

305 Lecture 2.6 - Complicated Truth Tables

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Plan

This lecture is about how to build more complicated truth tables than we have looked at so far.

Associated Reading

for all x , chapters 10 and 11.

The Example

We are going to work out the truth table for this sentence:

$$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$$

How Many Rows

- How many rows should there be in the truth table?

How Many Rows

- How many rows should there be in the truth table?
- There are three (3) atomic sentences, so there should be $2^3 = 8$ rows.

Laying Out the Rows

- The convention for these is a bit odd.
- Here's one way to think about it.
- For the left-most column you fill the first half of the rows with \mathbb{T} and then the second half of the rows with \mathbb{F} .

First Column

A	B	C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$
T			
T			
T			
T			
F			
F			
F			
F			

Second Column

- Then the second column has one quarter \mathbb{T} , followed by one quarter \mathbb{F} , followed by one quarter \mathbb{T} , followed by one quarter \mathbb{F} .
- In this case that means we alternate every two rows.

Second Column

A	B	C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$
T	T		
T	T		
T	F		
T	F		
F	T		
F	T		
F	F		
F	F		

Third Column

- From now on you do half as many rows between changes.
- In this table we did 4 rows with one value then 4 of another for column 1, 2 with one value then 2 with another for column 2, and now alternate every row for column 3.
- It's helpful to know the full algorithm in case you ever have to do this with 5 or more variables.
- But I won't do that in this course.

Third Column

A	B	C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$
T	T	T	
T	T	F	
T	F	T	
T	F	F	
F	T	T	
F	T	F	
F	F	T	
F	F	F	

Parsing the Sentence

Now we need to go back to our sentence.

$$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$$

What is its **main connective**?

Parsing the Sentence

Now we need to go back to our sentence.

$$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$$

What is its **main connective**?

- It's the first \rightarrow . The sentence is of the form $D \rightarrow E$, where D is $(A \vee \neg B)$ and E is $(B \rightarrow (A \wedge C))$

Building Up

So eventually, we will have the truth value for the whole sentence under the first \rightarrow .

- But that's some distance away.
- While that's where we want to get to, we have to build from the inside out.
- The first thing to do is to repeat the values for the atomic sentences.

Atomic Replicator

A	B	C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$				
T	T	T	T	T	T	T	T
T	T	F	T	T	T	T	F
T	F	T	T	F	F	T	T
T	F	F	T	F	F	T	F
F	T	T	F	T	T	F	T
F	T	F	F	T	T	F	F
F	F	T	F	F	F	F	T
F	F	F	F	F	F	F	F

What can we fill in immediately?

Next Steps

- We have enough on the table to include the values for $\neg B$.
- And we have enough on the table to include the values for the $A \wedge C$ on the far right.
- We'll do these in order.

Negation

Everywhere B is \top , $\neg B$ is \bot , so let's include all of those.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
$\top \top \top$	\top	\bot	\top	\top	\top	\top
$\top \top \bot$	\top	\bot	\top	\top	\top	\bot
$\top \bot \top$	\top		\bot	\bot	\top	\top
$\top \bot \bot$	\top		\bot	\bot	\top	\bot
$\bot \top \top$	\bot	\bot	\top	\top	\bot	\top
$\bot \top \bot$	\bot	\bot	\top	\top	\bot	\bot
$\bot \bot \top$	\bot		\bot	\bot	\bot	\top
$\bot \bot \bot$	\bot		\bot	\bot	\bot	\bot

Negation (cont)

Everywhere B is \mathbb{F} , $\neg B$ is \mathbb{T} , so let's include all of those.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
$\mathbb{T} \mathbb{T} \mathbb{T}$	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{T}	\mathbb{T}
$\mathbb{T} \mathbb{T} \mathbb{F}$	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{T}	\mathbb{F}
$\mathbb{T} \mathbb{F} \mathbb{T}$	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{T}
$\mathbb{T} \mathbb{F} \mathbb{F}$	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{F}
$\mathbb{F} \mathbb{T} \mathbb{T}$	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{T}
$\mathbb{F} \mathbb{T} \mathbb{F}$	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{F}
$\mathbb{F} \mathbb{F} \mathbb{T}$	\mathbb{F}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{F}	\mathbb{T}
$\mathbb{F} \mathbb{F} \mathbb{F}$	\mathbb{F}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{F}	\mathbb{F}

