$\Box A \rightarrow \Box \Box A \text{ (in KTB)}$ 

Build a Tableau

To Check Whether it is Valid

## **Hypothesis**

 $\square A \rightarrow \square \ \square \ A$  is a theorem of KTB.

• So we can use all the rules, plus the special rules for B.

1. 1,  $\mathbb{F} \square A \rightarrow \square \square A$  Assumption

Start with it being false at 1.

$$\Box A \rightarrow \Box \Box A$$

You know the drill - left hand side true, right hand side false.

It's T, so true box sentences are true.

1.	1,	Assumption
2.	1, T □ A	$\rightarrow \mathbb{F}$ , 1
3.	1,	$\rightarrow \mathbb{F}$ , 1
4.	1,	T 🗆 2
5.	1.1.	□ <b>F</b> . 3

False  $\square$  sentences have to be false somewhere.

 $\Box A \to \Box \Box A$ 

1.	$1, \mathbb{F}  \Box A \rightarrow \Box \Box A \checkmark$	Assumption
2.	1, ⊤ □ A	$\rightarrow \mathbb{F}$ , 1
3.	1,	$\rightarrow \mathbb{F}$ , 1
4.	1, ⊤ A	T □ 2
5.	1.1,	□ <b>F</b> , 3
6.	1.1, T A	□ <b>T</b> , 2

True  $\square$  sentences have to be true everywhere accessible.

 $\Box A \to \Box \Box A$ 

1.	1,	Assumption
2.	1, ⊤ □ A	<b>→F</b> , 1
3.	1,	<b>→F</b> , 1
4.	1, ⊤ A	T 🗆 2
5.	1.1,	□ <b>F</b> ,3
6.	1.1, <b>⊤</b> A	□ <b>T</b> , 2
7.	1.1.1,	□ <b>F</b> , 5

Another false  $\square$  sentence that needs to be made false.

1.	1, $\mathbb{F} \square A \rightarrow \square \square A \checkmark$	Assumption
2.	1, T □ A	$\rightarrow \mathbb{F}$ , 1
3.	1,	$\rightarrow \mathbb{F}$ , 1
4.	1, ⊤ A	T 🗆 2
5.	1.1,	□ <b>F</b> , 3
6.	1.1, ⊤ A	□ <b>T</b> , 2
7.	1.1.1,	□ <b>F</b> , 5

And that's it. We've applied all the rules. The extra rules for T and B don't change much here from the basic tableau in K.

## **A Model**

- Three worlds, w<sub>1</sub>, w<sub>1,1</sub>, w<sub>1,1,1</sub>.
- The accessibility relations are  $w_1 Rw_{1.1}, w_{1.1} Rw_{1.1.1}, w_1 Rw_1, w_{1.1} Rw_{1.1}, w_{1.11} Rw_{1.11}, w_{1.1} Rw_1$  and  $w_{1.1.1} Rw_{1.1}$
- The first two are from the tree, the next three from T and the last two from B.
- A is true at w<sub>1</sub> and w<sub>1,1</sub> and false at w<sub>1,1,1</sub>.
- So □A will be true only at w<sub>1</sub>.
- So □ □ A will be false at w<sub>1</sub>, as required.