

# A backward error analysis framework for GMRES

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# What is GMRES?

Throughout the presentation, we focus on the Generalized Minimal RESidual (GMRES) algorithm.

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**Algorithm:** GMRES( $A, b, x_0, \tau$ )

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5:    $w_k = Av_k$   
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7:   for  $i = 1, \dots, k$  do  
8:      $h_{i,k} = v_i^T w_k$   
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- Chooses the vector  $x_k$  in  $\text{span}\{V_k\}$  that **minimizes**  $\|Ax_k - b\|$ .
- **Reiterate** until  $x_k$  is a satisfying approximant of  $x$ .

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# GMRES comes in many flavors

## Preconditioning

GMRES might converge too slowly. It is essential to use a preconditioner  $M$  that transforms  $Ax = b$  into an “easier” linear system to solve.

$$M^{-1}Ax = M^{-1}b \quad (\text{left}), \quad Au = b, \quad u = Mx \quad (\text{right})$$

**More possibilities:** split preconditioning, non-constant preconditioners (FGMRES).

*Example of  $M$ :* ILU, polynomial, block Jacobi, approximate inverse, an iterative method, ...

## Restart

**Principle:** under a chosen restart criterion, stop the iteration, erase  $V_k$ , restart GMRES with the initial guess  $x_0 = x_k$ .

The cost in memory and execution time of an iteration grows as we iterate  $\Rightarrow$  Restart cumulates more iterations while bounding the cost.

## Orthogonalization

The Arnoldi process can be constructed with any orthogonalization procedures: Householder QR, CGS, MGS, CGS2, ...

**Warning:** Different tradeoffs between numerical stability and performance!

# What is a backward error analysis?

## Backward and forward errors

Even for  $k = n$ , GMRES computed in finite precision won't deliver the exact solution. We quantify the quality of the computed solution  $\hat{x}_k$  by the quantities

$$bwd = \frac{\|A\hat{x}_k - b\|}{\|A\|\|\hat{x}_k\| + \|b\|}, \quad fwd = \frac{\|x - \hat{x}_k\|}{\|x\|}.$$

*"The process of bounding the backward error of a computed solution is called backward error analysis"* N. J. Higham, Accuracy and Stability of Numerical Algorithms.

## Why we care?

- Formal proof that the computed solution will always be correct.
- Reveals the sensitivity to rounding errors of the different operations.
- Is needed to derive a backward error analysis of an algorithm using GMRES.

## Existing backward error analysis of GMRES

Bounding the backward and forward error of GMRES is NOT EASY:

- GMRES is a complex algorithm made of different sub-algorithms  
→ we need a **backward error analysis on every sub-algorithm**.
- GMRES is an iterative process, **bounds** on the errors are only  
**valid from a certain  $k$**  → we need to answer the question: at which  $k$  the errors are satisfying.

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2007-2008	Flexible MGS GMRES [document] “A Note on GMRES Preconditioned by a Perturbed $LDL^T$ Decomposition with Static Pivoting” by M. Arioli, I. S. Duff, S. Gratton, and S. Pralet, SIAM SISC. [document] “Using FGMRES to obtain backward stability in mixed precision” by M. Arioli and I. S. Duff, ETNA.

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⇒ An almost infinite number of variants...

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- They are very smart, long, and hard ⇒ Understanding and **adapting them is a challenge**.

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- They are very smart, long, and hard ⇒ Understanding and **adapting** them is a challenge.

Consequences:

- **A few GMRES variants have error bounds** on their computed solution.
- Bounding errors of a new variant is **inconvenient** and **tedious**.

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⇒ We aim to propose a modular and generic backward error analysis tool for GMRES.

# Generic GMRES: an abstract algorithm

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**Algorithm:** GEN-GMRES( $A, b, M_l, Z_k$ )

---

- 1: Compute  $C_k = \tilde{A}Z_k$  where  $\tilde{A} = M_l^{-1}A$ .
  - 2: Compute  $\tilde{b} = M_l^{-1}b$ .
  - 3: Solve  $y_k = \operatorname{argmin}_y \|\tilde{b} - C_k y\|$ .
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**Principle:** Finding  $x_k \in \text{span}\{Z_k\}$  minimizing the left-preconditioned residual  $\|\tilde{b} - \tilde{A}x\|$ .

