

## CS70–Fall 2011 — Solutions to Homework 2

October 11, 2016

### 1. Practice Proving Propositions

- (a) **Claim:** For all natural numbers  $n$ , if  $n$  is even then  $n^2 + 2011$  is odd.

**Proof:** *by Induction*

Let

$$P(n) = n^2 + 2011 \quad \forall n \in \mathbb{N}$$

**Base Case:**

$$P(0) = 0^2 + 2011 = 2011$$

which is odd. So the base case is true.

**Inductive Step:**

Assume  $P(n)$  is true. Taking  $P(n + 2)$

$$P(n + 2) = (n + 2)^2 + 2011$$

$$P(n + 2) = n^2 + 4n + 2015$$

Take  $K(q) = q^2 + 4q + 2015$ . There is only one possibility

- i.  $q$  is even. If  $q$  is even,  $q^2$  is also even. Moreover  $4q$  becomes even. So we have

$$K(q) = \text{Even} + \text{Even} + \text{Odd}$$

$$K(q) = \text{Even} + \text{Odd}$$

$$K(q) = \text{Odd} \checkmark$$

So,  $P(n+2) = n^2 + 4n + 2015$  is odd. So,  $P(n) \rightarrow P(n+2)$  is true. The claim is correct. ✓

(b) **Claim:** For all natural numbers  $n$ ,  $n^2 + 5n + 1$  is odd.

**Proof:** *by Induction*

Let

$$P(n) = n^2 + 5n + 1 \quad \forall n \in \mathbb{N}$$

**Base Case:**

$$P(0) = 0^2 + 5(0) + 1 = 1$$

which is odd. So the base case is true.

**Inductive Step:**

Assume  $P(n)$  is true. Taking  $P(n+1)$

$$P(n+1) = (n+1)^2 + 5(n+1) + 1$$

$$P(n+1) = n^2 + 7n + 7$$

Take  $K(q) = q^2 + 7q + 7$ . There are two possibilities

- i.  $q$  is odd. If  $q$  is odd,  $q^2$  is also odd. Moreover  $7q$  becomes odd. So we have

$$K(q) = \text{Odd} + \text{Odd} + \text{Odd}$$

**Axiom:** Adding an Odd number with an Odd number gives an even number.

$$K(q) = \text{Even} + \text{Odd}$$

**Axiom:** Adding an Even number with an Odd number gives an Odd number.

$$K(q) = \text{Odd} \quad \checkmark$$

- ii.  $q$  is even. If  $q$  is even,  $q^2$  is also even. Moreover  $7q$  becomes even. So we have

$$K(q) = \text{Even} + \text{Even} + \text{Odd}$$

$$K(q) = \text{Even} + \text{Odd}$$

$$K(q) = \text{Odd} \quad \checkmark$$

So,  $P(n+1) = n^2 + 7n + 7$  is odd. So,  $P(n) \rightarrow P(n+1)$  is true.  
The claim is correct. ✓

- (c) **Claim:** For all real numbers  $a, b$ , if  $a + b \geq 2011$  then  $a > 1005$  or  $b > 1005$ .

**Proof:** *Direct Proof*

Let

$$P(a, b) = (a + b \geq 2011)$$

and

$$Q(a, b) = (a > 1005 \text{ or } b > 1005)$$

We have to prove that  $P(a, b) \rightarrow Q(a, b)$  is true. Just take  $a$  to be 1005, therefore for the implication to be true  $b$  has to be greater than 1005, because otherwise  $Q(a, b)$  will be false, and so will the implication.

$$a + b \geq 2011$$

$$1005 + b \geq 2011$$

$$b \geq 1006$$

$$b > 1005$$

which implies that  $Q(a, b)$  is true. ✓

- (d) **Claim:** For all real numbers  $r$ , if  $r$  is irrational then  $r/4$  is irrational.

