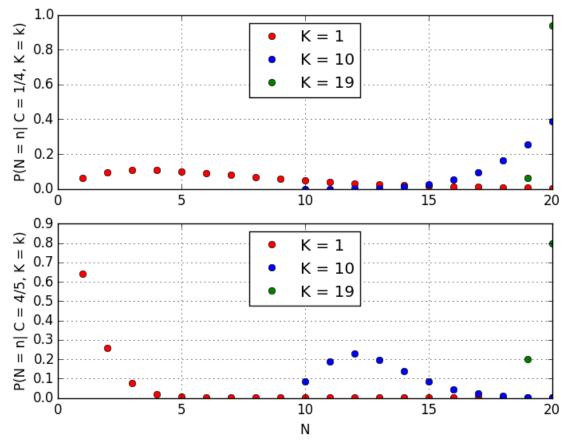
Therefore:

$$P_{N|C,K}(n|c,k) = \begin{cases} \frac{1/30\binom{n}{k}(1/4)^k(3/4)^{n-k}}{1/30(1/4)^k \sum_{n=1}^{20} \binom{n}{k}(3/4)^{n-k}} & \text{for } c = 1/4 \\ \frac{1/60\binom{n}{k}(4/5)^k(1/5)^{n-k}}{1/60(4/5)^k \sum_{n=1}^{20} \binom{n}{k}(1/5)^{n-k}} & \text{for } c = 4/5 \end{cases}$$

$$= \begin{cases} \frac{\binom{n}{k}(3/4)^{n-k}}{\sum_{n=1}^{20} \binom{n}{k}(3/4)^{n-k}} & \text{for } c = 1/4 \\ \frac{\binom{n}{k}(1/5)^{n-k}}{\sum_{n=1}^{20} \binom{n}{k}(1/5)^{n-k}} & \text{for } c = 4/5 \end{cases}$$

D. PLOTTING THE DISTRIBUTION $P_{N|K,C}$

Below are plots of the distribution of N given K=1,K=10,K=19 and C=1/4 on one graph and K=1,K=10,K=19 and C=4/5 on another graph.



These results make intuitive sense in the context of the example posed in question b. First, comparing the PMFs when C = 4/5 and C = 1/4, it is clear that for any K = k lower values of N are more likely when C = 4/5 than when C = 1/4. More specifically,

it is apparent from the plots that given any K = k, for any value x such that $k \le x < 20$

$$P(N \le x | C = 1/4) < P(N \le x | C = 4/5)$$

In the context of the problem, that makes sense- if you get heads with probability 4/5, you'd expect to need fewer flips to get k heads than you would if you had the coin with 1/4 probability of heads.

3. Router

A. MEAN AND VARIANCE OF NUMBER OF PACKETS

Let N be the number of packets that arrive at the router in a second. Note $N \sim \operatorname{Pois}(\lambda)$. Let X be the number of packets routed through connection 1. Then $X \sim \operatorname{B}(N,p)$. Now we can find the mean and variance of X:

$$E(X) = E(E(X|N))$$

$$= E(N \cdot p) = p \cdot E(N) = p \cdot \lambda$$

$$Var(X) = E(Var(X|N)) + Var(E(X|N))$$

$$= E(N \cdot p \cdot (1-p)) + Var(N \cdot p)$$

$$= p \cdot (1-p) \cdot E(N) + p^2 \cdot Var(N)$$

$$= p \cdot (1-p) \cdot \lambda + p^2 \cdot \lambda$$

$$= p \cdot \lambda$$

Thus, $E(X) = Var(X) = p \cdot \lambda$.

B. PMF OF NUMBER OF PACKETS

The PMF $P_X(x)$ is

$$P_X(x) = \sum_{n=0}^{\infty} P(X = x | N = n) \cdot P(N = n)$$

Note since $N \geq X$ by definition, this is equivalent to

$$P_X(x) = \sum_{n=x}^{\infty} \left[\binom{n}{x} p^x (1-p)^{n-x} \cdot \frac{\lambda^n}{n!} e^{-\lambda} \right]$$

$$= p^x e^{-\lambda} \cdot \sum_{n=x}^{\infty} \left[\frac{n!}{(n-x)!x!} \cdot (1-p)^{n-x} \cdot \frac{\lambda^n}{n!} \right]$$

$$= \frac{p^x e^{-\lambda}}{x!} \cdot \sum_{n=x}^{\infty} \left[\frac{(1-p)^{n-x}}{(n-x)!} \cdot \lambda^n \right]$$

$$= \frac{p^x e^{-\lambda}}{x!} \cdot \sum_{n=x}^{\infty} \left[\frac{(1-p)^{n-x}}{(n-x)!} \cdot \lambda^{n-x} \cdot \lambda^x \right]$$

