big data Lecture 3: map reduce

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1 motivation: text indexing

- if you have N documents and want to make an index mapping all words to the document it appeared in on a single machine it would take $\Omega(N)$ where N is the number of documents (that is you would at least have to look at all documents once and do something with them)
- but this problem is parallel, can just look at all documents on there own and combine results
- a parallel implementation with M computers would be $\Omega(N/M)$
- so that is a good idea but need a way to distribute the work and collect results that is where map reduce comes in

map reduce

- distributed programs are hard to write well, so if we restrict how we program it becomes easier to get parallelism
- so we are again getting power by restricting what we can do
- map and reduce are common operations in functional programming
- a map function looks like like map(function f, values[$x_1 \cdots x_n$]) \rightarrow [$f(x_1) \cdots f(x_n)$] so that is it takes a function and applied it to some list of values and outputs a list with that function applied to all values
- a reduce function works like reduce(function g, values[$x_1 \cdots x_n$]) $\rightarrow g(x_1, reduce(g[x_2 \cdots x_n]))$
- a reduce function takes maps a list to a value. it does this by recursively applying g to pairs of items

Define functions "sum" and "square"

o sum: x, y \rightarrow x + y sum: [] \rightarrow 0

o square: $X \rightarrow X * X$

- reduce(sum, map(square, [x₁, x₂, ..., x_n]))
- the above example, explains a simple example. the map functions maps all elements in a list to its squared values
- the reduce function holds a variable called sum that is initialized at zero at recursively calls sum on that variable and the next list element

working with map reduce

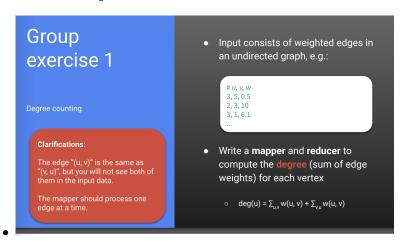
- the programer must write a mapper and reducer function. the more simple the better
- the mapper consumers key value pairs as input and outputs intermediate key value pairs
- the reducer consumes a single key and list of values and produces values for each key. note that in map reduce unlike in functional programming the reducer is not applied recursively

workflow

- map phase
 - distribute data to mappers
 - run mappers on chunks of data to get intermediate key value pairs
- sort shuffle phase
 - assign a parton of the intermediate results to each reducer by key
 - move data from mappers to reducers
- reduce phase
 - run the reducer call on chunk keys
 - collect the output

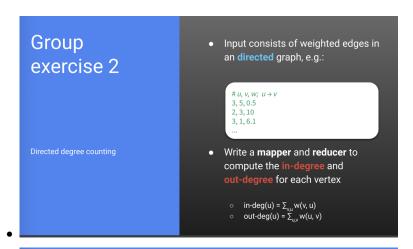
- make the mapper simple
- let the mr frame work route the intermediate results
- keep the reducer simple if possible

slido example





- this one is pretty straight forward
- \bullet tips for map reduce
 - 1. do not use floating point keys they don't hash well
 - 2. keep map and reduce simple (the sorting is what is optimized so let that do as much work as possible)
 - 3. compare your algorithm to a simple implementation



Solution: Directed degree counting def mapper(u, v, w): def reducer(key, weights): emit (u, 'out'), w # u, direction ← key emit sum(weights) The logic is nearly the same as in problem 1 We can pack information into the intermediate keys to keep things simple! Sort/shuffle phase does the hard work for us.

• the logic is the same as the last problem the real point is think about changing your keys before making your map program more complex

map reduce in practice

- can intermediate outputs be randomly assigned to reducers
- no we need to make sure all intermediate outputs with the same key go to the same reducer
- item this leads to key skew
- so suppose there is are 4 keys and 4 reducers but one key shows up 95% of the time, then that one reduce worker is doing 95% of the work this is called key skew or data skew

combiners

• key skew causes high latency

- lots of keys also means a lot of communication
- we can make the reducer's job easier with combiners which reduce intermediate key value pairs within the mapper worker, before shuffeling the data

Heuristics for using map reduce well

- have fewer mappers than inputs
- having fewer reducers than intermediate keys
- combiners can help but sometimes a more complex map is better
- sometimes doing some other sorting stuff can reduce communication

criticism of map reduce

- map reduce has criticism because it is
 - 1. too low level
 - 2. not well implementation
 - 3. not novel
 - 4. missing key DBMS features
 - 5. and not compatible with DBMS tools
- i mean i read the paper 1 4 and 5 are for sure, and i would say probably 2 (due to key skew)

too low level

- map reduce does not have schemas, so a programer has to infer them as they work, and there is no garuntee of data consistency
- map reduce does not use high level access language like sql which people think is better
- there are tools that can be used in tandom with map reduce that can address this

poor implementation

- map reduce does not have the ability to index
- is this a major issue? depends on the context

not novel

• map reduce uses old ideas so was not novel but that does not really matter

missing features

- many features from DBMS are not in mpa reduce
- why are they missing? is this a fair comparison? I am not sure

lack of DBMS compatability

 since this was published this has shifted a lot so not really an issue any more

why was map reduce so sucessfull

- map and reduce are simple abstractions that are powerful
- many jobs are only one shot (ie only need to be done once) so it may not be worth the time to build a database infrastructure

map reduce is not a database managment system

• but there is some overlap, simple queries can be done in map reduce but not hard queries

map reduce is not a general comutation engine

- maps are reducers are really flexible so a lot can be done with map reduce
- it can not however work with iterative or recursive algorithms
- not can it work with non deterministic functions

no transactions

- data base management systems use transactions to make sure data is consistency
- map reduce does not have transactions but still avoids these issues by having data be immutable, and require programs to be deterministic
- it has error tolerence if a node fails we can re assign idle workers to do its task

real issues with map reduce

- key skew so there are latency and scheduling issues
- we store intermediate results in many files
- not all tasks fit neatly into map reduce including non deterministic tasks, iterative algorithms or recursive algorithms
- cant do visualizations

what is the role of map reduce today

- map reduce is good for large batch jobs that run once or infrequently
- like data transformations or feature extaction

why is map reduce studied

- it is important for the development of big data tools
- map and reduce functions are a good way to break down problems
- the hadoop eco system is much bigger than map reduce
- there are still legacy code bases run in map reduce