

Using Ants as a Genetic Crossover Operator in GLS to Solve STSP

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Abstract— Ant Colony Algorithm (ACA) and Genetic Local Search (GLS) are two optimization algorithms that have been successfully applied to the Traveling Salesman Problem (TSP). In this paper we define new crossover operator then redefine ACA's ants as operate according to defined crossover operator then put forward our GLS that uses these ants to solve Symmetric TSP (STSP) instances.

Keywords—ACA; GLS; Heuristic crossover; Local search

I. INTRODUCTION

GLSs are metaheuristics that not only exploit Genetic Algorithm (GA) but also use profits of local search and have been successfully applied to TSP [8-9-10-11-12].

ACA is another metaheuristic that was firstly invented by three Italy scholars. This algorithm is inspired by the behavior of real ants that enables them to find the shortest paths between food sources and their nest. Many versions of ACA have been successfully applied to TSP [13-14-16-17-18].

Each of GLS and ACA has been successfully applied to TSP but each of ACA with GA family can be better to solve TSP [21]. Reference [15] embeds ACA with GA to solve TSP. In this paper we will combine ACA with GLS to solve TSP. we will use ACA's ants as crossover operator in GLS. Many heuristic crossovers have been invented [1-2-3-4-5-6-7] also here we presented a heuristic crossover. In this paper we design ant that operates as our heuristic crossover then we use it in our GLS.

This paper organized as follows. Next section introduces ACA and GLS. In section III we present our ant based GLS. In section IV we put forward our heuristic crossover and explain how our designed ant operates as our heuristic crossover. We state our local searches in section V. Section VI presents results of experiments. Section VII summarizes this paper.

II. REVIEW OF ACA AND GLS

A. ACA

ACA is inspired by the behavior of real ants that enables them to find the shortest paths between food sources and their nest. In the process of finding the shortest path the ants are guided by exploiting pheromone that each ant deposits on the ground when walking. To solve TSP, ACA executes multiple iterations. During each iteration, m ants located on different nodes beginning to build a tour in n ($=$ problem size) times. In each time each of ants selects next node according to transition rule then update pheromone locally. After each ant's tours are completed global pheromone update is

executed [13-14]. General pseudo code for ACA is shown in Fig. 1.

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1) for i=1 to Number-of-Iteration
2)   for k=1 to N(=number of city)
3)     while (tourk is incomplete)
4)       ant[k] select next city (next edge) according to transition rule
5)       update pheromone locally
6)     update pheromone globally
    
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Figure 1. ACA pseudo code

Ant Colony System (ACS) [14] is one of ACA versions that transition rule, local update and global update are defined as follow: transition rule uses (1)

$$S = \begin{cases} \arg \max_{u \in J_k(r)} \{ [\tau(r, u)] \cdot [\eta(r, u)]^\beta \}, & \text{if } q \leq q_0 \\ \text{use (2)} & , \text{ otherwise} \end{cases} \quad (1)$$

$$p_k(r, s) = \begin{cases} \frac{[\tau(r, s)] \cdot [\eta(r, s)]^\beta}{\sum_{u \in J_k(r)} [\tau(r, u)] \cdot [\eta(r, u)]^\beta} & \text{if } s \in J_k(r) \\ 0 & \text{otherwise} \end{cases} \quad (2)$$

When ant k that is located on node r selects city s as next city, pheromone local update is executed as (3)

$$\tau(r, s) \leftarrow (1 - \rho) \cdot \tau(r, s) + \rho \cdot \tau_0 \quad (3)$$

In ACS global updating rule uses best tour that is produced by one of ants in steps 3 to 4 of Fig. 1. Global updating rule is as (4)

$$\tau(r, s) \leftarrow (1 - \alpha) \cdot \tau(r, s) + \alpha \cdot (\text{best tour length})^{-1} \quad (4)$$

In (1), (2), (3) and (4) $J_k(r)$ is the set of cities that remain to be visited by ant k positioned on city r and $\eta(r, u) = 1/\text{distance}(r, u)$ and q is a random number uniformly distributed in $[0 \dots 1]$, q_0 ($0 \leq q_0 \leq 1$), α , β and ρ are parameters and τ is pheromone array.

Max-Min Ant System (MMAS) is another version of ACA that limits pheromone values between τ_{\min} and τ_{\max} [16-17-18].

B. GLS

GLS is another metaheuristic that is obtained from combination of GA and Local Search (LS) and can be implemented as Fig. 2. Demonstrated "while" loop in line 2 repeated until no better child produced. If one of produced children is better than one of population individual then "while" loop will be continued.

