

## Half - Subtractor

=> Half Subtractor is a combinational circuit which is used to perform subtraction of two bits. It has two inputs 'A' & 'B' (minuend & subtrahend) and two outputs 'D' (Difference) and 'Bo' (borrow).

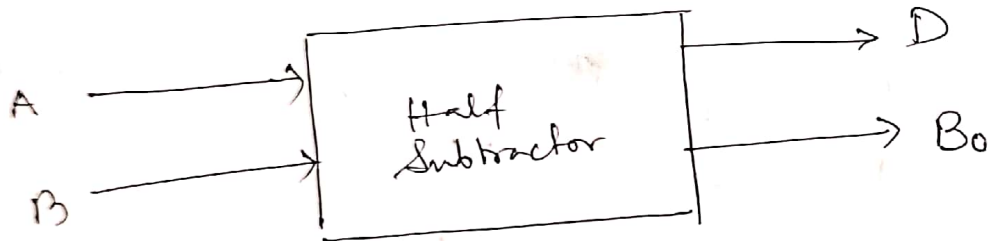


Fig. - Logic Symbol

Inputs		Outputs	
A	B	D	Bo
0	0	0	0
0	1	1	1
1	0	1	0
1	1	0	0

=> From the above truth table, it is visible that the difference (D) is '0' if  $A = B$  and  $D = 1$ , if  $A \neq B$ .

=> The borrow output 'Bo' is '1', whenever  $A = 0$  and  $B = 1$  or 'A' is lesser than 'B'.

$\Rightarrow$  So, the boolean expression for difference and Borrow output can be written as -

$$D = \bar{A}B + A\bar{B} = A \oplus B$$

$$B_0 = \bar{A}B$$

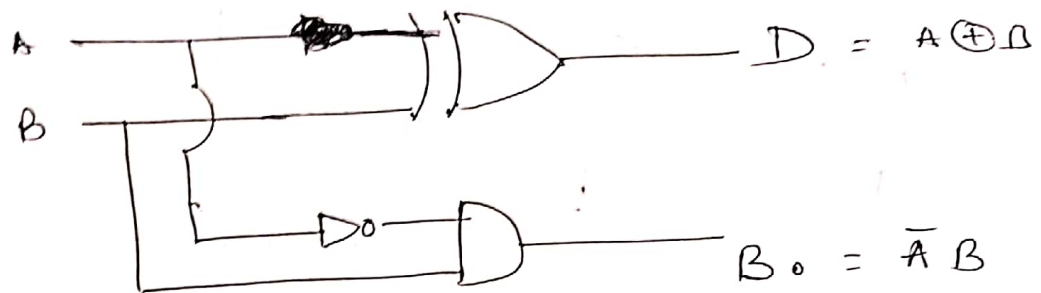


Fig. - Half Subtractor using Logic Gates

### Half Subtractor using Universal Logic:-

① using NAND Gate:-

$$D = A \oplus B = \bar{A}B + A\bar{B}$$

$$= \bar{A}B + B\bar{B} + A\bar{B} + A\bar{A}$$

$$= B(\bar{A} + \bar{B}) + A(\bar{A} + \bar{B})$$

$$= \cancel{(\bar{A} + \bar{B})}(\cdot)$$

$$= B \cdot \overline{AB} + A \cdot \overline{AB}$$

$$= \overline{B \cdot AB} \cdot \overline{A \cdot AB} \quad [\text{using De-Morgan's}]$$

