- There are three types of DFS traversals in tree \rightarrow pre, in and post
- and only one type of BFS \rightarrow level order traversal

```
class Solution {
public:
    vector<vector<int>> levelOrder(TreeNode* root) {
        vector<vector<int>> ans;
        if (root == nullptr) {
            return ans;
        }
        queue<TreeNode*> q;
        q.push(root);
        while (!q.empty()) {
            // Get the size of the current level
            int size = q.size();
            vector<int> level;
            for (int i = 0; i < size; i++) {
                TreeNode* node = q.front();
                q.pop();
                level.push_back(node->val);
                // Enqueue the child nodes if they exist
                if (node->left != nullptr) {
                    q.push(node->left);
                }
                if (node->right != nullptr) {
                    q.push(node->right);
                }
            }
            ans.push_back(level);
        }
        return ans;
    }
};
```

Height / Max Depth of BT

```
int maxDepth(TreeNode* root) {
  if (root == nullptr) return 0;
```

```
int leftDepth = maxDepth(root->left);
int rightDepth = maxDepth(root->right);
return max(leftDepth, rightDepth) + 1;
}
```

Check if tree is balanced or not

```
class Solution {
public:
    bool isBalanced(Node* root) {
        return dfsHeight(root) != -1;
   }
    int dfsHeight(Node* root) {
        if (root == NULL) return 0;
        int leftHeight = dfsHeight(root->left);
        if (leftHeight == -1)
            return -1;
        int rightHeight = dfsHeight(root->right);
        if (rightHeight == -1)
            return -1;
        if (abs(leftHeight - rightHeight) > 1)
            return -1;
        return max(leftHeight, rightHeight) + 1;
    }
};
```

Diameter of a tree

```
class Solution {
public:
    int diameterOfBinaryTree(Node* root) {
        int diameter = 0;
        height(root, diameter);

        return diameter;
    }

private:
    int height(Node* node, int& diameter) {
        if (!node) {
            return 0;
        }
}
```

```
int lh = height(node->left, diameter);
int rh = height(node->right, diameter);

diameter = max(diameter, lh + rh);

return 1 + max(lh, rh);
};
```

Max path sum

```
class Solution {
public:
    int findMaxPathSum(Node* root, int &maxi) {
        if (root == nullptr) {
            return 0;
        }
        int leftMaxPath = max(0, findMaxPathSum(root->left, maxi));
        // to make sure we don't take any -ve sum ahead
        int rightMaxPath = max(0, findMaxPathSum(root->right, maxi));
        maxi = max(maxi, leftMaxPath + rightMaxPath + root->data);
        return max(leftMaxPath, rightMaxPath) + root->data;
   }
    int maxPathSum(Node* root) {
        int maxi = INT_MIN;
        findMaxPathSum(root, maxi);
        return maxi;
   }
};
```

Check if two trees are identical

```
class Solution {
public:
    bool isIdentical(Node* node1, Node* node2){

    if( node1== NULL || node2==NULL){
      // this ensure same structure
      return node1 == node2;
    }

    return ((node1->data == node2->data)
```

```
// this ensures same values in them
    && isIdentical(node1->left, node2->left)
        && isIdentical(node1->right, node2->right));
}
```

Zig-Zag / Spiral Traversal

```
vector<vector<int>> zigzagLevelOrder(TreeNode* root) {
    vector<vector<int>> result;
    if (root == nullptr) return result;
    queue<TreeNode*> q;
    q.push(root);
    bool leftToRight = true;
    while (!q.empty()) {
        int levelSize = q.size();
        vector<int> level;
        for (int i = 0; i < levelSize; ++i) {</pre>
            TreeNode* node = q.front();
            q.pop();
            level.push_back(node->val);
            if (node->left) q.push(node->left);
            if (node->right) q.push(node->right);
        }
        // Reverse the level if we're going right to left
        if (!leftToRight) {
            reverse(level.begin(), level.end());
        }
        result.push_back(level);
        leftToRight = !leftToRight; // Toggle direction
    }
    return result;
}
```