Conjunction

If A, C are both \top , so is $A \wedge C$. So let's include those.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	F	T	T	T	T
T T F	T	F	T	T	T	F
T F T	T	T	F	F	T	T
T F F	T	T	F	F	T	F
F T T	F	F	T	T	F	T
F T F	F	F	T	T	F	F
F F T	F	T	F	F	F	T
F F F	F	T	F	F	F	F

Conjunction (cont)

And $A \wedge C$ is false everywhere else.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	F	T	T	T	T
T T F	T	F	T	T	T	F
T F T	T	T	F	F	T	T
T F F	T	T	F	F	T	F
F T T	F	F	T	T	F	F
F T F	F	F	T	T	F	F
F F T	F	T	F	F	F	T
F F F	F	T	F	F	F	F

Complex Disjunction

- The next step is combining the values of A and $\neg B$ to get the value of $A \vee \neg B$.
- The main thing to remember here is what your inputs are.
- In this case it's not too confusing; it's the values immediately to either side of the \vee .
- But that won't be the general case.

Disjunction

When A is \mathbb{T} , so is $A \vee \neg B$.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T T T
T T F	T	T	F	T	T	T F F
T F T	T	T	T	F	F	T T T
T F F	T	T	T	F	F	T F F
F T T	F		F	T	T	F F T
F T F	F		F	T	T	F F F
F F T	F		T	F	F	F F T
F F F	F		T	F	F	F F F

Disjunction (cont)

And when $\neg B$ is \top , so is $A \vee \neg B$.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T T T
T T F	T	T	F	T	T	F F F
T F T	T	T	T	F	F	T T T
T F F	T	T	T	F	F	T F F
F T T	F		F	T	T	F F T
F T F	F		F	T	T	F F F
F F T	F	T	T	F	F	F F T
F F F	F	T	T	F	F	F F F

Disjunction (part III)

Otherwise, $A \vee \neg B$ is \mathbb{F} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T T T
T T F	T	T	F	T	T	F F F
T F T	T	T	T	F	F	T T T
T F F	T	T	T	F	F	T F F
F T T	F	F	F	T	T	F F T
F T F	F	F	F	T	T	F F F
F F T	F	T	T	F	F	F F T
F F F	F	T	T	F	F	F F F

Conditional

- Now we have to do $B \rightarrow (A \wedge C)$.
- We have to remember the table for \rightarrow - TFFT.
- And we have to remember that what's on the right-hand side of this conditional is a complex sentence: $A \wedge C$.

Conditional Terminology

- It's helpful to recall our distinctive terminology for conditionals
- We'll often call the left-hand side of a conditional the **antecedent**. ('Ante' for before, if that helps.)

Conditional Terminology

- It's helpful to recall our distinctive terminology for conditionals
- We'll often call the left-hand side of a conditional the **antecedent**. ('Ante' for before, if that helps.)
- And we'll call the right-hand side the **consequent** (i.e., what comes after).
- We won't have fancy distinct terminology for the left-hand and right-hand sides of other sentences, because they are symmetric.

Row 1

B is \top , $A \wedge C$ is \top , so this is $\top \rightarrow \top$, i.e., \top .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
$\top \top \top$	\top	\top	\neg	\top	\rightarrow	\top
$\top \top \neg \top$	\top	\top	\neg	\top	\rightarrow	\top
$\top \neg \top \top$	\top	\top	\neg	\top	\rightarrow	\top
$\top \neg \top \neg \top$	\top	\top	\neg	\top	\rightarrow	\top
$\neg \top \top \top$	\neg	\top	\neg	\top	\rightarrow	\top
$\neg \top \top \neg \top$	\neg	\top	\neg	\top	\rightarrow	\top
$\neg \neg \top \top$	\neg	\neg	\top	\rightarrow	\neg	\top
$\neg \neg \top \neg \top$	\neg	\neg	\top	\rightarrow	\neg	\top

