



# Royal University Phnom Penh IT Department

Subject : Data Structure & Algorithm

Member Of Group 6

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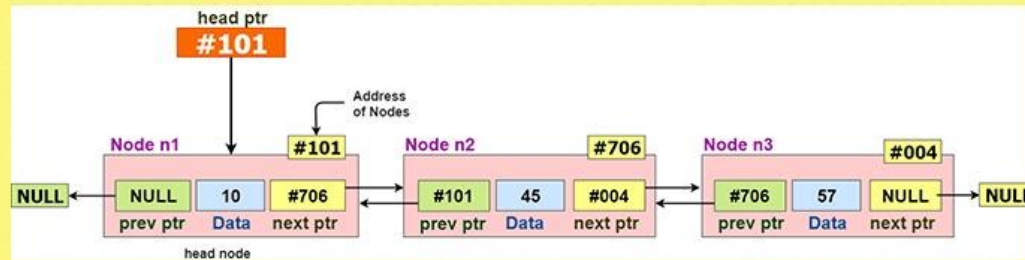
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## Topic

### DOUBLY LINKED LIST

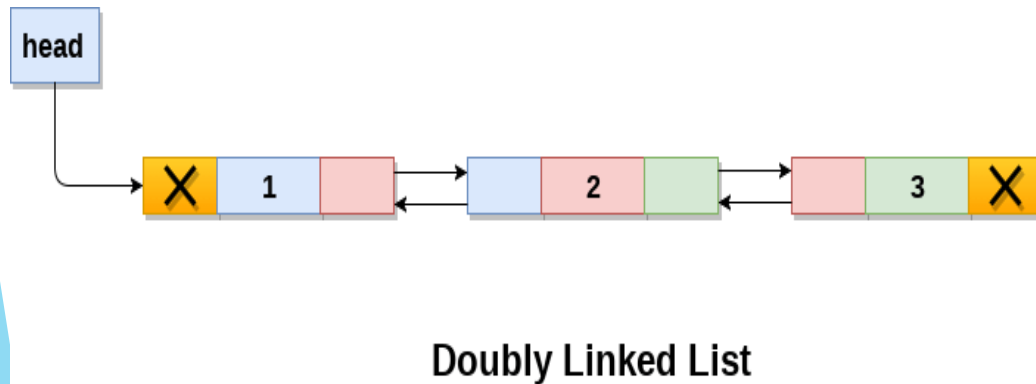
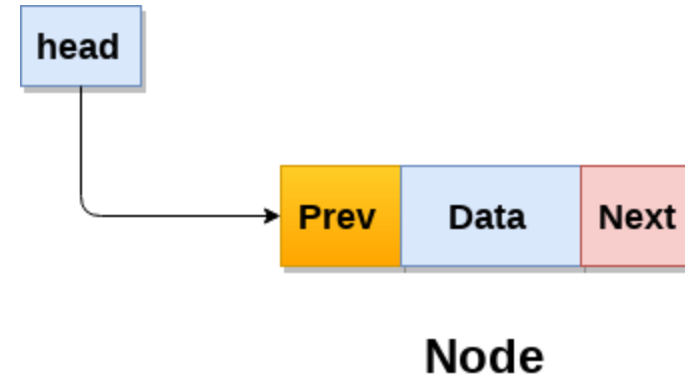


✂ SIMPLE SNIPPETS

- Tanmay Sakpal

# Introduction

Doubly linked list is a complex type of linked list in which a node contains a pointer to the previous as well as the next node in the sequence. Therefore, in a doubly linked list, a node consists of three parts: node data, pointer to the next node in sequence (next pointer), pointer to the previous node (previous pointer). A sample node in a doubly linked list is shown in the figure.



A doubly linked list containing three nodes having numbers from 1 to 3 in their data part, is shown in the following image.

## Advantages:

- Provides easy traversal in both directions.
- Insertion and deletion of nodes can be done easily from both ends.
- Improves the efficiency of algorithms where bidirectional traversal is required.

