

## How to determine the truth of sentences involving $\forall$

1. Give every object a name.
2. For each name in turn, create a new sentence like this: delete the quantifier and replace all instances of the variable it binds with that name
3. If ALL of the new sentences are true, so is the original.

## Scope

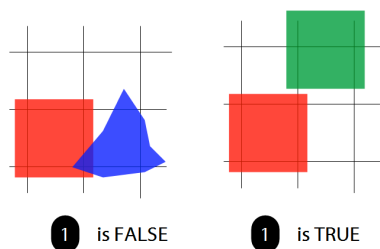
Underlining shows the scope of the quantifiers

"All squares are blue"  
 $\forall x ( \text{Square}(x) \rightarrow \text{Blue}(x) )$

"If everything is square, everything is blue"  
 $\forall x \text{ Square}(x) \rightarrow \forall x \text{ Blue}(x)$

## Multiple quantifiers

1  $\exists x \exists y \text{ Above}(x,y)$   
"Something is above something"



## Quantifiers bind variables

Read §9.2 of Barwise & Etchemendy

## The Soundness & Completeness Theorems

' $A \vdash B$ ' means there is a proof of B using premises A

' $\vdash B$ ' means there is a proof of B using no premises

' $A \models B$ ' means B is a logical consequence of A

' $\models B$ ' means B is a tautology

' $A \models_T B$ ' means B is a logical consequence of A just in virtue of the meanings of truth-functions (the textbook LPL calls this 'tautological consequence')

Soundness: If  $A \vdash B$  then  $A \models B$

I.e. if you can prove it in Fitch, it's valid

Completeness: If  $A \models B$  then  $A \vdash B$

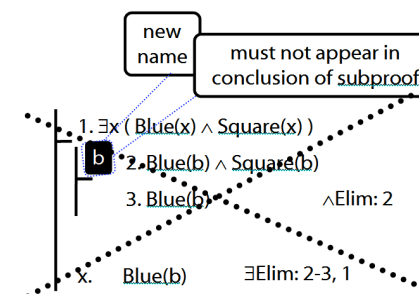
I.e. if it's valid just in virtue of the meanings

## Proof example: $\exists$ Elim

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1.  $\exists x ( \text{Blue}(x) \wedge \text{Square}(x) )$ 
  |
  | b
  | 2.  $\text{Blue}(b) \wedge \text{Square}(b)$ 
  | 3.  $\text{Blue}(b)$   $\wedge$ Elim: 2
  | 4.  $\exists x \text{ Blue}(x)$   $\exists$ Intro: 3
x.  $\exists x \text{ Blue}(x)$   $\exists$ Elim: 2-4, 1
  
```

## Invalid use of $\exists$ Elim



## Two objects are broken

$\exists x \exists y ( \text{Broken}(x) \wedge \text{Broken}(y) \wedge \neg(x=y) )$

To translate sentences involving number into FOL, use identity.

## Reflexive, symmetric and transitive

A *reflexive* relation is one that everything bears to itself.

(E.g. SameShape)

$\forall x R(x,x)$



A *symmetric* relation is one such that if x bears it to y, then y bears it to x. (E.g. Adjacent(x,y).)

$\forall x \forall y ( R(x,y) \rightarrow R(y,x) )$



A *transitive* relation is one such that if x bears it to y and y bears it to z then x bears it to z. (E.g. LeftOf is transitive; DifferentShape is *not* transitive)

$\forall x \forall y \forall z ( [ R(x,y) \wedge R(y,z) ] \rightarrow R(x,z) )$



## Exercises 06

*For your sixth seminar*

*Only for fast groups*

11.2, 11.4

11.8, 11.9, \*11.11

13.12, 13.14, 13.16 (quantifier proofs)

13.28

*Despite what book says, DON'T even THINK  
about using TAUT CON.*