$$= \frac{p^x e^{-\lambda} \lambda^x}{x!} \cdot \sum_{n=x}^{\infty} \left[\frac{[(1-p)\lambda]^{n-x}}{(n-x)!} \right]$$

$$= \frac{p^x e^{-\lambda} \lambda^x}{x!} \cdot \sum_{n=0}^{\infty} \left[\frac{[(1-p)\lambda]^n}{n!} \right]$$

$$= \frac{p^x e^{-\lambda} \lambda^x}{x!} e^{(1-p)\cdot\lambda}$$

$$= \frac{(p\lambda)^x}{x!} e^{\lambda-p\lambda-\lambda} = \frac{(p\lambda)^x}{x!} e^{-p\lambda}$$

Thus, $P_X(x) = \frac{(p\lambda)^x}{x!} e^{-p\lambda}$ and $X \sim \text{Pois}(p\lambda)$.

4. Cheap GPS

A. FINDING PRECISION Δ_1

First, recall Jennifer has defined the precision of the location estimate as the smallest Δ such that the probability of the error being larger than Δ is smaller than 1%. Thus, using X_i as an estimator of d_i , we want to find Δ such that:

$$P(|X_i - d_i| > \Delta) < .01$$
$$P(|Z_i| > \Delta) < .01$$

Using Chebyshev's inequality, we know

$$P(|Z_i - E(Z_i)| > a) \le \frac{Var(Z_i)}{a^2}$$

Thus

$$P(|Z_i - E(Z_i)| = |Z_i - 0| = |Z_i| > \Delta) < \frac{1}{\Delta^2}$$

So $P(|Z_i| > \Delta) < 0.1$ when $\Delta \ge 10$. Hence, the precision Δ_1 of using X_i as an estimate of d_i is $\Delta_1 = 10$.

B. FINDING PRECISION Δ_2

First,

$$Y_{i} = \frac{1}{m} \sum_{j=i-(m-1)}^{i} X_{i}$$

$$= \frac{1}{m} \sum_{j=i-(m-1)}^{i} (d_{j} + Z_{j})$$

$$= \frac{1}{m} \sum_{j=i-(m-1)}^{i} (d_{j}) + \frac{1}{m} \sum_{j=i-(m-1)}^{i} (Z_{j})$$

Let's consider just

$$\sum_{j=i-(m-1)}^{i} (d_j)$$

First, since Mary never moves backward and has a maximum speed of 2 meters per second,

$$d_i - 2(n) \le d_{i-n} \le d_i$$

Thus, we can place the following bound on the sum of the d_i s:

$$\sum_{j=i-(m-1)}^{i} (d_i - 2(i-j)) \le \sum_{j=i-(m-1)}^{i} d_j \le \sum_{j=i-(m-1)}^{i} d_i$$

$$m \cdot d_i - 2 \sum_{j=i-(m-1)}^{i} (i-j) \le \sum_{j=i-(m-1)}^{i} d_j \le \sum_{j=i-(m-1)}^{i} d_i$$

$$m \cdot d_i - 2 \sum_{j=0}^{m-1} (j) \le \sum_{j=i-(m-1)}^{i} d_j \le \sum_{j=i-(m-1)}^{i} d_i$$

$$m \cdot d_i - 2 \frac{(m-1)m}{2} \le \sum_{j=i-(m-1)}^{i} d_j \le \sum_{j=i-(m-1)}^{i} d_i$$

$$m \cdot d_i - (m-1)m \le \sum_{j=i-(m-1)}^{i} d_j \le \sum_{j=i-(m-1)}^{i} d_i$$

Now recall

$$Y_i = \frac{1}{m} \sum_{j=i-(m-1)}^{i} d_j + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j$$

Therefore, multiplying by $\frac{1}{m}$ and adding $\frac{1}{m}\sum_{j=i-(m-1)}^{i}Z_{j}$ throughout bounds Y_{i} :

$$\frac{1}{m}\left(m \cdot d_i - (m-1)m\right) + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j \le Y_i \le \frac{1}{m} \sum_{j=i-(m-1)}^{i} d_i + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j$$

Simplifying, we get

$$d_{i} - (m-1) + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_{j} \leq Y_{i} \leq \frac{1}{m} \cdot m \cdot d_{i} + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_{j}$$
$$-(m-1) + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_{j} \leq Y_{i} - d_{i} \leq \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_{j}$$

This implies

$$|Y_i - d_i| \le \max \left\{ \left| -(m-1) + \frac{1}{m} \sum_{j=i-(m-1)}^i Z_j \right|, \left| \frac{1}{m} \sum_{j=i-(m-1)}^i Z_j \right| \right\}$$

