

Relational Algebra:

Selection (σ):

Selects tuples that satisfy certain conditions.

Operators: $=, \neq, \geq, <, >, \leq$ Connectors: *and, or, not*

$\sigma_{\text{sentiment} > 0.9}$ (Chirp) — Selects Chirps with sentiment less than 0.9.

$\sigma_{\text{last_name} = \text{'Trump'}}$ (Bird) — Selects Birds whose last name is Trump.

Projection (π):

Projects a subset of a table's columns.

$\pi_{\text{tag, first_name}}$ (Bird) — Projects the tag and first name of all Birds

*Cross-Product (\times):

Combines two relations with every possible combination of tuples.

R		R CROSS S			
A	1	A	1	A	1
B	2	A	1	C	2
D	3	A	1	D	3
F	4	A	1	E	4
E	5	B	2	A	1
		B	2	C	2
		B	2	D	3
		B	2	E	4
		D	3	A	1
		D	3	C	2
		D	3	D	3
		D	3	E	4

*Difference ($-$):

Selects tuples that are present in one relation but not the other.

R		R DIFFERENCE S	
A	1		
B	2		
D	3		
F	4		
E	5		

S		S DIFFERENCE R	
A	1		
C	2		
D	3		
E	4		

*Union (\cup):

Selects tuples that are present in both relations.

R		R UNION S	
A	1	A	1
B	2	B	2
D	3	C	2
F	4	D	3
E	5	E	5

S			
A	1		
C	2		
D	3		
E	4		

Natural Join (\bowtie):

Combines two relations by finding a common attribute between them

Employee			Dept		Employee \bowtie Dept			
Name	Empld	DeptName	DeptName	Manager	Name	Empld	DeptName	Manager
Harry	3415	Finance	Finance	George	Harry	3415	Finance	George
Sally	2241	Sales	Sales	Harriet	Sally	2241	Sales	Harriet
George	3401	Finance	Production	Charles	George	3401	Finance	George
Harriet	2202	Sales			Harriet	2202	Sales	Harriet

Conditional Join (\ltimes_c):

Combines two relations similar to cross product but with a condition

Car		Boat		Car $\bowtie_{\text{CarPrice} > \text{BoatPrice}}$ Boat			
CarModel	CarPrice	BoatModel	BoatPrice	CarModel	CarPrice	BoatModel	BoatPrice
CarA	20,000	Boat1	10,000	CarA	20,000	Boat1	10,000
CarB	30,000	Boat2	40,000	CarB	30,000	Boat1	10,000
CarC	50,000	Boat3	60,000	CarC	50,000	Boat1	10,000
				CarC	50,000	Boat2	40,000

Division (\div):

Reduces a relation by performing the opposite of a cartesian product

Completed		DBProject		Completed \div DBProject	
Student	Task	Task		Student	
Fred	Database1	Database1		Fred	
Fred	Database2	Database2		Sarah	
Fred	Compiler1				
Eugene	Database1				
Eugene	Compiler1				
Sarah	Database1				
Sarah	Database2				

*Must be union compatible: 1) Same number of columns 2) Corresponding columns are of the same variable type

Relational Calculus [Examples]:

Sailors(sid, sname, rating, age), Reserves(sid, bid, date), Boats(bid, bname, color)

- Find sailors with a rating > 7
 $\{s \mid s \in \text{Sailors} \wedge s.\text{rating} > 7\}$
- Find names of sailors who've reserved a red boat
 $\{t(sname) \mid \exists s \in \text{Sailors}(t.sname = s.sname \wedge \exists r \in \text{Reserves}(r.sid = s.sid \wedge \exists b \in \text{Boats}(b.bid = r.bid \wedge b.color = \text{'red'})))\}$
- Find the names of sailors who've reserved all "Interlake" boats
 $\{t(sname) \mid \exists s \in \text{Sailors}(t.sname = s.sname \wedge \forall b \in \text{Boats}(b.bname = \text{'Interlake'} \rightarrow (\exists r \in \text{Reserves}(r.sid = s.sid \wedge b.bid = r.bid))))\}$

MySQL:

Queries:

SELECT

- DISTINCT
- T.attr, T.attr as attribute
- COUNT(*)
- MAX(T.attr), MIN(T.attr), AVG(T.attr)

FROM

- Table T
- Table2 T2

WHERE

- =, !=, >, <, >=, <=
- [NOT] IN
- [NOT] EXISTS
- ANY
- ALL
- LIKE _ (1 char), % (n chars)

GROUP BY

- attribute

HAVING

- condition on grouping

Insertions:

```
INSERT INTO table_name (field1, field2, ...fieldN)
VALUES
(value1, value2, ...valueN);
```

Deletions:

```
DELETE FROM table_name [WHERE clause]
```

Views:

```
CREATE VIEW view_name(attr1, attr2, ...) AS
SELECT []
FROM []
WHERE []
```

Procedures:

CALL Procedure(params)

```
CREATE PROCEDURE NewChirp(
    new_btag VARCHAR(30),
    content VARCHAR(255))
BEGIN
    DECLARE new_cno INT(11);
    SET new_cno = (SELECT MAX(cno)+1 FROM Chirp
        WHERE btag = new_btag);
    INSERT INTO Chirp(btag, cno, content)
    VALUES (new_btag, new_cno, content);
END
```

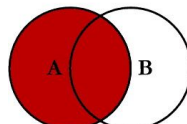
Alter Table:

```
DROP COLUMN col_name
DROP PRIMARY KEY
DROP FOREIGN KEY fk_name
```

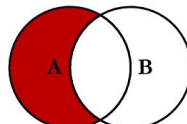
Update Table:

```
UPDATE [LOW_PRIORITY] [IGNORE] table
SET column1 = expression1,
    column2 = expression2,
    ...
[WHERE conditions]
[ORDER BY expression [ASC | DESC]]
[LIMIT number_rows]
```

SQL JOINS

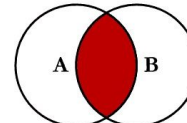


```
SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
ON A.Key = B.Key
```

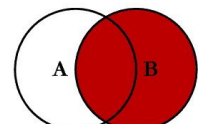
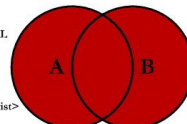


```
SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
ON A.Key = B.Key
WHERE B.Key IS NULL
```

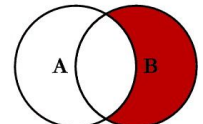
```
SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
ON A.Key = B.Key
```



```
SELECT <select_list>
FROM TableA A
INNER JOIN TableB B
ON A.Key = B.Key
```



```
SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
ON A.Key = B.Key
```



```
SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
ON A.Key = B.Key
WHERE A.Key IS NULL
```



```
SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
ON A.Key = B.Key
WHERE A.Key IS NULL
OR B.Key IS NULL
```

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