Designing an Index for ZooDB

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Outline

- 1 Introduction
- 2 The new Index Implementation
- 3 Benchmarks

ZooDB

- an open source object database written in Java
- JDO standard compliant
- 4 times faster than competitor db4o
- zoodb.org

Key-Value datastructure that allows ordered iteration.

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```
Example:
```

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ZooJdoHelper.createIndex(pm, Person.class, "name",
false);
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$$\begin{array}{c} \mathsf{OID} \\ \mathsf{Object}\text{-}\mathsf{ID} \\ \mathsf{Diskpos} \end{array} \rightarrow$$

 $\begin{array}{c} \mathsf{FSM} \\ \mathsf{Page}\text{-}\mathsf{ID} \to \mathsf{TxID} \end{array}$

B+ Tree

- node fills one disk page
- inner node contains keys and children pointer, leaves contain keys and values
- key unique vs. key-value unique

Example: insert

 $3908 \log + 2665 \log of test$

Goals

- fast B+ tree index
- buffer manager to allow caching
- prefix sharing

Challenges - General

- edge cases
- runtime dominated by disk access
 - change nodes infrequently
 - fewer nodes is better
- only parent to child pointer
- determine best when its time to split/redistribute
- key unique vs. key-value unique
- buffer manager lookup takes time
- prefix-sharing encoding/decoding takes time
- prefix-sharing rebalancing takes time
- general optimizations
 - avoid polymorphism
 - bit-level operations

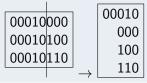
Exploit common prefix of keys in nodes

00010000 00010100 00010110

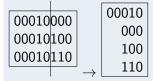
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Exploit common prefix of keys in nodes



- variable number of key-value entries per page
- prefix determines
 - if 2 nodes can be merged without overflow
 - the number of entries that can be redistributed from one node to the other

Our B+ Tree

- 64 bit keys and 64 bit values
- key unique and key non-unique variants
- keys encoded based on common bit prefix

Class Diagram

Make it and add it here

Operations

- Search Similar to normal B+ Tree
- Insert overflow
 - attempt to redistribute values to left sibling before creating a new node
- Delete underflow
 - · check if possible to merge with left or right neighbour
 - check if possible to split current node between left and right
 - redistribute from left or right

Benchmarks

Fine grained

- insert, remove, write
- duration, number of nodes
- prefix-sharing vs. no prefix-sharing

Whole system

- test harness
- partial PolePosition benchmark
- StackOverflow

Summary

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Outlook

• ...