

Example: "Brute force" for GCD

(greatest common divisor)

$$\text{gcd}(a, b) = \max_{d \in \mathbb{Z}} \{ d \mid a \text{ and } d \mid b \}.$$

$$\exists q \in \mathbb{Z} \text{ s.t. } a = d \cdot q$$

↑  
"there exists"

$$\text{Es., } \text{gcd}(6, 9) = 3$$

$\begin{array}{cc} 2 \cdot 3 & 3 \cdot 3 \\ \underline{\quad} & \underline{\quad} \end{array}$

Idea: start w/ largest candidate ( $\min(a, b)$ )  
↓  
count backwards until the first common divisor is found.

```
int main() {  
    int a, b;  
    int min;  
    cin >> a >> b;  
    if (a < b)  
        min = a;  
    else  
        min = b;  
    int d = min; // d = candidate for gcd  
    while (a % d != 0 || b % d != 0)  
        d--; // keep looking  
    cout << "gcd = " << d;
```

return 0;

Exercise! Find the exponent of 2 in the factorization of an integer.

E.g. if input = 40 = 5 · 2<sup>③</sup> ← answer is 3.

$$9 = 3^2 · 2^{\textcircled{0}}$$

$$6 = 3 · 2^{\textcircled{1}}$$

Again, we'll kind of "brute force" the answer.

Idea: start w/ 2<sup>0</sup> and keep multiplying by 2 until division has a remainder.

$$\underline{1 \leq k} \quad == 2^k$$

$$100 = 10^2 \quad (\text{base } 10)$$

$$100 = 2^2 \quad (\text{base } 2)$$

Alternate approach:

count the # of times 2 divides the input.

$$n = r \cdot 2^k, \quad r \text{ odd (so } k == \text{answer)}$$

$$\begin{array}{ccc} (n/2) / 2 / 2 \\ \uparrow \quad \uparrow \quad \uparrow \end{array}$$

count # of divisions to find k.

$$40 / 2 = \underline{20}$$

$$20 / 2 = 10$$

$$10 / 2 = 5 \text{ odd so } k=3.$$

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```
int n;  
cin >> n;  
int k = 0; // count the times n is  
            divided by 2.  
while (n % 2 == 0) {  
    n = n / 2;  
    k++; // k = k + 1  
}  
cout << k;
```