

Dependent Types to Push Corners of the Property-based Testing

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I hope you'll agree that

- Dependent types are cool

```
data AtIndex : Fin n → Vect n a → a → Type where
  Here  : AtIndex FZ (x :: xs) x
  There : AtIndex idx xs y → AtIndex (FS idx) (x::xs) y
```

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  There : AtIndex idx xs y → AtIndex (FS idx) (x::xs) y
```

- Property-based testing is cool too

```
indexInserted : Property
indexInserted = property $ do
  (n, x)      <- forAll [| (nat (linear 0 100), anyInt64) |]
  (xs, idx)   <- forAll [| (vect n anyInt64, fin (linearFin n)) |]
  f           <- apply <$> forAll (function anyInt32)
  insertAt idx (f x) (map f xs) == map f (insertAt idx x xs)
```

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- We get good things
 - less incorrect implementations
 - more confidence in the code
 - less unneeded code/checks
 - help from compilers and tools
 - ...

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 - help from compilers and tools
 - ...
 - better and more powerful tests!

Type-driven property-based testing?

Being given two sorted lists, function merge produces a sorted list

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- Special generator of sorted lists for `List`?

```
sortedList : Gen (List Nat)
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```
sortedList = anyList <&> foldr (\x, xs => x :: map (x+) xs) []
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- Special wrapper for sorted lists?

```
data SortedList = SL (List Nat)
sortedList : Gen SortedList
```

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- Special wrapper for sorted lists?

```
data SortedList = SL (List Nat)
sortedList : Gen SortedList
```

- Describe intent in type, derive generator!

Type-driven property-based testing!

```
mutual
data SortedList : Type where
  Nil  : SortedList
  (::)  : (x : Nat) → (xs : SortedList) →
        (0 _ : So $ canPrepend x xs) ⇒ SortedList
canPrepend : Nat → SortedList → Bool
canPrepend n = \case [] ⇒ True; x::xs ⇒ n < x
anySortedList : Gen SortedList -- derived complete generator
```

Type-driven property-based testing!

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canPrepend : Nat → SortedList → Bool
canPrepend n = \case [] ⇒ True; x::xs ⇒ n < x

anySortedList : Gen SortedList -- derived complete generator

toList : SortedList → List Nat

mergeSorted : Property
mergeSorted = property $ do
  (xs, ys) <- forAll [| (anySortedList, anySortedList) |]
  assert $ sorted (merge (toList xs) (toList ys))
```

Type-driven property-based testing!

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- Completeness and good distribution by library

Type-driven property-based testing!

- Surely, you can still make a mistake
- ...but in a declarative specification
- Completeness and good distribution by library
- ...but fine tuning is hard

Let's try

- New language implementation
- Restricted dialect of TypeScript, static strict typing
- Interpreter + JIT; AOT compiler

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- Restricted dialect of TypeScript, static strict typing
- Interpreter + JIT; AOT compiler
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 - semantically correct programs should be successfully compilable
 - among them, halting programs should run and be interpretable without unexpected crashes

Let's try

- New language implementation
- Restricted dialect of TypeScript, static strict typing
- Interpreter + JIT; AOT compiler
- Properties
 - semantically correct programs should be successfully compilable
 - among them, halting programs should run and be interpretable without unexpected crashes
 - all these running modes should produce the same result

How specification can look like

```
data Stmts : (functions : List (Name, FunSig)) →  
              (varsBefore : List (Name, Type)) →  
              (varsAfter  : List (Name, Type)) → Type where
```

```
(.) : (ty : Type) → (n : Name) →  
      Stmts funs vars ((n, ty)::vars)
```

```
(#=) : (n : Name) → (0 lk : n `IsIn` vars) ⇒  
       (v : Expr funs vars (found lk)) →  
       Stmts funs vars vars
```

```
If : (cond : Expr funs vars Bool) →  
      Stmts funs vars vThen → Stmts funs vars vElse →  
      Stmts funs vars vars
```

```
(>>) : Stmts funs preV midV → Stmts funs midV postV →  
       Stmts funs preV postV
```

How specification can look like

```
record FunSig where
  constructor (⇒)
  From : List Type
  To    : Type
```

```
data Expr : List (Name, FunSig) → List (Name, Type) →
           Type → Type where
```

```
C : (x : ty) → Expr funs vars ty
```

```
V : (n : Name) → (0 lk : n `IsIn` vars) ⇒
  Expr funs vars (found lk)
```

```
F : (n : Name) → (0 lk : n `IsIn` funs) ⇒
  All (Expr funs vars) (found lk).From →
  Expr funs vars (found lk).To
```

Semantically correct program

