

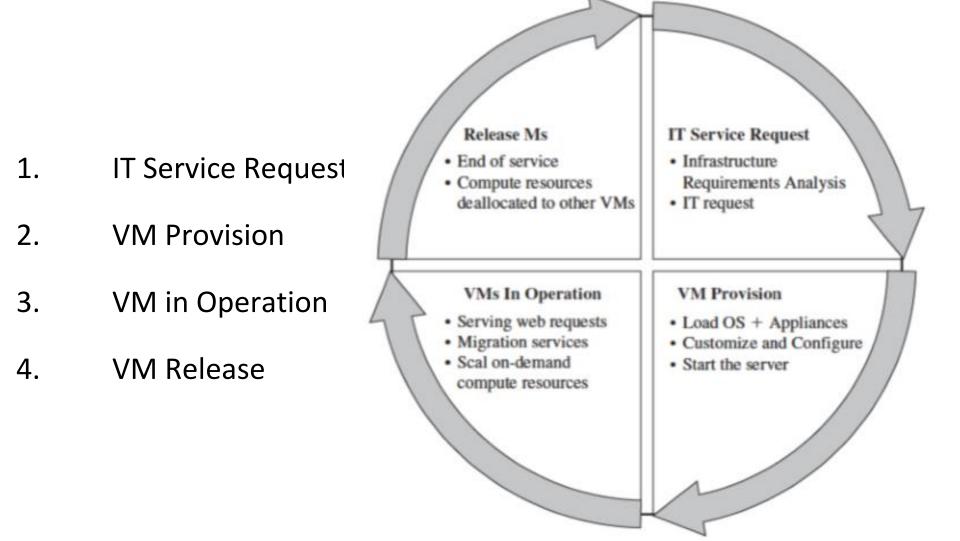


# Cloud Computing Session 11-12 VM Provisioning & Migration

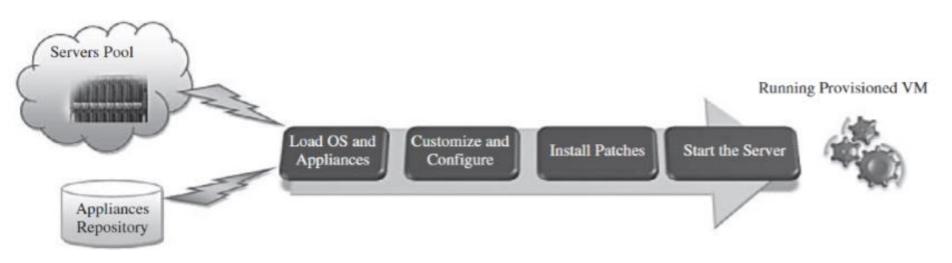
# **Agenda**

- VM Provisioning
- VM Migrations
  - Cold / Regular
  - Hot / Live
    - Flavors

# Virtual Machine Provisioning Life Cycle



# VM Provisioning Process



- Select a server from a pool of available servers along with the appropriate
   OS template
- 2. Load the appropriate software.
- 3. Customize and configure the machine (e.g., IP address, Gateway) to an associated network and storage resources.
- 4. Finally, the virtual server is ready to start with its newly loaded S/W.

# **VM** Provisioning

- Manually installing an OS,
- Using a preconfigured VM template,
- Cloning an existing VM
- Importing a physical server or a Server from another hosting platform.
- Physical servers can also be virtualized and provisioned using P2V (Physical to Virtual)

# VM Provisioning using templates

- Using template reduces the time required to create a new virtual machine.
- Administrators can create different templates for different purposes.
   For example –
- Vagrant provision tool using VagrantFile (template file)
- Heat Orchestration Tool of openstack (Heat template in YAML format)
- Instance creation in cloud, Load balancer in cloud

Enables the administrator to quickly provision a correctly configured virtual server on demand.

# VM Migration

#### What is it?

The process of moving a virtual machine from one host server or storage location to another.