- 1) initialize population with a construction heuristic
- 2) **while** population is changed
- 3) **for** i=1 to Generation-Size
- 4) Select father and mother from population
- 5) child \leftarrow operate crossover on father and mother
- 6) operate local search on child
- 7) add child to population
- 8) reduce population
- 9) return best individual of population

Figure 2. General definition of GLS

[8-9-10-11-12] use GLS to solve TSP. In these papers procedure like Fig 2 is applied on initial population. In [11] 2-opt tour improvement is applied on initial population.

III. OUR ANT BASED GLS

We use ACA's ants as crossover operator in our GLS. As crossover operator each ant takes two parents and produces child. Fig 3 shows our ant based GLS.

- 1) initialize population with a construction heuristic
- 2) use Classify_based_LS to improve population individuals
- 3) **while** population is changed
- 4) **for** k=1 to Generation-Size
- 5) Select father and mother from population
- 6) **for** i=1 to N(=number of city)
- 7) child_k[i] \leftarrow ant_k go to next city
- 8) update $\tau(\text{child}_k[i-1], \text{child}_k[i])$ according to (3) locally
- 9) use our 2-opt and 3-opt move based local search on child_k
- 10) add child_k to population
- 11) reduce population
- 12) use population's best individual to update pheromone globally use (4)
- 13) return best individual of population

Figure 3. Our ant based GLS

Each ant generates child in N(=number of cities) steps. In each step each ant selects next city and updates pheromone, according to (3), locally. After N steps ants produce their

children and our local searches are applied on ant's tour and these tours are added to population and "for" loop in line 5 of Fig. 3 is terminated. If one of produced children is better than one of population individuals then "while" loop will be continued.

IV. OUR ANT

A. Our Crossover

Our crossover operator selects a random city "c" and copies it to child at first then puts three (or more) pointers on each of father and mother (one of pointer in each parent must be pointed to right neighbor of selected city) then probes witch of pointed nodes is nearest to "c" then it is copied to "c" and its pointer goes one forward. Please consider that when witch of pointers reach to beginning position of other pointers then it must be removed. For example suppose that distance array of graph will be as Fig. 4, Fig. 5 show two steps of our crossover operation.

This crossover uses pointers to operate so we call it pointer based crossover (PBX).

	1	2	3	4	5	6	7	8
1	0	12	19	31	22	17	23	12
2	12	0	15	37	21	28	35	22
3	19	15	0	50	36	35	35	21
4	31	37	50	0	20	21	37	38
5	22	21	36	20	0	25	40	33
6	17	28	35	21	25	0	16	18
7	23	35	35	37	40	16	0	14
8	12	22	21	38	33	18	14	0

Figure 4. Distance array of graph (dis)

<p>Selected node (c): 4 (is selected randomly)</p> <p>Father: 4 5 7 3 1 2 6 8</p> <p style="text-align: center;">↑ ↑ ↑</p> <p style="text-align: center;">F_pointer 1 F_pointer 2 F_pointer 3</p> <p>Mother: 3 1 7 5 6 4 2 8</p> <p style="text-align: center;">↑ ↑ ↑</p> <p style="text-align: center;">M_pointer 1 M_pointer 2 M_pointer 3</p> <p>child: 4</p> <p>dis[4,5] is shorter than dis[4,7], dis[4,1] and dis[4,6] so F_pointer1 must go one forward</p>	=>	<p>Selected node (c): 5</p> <p>Father: 4 5 7 3 1 2 6 8</p> <p style="text-align: center;">↑ ↑</p> <p style="text-align: center;">F_pointer 2 F_pointer 3</p> <p>Mother: 3 1 7 5 6 4 2 8</p> <p style="text-align: center;">↑ ↑ ↑</p> <p style="text-align: center;">M_pointer 1 M_pointer 2 M_pointer 3</p> <p>child: 4 5</p> <p>F_pointer1 reaches to F_pointer2 so must be removed</p>
<p>Selected node (c): 5 (winner node of previous step)</p> <p>Father: 4 5 7 3 1 2 6 8</p> <p style="text-align: center;">↑ ↑</p> <p style="text-align: center;">F_pointer 2 F_pointer 3</p> <p>Mother: 3 1 7 5 6 4 2 8</p> <p style="text-align: center;">↑ ↑ ↑</p> <p style="text-align: center;">M_pointer 1 M_pointer 2 M_pointer 3</p> <p>child: 4 5</p> <p>dis[5,1] is shorter than dis[5,7] and dis[5,6] so F_pointer3 must go one forward</p>	=>	<p>Selected node (c): 1</p> <p>Father: 4 5 7 3 1 2 6 8</p> <p style="text-align: center;">↑ ↑</p> <p style="text-align: center;">F_pointer 2 F_pointer 3</p> <p>Mother: 3 1 7 5 6 4 2 8</p> <p style="text-align: center;">↑ ↑ ↑</p> <p style="text-align: center;">M_pointer 1 M_pointer 2 M_pointer 3</p> <p>child: 4 5 2</p>