$$B_0 = \bar{A}B = \bar{A}B + B\bar{B}$$

$$= B(\bar{A} + \bar{B}) = B \cdot \overline{AB}$$

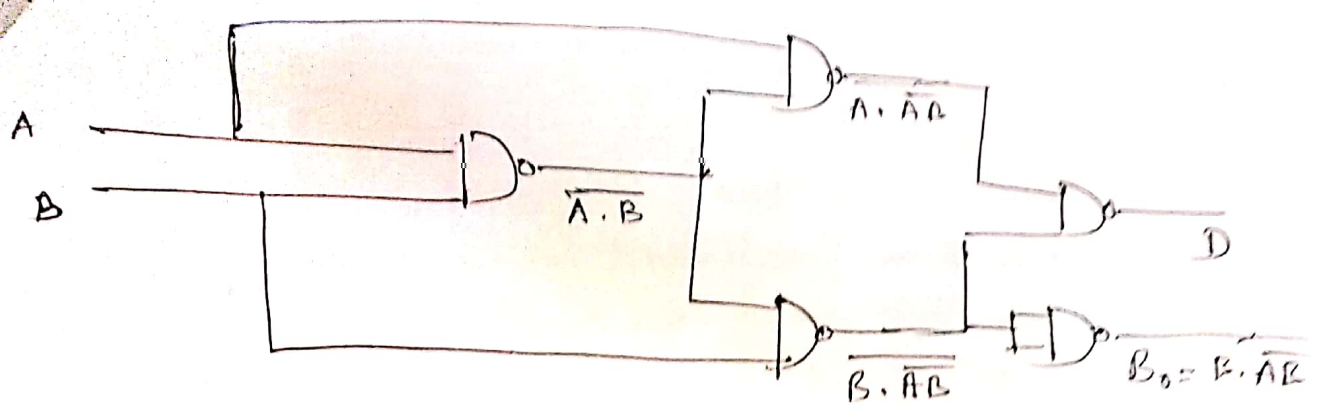


Fig → Half Subtractor using NAND Logic

② Using NOR-Logic:-

$$D = A \oplus B = \bar{A}B + A\bar{B}$$

$$= \bar{A}B + B\bar{B} + A\bar{B} + A\bar{A}$$

$$[\because B\bar{B} = 0 = A\bar{A}]$$

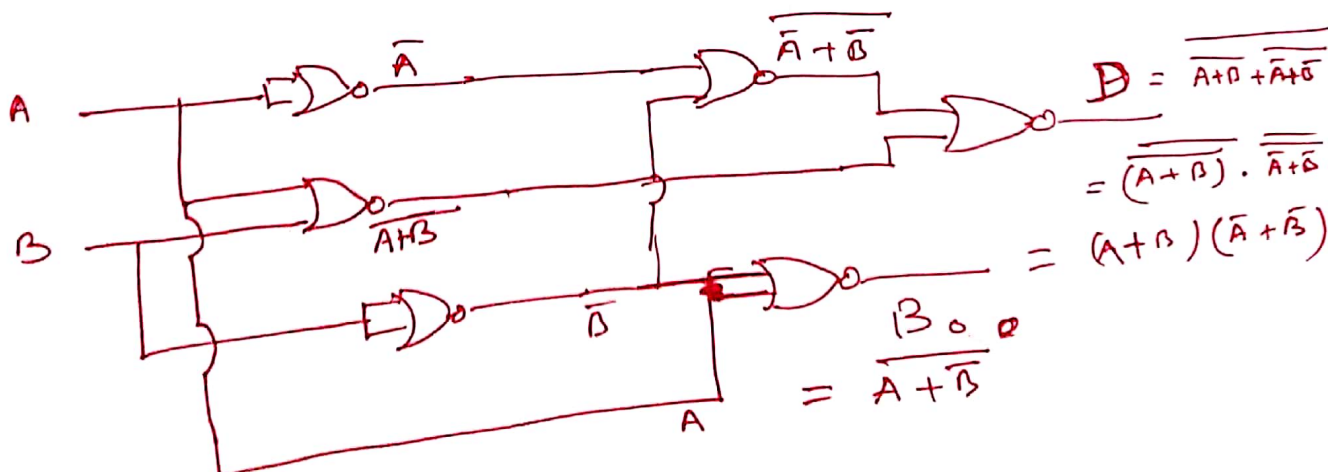
$$= B(\bar{A} + \bar{B}) + A(\bar{A} + \bar{B})$$

$$= (\bar{A} + \bar{B})(A + B)$$

$$B_0 = \bar{A}B = \bar{A}B + B\bar{B}$$

$$= B(\bar{A} + \bar{B})$$

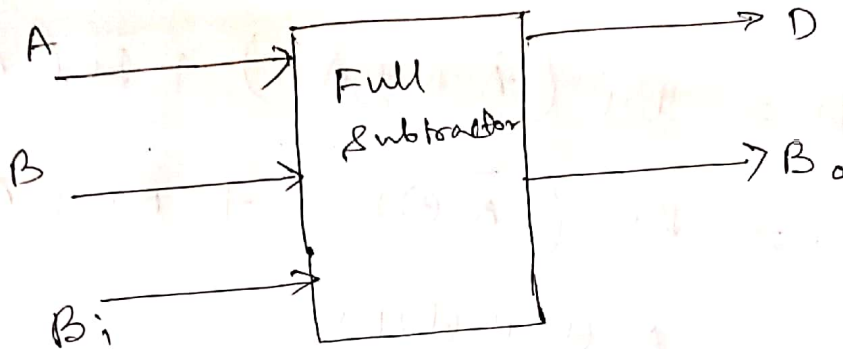
$$B_0 = \bar{A}B = \overline{\overline{\bar{A}B}} = \overline{\overline{\bar{A}} + \overline{\bar{B}}} = \overline{A + B}$$



