305 Lecture 4.1 - Using Truth Trees

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Showing how we use truth trees to check for logical truth and validity.



Boxes and Diamonds, section 2.4.

1.
$$\mathbb{T} \neg A$$
2. $\mathbb{F} A$ $\neg \mathbb{T}, 1$
2. $\mathbb{T} A \wedge B$
2. $\mathbb{T} A \wedge B$
3. $\mathbb{T} B \wedge \mathbb{T}, 1$
2. $\mathbb{F} A \wedge \mathbb{F} B \wedge \mathbb{F}, 1$
3. $\mathbb{T} B \wedge \mathbb{T}, 1$
4. $\mathbb{T} A \vee B$
5. $\mathbb{T} A \wedge \mathbb{T}, 1$
6. $\mathbb{T} A \vee B$
7. $\mathbb{T} A \wedge \mathbb{T} B \wedge \mathbb{T}, 1$
7. $\mathbb{T} A \to \mathbb{T} B \wedge \mathbb{T}, 1$
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7. $\mathbb{T} A \to \mathbb{T}, 1$
7. $\mathbb{T} A \to \mathbb{T}, 1$

Closure

- A tree without branches closes if there is a sentence on the tree that is marked as both true and false.
- Here is an example. The first line is stipulated, the rest are derived.

Interpreting Closure

- If a tree is closed, it means the initial assumptions can't be true.
- So this tree means that the initial assumption A ∧ ¬A can't be true.

Closure with Branches

- A branching tree closes if every branch closes.
- The next slide has an example, with in this case the top two lines stipulated.

4.

5.

6.

7.

 $\mathbb{T}\mathsf{A}$

Χ

FΑ

 $\mathbb{F} \mathsf{B}$

 $\mathbb{T}\mathsf{B}$

Χ



∨T, 2









Closure with Branches

- A tree with any open branches is open, i.e., not closed.
- If a tree has some open branches and some closed branches, it is open.
- All that matters is if all branches are closed.
- The next slide is an example of an open tree.

4. $\mathbb{T}\mathsf{B}$

5.

6.

 $\mathbb{T}\mathsf{A}$

Χ

 $\mathbb{T} \neg A \wedge B$

 $\mathbb{F}\mathsf{A}$

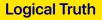
 $\mathbb{T}\mathsf{B}$

 $\wedge \mathbb{T}$, 1

 $\wedge \mathbb{T}$, 1

¬₹, 3

∨**T**, 2



Here is the algorithm for seeing whether a sentence is a logical truth.

1. Start a tree by saying at line 1 that the sentence is **False**.

Logical Truth

Here is the algorithm for seeing whether a sentence is a logical truth.

- 1. Start a tree by saying at line 1 that the sentence is **False**.
- 2. If the tree closes, it is a logical truth.

Logical Truth

Here is the algorithm for seeing whether a sentence is a logical truth.

- 1. Start a tree by saying at line 1 that the sentence is **False**.
- 2. If the tree closes, it is a logical truth.
- 3. If the tree does not close, it is not a logical truth (of propositional logic).

Looking for Counterexamples (Tables)

- Truth tables didn't just tell us that something failed to be a logical truth (of propositional logic).
- · It told us where the failure was.
- You didn't just know that there was an F in the relevant column, you knew which row it was on.
- And that told you where to look for counterexamples.

Looking for Counterexamples (Trees)

- The same thing happens with trees.
- By reading off the open branch, you can see where the sentence fails.
- The trick is to focus on the **atomic** sentences on the branch.
- · These are the ones with no connectives at all.

1.
$$\mathbb{F} A \rightarrow (A \wedge B)$$

2. $\mathbb{T} A \rightarrow \mathbb{T}, 1$
3. $\mathbb{F} A \wedge B \rightarrow \mathbb{T}, 1$
4. $\mathbb{F} A \mathbb{F} B \wedge \mathbb{F}, 3$

- The right hand branch doesn't close.
- The atomics on that branch are that A is T and B is F.

Χ

• So that's the line on the truth table where A \rightarrow (A \land B) is \digamma .

Here is the algorithm for seeing whether an argument is valid (in propositional logic).

1. Start a tree with one line for each premise, saying that the premise is **True**.

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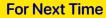
- 1. Start a tree with one line for each premise, saying that the premise is **True**.
- 2. Then have a line that says the conclusion is False.

Here is the algorithm for seeing whether an argument is valid (in propositional logic).

- 1. Start a tree with one line for each premise, saying that the premise is **True**.
- 2. Then have a line that says the conclusion is **False**.
- 3. If the tree closes, the argument is valid.

Here is the algorithm for seeing whether an argument is valid (in propositional logic).

- 1. Start a tree with one line for each premise, saying that the premise is **True**.
- 2. Then have a line that says the conclusion is False.
- 3. If the tree closes, the argument is valid.
- 4. If the tree does not close, the argumnt is not valid (in propositional logic).



We will illustrate this algorithm.