Image courtesy: Britannica

# **VM Migration**

When migration should be started? (When)

Selecting the candidate VM( Which)

Selecting the target PM (Where)

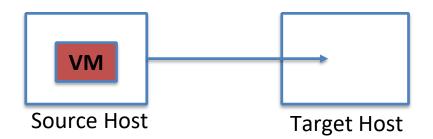
Transfer the VM (Methods)

Image courtesy: researchgate.net

# **VM Migration Flavours**

#### Main Classification

- 1. Hot/live migration
- 2. Cold/regular migration



# **Facilitating VM Migration**

All key machines' components, such as CPU, storage disks, networking, and memory, are completely virtualized, thereby facilitating the entire state of a virtual machine to be captured by a set of easily moved data files.

# **Why Migration**

### Load Balancing

Move VMs from highly loaded servers to lightly loaded servers

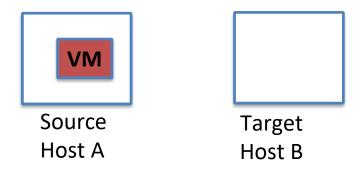
#### Server maintenance

When server needs to be upgraded

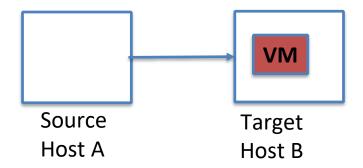
### Energy savings

 Move out VMs before shutting down servers to reduce energy usage

# Migrating a VM from Host A to Host B



Before Migrating a VM from host A to host B



After Migration of VM from host A to host B

# **Cold / Regular Migration**

- The migration of a powered-off virtual machine
- Done in the following steps:
- 1a. If the option to move to a different datastore was chosen,
  - the configuration files, including the NVRAM file (BIOS settings), and log files are moved from the source host to the destination host's associated storage area.
- 1b. If the option to move the virtual machine's disks, these are also moved.
- 2. The virtual machine is registered with the new host.
- 3. After the migration is completed, the old version of the virtual machine is deleted from the source host.

# **Why Live Migration**

 Live migration involves migrating a running/alive virtual machine from one host to the other

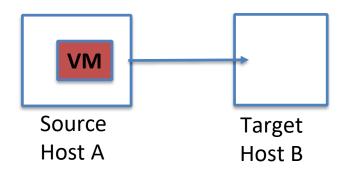
#### Why live?

- To keep long-running jobs alive
- To keep network connections alive
- Broadly, to avoid disruptions to users of VM
- No noticeable interruption to the VM (Ideally)

#### Why VM?

- Why not migrate individual processes?
- Process migration may leave residual dependencies (state) at source host
- E.g. System call redirection, shared memory, open files, interprocess communication, etc.

# **Live Migration: High Level Steps**



#### Migrating a VM from host A to host B

- 1. Setup target host B. Reserve resources for the VM
- 2. Push: Push some memory of VM from A to B
- **3. Stop-and-copy**: Stop the VM at A, copy CPU context, and some memory
- 4. Pull: Start VM at host B, pull any further memory required from A
- 5. Clean up state from host A, migration complete
  - Total migration time: Time for steps 2,3,4
  - Service downtime: Time for step 3

Slide text courtesy: IITB

# **Performance Goals in Live Migration**

- Minimizing Downtime
- Reducing total migration time
- Avoiding interference with normal system activity
- Minimizing network activity

#### **Evaluate metrics:**

- Impact on application performance
- Network bandwidth
- Total pages transferred

# **Challenges in Live Migration**

- VMs have lots of state in memory
- Some VMs have soft real-time requirements :
  - For example: web servers, databases and game servers, etc.
  - Need to minimize down-time

# **Types of Live Migration**

- 1. Pure Stop and Copy
- 2. Pre-Copy
- 3. Post-Copy
- 4. Hybrid-approach of pre-copy and post-copy

### **Pure Stop-and-Copy**

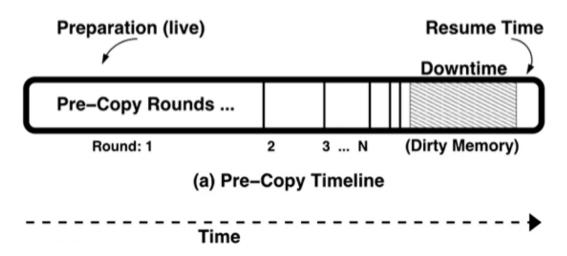
- Freeze VM at source
- Copy the VM's memory contents to target
- Restart VM at target
  - Long downtime
  - Minimal total migration time = downtime

