

Programming and Modelling

Weekly Practice Exercises

Week 4: Automated Reasoning about Programs

Work on your solutions in groups of 2-3 students. In the next programming café the TAs will go over the questions and your answers together as a class.

Exercises

- (1) This week you will continue reading through the “Course Notes” online pages of the Logika website. The online course notes are available at:

<https://logika.v3.sireum.org/dschmidt/index.html>

Remark: There might seem like a lot of sections, but many of them are very short, just a few paragraphs.

- Section 3.4. “Logika Proof Syntax”
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- Section 3.5. “The premise Justification”
- Section 3.6. “The assume Justification”
- Section 3.7. “Logical Operator Based Justifications”
- Section 3.8. “And-Introduction and And-Elimination”
- Section 3.9. “Or-Introduction and Or-Elimination”
- Section 3.10. “Implies-Introduction and Implies-Elimination”
- Section 3.11. “Negation”
- Section 3.12. “Negation Introduction”

- (2) In the lecture we discussed the automated theorem proof checker tool Sireum Logika.

- (a) Install the Integrated Verification Environment (IVE) of Sireum Logika (version 3). Follow the instructions in the lecture slides. The download page is:

<https://logika.v3.sireum.org/doc/01-getting-started/index.html>

- (b) Following the instructions in the lecture slides, create a new Logika project and file, and verify the proof given in the lecture slides for the sequent:

$\vdash a \rightarrow a$

- (3) Consider the following Logika proof:

```
s, q, r ⊢ r ∧ (q ∧ s)
{
  1. p           assume
  2. q           premise
  3. r           premise
  4. q ∨ p       vi 2 1
  5. r ∧ (q ∧ p) ∧i 3 4
}
```

When I type this proof into Sireum Logika, I get a number of errors.

- (a) Create a new Logika file and type in this proof to reproduce the errors.
 - (b) Create your own correct proof using graphical notation.
 - (c) Translate your graphical proof into Logika's textual notation, and verify your proof in Logika.
- (4) In the lectures we went through a number of examples for translating graphical proofs into Logika textual proofs.
- (a) Verify the “*Idempotence of conjunction 1*” textual Logika proof from the lecture slides using Logika.
 - (b) Translate the following graphical proof of “*Idempotence of disjunction 1*” into the textual Logika notation, and verify it using Logika:

$$\frac{\frac{1}{[\mathbf{A}]} \vee \mathbf{I}_1}{\mathbf{A} \vee \mathbf{A}} \rightarrow \mathbf{I}.1$$

- (5) Prove the following sequents, and then translate your proofs into textual Logika notation and verify them using Logika:
- (a) $\mathbf{A} \wedge \mathbf{B} \vdash \mathbf{A} \vee \mathbf{B}$
 - (b) $\mathbf{P} \vee \mathbf{P} \vdash \mathbf{P}$
- Hint:* Try applying the or-elimination rule. Start with what you might assume in each of the branches of the disjunction. The or-elimination rule discharges these assumptions
- (6) In the lecture we discussed proof rules involving *negation*. Using the graphical notation from lectures, prove the following sequents:
- (a) $\mathbf{P} \rightarrow \mathbf{Q}, \neg \mathbf{Q} \vdash \neg \mathbf{P}$
 - (b) $\neg \mathbf{Q} \vdash \mathbf{Q} \rightarrow \mathbf{S}$
- (7) Translate each of your proofs from exercise (6) above into Logika textual notation, and use Logika to verify your proofs.
- (8) [Optional challenge] Below are some additional sequents for you to practice creating proofs.
- (a) $\vdash ((\mathbf{A} \wedge \mathbf{B}) \wedge \mathbf{C}) \rightarrow (\mathbf{A} \wedge (\mathbf{B} \wedge \mathbf{C}))$
 - (b) $\mathbf{P} \rightarrow \mathbf{Q}, \mathbf{Q} \rightarrow \mathbf{R} \vdash \mathbf{P} \rightarrow \mathbf{R}$
 - (c) $\vdash (\mathbf{P} \rightarrow \mathbf{Q}) \rightarrow (\mathbf{Q} \rightarrow \mathbf{R}) \rightarrow (\mathbf{P} \rightarrow \mathbf{R})$
 - (d) $(\mathbf{A} \vee \mathbf{B}), \mathbf{C} \vdash (\mathbf{A} \wedge \mathbf{C}) \vee (\mathbf{B} \wedge \mathbf{C})$

Hint: Look carefully at the or-introduction rule: if you have proved some expression $\mathbf{E1}$ then you can extend this to $\mathbf{E1} \vee \mathbf{E2}$, for any arbitrary expression $\mathbf{E2}$ that could be useful in your proof.