Boundary Traversal :

```
// Utility function to check if a node is a leaf.
bool isLeaf(Node* root) {
   return !root->left && !root->right;
}
```

```
// Function to add the left boundary nodes (excluding leaf nodes).
void addLeftBoundary(Node* root, vector<int>& res) {
    Node* cur = root->left;
    while (cur) {
        if (!isLeaf(cur)) res.push_back(cur->data);
        if (cur->left) cur = cur->left;
        else cur = cur->right;
   }
}
// Function to add the right boundary nodes (excluding leaf nodes).
void addRightBoundary(Node* root, vector<int>& res) {
    Node* cur = root->right;
    vector<int> tmp; // Temporary vector to store nodes in reverse order.
    while (cur) {
        if (!isLeaf(cur)) tmp.push_back(cur->data);
        if (cur->right) cur = cur->right;
        else cur = cur->left;
    // Add the right boundary nodes in reverse order.
    for (int i = tmp.size() - 1; i >= 0; --i) {
        res.push_back(tmp[i]);
    }
}
// Function to add all the leaf nodes.
void addLeaves(Node* root, vector<int>& res) {
// any traversal is fine because we are printing
// left subtree before right subtree
    if (isLeaf(root)) {
        res.push_back(root->data);
        return;
    if (root->left) addLeaves(root->left, res);
    if (root->right) addLeaves(root->right, res);
}
// Main function to perform boundary traversal.
vector<int> printBoundary(Node* root) {
    vector<int> res:
    if (!root) return res; // Return empty vector if the tree is empty.
    // Add the root if it's not a leaf node.
    if (!isLeaf(root)) res.push_back(root->data);
    // Add the left boundary.
    addLeftBoundary(root, res);
```

```
// Add all the leaf nodes.
addLeaves(root, res);

// Add the right boundary.
addRightBoundary(root, res);

return res;
}
```

Vertical Order Traversal

```
class Solution {
public:
   // Function to perform vertical order traversal
   // and return a 2D vector of node values
   // if collision occurs return them in a sorted order,
    // considering duplicacy is allowed
   vector<vector<int>> findVertical(Node* root){
        // vertical lvl, horizontal lvl , node value
        map<int, map<int, multiset<int>>> nodes;
        // Queue for BFS traversal, each
        // element is a pair containing node
        // and its vertical and level information
        queue<pair<Node*, pair<int, int>>> todo;
        // Push the root node with initial vertical
        // and level values (0, 0)
        todo.push({root, {0, 0}});
        while(!todo.empty()){
            auto p = todo.front();
            todo.pop();
            Node* temp = p.first;
            // Extract the vertical and level information
            // x -> vertical
            int x = p.second.first;
            // y -> level
            int y = p.second.second;
            // Insert the node value into the
            // corresponding vertical and level
            // in the map
            nodes[x][y].insert(temp->data);
```

```
// Process left child
    if(temp->left){
        todo.push({
            temp->left,
            {
                // Move left in
                // terms of vertical
                x-1,
                // Move down in
                // terms of level
                y+1
            }
        });
    }
    // Process right child
    if(temp->right){
        todo.push({
            temp->right,
                // Move right in
                // terms of vertical
                x+1,
                // Move down in
                // terms of level
                y+1
            }
        });
    }
}
// Prepare the final result vector
// by combining values from the map
vector<vector<int>> ans;
for(auto p: nodes){
    vector<int> col;
    for(auto q: p.second){
        // Insert node values
        // into the column vector
        col.insert(col.end(), q.second.begin(), q.second.end());
    }
    // Add the column vector
    // to the final result
    ans.push_back(col);
return ans;
```

```
}
};
```

Top view

```
void topView(Node* root) {
    if (root == nullptr) return;
    // Map to store the first node at each horizontal distance.
    map<int, int> topNodes; // no duplicates allowed first come first served
    // It stores pairs of (node, horizontal distance).
    queue<pair<Node*, int>> q;
    q.push({root, 0});
    while (!q.empty()) {
        auto it = q.front();
        q.pop();
        Node* node = it.first;
        int horizontalDist = it.second;
        // If the horizontal distance is encountered for
        // the first time, add the node.
        if (topNodes.find(horizontalDist) == topNodes.end()) {
            topNodes[horizontalDist] = node->data;
        }
        // Move to the left and right children with
        // updated horizontal distances.
        if (node->left) {
            q.push({node->left, horizontalDist - 1});
        }
        if (node->right) {
            q.push({node->right, horizontalDist + 1});
        }
    }
    // Print the top view.
    for (auto it : topNodes) {
        cout << it.second << " ";</pre>
    }
    cout << endl;</pre>
}
```