### **Pre-copy Migration**

- 1. Pre-migration process
- 2. Reservation process
- 3. Iterative copy
- 4. Stop and copy
- 5. Commitment

### **Pre-copy Migration**

# Pre-copy migration



DON'T freeze VM at source Let the VM continue to run Copy VM's pseudo-physical memory contents to target over multiple iterations

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### **Pre-Copy Technique**

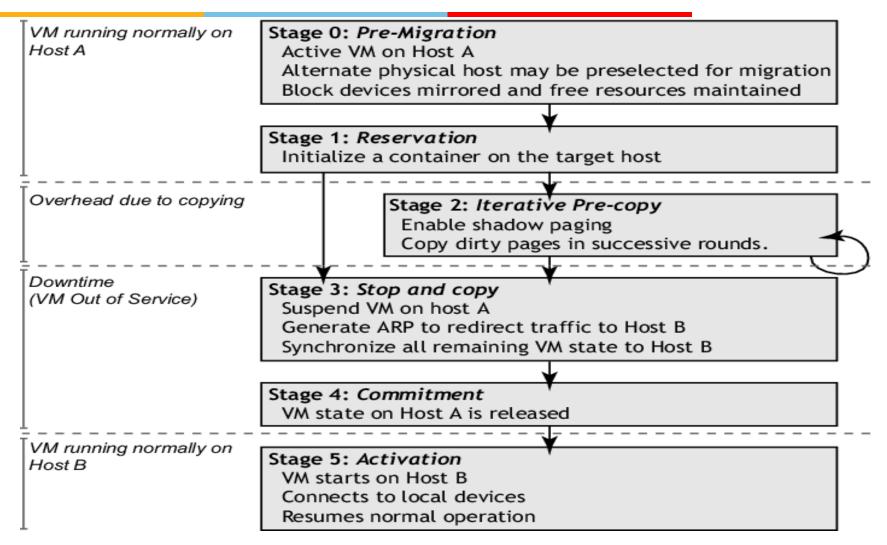


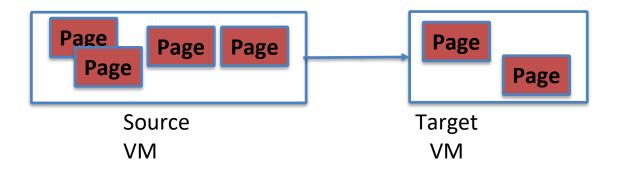
Image courtesy: Live Migration of Virtual Machines", Christopher Clark, Keir Fraser, Steven Hand, Jacob Gorm Hansen<sup>†</sup>, Eric Jul<sup>†</sup>, Christian Limpach, Ian Pratt, Andrew Warfield

### **Pre-copy Migration**

- First iteration → copy all pages.
- Each subsequent iteration → copy pages that were dirtied by the VM during the previous iteration
- Do a short stop-and-copy when number of dirty pages is "small enough".
- But what if number of dirty pages never converges to a small enough number?
  - After a fixed number of iterations, give up and stop-andcopy.

# How do we track dirtied pages?

- Mark the VM's memory pages as read-only after each iteration.
- Trap write operations via hypervisor and track dirtied pages.
- Reset after each iteration
- Works well as long as writes are infrequent



# **Steps of Tracking Dirty Pages**

Shadow page table constructed on demand for every round

- Dirty bitmap maintained for every round
- Any page access by guest, page fault to hypervisor, shadow page table updated
- PTE marked as read-only by default in shadow
- If valid write access, shadow PTE marked writeable, page marked dirty in bitmap
- At end of round, dirty pages are marked for transfer in control software
- Last set of dirty pages copied in stop-and-copy
- Shadow page table and dirty bitmap reinitialized after every round
- Guest page table in target host changed based on new physical addresses

