

## Q1

### Explaining The Algorithm

I made a dynamic programming algorithm. I proposed a subproblem which is  $\text{minimumPenalty}(i,a)$ . input  $i$  is the hotel rank and input  $a$  is the hotels distance array.

Subproblem is finding minimum penalty from starting hotel to hotel  $i$ . So I use  $\text{minimumPenalty}$  function recursively. Also its base case is when  $i$  is 0 so it returns 0 and empty path. Because there is no hotel. Firstly, I calculate the previous day stopped hotel and calculate the penalty after sum with today penalty and I choose minimum one. I store the hotel path in an array.

### Time Complexity Analyze $\Rightarrow O(n^2)$

There are  $n$  hotels and  $n$  subproblems. For one hotel, subproblem takes  $O(i)$  time. ( $i$  is the rank of the hotel). There are  $n$  hotels so time complexity is  $O(n^2)$

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$$\sum_{i=1}^n O(i) = O(1) + O(2) + O(3) + \dots + O(n)$$
$$O \sum_{i=1}^n (i) = O(1 + 2 + 3 + \dots + n) = O\left(\frac{n \cdot (n+1)}{2}\right) = O(n^2)$$

## Q2

### Explaining The Algorithm

I made a dynamic programming algorithm. I proposed a subproblem which is finding a result for string  $s[0 \dots i-1]$  and then storing it in the  $\text{resultsOfSubproblems}[i]$ . If result is true for the sub problem  $s[0 \dots i-1]$ , I store true in  $\text{resultsOfSubproblems}[i]$ . Otherwise I store false. I traverse the string and search the possible words (which is false in the  $\text{resultsOfSubproblems}$  list) in the dictionary and if I find it I make it true. When I reach end of the string and if last index of the  $\text{resultsOfSubproblems}$ , then I return true otherwise I return false.

**Time Complexity Analyze  $\Rightarrow O(n^2)$**

First for iteration takes  $O(n)$  time for  $n$  length string. Also the second for iteration takes  $n$  time for  $n$  length string.  $O(n * n) \Rightarrow$  Time complexity is  $= O(n^2)$

### Q3

#### Explaining The Algorithm

Firstly I take array of array as a parameter in the mergeTwoArrays method. If there is one array in the array then I return it (base case of recursive method). Otherwise I divide the array by two and send both two array same method recursively. Now there are two array and I send these arrays to mergeTwoArrays method for merging these two arrays and return this function output as a result of all program. In mergeTwoArrays method I traverse both two arrays and compare each indices for these two arrays. If first array element is smaller than the second array (in the same indices) then I store the first array element in the result array and increment the only first and result array indices. If I reach end one of these two arrays then I continue with the other array and store the element of this array to result array. Lastly I return the result array.

**Time Complexity Analyze  $\Rightarrow O(k * \log k n)$**

mergeTwoArrays method takes  $O(n+n)$  time. Because we traverse both two arrays which have  $n$  size. So time complexity of mergeTwoArrays is  $O(2n) = O(n)$ . Totally we do  $O(\log k)$  merging process. In mergeTwoArrays method time complexity is  $O(\log k)$  because we divide by two the problem. And we call  $O(nk)$ ,  $O(\log k)$  times. So, Total time complexity is this:  $O(k \log k n)$

### Q4

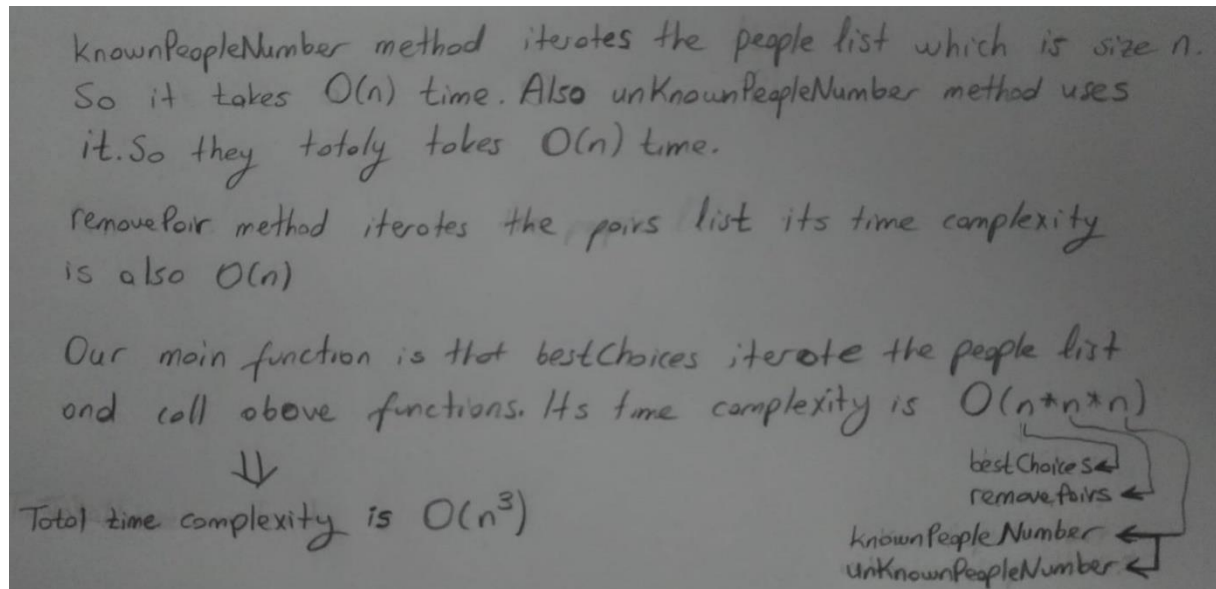
#### Explaining The Algorithm

I traverse the people list and if I see a person which is know less than 5 people or not know less than 5 people, then I remove this person from the people list and also I remove the pairs which has this person in the pairs list. Because we can't choose this person for the party.

After that if I remove a person and pairs then I iterate again the people list because this persons can effect the previous persons which have relation with this person.

I all rules is provided so we don't need to remove a person and pair we only increment the people list index and continue to iterate the list.

#### Time Complexity Analyze => $O(n^3)$



## Q5

### Explaining The Algorithm

getIndexesAndEquality method is an utility method for getting the infos from the constraints string.

Our main method is isSatisfied and in this method I iterate the constraintsList and I get the infos of constraint. With this infos I control this constraint is satisfied or not in the elementList.

If I see a constraint which is not satisfied then I return false otherwise I continue the iterating the constraintsList.

If all the constraints are satisfied then return true.

#### Time Complexity Analyze => $O(m)$

We iterate all the constraintList and its size is  $m$ . So it takes  $O(m)$  time. We don't need to use elementList for calculating the time complexity. Because we directly access the element with using the indices of the constraints.