## Full Subtractor

=> A full Subtractor is a Combinational Circuit that performs Subtraction involving three bits, namely minuend, Subtrahend & bit and the borrow from the previous stage.

Logic Symbol :-



$B_i$  = Borrow input from previous stage.

Truth Table:-

Inputs			Outputs	
A	B	$B_i$	D	$B_o$
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	0	1
1	0	0	1	0
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1



From the truth table;

$$D = \bar{A} \bar{B} B_i + \bar{A} B \bar{B}_i + A \bar{B} \bar{B}_i + A B B_i$$

Simplifying the above expression  $\rightarrow$

$$\begin{aligned} D &= B_i (\bar{A} \bar{B} + A B) + \bar{B}_i (\bar{A} B + A \bar{B}) \\ &= B_i (A \oplus B) + \bar{B}_i (A \oplus B) \\ &= A \oplus B \oplus B_i \end{aligned}$$

\* Similarly for  $B_0$ , the SOP expression can be written as  $\rightarrow$

$$B_0 = \bar{A} \bar{B} B_i + \bar{A} B \bar{B}_i + \bar{A} B B_i + A B B_i$$

The expression can be simplified as K-map.

K-map for  $B_0$  (SOP expression:  $\bar{A} \bar{B} B_i + \bar{A} B \bar{B}_i + \bar{A} B B_i + A B B_i$ )

	$B_i$	00	01	11	10
$A$	0	0	1	1	1
	1	0	1	1	0

Groupings:  $I_1$  (cells 01, 11),  $I_2$  (cells 01, 11),  $I_3$  (cells 11, 10)

$$B_0 = \bar{A}B + \bar{A}B_i + BB_i$$

Using the above simplified expressions, the full-subtractor can be realized as shown below: -