**Proof:** *by Contraposition*

Let  $P(r)$  be the statement that  $r$  is irrational, and  $Q(r)$  be the statement that  $r/4$  is irrational. Thus, we have to prove

$$P(r) \rightarrow Q(r)$$

By Contraposition,

$$\neg Q(r) \rightarrow \neg P(r)$$

.

Let  $X(r)$  be the proposition that  $r/4$  is rational, and  $Y(r)$  be the proposition that  $r$  is rational. So we have to prove,

$$X(r) \rightarrow Y(r)$$

Given that  $r/4$  is a rational number, we have

$$\left(\frac{r}{4}\right) = \left(\frac{p}{q}\right)$$

where  $p, q \in \mathbb{Z}$

$$r = \left(\frac{4p}{q}\right)$$

**Axiom:** *Given  $p$  is an integer,  $np$  is also integer where  $n \in \mathbb{Z}$*   
Thus,

$$r = \left(\frac{m}{q}\right)$$

where  $m = 4 \cdot p$ , is rational. The claim is correct. ✓

(e) **Claim:** For all natural numbers  $n$ ,  $10n^2 > n!$ .

**Proof:** *by Counterexample*

Let  $P(n)$  be the proposition  $10n^2 > n!$  where  $n \in \mathbb{N}$ . We could simply iterate through a couple of elements in  $\mathbb{N}$  until we find a  $n \in \mathbb{N}$  for which  $P(n)$  doesn't hold.

- i.  $P(1) = (10 > 1)$  ✓
- ii.  $P(2) = (40 > 2)$  ✓
- iii.  $P(3) = (90 > 6)$  ✓
- iv.  $P(4) = (160 > 24)$  ✓
- v.  $P(5) = (250 > 120)$  ✓
- vi.  $P(6) = (360 > 720)$  ✗

At  $n = 6$ , we have a counterexample to the claim  $P(n)$ . Thus the claim is incorrect. ✗

## 2. Interesting Induction

(a) For  $n \in \mathbb{N}$  with  $n \geq 2$ , define  $s_n$  by

$$s_n = \left(1 - \frac{1}{2}\right) \times \left(1 - \frac{1}{3}\right) \times \cdots \times \left(1 - \frac{1}{n}\right).$$

**Claim:**  $s_n = 1/n$  for every natural number  $n \geq 2$ .

**Proof:** *by Induction*

Given

$$s_n = \left(1 - \frac{1}{2}\right) \times \left(1 - \frac{1}{3}\right) \times \cdots \times \left(1 - \frac{1}{n}\right).$$

**Base Case:**

$$s_2 = \left(1 - \frac{1}{2}\right) = \left(\frac{1}{2}\right) \quad \checkmark$$

**Inductive Step:**

Assume  $s_n$  to be true.

$$s_{n+1} = \left(1 - \frac{1}{2}\right) \times \left(1 - \frac{1}{3}\right) \times \cdots \times \left(1 - \frac{1}{n}\right) \times \left(1 - \frac{1}{n+1}\right).$$

$$s_{n+1} = \left(\frac{1}{2}\right) \times \left(\frac{2}{3}\right) \times \cdots \times \left(\frac{n-1}{n}\right) \times \left(\frac{n}{n+1}\right).$$

$$s_{n+1} = \left(\frac{1 \cdot 2 \cdots (n-1) \cdot n}{2 \cdot 3 \cdots n \cdot (n+1)}\right).$$

$$s_{n+1} = \left(\frac{n!}{(n+1)!}\right)$$

$$s_{n+1} = \left(\frac{1}{n+1}\right) \checkmark$$

The claim is correct.  $\checkmark$ .

(b) Let  $a_n = 3^{n+2} + 4^{2n+1}$ . **Claim:** 13 divides  $a_n$  for every  $n \in \mathbb{N}$ .