# **Pre-Copy Performance**

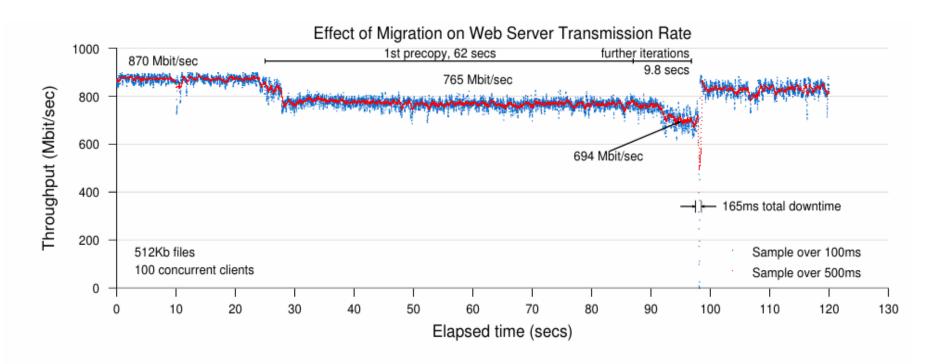
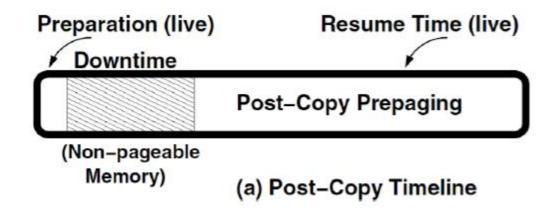


Figure 8: Results of migrating a running web server VM.

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### **Post-copy Migration**



- 1. Prepare target
- 2. Stop source VM, copy CPU context and minimum memory to target
- Start VM at target, but without its memory!
- 4. Pull memory from source via demand paging
- 5. Memory access at target causes page fault, page fetched from source VM

### **Post-copy Migration**

#### Advantage:

Each page transferred over the network only once.

Deterministic total migration time

#### Disadvantage:

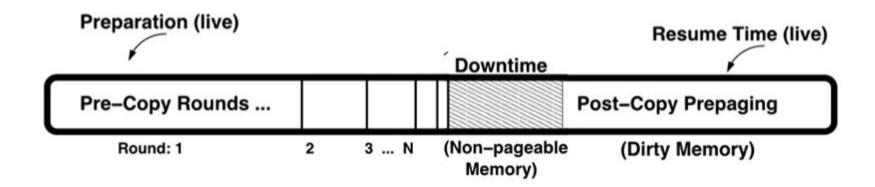
Cold start penalty at the destination

If migration fails, then VM is lost.

### **Optimizations on Post-copy Migration**

- Active pushing: Source proactively pushes important pages, in addition to pulling pages via page faults
- Pre-paging: A "bubble" of pages around faulted page and proactively pushed are pulled, in anticipation of future accesses
- **Dynamic self-ballooning:** VM periodically frees up unnecessary memory and gives it back to hypervisor
  - Reduces memory footprint, speeds up page transfer
  - Performed carefully without hurting application performance
  - Can be used to optimize pre-copy migration as well
- Hybrid: one pre-copy round, followed by post copy

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Combines the benefits & drawbacks of both

- 1. Perform one or more rounds of live pre-copy rounds
- 2. Pause VM and transfer execution state
- 3. Use post-copy to transfer any remaining dirty pages from source

# **Target machine Failure Handling**

#### Pre-copy

- Simply aborts the migration, restart with another target
- Latest state is on source only, so can recover

#### Post copy

- Source has stale memory, target has updated memory
- If target crashes, cannot recover application data (unless some replication is performed)

### **Source Machine Failure Handling**

#### Pre-copy

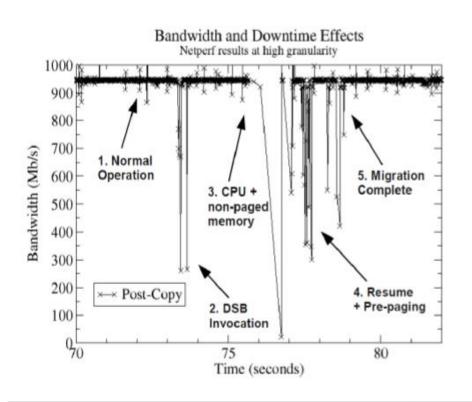
- Latest state was on source only
- Simply aborts the migration, start fresh with the service.
- If replication would have been there, service can be redirected to it and migration starts fresh from there.