Bottom view : If collision occurs take the rightmost guy

• same as top view but here we need to keep overwriting and what about collision, no extra handling as left is traversed first and then right

```
vector<int> bottomView(Node* root) {
   vector<int> ans;
   if (root == nullptr) return ans;
   // Map to store the last node at each horizontal distance (line).
    map<int, int> mpp; // why not unordered_map ?
    // => we need sorting, left most level first and then right
    // It stores pairs of (node, horizontal distance).
    queue<pair<Node*, int>> q;
    q.push({root, 0});
    while (!q.empty()) {
        auto it = q.front();
        q.pop();
        Node* node = it.first;
        int line = it.second;
        // Update the map with the current node for the given horizontal distance.
        mpp[line] = node->data;
        // Move to the left child with horizontal distance - 1.
        if (node->left != nullptr) {
            q.push({node->left, line - 1});
        }
        // Move to the right child with horizontal distance + 1.
        if (node->right != nullptr) {
            q.push({node->right, line + 1});
        }
   }
   // Store the final bottom view in the result vector.
   for (auto it : mpp) {
        ans.push_back(it.second);
   }
    return ans;
}
```

• Given a Binary Tree, return the left view of it

```
void recursionLeft(Node* root, int level, vector<int>& res){
   if(root == NULL){
```

```
return;
        }
        // Check if the size of the result vector
        // is equal to the current level as each level
        // can have only one node visible
        if(res.size() == level){
            res.push_back(root->data);
       }
        recursionLeft(root->left, level + 1, res);
        // First give priority to left if we want right view
        // then give it second priority
        recursionLeft(root->right, level + 1, res);
    }
vector<int> leftsideView(Node* root)
        vector<int> res;
        recursionLeft(root, 0, res);
        return res;
   }
```

Is symmetric ?

```
class Solution {
public:
    bool isSymmetric(TreeNode* root) {
        // If the tree is empty, it's symmetric.
        return root == NULL || isSymmetricHelp(root->left, root->right);
   }
private:
    bool isSymmetricHelp(TreeNode* left, TreeNode* right) {
        // If both subtrees are null, they are symmetric.
        if (left == NULL || right == NULL) return left == right; //**
        // If the values of the nodes don't match, the tree is not symmetric.
        if (left->val != right->val) return false;
        // Recursively check the left subtree of the left child
        // with the right subtree of the right child,
        // and the right subtree of the left child with the
        // left subtree of the right child.
        return isSymmetricHelp(left->left, right->right) &&
               isSymmetricHelp(left->right, right->left);
```

```
};
```

Print Path from root to Node

```
class Solution {
public:
    bool getPath(TreeNode* root, vector<int> &arr, int x) {
        // Base case: if the node is NULL, return false
        if (!root) return false;
        // Push the current node's value to the path
        arr.push_back(root->val);
        // If the current node's value is the target value, return true
        if (root->val == x) return true;
        // Recur for the left or right subtree
        if (getPath(root->left, arr, x) ||
        getPath(root->right, arr, x)) return true;
        // If the target value is not found in either subtree, backtrack
        arr.pop_back();
        return false;
    }
    vector<int> solve(TreeNode* A, int B) {
        vector<int> arr;
        // If the tree is empty, return an empty vector
        if (A == NULL) return arr;
        // Get the path from the root to the node with value B
        getPath(A, arr, B);
        return arr;
    }
};
```

Lowest Common Ancestor

 $\underline{\mathsf{Method}\ 1}$: we find paths of both roots, then compare both and find the last similar guy

Method 2 :

```
Case 1 : If we find a node return it else return NULL

Case 2 : If a node receives a node and NULL return node

Case 3 : If a node receives two non-nulls return the current node
```

```
class Solution {
public:
    TreeNode* lowestCommonAncestor(TreeNode* root, TreeNode* p, TreeNode* q) {
        // Base case: if root is NULL or root is either p or q, return root
        if (root == NULL || root == p || root == q) {
            return root;
        }
        // Recursive calls for left and right subtrees
        TreeNode* left = lowestCommonAncestor(root->left, p, q);
        TreeNode* right = lowestCommonAncestor(root->right, p, q);
        // If one subtree is NULL, return the non-NULL subtree.
        // Here we are not checking root but the return values of children
        if (left == NULL) {
            return right;
        } else if (right == NULL) {
            return left;
        } else {
            // If both left and right are non-NULL, root is the LCA
            return root;
        }
    }
};
```

