

Generalizing the inverse FFT off the unit circle

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1 Introduction

I would be discussing about inverse chirp Z-transform (ICZT) algorithm which is a generalization of inverse fast Fourier transform (IFFT) off the unit circle in the complex plane.

1.1 What is chirp Z-transform?

The chirp Z-transform(CZT) is the generalization of the discrete Fourier transform (DFT) or Fast Fourier transform (FFT). The differences between them are :-

- DFT samples the Z plane at uniformly-spaced points along the unit circle.
- CZT samples the Z plane not only from the unit circle but from any point on the Z-plane.
- The DFT transforms a sequence of N complex numbers $\{x_n\}$: $=x_0, x_1, x_2, \dots, x_{N-1}$ into another sequence given by $\{X_n\}$: $=X_0, X_1, X_2, \dots, X_{N-1}$, which is defined by:-

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-\frac{i2\pi}{N} kn}$$

for $k = 0, 1, 2, 3, 4, \dots, N-1$.

- The CZT transforms a sequence of N complex numbers $\{x_n\}$: $=x_0, x_1, x_2, \dots, x_{N-1}$ into another sequence given by $\{X_n\}$: $=X_0, X_1, X_2, \dots, X_{M-1}$, which is defined by:-

$$X(k) = \sum_{n=0}^{N-1} x(n) z_k^{-n}$$

$$z_k = A.W^{-k}, \text{ for } k = 0, 1, 2, 3, 4, \dots, M-1.$$

where A is the complex starting point, W is the complex ratio between points, and M is the number of points to calculate.

CZT distributes the samples along a logarithmic spiral contour (called chirp contour) that is defined by the formula $A^{-j}.W^{jk}$ where j denotes a zero-based input sample index and k denotes a zero-based output sample index.

An efficient algorithm for computing CZT has the complexity $O(n \log(n))$ which is same as the complexity of FFT.

1.2 What is inverse chirp Z-transform?

The ICZT is the inverse of the CZT. It is the generalised version of IDFT/IFFT, just like CZT is the generalised version of DFT/FFT. The ICZT maps the output of the CZT back to the input.

The IDFT transforms the sequence $\{X_n\} = x_0, x_1, x_2, \dots, x_{N-1}$ back to the sequence $\{x_n\} = x_0, x_1, x_2, \dots, x_{N-1}$ by the formula:-

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) e^{\frac{i2\pi}{N} kn}$$

Similarly, ICZT transforms the sequence $\{X_n\} = X_0, X_1, X_2, \dots, X_{M-1}$ back to the sequence $\{x_n\} = x_0, x_1, x_2, \dots, x_{N-1}$.

2 Algorithms for calculating ICZT and CZT

As stated before, CZT is defined as:

$$X(k) = \sum_{n=0}^{N-1} X(n) A^{-n} W^{nk}$$

Let $A = \text{diag}(A^0, A^{-1}, A^{-2}, A^{-3}, \dots, A^{N-1})$

The CZT can be written as $\mathbf{X} = \mathbf{W} \mathbf{A} \mathbf{x}$ where \mathbf{W} is a Vandermonde matrix defined as:-

$$\mathbf{W} = \begin{bmatrix} W^{0.0} & W^{0.1} & \dots & W^{(N-1).0} \\ W^{0.1} & W^{1.1} & \dots & W^{(N-1).1} \\ \vdots & \vdots & \dots & \vdots \\ W^{0.(M-1)} & W^{1.(M-1)} & \dots & W^{(N-1).(M-1)} \end{bmatrix}$$

Note, we can express \mathbf{W} as a product of a diagonal matrix, a Toeplitz matrix and another diagonal matrix. Using Bluestein's substitution, we can write $nk = \frac{n^2 + k^2 - (n-k)^2}{2}$

$$\begin{aligned} \text{Thus, } X(k) &= \sum_{n=0}^{N-1} X(n) A^{-n} W^{\frac{n^2}{2}} W^{\frac{k^2}{2}} W^{-\frac{(n-k)^2}{2}} \\ &= W^{\frac{k^2}{2}} \sum_{n=0}^{N-1} X(n) A^{-n} W^{\frac{n^2}{2}} W^{-\frac{(n-k)^2}{2}} \end{aligned}$$

We can write

$$\begin{aligned} \mathbf{P} &= \text{diag}(W^{\frac{0^2}{2}}, W^{\frac{1^2}{2}}, W^{\frac{2^2}{2}}, \dots, W^{\frac{(M-1)^2}{2}}) \\ \mathbf{Q} &= \text{diag}(W^{\frac{0^2}{2}}, W^{\frac{1^2}{2}}, W^{\frac{2^2}{2}}, \dots, W^{\frac{(N-1)^2}{2}}) \end{aligned}$$

and let \mathbf{V} be a M by N Toeplitz matrix with element in i th column and j th row as $W^{-\frac{(i-j)^2}{2}}$

As a result, we can write $\mathbf{X} = \mathbf{P} \mathbf{V} \mathbf{Q} \mathbf{A} \mathbf{x}$.

To compute the product of $\mathbf{P}, \mathbf{Q}, \mathbf{A}$, we require $O(n)$ time as they all are diagonal matrices. And the product of a Toeplitz matrix with a vector can be calculated in $O(n \log(n))$ time. Thus, the overall complexity for calculating CZT is $O(n \log(n))$.

We can derive the formula for ICZT using the above expression if $M = N$.

$$\mathbf{x} = \mathbf{A}^{-1} \mathbf{Q}^{-1} \mathbf{V}^{-1} \mathbf{P}^{-1} \mathbf{X}$$

The time complexity of inverting all the diagonal matrices is $O(n)$. We only need to find an efficient algorithm for inverting \mathbf{V} .

After applying Gohberg–Semencul formula, we can find \mathbf{V}^{-1} in $O(n \log(n))$ time.

3 References

<https://www.nature.com/articles/s41598-019-50234-9>

<https://www.nature.com/articles/s41598-020-60878-7>

Mersereau, R. An algorithm for performing an inverse chirp z-transform. IEEE Transactions on Acoustics, Speech and Signal Processing 22, 387–388 (1974).