Figure 5. Two first steps of our crossover

B. Our Ant

Our ant is same to pointed crossover; it takes father and mother and in each step adds one city to child according to transition rule that is defined in (5)

$$S = \begin{cases} \arg \max_{u \in PC} \{[\tau(r, u)] \cdot [\eta(r, u)]^\beta\}, & \text{if } q \leq q_0 \\ \text{use (6)} & , \text{ otherwise} \end{cases} \quad (5)$$

$$p_k(r, s) = \begin{cases} \frac{[\tau(r, s)] \cdot [\eta(r, s)]^\beta}{\sum_{u \in PC} [\tau(r, u)] \cdot [\eta(r, u)]^\beta} & \text{if } s \in PC \\ 0 & \text{otherwise} \end{cases} \quad (6)$$

PC is union of sets of pointed city in father and mother.

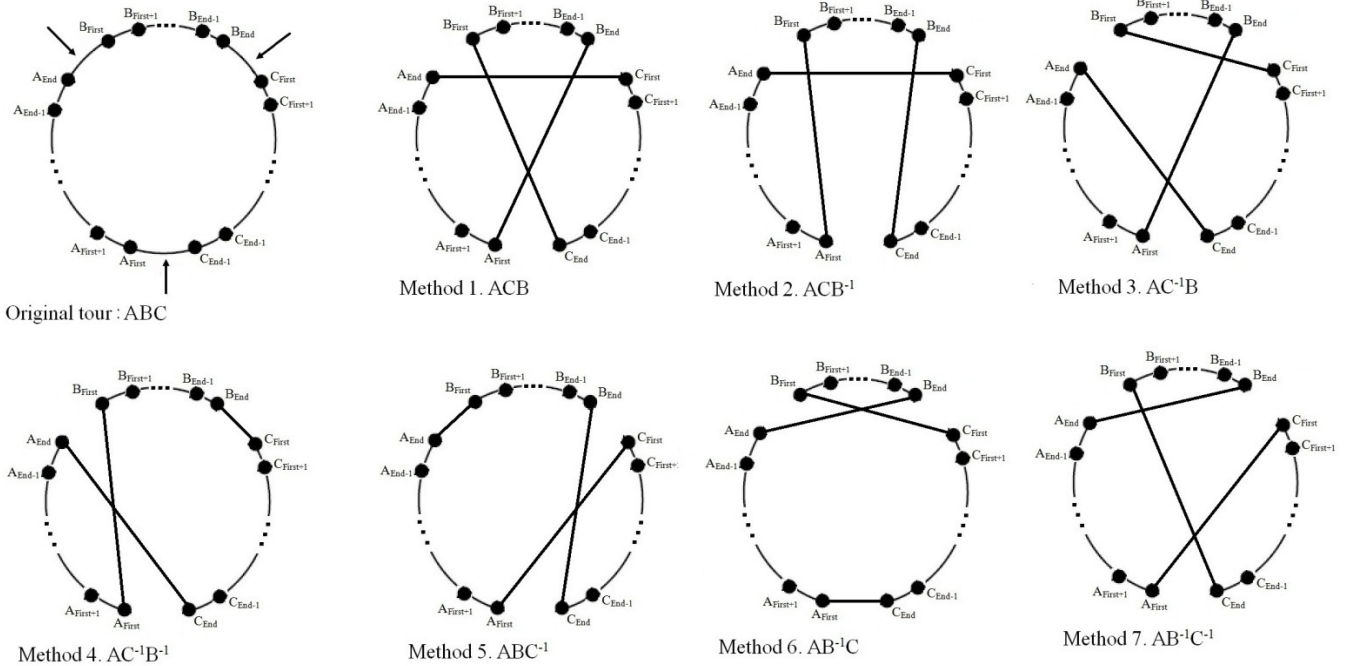
V. OUR LOCAL SEARCHES

A. 2opt_move_based_LS

2-opt move deletes two edges and reconnects two subpaths in other way [19] (see Fig. 6). In Fig. 5 can be easily seen that 2opt move selects a subpath from end and reverses it and the cost of obtained tour can be calculated with (1):

$$\begin{aligned} \text{Cost}_{\text{Obtained tour}} = & \text{Cost}_{\text{Original tour}} \\ & - \text{Distance}(A_{\text{End}}, B_{\text{First}}) - \text{Distance}(B_{\text{End}}, A_{\text{First}}) \\ & + \text{Distance}(A_{\text{End}}, B_{\text{End}}) + \text{Distance}(B_{\text{End}}, A_{\text{First}}) \end{aligned} \quad (7)$$