- Little assumptions on the operations.
- $Z_k$  can be any basis of rank  $k$ .
  - Can be seen as a **subspace projection method** solving the left-preconditioned system in  $\text{span}\{Z_k\}$ , where the left-preconditioner  $M_l$ , the basis  $Z_k$ , and the least squares solver are **not specified**.
  - Do not assume Arnoldi process.
  - Not presented as an iterative process.

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Specialization to:

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**Algorithm:** MGS GMRES

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- 1: Compute  $C_k = A\hat{V}_k$ , where  $M_l = I$  and  $\hat{V}_k$  is the computed Arnoldi basis.
  - 2:
  - 3: Solve  $y_k = \operatorname{argmin}_y \|b - A\hat{V}_k y\|$  by MGS Arnoldi.
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Specialization to:

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**Algorithm:** MGS GMRES with left- LU preconditioner

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- 1: Compute  $C_k = \tilde{A}\hat{V}_k$ , where  $\tilde{A} = U \setminus L \setminus A$  and  $\hat{V}_k$  is the Arnoldi basis.
  - 2: Compute  $\tilde{b} = U \setminus L \setminus b$ .
  - 3: Solve  $y_k = \operatorname{argmin}_y \|\tilde{b} - \tilde{A}\hat{V}_k y\|$  by MGS Arnoldi.
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Specialization to:

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**Algorithm:** CGS2 GMRES with flexible LU preconditioner

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- 1: Compute  $C_k = AZ_k$ , where  $M_l = I$  and  $Z_k = U \backslash L \backslash \hat{V}_k$ .
  - 2:
  - 3: Solve  $y_k = \operatorname{argmin}_y \|b - AZ_k y\|$  by CGS2 Arnoldi.
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GEN-GMRES is an abstract generic algorithm that **can be specialized to many GMRES algorithms**  $\Rightarrow$  Any result on GEN-GMRES holds for its specializations.

**Our goal:** Make a backward error analysis of GEN-GMRES.

One analysis to rule them all!

# Generic rounding error model

The terms  $\epsilon_{\tilde{A}}$ ,  $\epsilon_b$ ,  $\epsilon_{LS}$ , and  $\epsilon_z$  quantify the accuracies of every operation and are unspecified. They are only **specified for a given specialization** of GEN-GMRES.

## Matrix–matrix product with the basis (step 2)

$$\text{fl}(\tilde{A}Z_k) = \tilde{A}Z_k + \Delta_{\tilde{A}Z_k}, \quad \|\Delta_{\tilde{A}Z_k}\| \leq \epsilon_{\tilde{A}} \|\tilde{A}Z_k\|.$$

## Preconditioned RHS (step 3)

$$\text{fl}(M_l^{-1}b) = \tilde{b} + \Delta\tilde{b}, \quad \|\Delta\tilde{b}\| \leq \epsilon_b \|\tilde{b}\|.$$

## Least squares solution (step 4)

$$\begin{aligned} \hat{y}_k &= \operatorname{argmin}_y \|\tilde{b} + \Delta b' - (\text{fl}(AZ_k) + \Delta'_{AZ_k})\| \\ \|\Delta\tilde{b}' + \Delta'_{AZ_k} e_j\| &\leq \epsilon_{LS} \|[\tilde{b}, \text{fl}(AZ_k)] e_j\| \end{aligned}$$

## Compute the $k$ th approximant (step 5)

$$\hat{x}_k = \text{fl}(Z_k \hat{y}_k) = (Z_k + \Delta Z_k) \hat{y}_k, \quad \|\Delta Z_k\| \leq \epsilon_z \|Z_k\|$$

# A key dimension(/iteration)

We need to define the special dimension(/iteration)  $k$  at which we can demonstrate that the computed solution has attained a satisfying error.