#### Post Copy

- Source has stale memory, target has updated memory
- Target may not have pulled all application data (unless some replication is performed)
- Service has to be restarted

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### **Pre-Copy Versus Post-Copy**



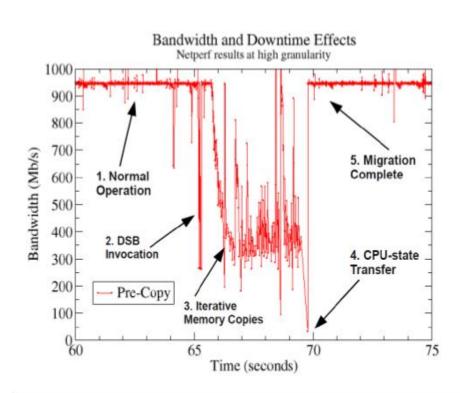
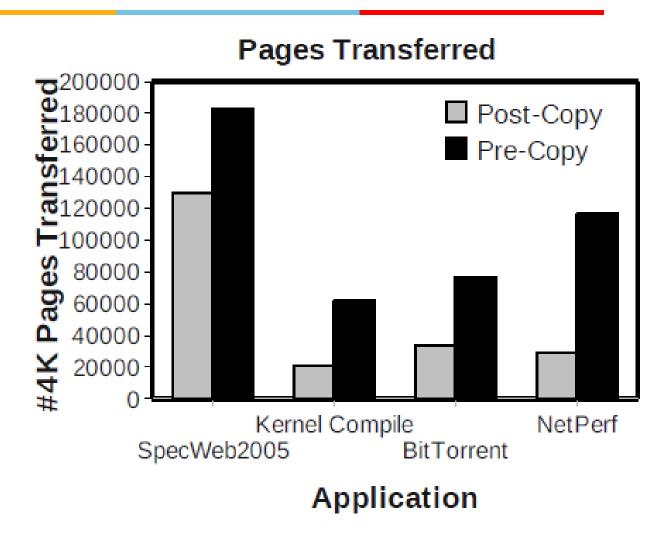
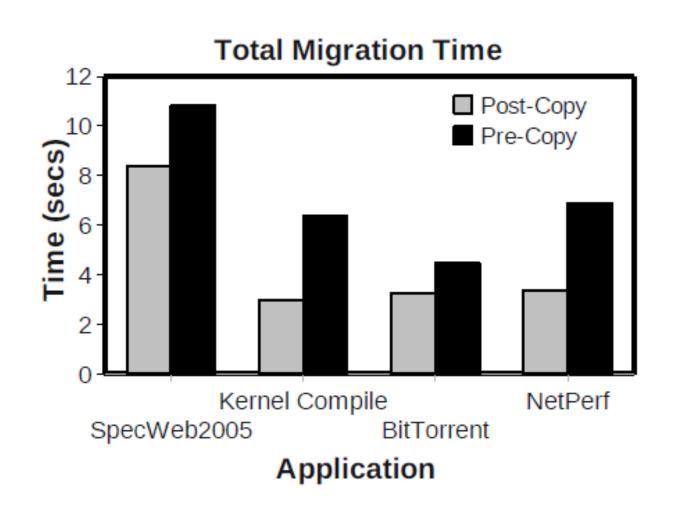
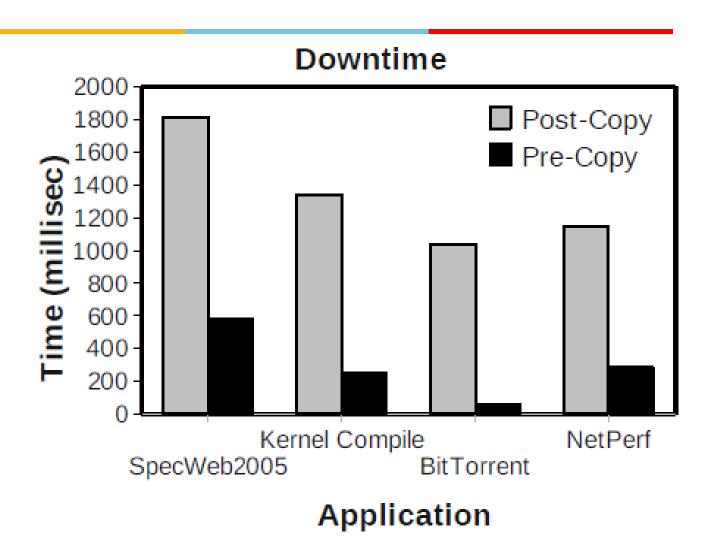