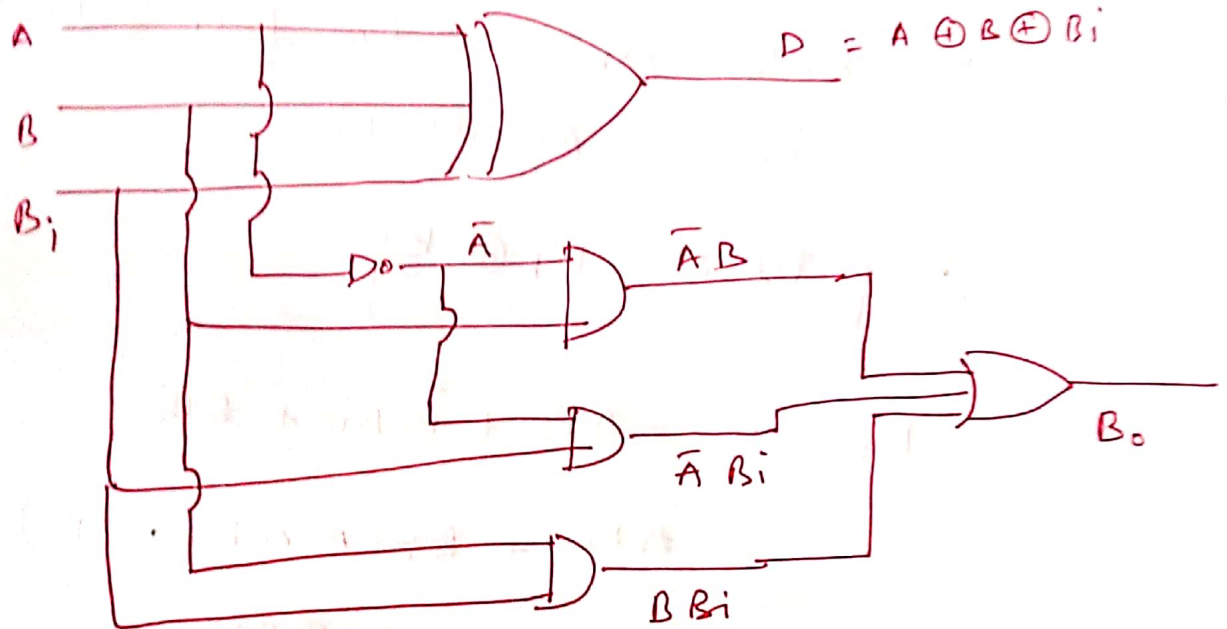
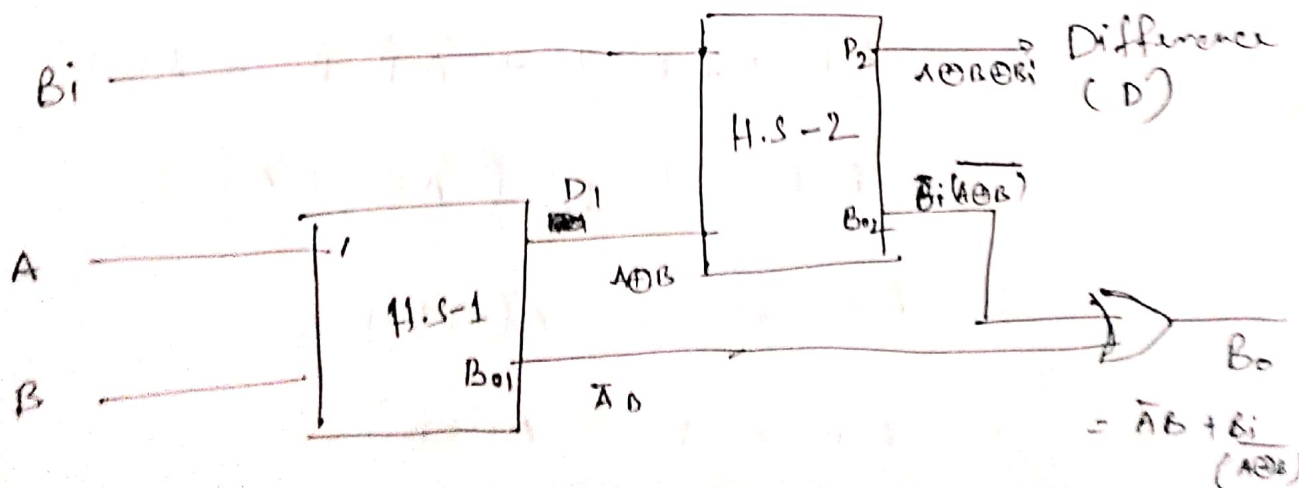


Fig. Full Adder Circuit using Logic Gates

\* The full-subtractor can be implemented using two half-subtractors and an OR Gate as shown below.



The boolean expression for half subtractor  $\rightarrow$

$$D_h = A \oplus B$$

$$B_{oh} = \bar{A}B$$

The boolean expression for full subtractor  $\rightarrow$

$$D_f = \underline{A \oplus B} \oplus B_i$$

$$\therefore \boxed{D_f = D_h \oplus B_i}$$

$$B_{of} = \bar{A}B + \bar{A}B_i + B B_i$$

$$= \bar{A}B + \cancel{B}(\bar{A} \cdot \bar{A}B_i(B + \bar{B})) \quad [\because B + \bar{B} = 1]$$
$$+ B B_i$$

$$= \bar{A}B + \bar{A}B B_i + \bar{A}\bar{B} B_i + B B_i$$

$$= \bar{A}B(1 + B_i) + \bar{A}\bar{B} B_i + B B_i$$

$$= \bar{A}B + \bar{A}\bar{B} B_i + B B_i \quad [\because 1 + B_i = 1]$$

$$= \bar{A}B + \bar{A}\bar{B} B_i + B B_i (A + \bar{A}) \quad [\because A + \bar{A} = 1]$$

$$= \bar{A}B + \bar{A}\bar{B} B_i + A B B_i + \bar{A} B B_i$$

$$= \bar{A}B(1 + B_i) + \bar{A}\bar{B} B_i + A B B_i$$

$$= \bar{A}B + B_i (A B + \bar{A}\bar{B})$$

$$= \bar{A}B + B_i (A \odot B)$$

$$= \bar{A}B + B_i (\overline{A \oplus B})$$