**Proof:** *by Induction*

Given

$$a_n = 3^{n+2} + 4^{2n+1}$$

**Base Case:**

$$a_0 = 3^2 + 4 = 13$$

$$13|a_0 = 13|13 \checkmark$$

**Inductive Step:** Assume  $a_n$  to be true.

$$a_{n+1} = 3^{n+3} + 4^{2n+3}$$

Using the Hint

$$a_{n+1} - 3a_n = (3^{n+3} + 4^{2n+3}) - (3^{n+2} + 4^{2n+1})$$

After solving algebraically

$$a_{n+1} - 3a_n = 52 \cdot 4^{2n}$$

$$13|(a_{n+1} - 3a_n) = 13|52 \cdot 4^{2n} \checkmark$$

This means that

$$(13|a_{n+1}) - (13|3a_n)$$

is true. Which tells us that  $13|a_{n+1}$  is true.  $\checkmark$

### 3. Proofs, Perhaps

- (a) This claim is false, and the proof is incorrect. The problem lies in the inductive step. The inductive hypothesis is

$$P(k) = (k^2 < k)$$

In Mathematical Induction, you assume  $P(k)$  is true and prove that  $P(k + 1)$  is true, and therefore prove that

$$P(k) \rightarrow P(k + 1)$$

is true  $\forall k \in D$  where  $D$  is some domain. However in this proof, you prove something (the inductive hypothesis) that you already had assumed to be true, which is incorrect.

- (b) This claim is false, and the proof is incorrect. In the Inductive step, you have

$$7^{k+1} = (7^k \cdot 7^k) / 7^{k-1}$$

However, when  $k = 0$

$$7^1 = (7^0 \cdot 7^0) / 7^{-1}$$

The claim only holds for  $k \in N$ , and  $-1 \notin N$ . Therefore, you can't substitute 1 in its place, and the proof breaks.

- (c) This claim is correct, and so is the proof.
- (d) This claim is false, and the proof is incorrect. The problem is in the claim that

$$\max[(a - 1), (b - 1)] = n$$

because this doesn't hold for  $a = b = 0$  because  $-1 \notin N$ .

4. Take the Tokens

**Claim:** For all natural numbers  $k$ , if the pile starts with  $4k + 1$  tokens, then Thuc has a winning strategy.

**Proof:** *by Induction*

Let  $P(k) = 4k + 1 \quad \forall k \in \mathbb{N}$

**Base Case:**

$$P(0) = 4(0) + 1 = 1$$

If there's only 1 token, then Tamara will end up picking it firstly, and thus Thuc wins.

**Inductive Step:**

Assume  $P(k)$  is true. Taking  $P(k + 1)$

$$P(k + 1) = 4(k + 1) + 1$$

$$P(k + 1) = (4k + 1) + 4$$

$$P(k + 1) = P(k) + 4$$

Tamara has to go first, and can pick up either 1, 2 or 3 tiles. Thus, Thuc can afterwards pick

$$4 - \text{No of tiles Tamara picked}$$

to get the number of tiles in the form of  $P(k) = 4k + 1$ , under which condition Thuc knows he will win. ✓



5. Rigorous Recursion

**Claim:** For all inputs  $n \in \mathbf{N}$ , the value returned by the program is  $G(n) = 3^n - 2^n$ .

**Proof:** *by Induction*

Given  $G(n) = 3^n - 2^n$

**Base Case:**

(a) *if* ( $n = 0$ )

$$G(0) = 3^0 - 2^0 = 0 \checkmark$$

(b) *if* ( $n = 1$ )

$$G(1) = 3^1 - 2^1 = 1 \checkmark$$

**Inductive Step:** Assume  $G(k)$  is true  $\forall 0 \leq k \leq i$ . Solving for  $G(k+1)$

$$\begin{aligned} 5G(k+1-1) - 6G(k+1-2) &= 5G(k) - 6G(k-1) \\ &= 5(3^k - 2^k) - 6(3^{k-1} - 2^{k-1}) \end{aligned}$$

After algrebic manipulation

$$= 3^{k+1} - 2^{k+1} \checkmark$$

6. Coloring Countries (Some of the content here including pictures has been taken from <http://mathonline.wikidot.com/6-colour-theorem-for-planar-graphs>)

(a) **Theorem:**

Every possible map can be colored with six colors, in such a way that no two neighboring countries have the same color.

