Designing an Index for ZooDB

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Outline

- 1 Introduction
- 2 The new Index Implementation
- 3 Benchmarks

ZooDB

- an open source object database written in Java
- JDO standard compliant
- 4 times faster than competitor db4o
- zoodb.org

Key-Value datastructure that allows ordered iteration.

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```
Example:
```

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ZooJdoHelper.createIndex(pm, Person.class, "name",
false);
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 $\begin{array}{l} \mathsf{Attribute} \\ \mathsf{Value} \to \mathsf{Object}\text{-}\mathsf{ID} \end{array}$

Key-Value datastructure that allows ordered iteration.

Example:

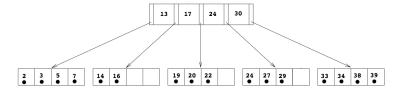
ZooJdoHelper.createIndex(pm, Person.class, "name",
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 $\begin{array}{c} \mathsf{Attribute} \\ \mathsf{Value} \to \mathsf{Object}\text{-}\mathsf{ID} \end{array}$

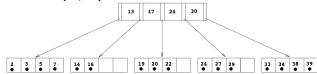
$$\begin{array}{c} \mathsf{OID} \\ \mathsf{Object}\text{-}\mathsf{ID} \\ \mathsf{Diskpos} \end{array} \rightarrow$$

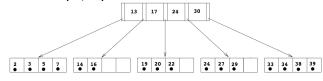
 $\begin{array}{c} \mathsf{FSM} \\ \mathsf{Page}\text{-}\mathsf{ID} \to \mathsf{TxID} \end{array}$

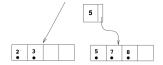
B+ Tree

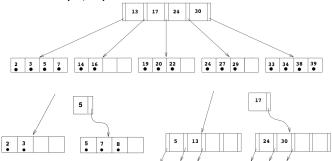


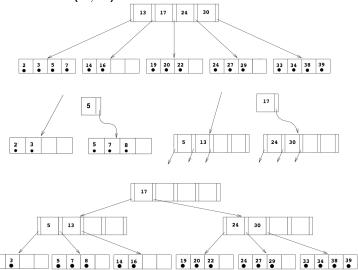
- node fills one disk page
- inner node contains keys and children pointer, leaves contain keys and values
- key unique vs. key-value unique











 $3908 \log + 2665 \log of test$

Goals

- fast B+ tree index
- buffer manager to allow caching
- prefix sharing

Challenges - General

- edge cases
- runtime dominated by disk access
 - change nodes infrequently
 - fewer nodes is better
- only parent to child pointer
- determine best when its time to split/redistribute
- key unique vs. key-value unique
- buffer manager lookup takes time
- prefix-sharing encoding/decoding takes time
- prefix-sharing rebalancing takes time
- general optimizations
 - avoid polymorphism
 - bit-level operations

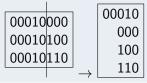
Exploit common prefix of keys in nodes

00010000 00010100 00010110

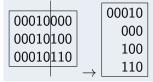
Exploit common prefix of keys in nodes



Exploit common prefix of keys in nodes



Exploit common prefix of keys in nodes



- variable number of key-value entries per page
- prefix determines
 - if 2 nodes can be merged without overflow
 - the number of entries that can be redistributed from one node to the other

Our B+ Tree

- 64 bit keys and 64 bit values
- key unique and key non-unique variants
- keys encoded based on common bit prefix

Class Diagram

Make it and add it here

Operations

- Search Similar to normal B+ Tree
- Insert overflow
 - attempt to redistribute values to left sibling before creating a new node
- Delete underflow
 - · check if possible to merge with left or right neighbour
 - check if possible to split current node between left and right
 - redistribute from left or right

Fine grained

- insert, remove, write
- duration, number of nodes
- prefix-sharing vs. no prefix-sharing

Whole system

- test harness
- partial PolePosition benchmark
- StackOverflow

Summary

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Outlook

• ...