

### Truth table for $\rightarrow$

Assuming that the rules of Fitch are such that it is impossible to prove an argument which is not logically valid, the truth-table for  $\rightarrow$  is fixed if we accept  $\rightarrow$ Elim and  $\rightarrow$ Intro.

How do the rules of proof for  $\rightarrow$  fix its truth table?

A	B	$A \rightarrow B$
T	T	
T	F	
F	T	
F	F	

### 'If' and ' $\rightarrow$ ' do not match

$\neg A$	Marnie will not miss her train
$A \rightarrow B$	If Marnie misses her train, she will arrive on time.

### 'If' and ' $\rightarrow$ ' do match

$\neg A \vee B$	America does not exist $\vee$ Baudrillard is wrong
If A, B	If America exists, Baudrillard is wrong
If A, B	If you love logic, things will fall into place
$\neg(A \wedge \neg B)$	Not both: you take logic and things don't fall into place

### Multiple quantifiers: simple examples

"Something is above something"  
 $\exists x \exists y \text{ Above}(x,y)$

"Everyone likes puffins"  
 $\forall y \forall x (\text{Puffin}(x) \rightarrow \text{Likes}(y,x))$

$\text{Puffin}(x)$  : x is a puffin

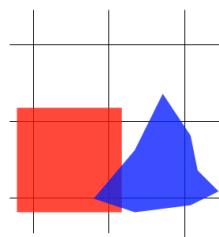
b : Steve

$\text{Likes}(x,y)$  : x likes y

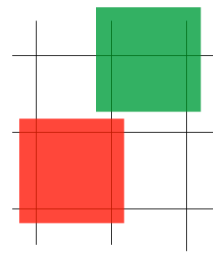
"Something makes someone want to die inside"  
 $\exists x \exists y \text{ WantToDieInside}(x,y)$

$\text{WantToDieInside}(x,y)$  : x makes y want to die inside

1  $\exists x \exists y \text{ Above}(x,y)$   
"Something is above something"



1 is FALSE



1 is TRUE

### The Soundness & Completeness Theorems

' $A \vdash B$ ' means there is a proof of B using premises A

' $\vdash B$ ' means there is a proof of B using no premises

' $A \models B$ ' means B is a logical consequence of A

' $\models B$ ' means B is a tautology

' $A \models_T B$ ' means B is a logical consequence of A just in virtue of the meanings of truth-functions (the textbook LPL calls this 'tautological consequence')

Soundness: If  $A \vdash B$  then  $A \models B$

I.e. if you can prove it in Fitch, it's valid

Completeness: If  $A \models_T B$  then  $A \vdash B$

I.e. if it's valid just in virtue of the meanings of truth functions, you can prove it in Fitch

### Proof Example using negation-Intro

1. P  
5.  $\neg\neg P$

### Substitution of logical equivalents

Two sentences are *logically equivalent* when each is a logical consequence of the other.

If \* and # are logically equivalent, then you take a complex formula and replace \* with # without changing the truth table of the complex formula.

Example. Because  $\neg( R \vee \neg P )$  is logically equivalent to  $\neg R \wedge P$ , it follows that:

$P \vee \neg( Q \wedge \neg( R \vee \neg P ) )$  is logically equivalent to  $P \vee \neg( Q \wedge ( \neg R \wedge P ) )$

And because  $\neg (Q \wedge (\neg R \wedge P))$  is logically equivalent to  $\neg Q \vee (R \vee \neg P)$ , it follows that  $P \vee \neg (Q \wedge (\neg R \wedge P))$  is logically equivalent to  $P \vee (\neg Q \vee (R \vee \neg P))$ .

### Example proof

1.  $SV(Q \wedge R)$

SVQ

## Fubar rules

 $\wedge F_{\text{ubar}}$ :
$$\begin{array}{|c} * \\ \vdots \\ * \wedge \# \end{array}$$

*Q1. What would be wrong with adding  $\Lambda F_{\text{ubar}}$  to Fitch?*

Q2. What would be wrong with having  $\wedge F_{\text{ubar}}$  in any system of proof?

## Tonk

## \*Intro

$$\begin{array}{c} P_i \\ \vdots \\ P_1 * P_2 \end{array}$$

\*Elim

$$\begin{array}{c} P_1 * P_2 \\ \vdots \\ P_i \end{array}$$

## Exercises 05

*For your sixth seminar*

### Not for fast groups

A. From the LPL textbook:

## 7.9 (truth functions)

6.17-20 (proof)

6.33, 6.40

DO NOT USE TAUT CON. EVER.

8.24–25 (proofs/counterexamples)

12.4–5, \*12.6–7 (counterexamples)

B. For each of the following sentences of FOL, give a logically equivalent sentence of idiomatic English using the specified interpretation. Your English sentences should be as concise as possible.

Domain : people and actions

D(x) : x is desirable

$V(x)$  :  $x$  is virtuous

A(x) : x is an action

$H(x)$  : x is a person

$P(x,y)$ :  $x$  performed  $y$

i.  $\forall x [ D(x) \rightarrow V(x) ]$

ii.  $\forall x [ [A(x) \wedge D(x)] \rightarrow V(x) ] ]$

iii.  $\exists x [ A(x) \wedge \neg [ D(x) \rightarrow V(x) ] ]$ 
$$*iv. \exists x \forall y [[H(x) \wedge A(y)] \wedge P(x,y)] \rightarrow V(y)]$$
$$\text{**v. } \neg \exists x [\exists y [H(x) \wedge P(x,y) \wedge A(y) \wedge \neg V(y)] \\ \wedge \neg \exists z [P(x,z) \wedge A(z) \wedge V(z)]]$$