Given a root, target and distance k. return the root values which are at a distance of k from target

```
parent_track[current->right] = current;
                queue.push(current->right);
           }
       }
   }
public:
   vector<int> distanceK(TreeNode* root, TreeNode* target, int k) {
        unordered_map<TreeNode*, TreeNode*> parent_track; //**
        markParents(root, parent_track, target);
       // Keeps track of visited nodes
        unordered_map<TreeNode*, bool> visited;
        queue<TreeNode*> queue;
        queue.push(target);
        // Tracks the current level
       // while traversing the tree
       int curr_level = 0;
       while (!queue.empty()) {
            int size = queue.size(); // simultaneous multiple BFS
            if (curr_level++ == k) {
                // Break if the current level
                // matches the required distance (k)
                break;
            }
            for (int i = 0; i < size; i++) {
                TreeNode* current = queue.front();
                queue.pop();
                if (current->left && !visited[current->left]) {
                    queue.push(current->left);
                    visited[current->left] = true;
                }
                if (current->right && !visited[current->right]) {
                    queue.push(current->right);
                    visited[current->right] = true;
                }
                if (parent_track[current] && !visited[parent_track[current]]) {
                    queue.push(parent_track[current]);
                    visited[parent_track[current]] = true;
                }
           }
       }
```

```
// Stores nodes at distance k from the target
vector<int> result;
while (!queue.empty()) {
    // Extract nodes at distance k from the queue
    TreeNode* current = queue.front();
    queue.pop();
    // Store node values in the result vector
    result.push_back(current->val);
}
return result;
}
```

Given a BinaryTree {has data, leftchild, rightchild}. One of its node is set on fire, find the minimum time taken for the entire tree to burn down.

- A simple BFS from the target would be sufficient to solve this problem
- But we can't backtrack to parents
- So, we will keep a track of parent node of each node

```
class Solution {
 private:
   void mapParents(Node* root, unordered_map<Node*, Node*> &parentMap,
                    Node* &targetNode, int target)
   {
        queue<Node*> q;
        q.push(root);
        while(!q.empty()) {
            Node* current = q.front();
            q.pop();
            if(current->data == target) { // Sub-Task
                targetNode = current;
            }
            if(current->left) {
                parentMap[current->left] = current;
                q.push(current->left);
            }
            if(current->right) {
                parentMap[current->right] = current;
                q.push(current->right);
            }
        }
```

```
}
public:
 int minTime(Node* root, int target)
     if(!root) return 0;
      unordered_map<Node*, Node*> parentMap;
      Node* targetNode = NULL;
     mapParents(root, parentMap, targetNode, target);
      queue<Node*> q;
      set<Node*> visited;
     q.push(targetNode);
     visited.insert(targetNode);
     int time = 0;
     while(!q.empty()) {
          int size = q.size();
          bool burned = false;
          for(int i = 0; i < size; i++) {
          // Why are we doing this ?
          // {SIMULTANEOUS BFS} other acc to this logic it would take more time
              Node* current = q.front();
              q.pop();
              // Burn left child
              if (current->left &&
              visited.find(current->left) == visited.end()) {
                  burned = true;
                  visited.insert(current->left);
                  q.push(current->left);
              }
              // Burn right child
              if (current->right &&
              visited.find(current->right) == visited.end()) {
                  burned = true;
                  visited.insert(current->right);
                  q.push(current->right);
              }
              // Burn parent
              if (parentMap[current] &&
              visited.find(parentMap[current]) == visited.end()) {
```