```
StdF : List (Name, FunSig)
StdF = [ ("+" , [Int, Int]  $\Rightarrow$  Int)
        ,("<" , [Int, Int]  $\Rightarrow$  Bool)
        ,("++", [Int]  $\Rightarrow$  Int)
        ,("||", [Bool, Bool]  $\Rightarrow$  Bool) ]

program : Stmts StdF [] ?
program = do
  Int. "x"
  "x"  $\#$ = C 5
  Int. "y"; Bool. "res"
  "y"  $\#$ = F "+" [V "x", C 1]
  If (F "<" [F "++" [V "x"], V "y"])
    (do "y"  $\#$ = C 0; "res"  $\#$ = C False)
    (do Int. "z"; "z"  $\#$ = F "+" [V "x", V "y"]
        Bool. "b"; "b"  $\#$ = F "<" [V "x", C 5]
        "res"  $\#$ = F "||" [V "b", F "<" [V "z", C 6]])
```


Semantically incorrect programs

```
failing "Mismatch between: Int and Bool"
```

```
bad : Stmts StdF [] ?
```

```
bad = do
```

```
  Int. "x"; "x" #= C 5
```

```
  Bool. "y"; "y" #= F "+" [V "x", C 1]
```

```
failing "Mismatch between: [] and [Int]"
```

```
bad : Stmts StdF [] ?
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```
bad = do
```

```
  Int. "x"; "x" #= C 5
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```
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```

```
failing "Mismatch between: Bool and Int"
```

```
bad : Stmts StdF [] ?
```

```
bad = do
```

```
  Int. "x"; "x" #= C 5
```

```
  Int. "y"; "y" #= F "+" [C True, V "x"]
```

Actual specification

- Using DepTyCheck library

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- Language subset
 - Halting programs
 - Loops, ifs, assignments, exceptions
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 - De Bruijn indices
 - “Continuation-passing style”-like data
 - Specialised polymorphic data types
 - Grouping type indices
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- Language subset
 - Halting programs
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 - Classes without arbitrary methods, numbers, arrays
- Specification “hacks”
 - De Bruijn indices
 - “Continuation-passing style”-like data
 - Specialised polymorphic data types
 - Grouping type indices
 - ...
- ~330 LOC of specification + same for harness
- Partially derived, partially hand-written generators

Actual specification

Testing...

Actual specification

```
ASSERTION FAILED: false
IN /.../optimizations/lse.cpp:851: GetEliminationCode
Backtrace [tid=2990599]:
#0 : 0x7f876a1776c0 PrintStack(std::ostream&)
#1 : 0x7f876a177562 debug::AssertionFail(...)
#2 : 0x7f8767185e10 compiler::Lse::GetEliminationCode(...)
#3 : 0x7f8767186a20 compiler::Lse::DeleteInstructions(...)
...
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...
```

Shrinking...

Actual specification

```
class C0 {  
  x0: boolean  
}  
  
function main() : void {  
  let x1: C0 = {x0: true}  
  while(x1.x0) {  
    x1.x0 = x1.x0  
    x1.x0 = false  
  }  
}
```

Actual specification

Testing...

Actual specification

```
TypeError: Unreachable statement. [<filename>:26:34]
```

Actual specification

```
TypeError: Unreachable statement. [<filename>:26:34]
```

Shrinking...

Actual specification

TypeError: Unreachable statement. [<filename>:3:30]

```
function main() : void {  
  let x1: Int = 1  
  while(([false, true])[x1]) {  
  }  
}
```

Actual specification

Testing...

Actual specification

```
ASSERTION FAILED: block->GetGraph() == GetGraph()
IN /.../optimizer/ir/graph_cloner.h:176: GetClone
Backtrace [tid=2902033]:
#0 : 0x7fe71892b820 PrintStack(std::ostream&)
#1 : 0x7fe71892b6c2 debug::AssertionFail(...)
#2 : 0x7fe71a61ae61 compiler::GraphCloner::GetClone(...)
#3 : 0x7fe71a61162a compiler::GraphCloner::CopyLoop(...)
#4 : 0x7fe71a611839 compiler::GraphCloner::CopyLoop(...)
#5 : 0x7fe71a611173 compiler::GraphCloner::CloneAnalyses(...)
#6 : 0x7fe71a610d1f compiler::GraphCloner::CloneGraph()
#7 : 0x7fe71a5b377c compiler::GraphChecker::GraphChecker(...)
...
```