**Proof:** *by Induction*

Let  $P(n)$  be the statement that the vertices in a planar graph can be coloured with 6 colors.

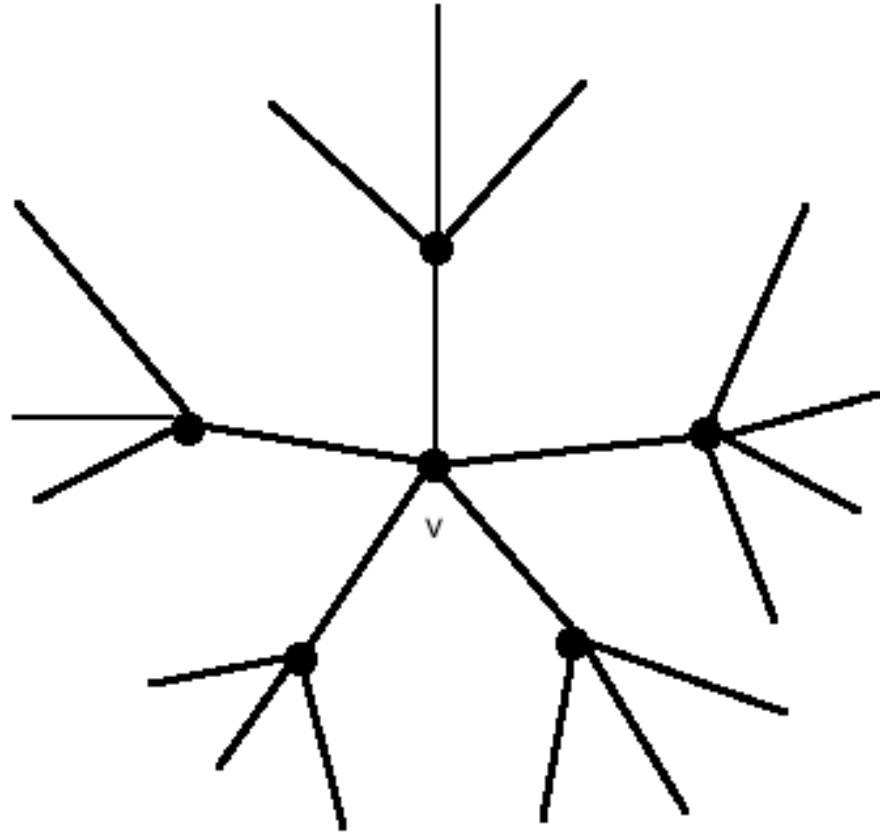
**Base Case:**

For  $1 \leq n \leq 6$ ,  $P(n)$  is correct because you can give each and every vertex a different color than its neighbor.

