

YaSat milestone 1

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file structure

```
.
├── makefile
├── benchmarks
│   ├── *.cnf
│   ├── SAT
│   │   ├── sanity
│   │   └── tiny
│   └── UNSAT
│       ├── sanity
│       └── tiny
├── doc
│   ├── project_description.pdf
│   └── report1.md
├── include
│   ├── dp11.hpp
│   ├── parser.h
│   └── state.hpp
├── lib
│   ├── MiniSat_v1.14_linux
│   └── sat_comparer
├── scripts
│   ├── check.cpp
│   └── test.sh
└── src
    ├── parser.cpp
    └── sat.cpp
```

Usage

- `makefile` is provided.
- All the source code can be compiled and run on `linux{id}.cs.nctu.edu.tw` by simply typing `make`. The resulting binary will be generated named `yasat`.

```
$ make
g++ -Iinclude/ -std=c++17 -Ofast -Wall -Wextra -Wshadow -Wold-style-cast -c -o obj/parser.o sr
g++ -Iinclude/ -std=c++17 -Ofast -Wall -Wextra -Wshadow -Wold-style-cast -c -o obj/sat.o src/s
g++ -Iinclude/ -std=c++17 -Ofast -Wall -Wextra -Wshadow -Wold-style-cast -o yasat obj/parser.o
```

```
$ la yasat
-rwxr-xr-x 1 pwhsieh gcs 48K Apr 19 20:24 yasat
```

Feature

In this project We are going to implement a SAT solver, and these are the features I've implemented in this milestone.

1. iterative DPLL
2. two literal watching

1. iterative DPLL

Since there's going to have all kinds of branching heuristics added to this project in the future, I thought it would be a good idea to maintain the recursive stack on my own.

DPLL structure

- This is the DPLL structure declaration. It takes `number of variables` and `clauses` as the constructor argument.

```
// #Declaration
struct DPLL {
    DPLL(int _num_vars, vector<vector<int>> &_clauses)
        : num_vars(_num_vars), num_clauses(_clauses.size()), clauses(_clauses) {}
    void init();
    void watch_not_false(int&, int&, int, int);
    bool watch_is_true(int, int);
    bool unit_propagate();
    bool backtrack();
    std::optional<std::vector<int>> solve();

    int num_vars;
    int num_clauses;
    std::queue<int> prop;
    std::stack<int> branch;
    std::stack<State, std::vector<State>> call_stack;
    std::vector<std::vector<int>> &clauses;
};
```

DPLL procedure

This is the main loop of the iterative DPLL, including `unit_propagate`, `backtrack`, and decision making.

```

std::optional<std::vector<int>> DPLL::solve() {
    // 0. init
    init();

    while (!call_stack.empty()) {
        // 1. unit propagate
        if (!unit_propagate()) {
            if (!backtrack()) return std::nullopt; //UNSAT
            continue;
        }

        // 2. check if curent state is satisfied already
        if (call_stack.top().done) break;

        // 3. if not, apply new decision
        call_stack.push(call_stack.top());
        int g = call_stack.top().pick_variable();
        call_stack.top().set_variable(g, prop);
        branch.push(-g);
    }

    // 4. SAT: return ans;
    auto res = std::vector<int>(num_vars);
    auto &state = call_stack.top();
    for (int i = 1; i <= num_vars; i++) {
        if (state.var(i) and *state.var(i)) {
            res[i-1] = i;
        } else {
            res[i-1] = -i;
        }
    }
    return res;
}

```

2. two literal watching

- During **unit propagation**, we only look at the first two literal in each clause. This so-called **two literal watching** method enable us to do quick unit propagation without iterating through all the clauses.
- This method drastically decrease the complexity of the unit propagation, since we at most look at all the clauses and literal **once**.
- **prop** record all the false literal. (ex: decision on **1 == true** , imply that **1 == false** is the false literal)

```

bool DPLL::unit_propagate() {
    auto &state = call_stack.top();
    while (!prop.empty()) {
        int false_literal = prop.front(); prop.pop();
    }
}

```

```

state.done = true;
for (int i = 0; i < num_clauses; i++) {
    if (clauses[i].size() == 1) break;

    // this clause is satisfied
    auto &[la, lb] = state.watch[i];
    if (la == -1 and lb == -1) continue;

    // one of the watched literal is already true
    if (watch_is_true(la, i) or watch_is_true(lb, i)) {
        la = lb = -1;
        continue;
    }

    // update watching literals that are not false
    watch_not_false(la, lb, i, false_literal);
    watch_not_false(lb, la, i, false_literal);

    // check if this became unit clause (implication)
    if (la == -1) std::swap(la, lb);
    if (la != -1 and lb == -1) {
        int tmp = clauses[i][la];
        if (state.var(tmp) and *state.var(tmp) != (tmp > 0)) {
            prop = {};
            return false; //conflict
        }
        state.set_variable(tmp, prop);
        la = -1;
    }
    if (la != -1 or lb != -1) state.done = false;
}
}
return true;
}

```

Results

- These result are verified by my own written `checker.cpp` , which is placed under `scripts/` .
- It checks the correctness of the SAT solver by taking the `.cnf` file and `.sat` file as input, and check if there's any conflict between them.

75 variable, 325 clauses

- These are some randomly generated 3-SAT dataset.
- 5 SAT and 5 UNSAT, 10 dataset in total.
- 10 cnf test cases take about *a second*.

make test will run the test.sh script to run the benchmark

```
$ time make test
./scripts/test.sh
[SAT]
[ACCEPTED]
[SAT]
[ACCEPTED]
[SAT]
[ACCEPTED]
[SAT]
[ACCEPTED]
[SAT]
[ACCEPTED]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
[UNSAT]
make test 1.03s user 0.04s system 91% cpu 1.176 total
```

Sudoku

- Out of curiosity, I use the encoded sudoku data from the previous project. And the result is actually quite good.
- 9×9 sudoku can be solved in 0.132 seconds, and 16×16 sudoku can be solved in 1.436 seconds.

9×9

```
$ time ./yasat benchmarks/9.cnf
[SAT]
./yasat benchmarks/9.cnf 0.03s user 0.00s system 23% cpu 0.132 total
```

16×16

```
$ time ./yasat benchmarks/16.cnf
[SAT]
./yasat benchmarks/16.cnf 1.42s user 0.01s system 99% cpu 1.436 total
```

Discussion

- The material provided in class has taught us how to implement a recursive version of `DPLL` , however, for better future proof, I decided to implement it in an *iterative way*.
- This is quite hard in the beginning because I can't quite figure out the right way to iterate through all the possible outcomes of the `cnf` while writing efficient code.
- Another obstacle I've encountered during this assignment is the implementation of `two-literal watching` . The whole concept is actually bug-prone and hard to debug.
- After all these hard work, fortunately, the result is great. Much better than I would have thought. It is capable of **solving `16 x 16` sudoku in about a second**. Nice!