## Key dimension

We define the key dimension  $k$  as the first  $k \leq n$  such that, for all  $\phi > 0$ , we have

$$\sigma_{\min}([\tilde{b}\phi, \tilde{A}Z_k]) \leq (\epsilon_{\tilde{A}} + \epsilon_b + \epsilon_{LS}) \|\tilde{b}\phi, \tilde{A}Z_k\|_F$$

and

$$\sigma_{\min}(\tilde{A}Z_k) \gg (\epsilon_{\tilde{A}} + \epsilon_b + \epsilon_{LS}) \|\tilde{A}Z_k\|_F.$$

The philosophy of these conditions is to capture the exact moment where  $\tilde{b}$  lies in the range of  $\tilde{A}Z_k$ , which is the moment where the basis  $Z_k$  contains the solution.

 “Modified Gram-Schmidt (mgs), least squares, and backward stability of MGS-GMRES” by C. C. Paige, M. Rozložník, and Z. Strakoš, 2006, SIAM SIMAX.

# Error bounds of GEN-GMRES

## Theorem

Consider the solution of a nonsingular linear system

$$Ax = b, \quad A \in \mathbb{R}^{n \times n}, \quad 0 \neq b \in \mathbb{R}^n,$$

with GEN-GMRES under the previous **error model**. If there exists a key dimension  $k$  as defined previously, then, GEN-GMRES produces a computed solution  $\hat{x}_k$  whose **backward** and **forward** error satisfies respectively

$$\frac{\|b - A\hat{x}_k\|}{\|b\| + \|A\|\|\hat{x}_k\|} \lesssim \Phi \kappa(M_l), \quad \frac{\|\hat{x}_k - x\|}{\|x\|} \lesssim \Phi \kappa(\tilde{A}),$$

where

$$\Phi \equiv \alpha \epsilon_{\tilde{A}} + \beta \epsilon_b + \beta \epsilon_{LS} + \lambda \epsilon_z$$

with

$$\alpha \equiv \sigma_{\min}^{-1}(Z_k) \frac{\|\tilde{A}Z_k\|}{\|\tilde{A}\|}, \quad \beta \equiv \max(1, \sigma_{\min}^{-1}(Z_k) \frac{\|\tilde{A}Z_k\|}{\|\tilde{A}\|}), \quad \lambda \equiv \sigma_{\min}^{-1}(Z_k) \|Z_k\|.$$

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Using the previous theorem requires some work:

- Show that your algorithm is a **specialization of GEN-GMRES**.
- **Determine**  $\epsilon_{\tilde{A}}$ ,  $\epsilon_b$ ,  $\epsilon_{LS}$ , and  $\epsilon_z$ . The difficulty of this step varies according to the existing literature of the sub-algorithms used.
- Show the existence of **the key dimension**. The difficulty also varies according to the existing literature.

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- Show the existence of **the key dimension**. The difficulty also varies according to the existing literature.

This Theorem is **backward compatible with the previous analyses**: Applying it on Householder GMRES, MGS GMRES, and Flexible MGS GMRES gives the same results as the existing analyses.

# Error model for restarted GEN-GMRES

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**Algorithm:** Restarted GEN-GMRES( $A, b, M_l$ )

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- 1: Initialize  $x_0$
  - 2: **repeat**
  - 3:   Compute  $r_i = Ax_i - b$ .
  - 4:   Solve  $Ad_i = r_i$  with GEN-GMRES.
  - 5:   Compute the approximant  $x_{i+1} = x_i + d_i$ .
  - 6: **until** convergence
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Residual computation (step 3)

$$\widehat{r}_i = b - A\widehat{x}_i + \Delta r_i, \quad |\Delta r_i| \leq \epsilon_R(|b| + |A|\|\widehat{x}_i\|).$$

Restart update (step 5)

$$\widehat{x}_{i+1} = \widehat{x}_i + \widehat{d}_i + \Delta x_i, \quad |\Delta x_i| \leq \epsilon_U \|\widehat{x}_{i+1}\|.$$