Figure 9. Impact of post-copy on NetPerf bandwidth.

**Figure 10.** Impact of pre-copy on NetPerf bandwidth.







- Post copy is better for memory-intensive applications
  - Longer downtime as compared to pre-copy,
  - But lower total migration time, fewer page transfers, lesser disruption to application
- Pre-copy suited for interactive application

### **Migrating Network Connections**

#### Within a LAN

- The migrated VM carries its IP address, MAC address, and all protocol state, including any open sockets
- Backward (re)learning delay at the network switches
- Switches needs to re-learn the new location of migrated VM's MAC address
- Solution:
  - Send an unsolicited ARP reply from the target host.
  - Intermediate switches will re-learn automatically.
- Few in-flight packets might get lost.

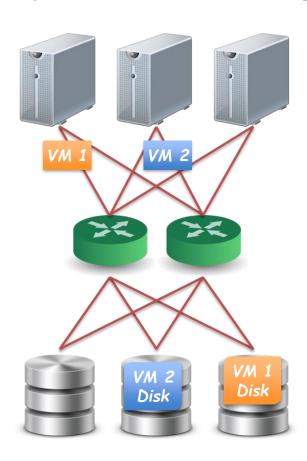
# **Migrating Network Connections**

#### **Across a WAN (Wide Area Network)**

- Source and destination subnets may have different IP addresses.
- Active network connections may need to be tunneled via VPN or similar mechanisms.

# **Storage Migration**

Many gigabytes of local disk image to migrate



**Network Attached Storage** 

# **Storage Migration**

#### For LAN

- Assume the storage is over the network and remains accessible from the new target machine.
- E.g. Network File System (NFS), or Network Block Device(NBD), or iSCSI etc.

#### For WAN

Disk image may need to be transferred.

- Can use pre-copy or post-copy for disk images
- Combined bandwidth saving optimizations such as compression, and/or de-duplication.

# **Summary**

- VM Provisioning
- VM Migration
  - Cold/Regular
  - Hot/Live
- VM Live Migration
  - Pure Stop-and-Copy
  - Pre-Copy
  - Post-Copy
  - Hybrid of Pre-Copy and Post-Copy

### References

- "Live Migration of Virtual Machines", Christopher Clark, Keir Fraser, Steven Hand, Jacob Gorm Hansen, Eric Jul, Christian Limpach, Ian Pratt, Andrew Warfield, NSDI'05: Proceedings of the 2nd conference on Symposium on Networked Systems Design & Implementation Volume, 2005, doi: <a href="https://dl.acm.org/doi/10.5555/1251203.1251223">https://dl.acm.org/doi/10.5555/1251203.1251223</a>
- "Post-Copy Based Live Virtual Machine Migration Using Adaptive Pre-Paging and Dynamic Self-Ballooning", Michael R. Hines and Kartik Gopalan, VEE '09: Proceedings of the 2009 ACM SIGPLAN/SIGOPS international conference on Virtual execution environments, 2009, doi: <a href="https://doi.org/10.1145/1508293.1508301">https://doi.org/10.1145/1508293.1508301</a>