Actual specification

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ASSERTION FAILED: block->GetGraph() == GetGraph()
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#7 : 0x7fe71a5b377c compiler::GraphChecker::GraphChecker(...)
...
```

Shrinking...

Actual specification

```
class C0 {  
  x0: boolean  
  
  f() : string {  
    return ""  
  }  
}
```

```
function main() : void {  
  let x2: C0 = {x0: true}  
  let fuel0 = 1  
  while(fuel0 > 0) {  
    do {  
      fuel0--  
      do {  
        fuel0--  
        let s = x2.f()  
      } while(true && (fuel0 > 0))  
    } while(true && (fuel0 > 0))  
  }  
}
```

Actual specification

- ...and so on
- 9 bugs in typechecker, JIT- and AOT-optimisers, code generator
- 8 more bugs during specification

What we've (re)learned

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- Same property can find many different bugs
- Subtle bugs are in interaction of many features
- PBT can find bug that are tricky to find manually
- PBT is cool!
- Dependent types may be cool for PBT!

What are generators anyway?

```
data Fin : Nat → Type
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```
finN      : (n : Nat) → Gen (Fin n)
```

```
finAny    : Gen (n ** Fin n)
```

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data LT : Nat → Nat → Type
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```
data Fin : Nat → Type
```

What's a signature for generator of `Fin`?

```
finN    : (n : Nat) → Gen (Fin n)
```

```
finAny  : Gen (n ** Fin n)
```

```
data LT : Nat → Nat → Type
```

```
ltAny    : Gen (1 ** r ** LT 1 r)
```

```
ltLeft   : (r : Nat) → Gen (1 ** LT 1 r)
```

```
ltRight  : (1 : Nat) → Gen (r ** LT 1 r)
```

```
ltGiven  : (1, r : Nat) → Gen (LT 1 r)
```

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- We can't afford saying “no”, recall

`finN : (n : Nat) → Gen (Fin n)`

Can generators be empty?

- QuickCheck says “yes”, limited retry
- Idris Hedgehog says “no”, disallowing even filtering
- We can't afford saying “no”, recall
`finN : (n : Nat) → Gen (Fin n)`
- Generation spends most of the time in retrying

Compile-time control of emptiness

```
data Emptiness = NonEmpty | MaybeEmptyDeep | MaybeEmpty
data Gen : Emptiness → Type → Type
```

Compile-time control of emptiness

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```

```
finN    : (n : Nat) → Gen MaybeEmpty (Fin n)
finSN   : (n : Nat) → Gen NonEmpty (Fin (S n))
finAny  : Gen NonEmpty (n ** Fin n)
```

Compile-time control of emptiness

```
data Gen : Emptiness → Type → Type where
  --- ... other stuff ... ---
  OneOf : alem `NoWeaker` em ⇒ NotImmediatelyEmpty alem ⇒
    GenAlternatives True alem a → Gen em a
```

Compile-time control of emptiness

```
data GenAlternatives : (notEmpty : Bool) → Emptiness → Type → Type
```

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  oneOf : alem `NoWeaker` em ⇒ AltsNonEmpty altsNe em ⇒
    GenAlternatives altsNe alem a → Gen em a
```


Compile-time control of emptiness

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data GenAlternatives : (notEmpty : Bool) → Emptiness → Type → Type
data AltsNonEmpty : Bool → Emptiness → Type
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oneOf : alem `NoWeaker` em ⇒ AltsNonEmpty altsNe em ⇒
  GenAlternatives altsNe alem a → Gen em a
```

```
g1, g2 : Gen MaybeEmpty Nat
g1 = oneOf [empty, elements [0, 1, 2], empty]
g2 = elements [0, 1, 2]
```

Intro
○○○○○

Practicality
○○○○○○○○○○○

Adventures of generators
○○○○

Adventures of derivation
●○○

The end
○

Practical derivation

Practical derivation

- Complete
- Total

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- Complete
- Total
- Good distribution

Practical derivation

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```
fins : Fuel → (n : Nat) → Gen MaybeEmpty (List (Fin n))
```

Practical derivation

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```
fins : Fuel → (n : Nat) → Gen MaybeEmpty (List (Fin n))
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- Whole-type derivation

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fins : Fuel → (n : Nat) → Gen MaybeEmpty (List (Fin n))
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- Whole-type derivation

```
program : Fuel → (fs : _) → (vs : _) →  
              Gen MaybeEmpty (vs' ** Stmts fs vs vs')
```


Practical derivation

- Complete
- Total
- Good distribution

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fins : Fuel → (n : Nat) → Gen MaybeEmpty (List (Fin n))
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- Whole-type derivation