## Disadvantages:

- Doubly linked list requires more memory as compared to Singly Linked List.
- Code implementation is more complex than a singly linked list.

# structure of a node in doubly linked list in C programming

```
struct node
{
    struct node *prev;
    int data;
    struct node *next;
}
```

The **prev** part of the first node and the **next** part of the last node will always contain null indicating end in each direction.



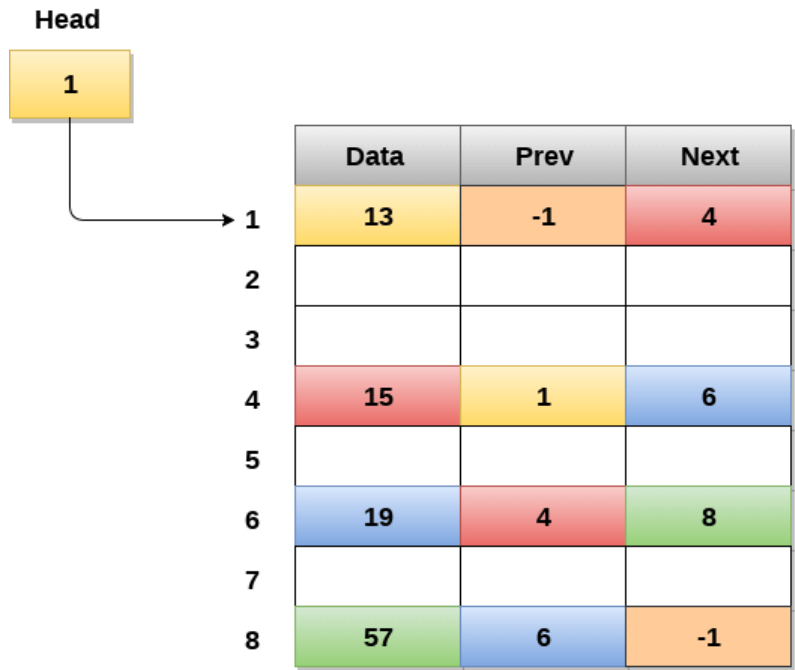
In a singly linked list, we could traverse only in one direction, because each node contains address of the next node and it doesn't have any record of its previous nodes. However, doubly linked list overcome this limitation of singly linked list. Due to the fact that, each node of the list contains the address of its previous node, we can find all the details about the previous node as well by using the previous address stored inside the previous part of each node.

# Memory Representation of a doubly linked list

Memory Representation of a doubly linked list is shown in the following image. Generally, doubly linked list consumes more space for every node and therefore, causes more expansive basic operations such as insertion and deletion. However, we can easily manipulate the elements of the list since the list maintains pointers in both the directions (forward and backward).

In the following image, the first element of the list that is i.e. 13 stored at address 1. The head pointer points to the starting address 1. Since this is the first element being added to the list therefore the prev of the list contains null. The next node of the list resides at address 4 therefore the first node contains 4 in its next pointer.

We can traverse the list in this way until we find any node containing null or -1 in its next part.



Memory Representation of a Doubly linked list

# Operations on doubly linked list

SN	Operation	Description
1	<u>Insertion at beginning</u>	Adding the node into the linked list at beginning.
2	<u>Insertion at end</u>	Adding the node into the linked list to the end.
3	<u>Insertion after specified node</u>	Adding the node into the linked list after the specified node.
4	<u>Deletion at beginning</u>	Removing the node from beginning of the list
5	<u>Deletion at the end</u>	Removing the node from end of the list.
6	Deletion of the node having given data	Removing the node which is present just after the node containing the given data.
7	<u>Searching</u>	Comparing each node data with the item to be searched and return the location of the item in the list if the item found else return null.
8	<u>Traversing</u>	Visiting each node of the list at least once in order to perform some specific operation like searching, sorting, display, etc.