

1) a) $T(n) = 9T(n/4) + n^2 \rightarrow \frac{a}{b^d} = \frac{9}{4^2} < 1 \rightarrow T(n) \in \Theta(n^d) = \Theta(n^2)$
 $\Rightarrow T(n) = \Theta(n^2)$

b) $T(n) = 3T(n/3) + \log n \rightarrow$ Derste gösterilen Master Theorem'de $\Theta(n^d)$ teriminin d 0 ile polinom olarak aldık. Bu sebeple bize derste gösterilen şekilde bir alt sınır belirleyebiliyoruz diye düşünürüm.

- $d=0$ için $P(n) = 3P(n/3) + 1$ biçiminde olur. $\frac{a}{b^d} = \frac{3}{3^0} > 1$
 $\rightarrow P(n) \in \Theta(n^{\log_b a}) = \Theta(n^{\log_3 3})$
 $P(n) = \Theta(n)$

* Sonuç olarak $T(n)$ için $d=0$ olduğu durumdan $\Omega(n)$ alt sınırı koyabiliriz
 $T(n) \in \Omega(n)$ Çünkü $P(n)$ 'de her aşamada fonksiyon çağrıları dışında sabit iş varken, $T(n)$ 'de $\log n$ 'lik bir iş yapıyor.

c) $T(n) = 3T(n/2) + n \rightarrow \frac{a}{b^d} = \frac{3}{2^1} > 1 \rightarrow T(n) \in \Theta(n^{\log_b a}) = \Theta(n^{\log_2 3})$
 $\rightarrow T(n) = \Theta(n^{\log_2 3})$

2) F1 $\rightarrow \boxed{O(n)}$

F2 $\rightarrow \sum_{i=0}^N \sum_{j=0}^i \sum_{k=0}^j 1 = \sum_{i=0}^N \sum_{j=0}^i j = \sum_{i=0}^N \left(\frac{i \cdot (i+1)}{2} \right) = \frac{1}{2} \cdot \left(\sum_{i=0}^N i^2 + \sum_{i=0}^N i \right)$
 $= \frac{1}{2} \left(\frac{N \cdot (N+1) \cdot (2N+1)}{6} + \frac{N \cdot (N+1)}{2} \right) = \boxed{O(N^3)}$

$$\begin{aligned}
 F3 \rightarrow T(N) &= N \cdot T(N-1) + N = N \cdot [(N-1) \cdot T(N-2) + N-1] + N \\
 &= N \cdot (N-1) \cdot T(N-2) + (N-1) \cdot N + N = N \cdot (N-1) \cdot (N-2) \cdot T(N-3) + N \cdot (N-1) \cdot (N-2) + N \cdot (N-1) + N \\
 &= N \cdot (N-1) \cdot \dots \cdot 2 \cdot T(1) + N! + \frac{N!}{2!} + \frac{N!}{3!} + \frac{N!}{4!} + \dots + \frac{N!}{(N-1)!} \\
 &= N! + N! + \frac{N!}{2!} + \dots + \frac{N!}{(N-1)!} = O(N!)
 \end{aligned}$$

$$F4 \rightarrow T(N) = 2T(N/2) + O(N) \rightarrow \frac{a}{b^d} = \frac{2}{2^1} = 1 \rightarrow T(N) \in \Theta(n^d \log n)$$

$$T(N) = O(N \log N)$$

3)

	$F(n)$	$g(n)$	
O	n^2	n^3	$\rightarrow \lim_{n \rightarrow \infty} \frac{n^2}{n^3} = 0 \rightarrow O$
Ω	$n \log n$	n	$\rightarrow \lim_{n \rightarrow \infty} \frac{n \log n}{n} = \infty \rightarrow \Omega$
Θ	1	$3 + \sin n$	$\rightarrow \lim_{n \rightarrow \infty} \frac{1}{3 + \sin n} = \text{sabit} \rightarrow \Theta$
Ω	3^n	2^n	$\rightarrow \lim_{n \rightarrow \infty} \frac{3^n}{2^n} = \infty \rightarrow \Omega$
Θ	4^{n+4}	2^{2n+2}	$\rightarrow \lim_{n \rightarrow \infty} \frac{4^{n+4}}{2^{2n+2}} = \lim_{n \rightarrow \infty} \frac{2^{2n+8}}{2^{2n+2}} = \text{sabit} \rightarrow \Theta$
O	$n \log n$	$n^{\frac{100}{100}}$	$\rightarrow \lim_{n \rightarrow \infty} \frac{n \log n}{n^{\frac{100}{100}}} = 0 \rightarrow O$
Θ	$\log(\log n)$	$\log n^3$	$\rightarrow \lim_{n \rightarrow \infty} \frac{\log(\log n)}{\log n^3} = \lim_{n \rightarrow \infty} \frac{\frac{1}{2} + \log n}{3 \log n} = \text{sabit} \rightarrow \Theta$
O	$n!$	$(n+1)!$	$\rightarrow \lim_{n \rightarrow \infty} \frac{n!}{(n+1)!} = 0 \rightarrow O$

4) a) $T(n) = 2^{n+1} + 3^{n-1} = 2 \cdot 2^n + \frac{3^n}{3} \rightarrow 3^n, 2^n$ 'den daha hızlı büyüyor

$$T(n) \in \Theta(3^n) \rightarrow c_1 \cdot 3^n \leq 2^{n+1} + 3^{n-1} \leq c_2 \cdot 3^n$$

$\Rightarrow c_1 = 1, c_2 = 2, n_0 = 1$ iken söyler

$$3 \leq 5 \leq 6$$

$$b) T(n) = 2n \log(n+2)^2 + (n+2)^2 \log \frac{n}{2} = 4n \log(n+2) + (n+2)^2 \log \frac{n}{2}$$

$n^2 \log n$ domine et lici $T(n) \in \Theta(n^2 \log n)$

$$c_1 \cdot n^2 \log n \leq 4n \log(n+2) + (n+2)^2 \log \frac{n}{2} \leq c_2 \cdot n^2 \log n$$

$$c_1 = 1, c_2 = 2, n_0 = 9$$

$$17.47 < 31.65 \leq 34.94$$

$$5) F(n) = \sum_{i=1}^n (i+1) \cdot 2^{i-1} = 2 \cdot 2^0 + 3 \cdot 2^1 + \dots + (n+1) \cdot 2^n$$

$$F(n) \in O(n \cdot 2^n)$$

$$F(n) = \sum_{i=1}^n 2^{i-1} + \sum_{i=1}^n i \cdot 2^{i-1} = 2^n - 1 + (n-1) \cdot 2^n + 1 = n \cdot 2^n$$

$$6) T(n) = T(n-2) + 2n$$

$$T(n) = T(n-4) + 2 \cdot (n-2) + 2n = T(n-6) + 2 \cdot (n-4) + 2 \cdot (n-2) + 2n$$

$$= T(0) + 2 \cdot 0 + 2 \cdot 2 + \dots + 2 \cdot (n-2) + 2 \cdot n = 2 \cdot \left(\frac{n}{2} \cdot \left(\frac{n}{2} + 1 \right) \right) = n \cdot \left(\frac{n}{2} + 1 \right)$$

$$= \frac{n^2}{2} + n \rightarrow T(n) = \frac{n^2}{2} + n \rightarrow T(n) \in O(n^2)$$