But, by the triangle inequality,

$$\left| -(m-1) + \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j \right| \le |m-1| + \left| \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j \right|$$

Thus, we have

$$|Y_i - d_i| \le |m - 1| + \left| \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j \right|$$

That implies (noting $m \ge 1$)

$$P(|Y_i - d_i| > \Delta) \le P\left(|m - 1| + \left|\frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j\right| > \Delta\right) = P\left(\left|\frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j\right| > \Delta - (m-1)\right)$$

Now let's let

$$X = \frac{1}{m} \sum_{j=i-(m-1)}^{i} Z_j$$

Then E(X) = 0 (since $E(Z_j) = 0$ and expectation is a linear operator), and (by the independence of the Z_j s)

$$Var(X) = \sum_{j=i-(m-1)}^{i} Var\left(\frac{1}{m}Z_j\right) = \sum_{j=i-(m-1)}^{i} \frac{1}{m^2} Var(Z_j) = 0$$

Now, setting

$$P(|X| > \Delta - (m-1)) < 0.01$$

Chebyshev's inequality implies

$$0.01 \ge \frac{Var(X)}{(\Delta - (m-1))^2}$$

Thus

$$(\Delta_2 - (m-1))^2 = \frac{100}{m} \implies \Delta_2 = 10 \cdot m^{-1/2} + m - 1$$

C. EVALUATING Δ_2 FOR m=4

When m=4,

$$\Delta_2 = 10 \cdot m^{-1/2} + m - 1 = 10 \cdot (4)^{-1/2} + 4 - 1 = 8$$

That's why Jennifer is using the running mean- it's an improved estimator of d_i (in the sense it is more precise, though it is unfortunately biased assuming Mary's speed is non-zero).

D. MINIMIZING Δ_2 FUSING THE RUNNING MEAN

First, let's express Δ_2 as a function of m:

$$\Delta_2(m) = 10 \cdot m^{-1/2} + m - 1$$

Then

$$\Delta_2(1) = 10$$
 $\Delta_2(2) \approx 8.071$
 $\Delta_2(3) \approx 7.7735$
 $\Delta_2(4) = 8$

Finally, note $d\Delta_2/dm = -5m^{-3/2} + 1 > 0$ for all $n \in \mathbb{N}, n \geq 4$. Thus, Δ_2 is strictly increasing from 4 to ∞ . Hence, the best precision Jennifer can achieve using a running mean is $\Delta_2(3) \approx 7.7735$.

5. Iterated Expectation for Random Vectors

A. PROOF FOR DISJOINT RANDOM VECTORS

We will show that for any disjoint subvectors indexed by $\mathcal{I}, \mathcal{J} \subseteq \{1, 2, ..., n\}, \mathcal{I} \cap \mathcal{J} = \emptyset$,

$$E(E(\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{I}})) = E(\mathbf{X}_{\mathcal{I}})$$

Note we will demonstrate this in the continuous case- the discrete case is similar (replacing integrals with sums).

First, let

$$\begin{split} h(\vec{x_{\mathcal{J}}}) &:= E(\mathbf{X}_{\mathcal{I}} | \mathbf{X}_{\mathcal{J}} = \vec{x_{\mathcal{J}}}) \\ &= \int_{\mathcal{I}} \vec{x_{\mathcal{J}}} \cdot f_{\mathbf{X}_{\mathcal{I}} | \mathbf{X}_{\mathcal{J}}} (\vec{x_{\mathcal{I}}} | \vec{x_{\mathcal{J}}}) \mathbf{d} \vec{x_{\mathcal{I}}} \end{split}$$

Before we proceed, let's clarify this abuse of notation:

$$\int_{\mathcal{I}} \text{ represents } \int_{x_{\mathcal{I}_1} = -\infty}^{\infty} \int_{x_{\mathcal{I}_2} = -\infty}^{\infty} \dots \int_{x_{\mathcal{I}_i} = -\infty}^{\infty}$$

$$\mathbf{d}\vec{x_{\mathcal{I}}}$$
 represents $\mathrm{d}x_{\mathcal{I}_i}\mathrm{d}x_{\mathcal{I}_{i-1}}...\mathrm{d}x_{\mathcal{I}_1}$

