

Designing an Index for ZooDB

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Outline

- 1 Introduction
- 2 Goals & Challenges
- 3 The new Index Implementation
- 4 Benchmarks



- an open source object database written in Java
- JDO standard compliant
- 4 times faster than competitor db4o
- zoodb.org

Database Index

Key-Value data structure

1. **fast** retrieval
2. **ordered** iteration
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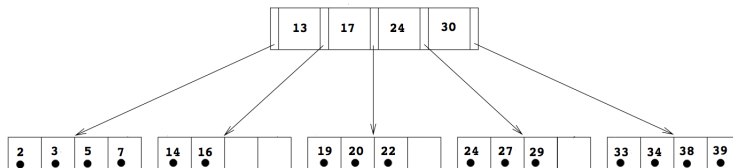
ObjectID Index
OID → Diskpos

Free Space Index
Page-ID → TxID

B+ Tree

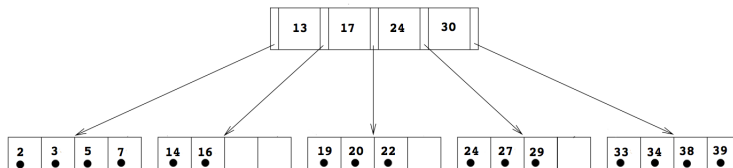
Images adapted from Database Management Systems by Ramakrishnan and Gehrke.

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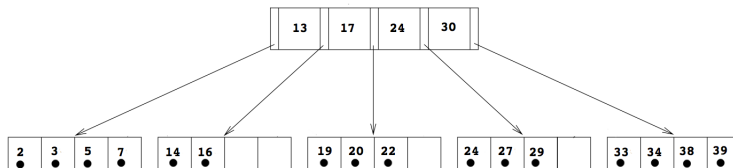
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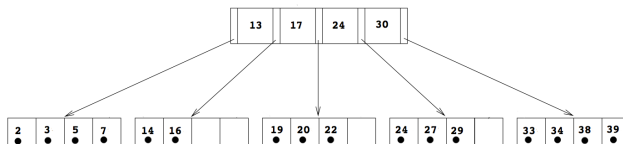
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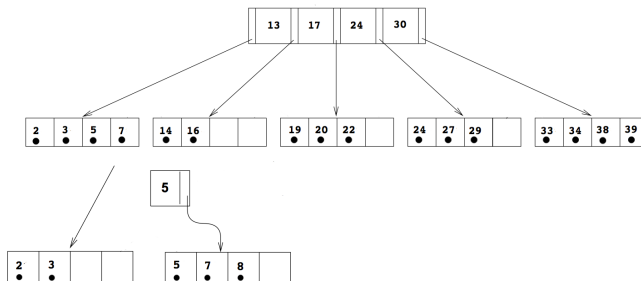


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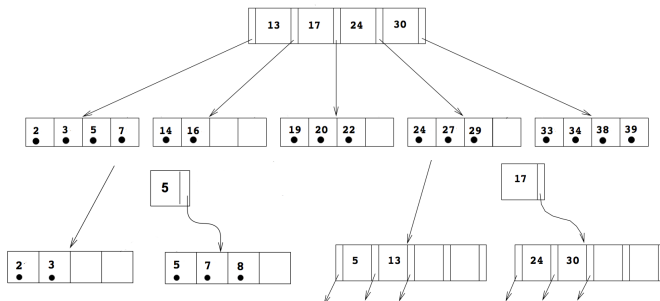
Example: insert (8, v)



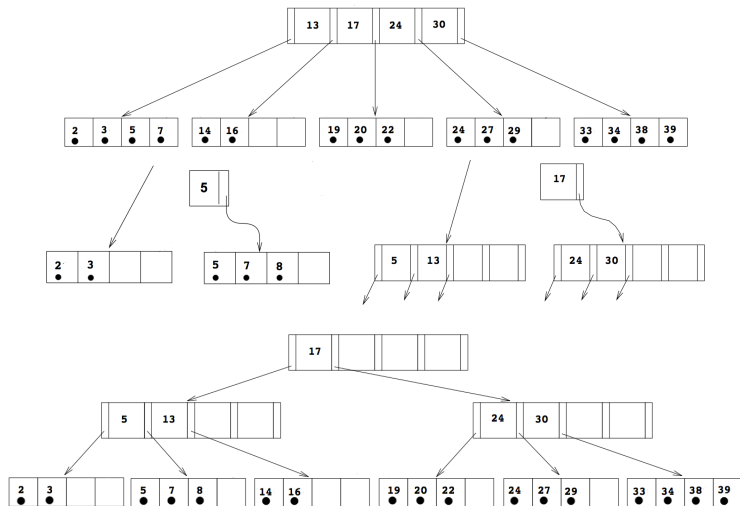
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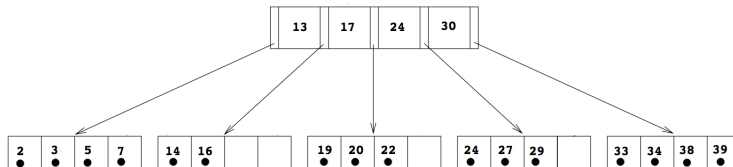