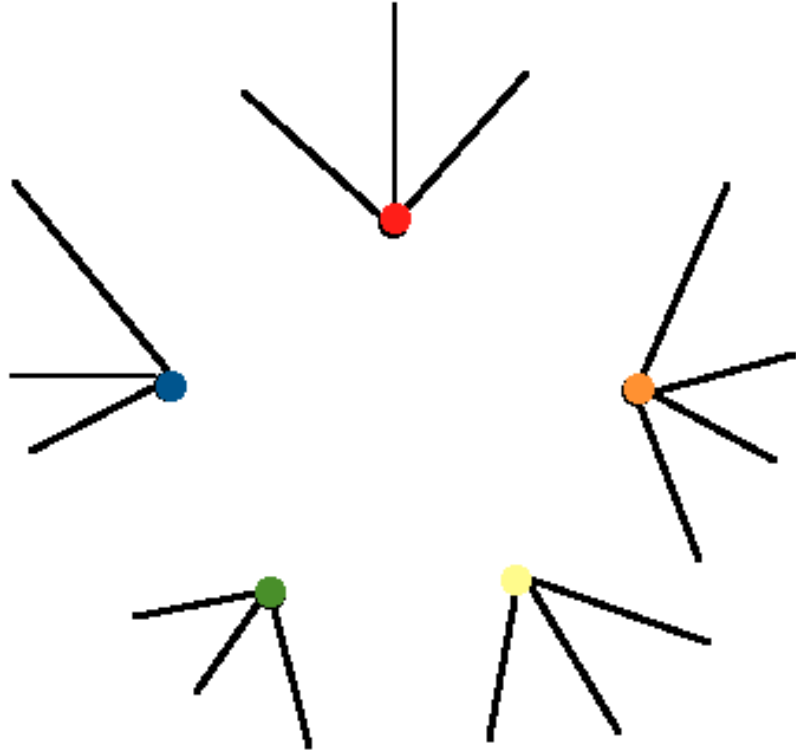
**Inductive Step:**

Assume that  $P(n)$  is true, which means that if we have a planar graph with  $n$  vertices, we can obtain a good coloring. We want to prove that if we have  $n + 1$  vertices, the coloring is still possible.

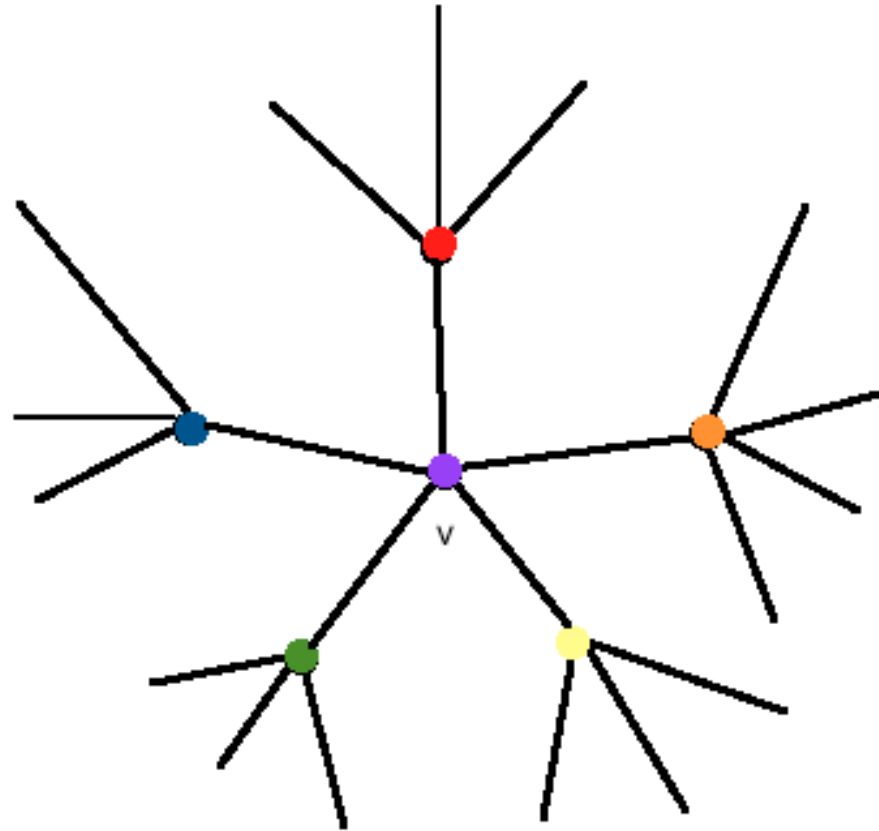
Now, assume  $G$  is a planar graph with  $n + 1$  vertices. The lemma tells us that the degree of the vertex  $\leq 5$ . Therefore, we would have a graph as shown in the figure below.



If we remove the vertex  $v$ , the number of vertices become  $n$ , for which we know that good coloring is possible, as in the following figure .



Adding the vertex  $v$  back, the figure becomes



Thus, a good coloring is possible. So  $P(n + 1)$  is true. ✓