# Error bounds of restarted GEN-GMRES

## Theorem

Consider the solution of a nonsingular linear system

$$Ax = b, \quad A \in \mathbb{R}^{n \times n}, \quad 0 \neq b \in \mathbb{R}^n,$$

with restarted GEN-GMRES under the previous **error models**. If, for each restart, **the conditions of the previous Theorem (for GEN-GMRES) are met**, then it exists an iteration  $i$  such that restarted GEN-GMRES produces a computed  $\hat{x}_i$  satisfying

$$\frac{\|b - A\hat{x}_i\|}{\|b\| + \|A\|\|\hat{x}_i\|} \leq \epsilon_R + \epsilon_U \quad \text{and} \quad \frac{\|\hat{x}_i - x\|}{\|x\|} \leq \epsilon_R \operatorname{cond}(A, x) + \epsilon_U,$$

provided that for all  $i$

$$\Phi_i \left( \frac{\|M_l\|\|\tilde{A}\|}{\|A\|} \kappa(A) + \kappa(M_l) \right) \ll 1 \quad (\text{backward}) \quad \text{and} \quad \Phi_i \kappa(\tilde{A}) \ll 1 \quad (\text{forward}),$$

where

$$\Phi_i \equiv \alpha_i \epsilon_{\tilde{A}} + \beta_i \epsilon_b + \beta_i \epsilon_{LS} + \lambda_i \epsilon_z$$

with

$$\alpha_i \equiv \sigma_{\min}^{-1}(Z_k^{(i)}) \frac{\|\tilde{A}Z_k^{(i)}\|}{\|\tilde{A}\|}, \quad \beta_i \equiv \max(1, \sigma_{\min}^{-1}(Z_k^{(i)}) \frac{\|\tilde{A}Z_k^{(i)}\|}{\|\tilde{A}\|}), \quad \lambda_i \equiv \sigma_{\min}^{-1}(Z_k^{(i)}) \|Z_k^{(i)}\|.$$

## What about mixed precision?

No mixed precision in this presentation so far!

胸怀怒火，怒火胸怀

# What about mixed precision?

Crazy number of mixed precision GMRES algorithms:

- Hartwig Anzt, Vincent Heuveline, and Björn Rocker, “*An Error Correction Solver for Linear Systems: Evaluation of Mixed Precision Implementations*”, 2011.
- Mario Arioli, Iain S. Duff, Serge Gratton, and Stéphane Pralet, “*A Note on GMRES Preconditioned by a Perturbed  $LDL^T$  Decomposition with Static Pivoting*”, 2007.
- Erin Carson and Nicholas J. Higham, “*A new analysis of iterative refinement and its application to accurate solution of ill-conditioned sparse linear systems*”, 2017.
- Erin Carson and Noaman Khan, “*Mixed Precision Iterative Refinement with Sparse Approximate Inverse Preconditioning*”, 2022.
- Neil Lindquist, Piotr Luszczek, and Jack Dongarra, “*Improving the performance of the GMRES method using mixed-precision techniques*”, 2020.
- Jennifer A. Loe, Christian A. Glusa, Ichitaro Yamazaki, Erik G. Boman, and Sivasankaran Rajamanickam, “*A Study of Mixed Precision Strategies for GMRES on GPUs*”, 2021.
- José Aliaga, Hartwig Anzt, Thomas Grützmacher, Enrique Quintana-Ortí, and Andrés Tomás, “*Compressed basis GMRES on high performance GPUs*”, 2020.
- ...

A lot of them **are not** covered by a backward error analysis!

## What about mixed precision?

Our framework has been designed to facilitate backward error analyses of mixed precision GMRES:

- Mixed precision at the **preconditioner level**: only need to study the accuracy of the product  $M_l A Z_k$ .
  - Mixed precision at the **orthogonalization level**: only need to study the accuracy of the orthogonalization process and evaluate the loss of orthogonality on the basis.
  - Mixed precision at the **restart level**: only need to consider at which precision the residual, the update and the GMRES solver are computed.
  - Mixed precision in **every of these parts** works as well.
- ⇒ **Goal:** Help **keep up backward error analysis coverage** of the increasing number of mixed precision GMRES algorithms.

# Conclusion

## Takeaways

- Many GMRES variants not covered by a backward error analysis.
- We propose a backward error analysis framework to efficiently derive error bounds on new variants.
- We can apply this framework on most existing mixed precision GMRES.

It is still an ongoing work. Preprint will be available soon.