Additionally, note that while $h(\vec{x_J}) = E(\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{J}} = \vec{x_J})$ is a function mapping $\mathbb{R}^j \to \mathbb{R}^j$ (where j is the length of $\vec{x_J}$), $h(\mathbf{X}_{\mathcal{J}}) = E(\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{J}})$ is a random vector (since a function of a random vector is itself a random vector). Then

$$\begin{split} E(E(\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{J}})) &= E(h(\mathbf{X}_{\mathcal{J}})) \\ &= \int_{\mathcal{J}} h(\vec{x_{\mathcal{J}}}) \cdot f_{\mathbf{X}_{\mathcal{J}}}(\vec{x_{\mathcal{J}}}) \mathbf{d}\vec{x_{\mathcal{J}}} \\ &= \int_{\mathcal{J}} \left(\int_{\mathcal{I}} \vec{x_{\mathcal{I}}} f_{\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{J}}}(\vec{x_{\mathcal{I}}}|\vec{x_{\mathcal{J}}}) \mathbf{d}\vec{x_{\mathcal{I}}} \right) \cdot f_{\mathbf{X}_{\mathcal{J}}}(\vec{x_{\mathcal{J}}}) \mathbf{d}\vec{x_{\mathcal{J}}} \\ &= \int_{\mathcal{I}} \vec{x_{\mathcal{I}}} \left(\int_{\mathcal{J}} f_{\mathbf{X}_{\mathcal{I}}|\mathbf{X}_{\mathcal{J}}}(\vec{x_{\mathcal{I}}}|\vec{x_{\mathcal{J}}}) \cdot f_{\mathbf{X}_{\mathcal{J}}}(\vec{x_{\mathcal{J}}}) \mathbf{d}\vec{x_{\mathcal{J}}} \right) \mathbf{d}\vec{x_{\mathcal{I}}} \\ &= \int_{\mathcal{I}} \vec{x_{\mathcal{I}}} f_{\mathbf{X}_{\mathcal{I}}}(\vec{x_{\mathcal{I}}}) \mathbf{d}\vec{x_{\mathcal{I}}} \\ &= E(\mathbf{X}_{\mathcal{T}}) \end{split}$$

B. MEAN AND VARIANCE OF K IN PROBLEM 2

First, from question 2 we know

$$P_{K|N,C} = \frac{P_{N,C,K}}{P_{N,C,}}$$
$$= {N \choose k} C^k (1 - C)^{N-k}$$

So

$$E(K|N,C) = \sum_{k=0}^{N} k \binom{N}{k} C^k (1-C)^{N-k}$$
$$= C \cdot N$$

But then

$$E(K) = E(E(K|N,C)) = E(CN)$$

$$= \sum_{n=1}^{20} \sum_{c \in \{1/4,4/5\}} c \cdot n \cdot P_{N,C}(n,c)$$

$$= \sum_{n=1}^{20} \left[\frac{1}{4} \cdot n \cdot \frac{1}{30} + \frac{4}{5} \cdot n \cdot \frac{1}{60} \right]$$

$$= \sum_{n=1}^{20} \frac{13}{600} \cdot n$$

$$= \frac{13 \cdot 20 \cdot 21}{600 \cdot 2} = 4.55$$

Next, let's find the variance of K. First, note

$$Var(K) = E(K^2) - (E(K))^2$$

Well, $E(K^2) = E(E(K^2|N, C))$, and

$$E(K^{2}|N,C) = \sum_{k=0}^{N} k^{2} {N \choose k} C^{k} (1-C)^{N-k}$$

But this is just the second moment of a binomial (though, significantly, a binomial parameterized by random variables). Recall the variance of B(N, C) is NC(1 - C). So

$$E(K^{2}|N,C) = Var(K|N,C) + (E(K|N,C))^{2}$$
$$= NC(1-C) + (NC)^{2}$$

Therefore,

$$E(K^{2}) = E(E(K^{2}|N,C))$$

$$= \sum_{n=1}^{20} \sum_{c \in \{1/4,4/5\}} \left((nc(1-c) + (nc)^{2}) P_{N,C}(n,c) \right)$$

$$= \sum_{n=1}^{20} \left[\left(n \cdot \frac{1}{4} \cdot \frac{3}{4} + \left(n \cdot \frac{1}{4} \right)^{2} \right) \frac{1}{30} + \left(n \cdot \frac{4}{5} \cdot \frac{1}{5} + \left(n \cdot \frac{4}{5} \right)^{2} \right) \frac{1}{60} \right]$$

$$= \sum_{n=1}^{20} \left[\left(\frac{3}{16} n + \frac{1}{16} n^{2} \right) \frac{1}{30} + \left(\frac{4}{25} n + \frac{16}{25} n^{2} \right) \frac{1}{60} \right]$$

$$= 38.465$$

So
$$Var(K) = E(K^2) - (E(K))^2 = 38.465 - (4.55)^2 = 17.7625$$
.