2opt_move_based_LS uses 2-opt-move showed in Fig. 6. 2-opt-move first selects one point at random and considers the subpath from selected point to end of tour then calculates



Predicted_Cost with (7), if Predicted_Cost is be small than tour's cost then subpath is inverted. 2opt_move_based_LS repeat 2-opt-move until Predicted_Cost is been greater than tour cost.

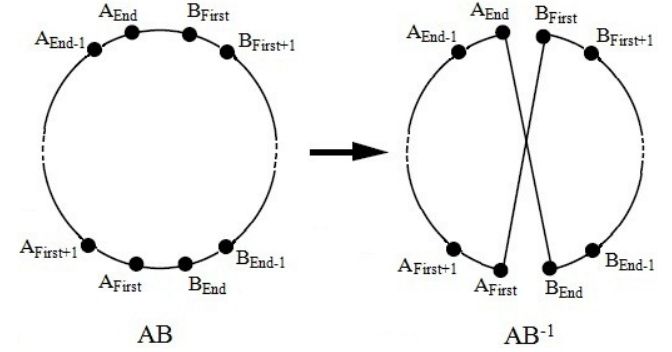


Figure 6. 2-opt move: we use X, X^{-1} replace of $(X_{\text{First}} X_{\text{First}+1} \dots X_{\text{End}-1} X_{\text{End}})$ and $(X_{\text{End}} X_{\text{End}-1} \dots X_{\text{First}+1} X_{\text{First}})$

B. 3opt_move_based_LS

3-opt move deletes three edges and reconnects three subpaths in other ways [19] (see Fig. 7). The cost of obtained tours can be calculated with (8):

$$\begin{aligned} \text{Cost}_{\text{Obtained tour}} = & \text{Cost}_{\text{Original tour}} \\ & - \text{Deleted edges weight} \\ & + \text{Added edges weight} \end{aligned} \quad (8)$$

For example in method 1 of Fig. 7 deleted and added edges sets are $\{(A_{\text{End}}, B_{\text{First}}), (B_{\text{End}}, C_{\text{First}}), (C_{\text{End}}, A_{\text{First}})\}$ and $\{(A_{\text{End}}, C_{\text{First}}), (C_{\text{End}}, B_{\text{First}}), (B_{\text{End}}, A_{\text{First}})\}$ respectively.

Figure 7. 3-opt move: seven distinct tour are obtained from original tour. We use X, X^{-1} replace of $(X_{\text{First}} X_{\text{First}+1} \dots X_{\text{End}-1} X_{\text{End}})$ and $(X_{\text{End}} X_{\text{End}-1} \dots X_{\text{First}+1} X_{\text{First}})$

$\alpha=.9$, $\beta=2$, $\rho=.1$ $q_0=.9$. TABLE II shows results of this experiment.

TABLE II. EXPERIMENTAL RESULTS OF ANT BASED GLS FOR SOME TSPLIB INSTANCES, GENERATION SIZE=500, POPULATION SIZE=50

	Best length (quality)	Average length (quality)	Worst length (quality)	Average Time (second)
eil51	427 (0.23%)	428.7 (0.63%)	433 (1.64%)	10.25
eil76	543 (0.93%)	548.07 (1.87%)	558 (3.72%)	16.13
kroA100	21331 (0.23%)	21903.67 (2.92%)	23106 (8.57%)	28.52
A280	2609 (1.16%)	2703.23 (4.82%)	2832 (9.81%)	133.99

In this table “Best length”, “Average length” and “Worst length” show the best, average, and worst tour lengths of runs, respectively. “Average Time” column gives the average running time in seconds. In “Best length”, “Average length” and “Worst length” columns the values in parentheses is result of calculating

$$\frac{\text{cost of solution found} - \text{known optimum cost}}{\text{known optimum cost}} \times 100$$

VII. CONCLUSION

In this paper we presented GLS that was using ants as crossover operator. These ants were operating as our presented heuristic crossover. Our heuristic crossover was using several pointers to operate so we called it pointer based crossover (PBX). We also presented Classify_based local search to improve initial population individual. Experimental results have shown that Classify_based local search can decrease cost of random tours up to 65%.

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