```
program : Fuel → (fs : _) → (vs : _) →  
              Gen MaybeEmpty (vs' ** Stmts fs vs vs')
```

- Efficient

Challenges

```
data Z : Type where
  Zxy : X n m → Y n k → Z
  Zyx : Y n m → X n k → Z
```

Challenges

```
data Z : Type where
  Zxy : X n m → Y n k → Z
  Zyx : Y n m → X n k → Z

zs : Fuel → Gen MaybeEmpty Z
zs = deriveGen
```

Challenges

```

genX'' : Gen0 (n ** m ** X n m)
genY'  : (n:_) → Gen0 (k ** Y n k)
genZxy : Gen0 Z
genZxy = do
  (n ** m ** x) <- genX''
  (k ** y) <- genY' n
  pure (Zxy x y)

```

-- or --

```

data Z : Type where
  Zxy : X n m → Y n k → Z
  Zyx : Y n m → X n k → Z

zs : Fuel → Gen MaybeEmpty Z
zs = deriveGen

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```

genX' : (n:_) → Gen0 (m ** X n m)
genY'' : Gen0 (n ** k ** Y n k)
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  (m ** x) <- genX' n
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```

Challenges

```
data Y : Nat → Nat → Type where
  MkY : (n, m : _) → Y n m
```

```
data Z : Type where
  Zxy : X n m → Y n k → Z
  Zyx : Y n m → X n k → Z
```

```
zs : Fuel → Gen MaybeEmpty Z
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genX'' : Gen0 (n ** m ** X n m)
genY'  : (n : _) → Gen0 (k ** Y n k)
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-- or --

```
genX' : (n : _) → Gen0 (m ** X n m)
genY'' : Gen0 (n ** k ** Y n k)
genZxy' : Gen0 Z
genZxy' = do
  (n ** k ** y) <- genY''
  (m ** x) <- genX' n
  pure (Zxy x y)
```

Challenges

```
data X : Nat → Nat → Type where
  X45 : X 4 5
  X5m : X 5 m
```

```
data Y : Nat → Nat → Type where
  MkY : (n, m : _) → Y n m
```

```
data Z : Type where
  Zxy : X n m → Y n k → Z
  Zyx : Y n m → X n k → Z
```

```
zs : Fuel → Gen MaybeEmpty Z
zs = deriveGen
```

```
genX'' : Gen0 (n ** m ** X n m)
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genY'' : Gen0 (n ** k ** Y n k)
genZxy' : Gen0 Z
genZxy' = do
  (n ** k ** y) <- genY''
  (m ** x) <- genX' n
  pure (Zxy x y)
```

Challenges

```
f : Nat → Nat
```

```
data Yn : (n : Nat) → Fin n → Type
```

```
data Z : Type where
```

```
  Z1 : (x : Fin (f n)) → Yn (f n) x → Z  -- n, (x ** y)
```

Challenges

$f : \text{Nat} \rightarrow \text{Nat}$

`data Yn : (n : Nat) → Fin n → Type`

`data Yf : (n : Nat) → Fin (f n) → Type`

`data Z : Type where`

`Z1 : (x : Fin (f n)) → Yn (f n) x → Z -- n, (x ** y)`
`Z2 : (x : Fin (f n)) → Yf n x → Z -- (n ** x ** y)`

Challenges

```
f : Nat → Nat
g : Fin n → Fin n
```

```
data Yn : (n : Nat) → Fin n → Type
data Yf : (n : Nat) → Fin (f n) → Type
```

```
data Z : Type where
  Z1 : (x : Fin (f n)) → Yn (f n) x → Z    -- n, (x ** y)
  Z2 : (x : Fin (f n)) → Yf n x → Z          -- (n ** x ** y)
  Z3 : (x : Fin (f n)) → Yf n (g x) → Z      -- n, x, y
```

Challenges

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f : Nat → Nat
g : Fin n → Fin n
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data Yf : (n : Nat) → Fin (f n) → Type
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data Z : Type where
  Z1 : (x : Fin (f n)) → Yn (f n) x → Z      -- n, (x ** y)
  Z2 : (x : Fin (f n)) → Yf n x → Z           -- (n ** x ** y)
  Z3 : (x : Fin (f n)) → Yf n (g x) → Z       -- n, x, y
  Z4 : (x : Fin n) → Yn n (g x) → Z          -- (n ** x), y
```

If you're interested

This presentation



DepTyCheck, examples



Code from the slides



Thank you!

This presentation



Code from the slides



DepTyCheck, examples



Questions?