Designing an Index for ZooDB

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Outline

- 1 Introduction
- 2 Goals & Challenges
- 3 The new Index Implementation
- 4 Benchmarks

- an open source object database written in Java
- JDO standard compliant
- 4 times faster than competitor db4o
- zoodb.org

Key-Value data structure that allows for fast retrieval and ordered iteration of records stored in a file.

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ZooJdoHelper.createIndex(pm, Person.class, "name",
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 $\begin{array}{l} \mathsf{Attribute} \; \mathsf{Index} \\ \mathsf{Value} \; \to \; \mathsf{Object}\text{-}\mathsf{ID} \end{array}$

Key-Value data structure that allows for fast retrieval and ordered iteration of records stored in a file.

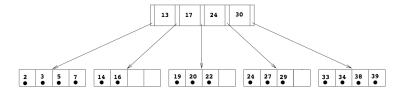
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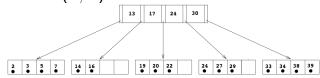
> Attribute Index $Value \rightarrow Object-ID$

ObjectID Index $\mathsf{OID} \to \mathsf{Diskpos}$ Free Space Index Page-ID \rightarrow TxID

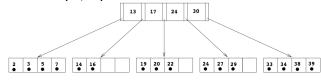
B+ Tree

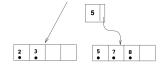


- node fills one disk page
- inner node contains keys and children pointer, leaves contain keys and values
- key unique vs. key-value unique

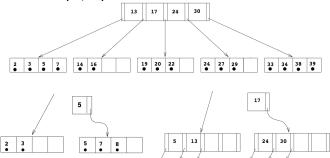


Example: insert (8, v)

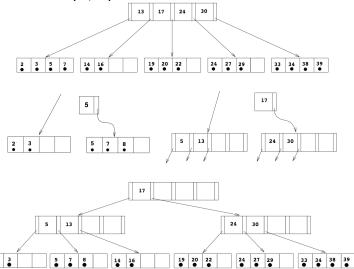




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Goals

- fast B+ tree index
- key unique and key-value unique
- buffer manager to allow caching
- prefix sharing

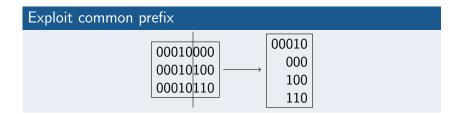
Prefix Sharing

Exploit common prefix

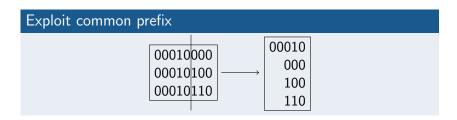
00010000 00010100 00010110

Exploit common prefix 00010000 00010100 00010110 00010110

Prefix Sharing



Prefix Sharing

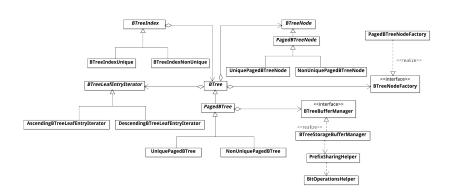


- variable number of key-value entries per page
- prefix determines
 - if 2 nodes can be merged without overflow
 - the number of entries that can be redistributed from one node to the other

Challenges

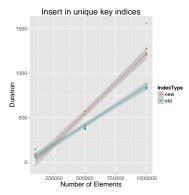
- runtime dominated by disk access
 - prefer few nodes
 - rarely modify nodes
- new features are costly
- Textbook algorithms need to be adapted
 - 1. not optimized for practical scenarios
 - 2. do not cover duplicates nor prefix sharing
- Low-level implementation optimizations

Index Implementation

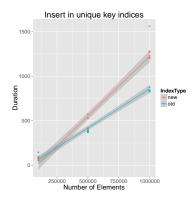


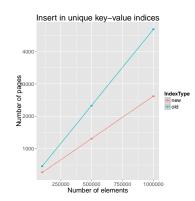
Operations

- Search Similar to normal B+ Tree
- Insert overflow
 - attempt to redistribute values to left sibling before creating a new node
- Delete underflow
 - check if possible to merge with left or right neighbour
 - check if possible to split current node between left and right
 - redistribute from left or right
- Write
 - · only write dirty nodes
 - prefix encoding

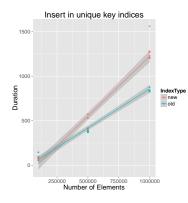


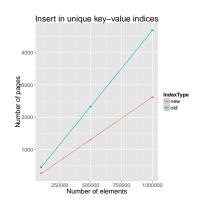
Microbenchmarks





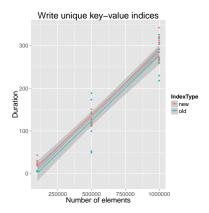
Microbenchmarks

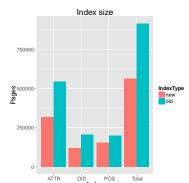




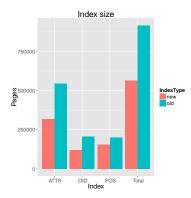
- in every microbenchmark the new index is significantly slower
- in most microbenchmarks there s a significantly lower number of nodes

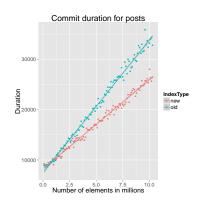
Microbenchmarks





StackOverflow Import





• ...

Outlook

• ..