Count total Nodes in a COMPLETE Binary Tree

```
class Solution {
public:
   int countNodes(TreeNode* root) {
        if (root == NULL) {
            return 0;
        }
        // Find the left height and
        // right height of the tree
        int lh = findHeightLeft(root);
        int rh = findHeightRight(root);
        // Check if the last level
        // is completely filled
        if (lh == rh) {
            // If so, use the formula for
            // total nodes in a perfect
            // binary tree ie. 2^h - 1
            return (1 << lh) - 1;
        }
        // If the last level is not completely
        // filled, recursively count nodes in
        // left and right subtrees
        return 1 + countNodes(root->left) + countNodes(root->right);
   }
   int findHeightLeft(TreeNode* node) {
```

```
int height = 0;
while (node) {
    height++;
    node = node->left;
}
return height;
}

int findHeightRight(TreeNode* node) {
    int height = 0;
    while (node) {
        height++;
        node = node->right;
    }
    return height;
}
```

Given a Binary Tree convert it into a string. Now delete the Tree. Reconstruct the tree using the string and return the new root.

```
class Solution {
public:
   // Encodes the tree into a single string
    string serialize(TreeNode* root) {
        if (!root) {
            return "";
        }
        string s = "";
        queue<TreeNode*> q;
        q.push(root);
        while (!q.empty()) {
            TreeNode* curNode = q.front();
            q.pop();
            // Check if the current node is
            // null and append "#" to the string
            if (curNode == nullptr) {
                s += "#,";
            } else {
                // Append the value of the
                // current node to the string
                s += to_string(curNode->val) + ",";
                // Push the left and right children
                // to the queue for further traversal
```

```
q.push(curNode->left);
            q.push(curNode->right);
        }
    }
    return s;
}
// Decode the encoded
// data to a tree
TreeNode* deserialize(string data) {
    if (data.empty()) {
        return nullptr;
    }
    // Use a stringstream to
    // tokenize the serialized data
    stringstream s(data);
    string str;
    // Read the root value
    // from the serialized data
    getline(s, str, ',');
    TreeNode* root = new TreeNode(stoi(str));
    queue<TreeNode*> q;
    q.push(root);
    // Perform level-order traversal
    // to reconstruct the tree
    while (!q.empty()) {
        TreeNode* node = q.front();
        q.pop();
        // Read the value of the left
        // child from the serialized data
        getline(s, str, ',');
        // If the value is not "#", create a new
        // left child and push it to the queue
        if (str != "#") {
            TreeNode* leftNode = new TreeNode(stoi(str));
            node->left = leftNode;
            q.push(leftNode);
        }
        // Read the value of the right child
        // from the serialized data
        getline(s, str, ',');
        // If the value is not "#", create a
```

```
// new right child and push it to the queue
if (str != "#") {
        TreeNode* rightNode = new TreeNode(stoi(str));
        node->right = rightNode;
        q.push(rightNode);
    }
}
return root;
}
```

Morris Traversal is a method to traverse a binary tree without using additional space like a stack or recursion. It is a space-efficient algorithm because it uses no extra space apart from a few pointers. The idea is to temporarily modify the tree structure during the traversal and restore it afterward.

1. If the left child is NULL:

2. If the current node does not have a left child, print the current node's data and move to the right child.

Explanation: There's no left subtree to explore, so you simply move to the right child.

3. If the left child is not NULL:

- 4. Find the inorder predecessor of the current node in the left subtree. The inorder predecessor is the rightmost node in the left subtree.
- 5. Subcase 1: If the right child of the predecessor is NULL:

Make the current node the right child of its inorder predecessor (temporary link) and move the current node to its left child.

Explanation: This temporary link allows you to return to the current node after exploring the left subtree.

6. Subcase 2: If the right child of the predecessor is not NULL:

This means the left subtree has already been visited, so remove the temporary link (restore the original tree), print the current node's data, and move to the right child.

Explanation: You've finished visiting the left subtree, so you can now process the current node and move to the right subtree.

```
void morrisTraversal(Node* root) {
   Node* current = root;

while (current != NULL) {
   if (current->left == NULL) { // case 1
      cout << current->data << " ";
      current = current->right;
   } else { // case 2
      Node* predecessor = current->left;
```

```
while (predecessor->right != NULL && predecessor->right != current) {
    predecessor = predecessor->right;
}

if (predecessor->right == NULL) { // sub-case 1
    predecessor->right = current;
    current = current->left;
} else { // subcase 2
    predecessor->right = NULL;
    cout << current->data << " ";
    current = current->right;
}
}
}
```

