

662. Maximum Width of Binary Tree

(/problems/maximum-width-of-binary-tree/)

Dec. 10, 2017 | 23.3K views

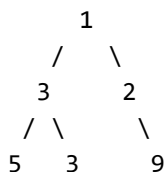
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Given a binary tree, write a function to get the maximum width of the given tree. The width of a tree is the maximum width among all levels. The binary tree has the same structure as a **full binary tree**, but some nodes are null.

The width of one level is defined as the length between the end-nodes (the leftmost and right most non-null nodes in the level, where the `null` nodes between the end-nodes are also counted into the length calculation.

Example 1:

Input:



Output: 4

Explanation: The maximum width existing in the third level with the length 4 (5,3,null,9).

Example 2:

Input:

```

      1
     /
    3
   / \
  5   3

```

Output: 2**Explanation:** The maximum width existing in the third level with the length 2 (5,3).**Example 3:****Input:**

```

      1
     / \
    3   2
   /
  5

```

Output: 2**Explanation:** The maximum width existing in the second level with the length 2 (3,2).**Example 4:****Input:**

```

      1
     / \
    3   2
   /   \
  5     9
 /     \
6       7

```

Output: 8**Explanation:** The maximum width existing in the fourth level with the length 8 (6,null,null,null,null,null,null,7).**Note:** Answer will in the range of 32-bit signed integer.**Approach Framework**

Explanation

As we need to reach every node in the given tree, we will have to traverse the tree, either with a depth-first search, or with a breadth-first search.

The main idea in this question is to give each node a `position` value. If we go down the left neighbor, then `position -> position * 2`; and if we go down the right neighbor, then `position -> position * 2 + 1`. This makes it so that when we look at the position values `L` and `R` of two nodes with the same depth, the width will be `R - L + 1`.

Approach #1: Breadth-First Search [Accepted]

Intuition and Algorithm

Traverse each node in breadth-first order, keeping track of that node's position. For each depth, the first node reached is the left-most, while the last node reached is the right-most.

Java

Python

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```

1  class Solution {
2      public int widthOfBinaryTree(TreeNode root) {
3          Queue<AnnotatedNode> queue = new LinkedList();
4          queue.add(new AnnotatedNode(root, 0, 0));
5          int curDepth = 0, left = 0, ans = 0;
6          while (!queue.isEmpty()) {
7              AnnotatedNode a = queue.poll();
8              if (a.node != null) {
9                  queue.add(new AnnotatedNode(a.node.left, a.depth + 1, a.pos * 2));
10                 queue.add(new AnnotatedNode(a.node.right, a.depth + 1, a.pos * 2 + 1));
11                 if (curDepth != a.depth) {
12                     curDepth = a.depth;
13                     left = a.pos;
14                 }
15                 ans = Math.max(ans, a.pos - left + 1);
16             }
17         }
18         return ans;
19     }
20 }
21
22 class AnnotatedNode {
23     TreeNode node;
24     int depth, pos;
25     AnnotatedNode(TreeNode n, int d, int p) {
26         node = n;
27         depth = d;

```

Complexity Analysis

- Time Complexity: $O(N)$ where N is the number of nodes in the input tree. We traverse every node.
- Space Complexity: $O(N)$, the size of our queue .

Approach #2: Depth-First Search [Accepted]

Intuition and Algorithm

Traverse each node in depth-first order, keeping track of that node's position. For each depth, the position of the first node reached of that depth will be kept in `left[depth]` .

Then, for each node, a candidate width is `pos - left[depth] + 1` . We take the maximum of the candidate answers.

Java

Python

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```

1  class Solution {
2      int ans;
3      Map<Integer, Integer> left;
4      public int widthOfBinaryTree(TreeNode root) {
5          ans = 0;
6          left = new HashMap();
7          dfs(root, 0, 0);
8          return ans;
9      }
10     public void dfs(TreeNode root, int depth, int pos) {
11         if (root == null) return;
12         left.computeIfAbsent(depth, x-> pos);
13         ans = Math.max(ans, pos - left.get(depth) + 1);
14         dfs(root.left, depth + 1, 2 * pos);
15         dfs(root.right, depth + 1, 2 * pos + 1);
16     }
17 }

```

Complexity Analysis

- Time Complexity: $O(N)$ where N is the number of nodes in the input tree. We traverse every node.
- Space Complexity: $O(N)$, the size of the implicit call stack in our DFS.

Analysis written by: @awice (<https://leetcode.com/awice>).

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 Preview

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dantegt (dantegt) ★ 7 🕒 December 11, 2017 10:05 AM



Was that really accepted? I wonder about the case when the input is a tree always going to the right. As far as I understand, your 'pos' integer variable would overflow with a depth > 32. I am asking it because it happened to me.

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Dird (dird) ★ 29 🕒 February 23, 2018 1:28 AM



What is (depth, x-> pos)?? Why not just (depth, pos)

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habibullah (habibullah) ★ 16 🕒 March 27, 2019 12:03 AM



The solution is not correct. It will overflow for some input. Converting the solution in c++ shows the overflow, overflow test cases doesn't have any max result in overflow area so java solution doesn't shows problem.

A correct approach could be to recenter the current level of tree to avoid overflow.

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shlykovich (shlykovich) ★ 14 🕒 July 26, 2019 10:10 AM



I didn't find the explanation very intuitive, so this is my attempt to make things clear.

First of all, if you have two positive integers a and b where $a < b$, than there are exactly $(b - a + 1)$ integers between them including a and b. For example if $a = 5$, $b = 10$ than there are $10 - 5 + 1 = 6$ integers.

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ritakiv (ritakiv) ★ 6 🕒 July 17, 2019 3:22 AM



For those who struggle with an overflow exception in c++, just use ulong type for pos value.

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wanders (wanders) ★ 10 🕒 May 24, 2019 10:51 PM



Similar BFS python solution using just position:

```
def widthOfBinaryTree(self, root):  
    if not root:  
        return 0
```

[Read More](#)

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vlearner (vlearner) ★ 45 🕒 May 19, 2019 9:14 AM



nice explanation!

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bxs3514 (bxs3514) ★ 5 🕒 April 22, 2019 2:31 AM



Reset the left index(position) to 0 for every level can avoid the overflow.

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fudonglai (fudonglai) ★ 380 🕒 February 28, 2019 3:33 PM



Now the solutions may cause overflow, using Python can avoid it.

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gogol1anonly (gogol1anonly) ★ 2 🕒 January 23, 2019 5:45 PM



Really good problem.

Love the twists!!!

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