Row 2

B is \mathbb{T} , $A \wedge C$ is \mathbb{F} , so this is $\mathbb{T} \rightarrow \mathbb{F}$, i.e., \mathbb{F} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T
T T F	T	T	F	T	F	F
T F T	T	T	T	F	T	T
T F F	T	T	T	F	T	F
F T T	F	F	F	T	F	T
F T F	F	F	F	T	F	F
F F T	F	T	T	F	F	T
F F F	F	T	T	F	F	F

Row 3

B is \mathbb{F} , $A \wedge C$ is \mathbb{T} , so this is $\mathbb{F} \rightarrow \mathbb{T}$, i.e., \mathbb{T} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T
T T F	T	T	F	T	T	F
T F T	T	T	T	F	T	T
T F F	T	T	T	F	T	F
F T T	F	F	F	T	F	T
F T F	F	F	F	T	F	F
F F T	F	T	T	F	F	T
F F F	F	T	T	F	F	F

Row 4

B is \mathbb{F} , $A \wedge C$ is \mathbb{F} , so this is $\mathbb{F} \rightarrow \mathbb{F}$, i.e., \mathbb{T} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T
T T F	T	T	F	T	T	F
T F T	T	T	T	F	T	T
T F F	T	T	T	F	T	F
F T T	F	F	F	T	F	T
F T F	F	F	F	T	F	F
F F T	F	T	T	F	F	T
F F F	F	T	T	F	F	F

Row 5

B is \mathbb{T} , $A \wedge C$ is \mathbb{F} , so this is $\mathbb{T} \rightarrow \mathbb{F}$, i.e., \mathbb{F} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
T T T	T	T	F	T	T	T
T T F	T	T	F	T	F	F
T F T	T	T	T	F	T	T
T F F	T	T	T	F	T	F
F T T	F	F	F	T	F	T
F T F	F	F	F	T	F	F
F F T	F	T	T	F	F	T
F F F	F	T	T	F	F	F

Row 6

B is \mathbb{T} , $A \wedge C$ is \mathbb{F} , so this is $\mathbb{T} \rightarrow \mathbb{F}$, i.e., \mathbb{F} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$					
$\mathbb{T} \mathbb{T} \mathbb{T}$	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{T}
$\mathbb{T} \mathbb{T} \mathbb{F}$	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{F}	\mathbb{F}
$\mathbb{T} \mathbb{F} \mathbb{T}$	\mathbb{T}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{T}
$\mathbb{T} \mathbb{F} \mathbb{F}$	\mathbb{T}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{T}	\mathbb{F}
$\mathbb{F} \mathbb{T} \mathbb{T}$	\mathbb{F}	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{F}	\mathbb{T}
$\mathbb{F} \mathbb{T} \mathbb{F}$	\mathbb{F}	\mathbb{F}	\mathbb{F}	\mathbb{T}	\mathbb{F}	\mathbb{F}
$\mathbb{F} \mathbb{F} \mathbb{T}$	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{T}
$\mathbb{F} \mathbb{F} \mathbb{F}$	\mathbb{F}	\mathbb{T}	\mathbb{T}	\mathbb{F}	\mathbb{F}	\mathbb{F}

Rows 7 and 8

B is \mathbb{F} , $A \wedge C$ is \mathbb{F} , so this is $\mathbb{F} \rightarrow \mathbb{F}$, i.e., \mathbb{F} .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$				
T T T	T	T	F	T	T T T T T
T T F	T	T	F	T	T F T F F
T F T	T	T	T	F	F T T T T
T F F	T	T	T	F	F T T F F
F T T	F	F	F	T	T F F F T
F T F	F	F	F	T	T F F F F
F F T	F	T	T	F	F T F F T
F F F	F	T	T	F	F T F F F

Almost Done

- Now we just need to put the two parts together.
- We have a conditional whose left-hand side, the antecedent, is $A \vee \neg B$.
- And the right-hand side, the consequent, is $B \rightarrow (A \wedge C)$.
- In each row we've computed the truth values for the antecedent and consequent.
- Now it's a matter of just looking up how they combine.
- Remember that the truth table for \rightarrow is TFFT.
- To help, I've started by putting in blue the columns for the antecedent and consequent.

Getting There

Putting the two relevant columns in blue to make them stand out.