In a BST, apart from the tree itself, the left and right subtrees are BST's aswell.

```
// search in BST
struct Node* search(struct Node* root, int key) {
    // Base cases: root is NULL or key is present at root
    if (root == NULL || root->data == key) // Takin care of NULL is imp
        return root;

// Key is greater than root's data
    if (root->data < key)
        return search(root->right, key);

// Key is smaller than root's data
    return search(root->left, key);
}
```

Find ceil in BST

```
int findCeil(Node* root, int input) {
  int ans = -1;

while (root != NULL) {
    // If the current node's value is equal to the input, that's the ceil.
    if (root->data == input) {
        return root->data; // no need to further traverse
    }

    // If the current node's value is greater than the input, it could be the cei
    if (root->data > input) {
        ans = root->data; // Update the answer
        root = root->left;
```

```
// Move to the left subtree to find a smaller candidate
} else {
    // If the current node's value is less than the input,
    // move to the right subtree.
    root = root->right;
}
return ans;
}
```

Insertion in a BST

```
TreeNode<int>* insertionInBST(TreeNode<int>* root, int val) {
    // If the tree is empty, return a new node
    if (!root) return new TreeNode<int>(val);
    TreeNode<int>* curr = root;
    while (true) {
        if (val <= curr->data) {
            if (curr->left) {
                curr = curr->left;
            } else {
                curr->left = new TreeNode<int>(val);
                break;
            }
        } else {
            if (curr->right) {
                curr = curr->right;
            } else {
                curr->right = new TreeNode<int>(val);
                break;
            }
        }
    }
    return root;
}
```

Given a BST and a key, if the key exists in the binary tree delete it else do nothing. return the root of the tree.

```
class Solution {
private:
    // Helper function to find the minimum value in the right subtree.
    int find_min(TreeNode* node) {
        while (node->left) {
```

```
node = node->left;
       }
        return node->val;
   }
public:
   TreeNode* deleteNode(TreeNode* root, int key) {
        // Only reached when either tree is empty or
       // the key doesn't exists in the tree
       if (!root) return root;
       // Navigate the tree based on the key value
       if (key < root->val) {
            root->left = deleteNode(root->left, key);
       }
       else if (key > root->val) {
            root->right = deleteNode(root->right, key);
       }
       else {
            // Node found and now we handle the three cases:
            // Case 1: Node has no left child
            if (!root->left) {
                TreeNode* temp = root->right;
                delete root;
                return temp;
            }
            // Case 2: Node has no right child
            else if (!root->right) {
                TreeNode* temp = root->left;
                delete root;
                return temp;
            }
            // Case 3: Node has two children
            else {
                // Get the in-order successor (smallest in the right subtree)
                int minValue = find_min(root->right);
                root->val = minValue;
                // we are not swapping we have just replaced the node value
                // with inorder successor
                // Delete the in-order successor
                root->right = deleteNode(root->right, minValue);
           }
       }
        return root;
```

```
};
```

Kth smallest element in BST

• inorder traversal of a BST is always sorted

```
class Solution {
private:
   int count = 0; // Counter to track the number of nodes visited
    int result = -1; // To store the K-th smallest element
    // Helper function to perform in-order traversal
   void inorder(TreeNode* root, int k) {
        if (!root) return;
        // Traverse the left subtree
        inorder(root->left, k);
        // Increment the counter when visiting a node
        count++;
        // If the counter equals k, we have found the K-th smallest element
        if (count == k) {
            result = root->val;
            return;
        }
        // Traverse the right subtree
        inorder(root->right, k);
   }
public:
    int kthSmallest(TreeNode* root, int k) {
        inorder(root, k);
        return result;
   }
};
```

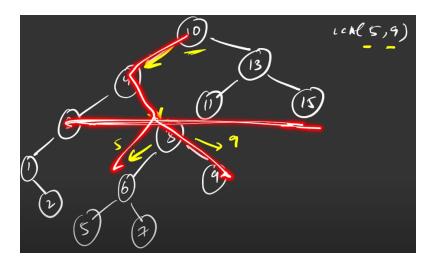
• If we want the kth largest element then modify k = n - k;

Given a BST validate if it is a valid BST

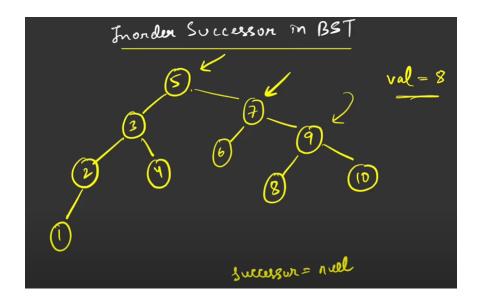
```
class Solution {
private:
   bool isIt(TreeNode* root, long long left_bound, long long right_bound) {
     if (!root) return true;
     if (root->val <= left_bound || root->val >= right_bound) {
```

• Why are we taking long long ? if INT_MAX or INT_MIN then when we check ≤ or ≥ case it is not a duplicacy!