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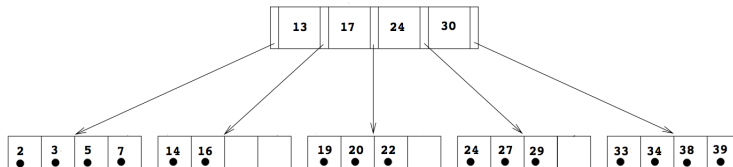
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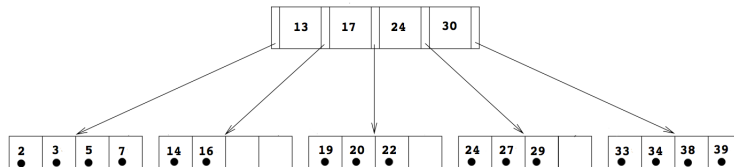
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- Insert, remove, search are logarithmic.

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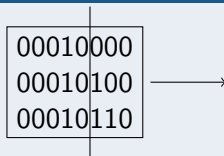
Prefix Sharing

Exploit common prefix

00010000
00010100
00010110

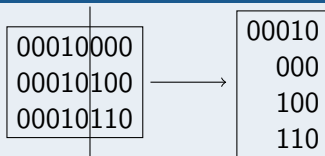
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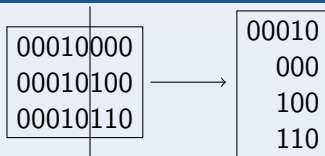
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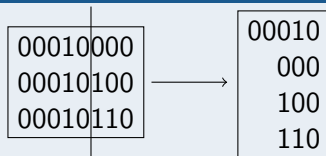
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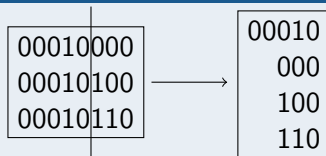
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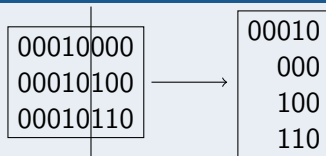
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 - the number redistributions

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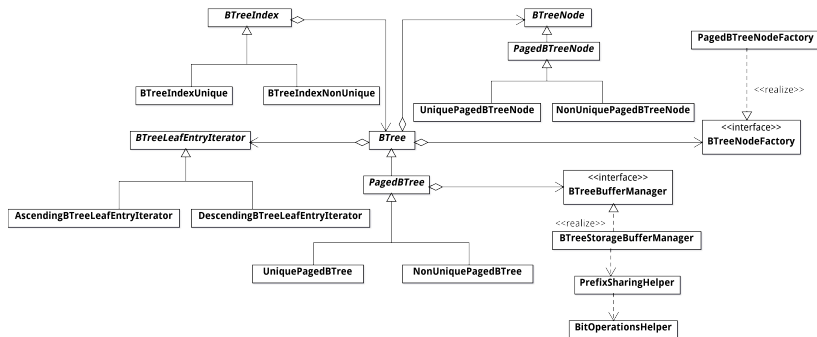
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 1. not optimized for practical scenarios
 2. do not cover duplicates nor prefix sharing
- low-level implementation optimizations

Index Implementation



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 - redistribute from left or right
- Write
 - only write dirty nodes
 - prefix encoding

Microbenchmarks

Duration

Operation	Baseline (No prefix sharing)	Prefix sharing
Search	1	0.9 - 1.1
Insert	1	1.6 - 2.8
Delete	1	1.45 - 2.9

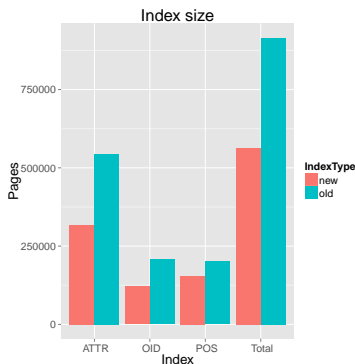
Size of B+ tree

Operation	Baseline (No prefix sharing)	Prefix sharing
Insert	1	0.5 - 1.1
Delete	1	0.5 - 0.75

StackOverflow Data Import

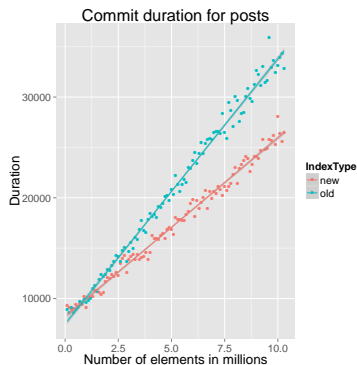
- Real-world workload consisting of importing StackOverflow data: users, posts, comments and votes
- 3 key unique attribute indexes and 9 key-value unique indexes

StackOverflow Import - Index Sizes



Index	Space saving (%)
Attribute	41.6
OID	41.5
POS	23.1
Total	38.5

StackOverflow Import - Commit times



Q&A

- Thank you for your attention!
- Questions ?