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$							
T T T	T	T	F	T	T	T	T	T
T T F	T	T	F	T	T	F	T	F
T F T	T	T	T	F	F	T	T	T
T F F	T	T	T	F	F	T	T	F
F T T	F	F	F	T	T	F	F	T
F T F	F	F	F	T	T	F	F	F
F F T	F	T	T	F	F	T	F	T
F F F	F	T	T	F	F	T	F	F

Row 1

That's $\top \rightarrow \top$, i.e., \top .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	T	F	T	F	F	F
T F T	T	T	T	F	F	T	T	T	T	T
T F F	T	T	T	F	F	T	T	F	F	F
F T T	F	F	F	T	T	F	F	F	T	T
F T F	F	F	F	T	T	F	F	F	F	F
F F T	F	T	T	F	F	T	F	F	T	T
F F F	F	T	T	F	F	T	F	F	F	F

Row 2

That's $\top \rightarrow \text{F}$, i.e., F .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	T	F	T	F	F	F
T F T	T	F	T	F	F	T	T	T	T	T
T F F	T	F	T	F	F	T	T	F	F	F
F T T	F	T	F	T	T	F	F	F	T	T
F T F	F	T	F	T	T	F	F	F	F	F
F F T	F	F	T	F	F	T	F	F	T	T
F F F	F	F	T	F	F	T	F	F	F	F

Row 3

That's $\top \rightarrow \top$, i.e., \top .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	F	T	F	T	F	F
T F T	T	T	T	F	T	F	T	T	T	T
T F F	T	T	T	F	F	T	T	F	F	F
F T T	F	F	F	T	T	F	F	F	F	T
F T F	F	F	F	T	T	F	F	F	F	F
F F T	F	T	T	F	F	T	F	F	F	T
F F F	F	T	T	F	F	T	F	F	F	F

Row 4

That's also $\top \rightarrow \top$, i.e., \top .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	F	T	F	T	F	F
T F T	T	T	T	F	F	T	T	T	T	T
T F F	T	T	T	F	F	T	T	F	F	F
F T T	F	F	F	T	T	F	F	F	T	T
F T F	F	F	F	T	T	F	F	F	F	F
F F T	F	T	T	F	F	T	F	F	T	T
F F F	F	T	T	F	F	T	F	F	F	F

Rows 5 and 6

That's $F \rightarrow F$, i.e., T .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	F	T	F	T	F	F
T F T	T	T	T	F	T	F	T	T	T	T
T F F	T	T	T	F	T	F	T	T	F	F
F T T	F	F	F	T	T	F	F	F	T	T
F T F	F	F	F	T	T	F	F	F	F	F
F F T	F	T	T	F	F	T	F	F	T	T
F F F	F	T	T	F	F	T	F	F	F	F

Rows 7 and 8

That's $\top \rightarrow \top$, i.e., \top .

A B C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$									
T T T	T	T	F	T	T	T	T	T	T	T
T T F	T	T	F	T	T	F	T	F	F	F
T F T	T	T	T	F	F	T	T	T	T	T
T F F	T	T	T	F	F	T	T	F	F	F
F T T	F	F	F	T	T	F	F	F	T	T
F T F	F	F	F	T	T	F	F	F	F	F
F F T	F	T	T	F	F	T	F	F	T	T
F F F	F	T	T	F	F	T	F	F	F	F

Summing Up

It's true everywhere except when A, B are both \mathbb{T} , and C is \mathbb{F} .

A	B	C	$(A \vee \neg B) \rightarrow (B \rightarrow (A \wedge C))$								
T	T	T	T	T	F	T	T	T	T	T	T
T	T	F	T	F	F	T	F	T	F	F	F
T	F	T	T	T	T	F	F	T	T	T	T
T	F	F	T	T	T	F	F	T	F	F	F
F	T	T	F	F	F	T	T	F	F	F	T
F	T	F	F	F	F	T	T	F	F	F	F
F	F	T	F	T	T	F	F	T	F	F	T
F	F	F	F	T	T	F	F	T	F	F	F

For Next Week

- We'll start on analysing arguments using truth tables.
- Remember to do the assignment by Friday 5pm.