LCA in BST



- If both of them are on the right or left we can simply move.
- The instant they we are not able to say both of them don't lie in left or right we got out LCA!
- It is not always the case that inorder successor or predecessor are children of the node.



• For 8 here, IS = 9 (It's parent)

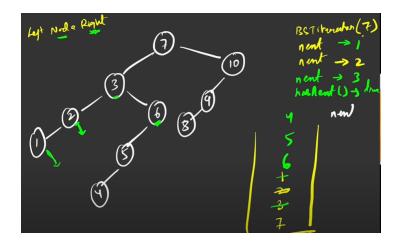
```
TreeNode* inorderSuccessor(TreeNode* root, TreeNode* p) {
    TreeNode* successor = NULL;

while (root != NULL) {
    if (p->val >= root->val) {
        root = root->right;
    } else {
        successor = root;
        root = root->left;
    }
}

return successor;
}
```

BST iterator

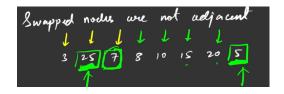
```
class BSTIterator {
   stack<TreeNode *> myStack;
public:
   BSTIterator(TreeNode *root) {
       pushAll(root);
    /** @return whether we have a next smallest number */
   bool hasNext() {
                                                    Ĭ
       return !myStack.empty();
    /** @return the next smallest number */
   int next() {
       TreeNode *tmpNode = myStack.top();
       myStack.pop();
       pushAll(tmpNode->right);
        return tmpNode->val;
private:
   void pushAll(TreeNode *node) {
       for (; node != NULL; myStack.push(node), node = node->left);
```



- So basically instead of storing entire inorder, jitna chahiye utna we are maintaining.
- ullet two sum in BST. Two pointer approach on inorder \Rightarrow BST iterator next and before.
- Instead of making a complete different class, make changes in class itself

```
class BSTIterator {
    stack<TreeNode *> myStack;
    // reverse -> true -> before
// reverse -> false -> next
    bool reverse - true;
public:
    BSTIterator(TreeNode *root, bool isReverse) {
        reverse = isReverse;
pushAll(roct);
    /** @return whether we have a next smallest number */
    bool hasNext() {
        return !myStack.empty();
     /** @return the next smallest number */
    int next() {
        TreeNode *tmpNode = myStack.top();
         myStack.pop();
         if(!reverse) pushAll(tmpNode->right);
         else pushAll(tmpNode->left);
         return tmpNode->val;
private:
    void pushAll(TreeNode *node) {
         for(;node != NULL; ) {
              myStack.push(node);
              if(reverse -- true) {
                  node = node->right;
              } else {
                  node = node->left;
```

If two elements in a BST are swapped, recover it



```
Swapped Nodes are adjecent

1 1 1 1 1 1 1 1 1

3 5 13 7 10 15 20 25
```

```
class Solution {
private:
     TreeNode* first:
     TreeNode* prev;
TreeNode* middle;
TreeNode* last;
private:
     void inorder(TreeNode* root) {
           if(root == NULL) return;
           inorder(root->left);
           if (prev != NULL && (root->val < prev->val))
                // If this is first violation, mark these two nodes as // 'first' and 'middle' \,
                if ( first == NULL )
                     first = prev;
middle = root;
                // If this is second violation, mark this node as last
                else
                     last = root;
           // Mark this node as previous
           inorder(root->right);
public:
     void recoverTree(TreeNode* root) {
   first = middle = last = NULL;
   prev = new TreeNode(INT_MIN);
           inorder(root);
           if(first && last) swap(first->val, last->val);
else if(first && middle) swap(first->val, middle->val);
```

Given a BT return the size of largest BST in it

