305 Lecture 11.3 - Two Basic Results

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• To go over two fundamental results in modal logic.

Associated Reading

• Boxes and Diamonds, section 3.4.

These two claims are equivalent.

- 1. □A
- 2. ¬ ♦ ¬A

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- 2. ¬ ♦ ¬A

From 1 to 2: If \Box A is true at x, then A is true for all y such that xRy. That means there is no y such that xRy and A is not true. That means there is no y such that xRy and \neg A is true. That means \Diamond \neg A is not true at w. That means \neg \Diamond \neg A is true at x.

These two claims are equivalent.

- 1. □A
- 2. ¬ ♦ ¬A

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From 2 to 1: If $\neg \diamondsuit \neg A$ is true at x, then $\diamondsuit \neg A$ is not true at w. So there is no world y such that xRy and $\neg A$ is true at y. So at all worlds y such that xRy, $\neg A$ is not true. So at all worlds y such that xRy, A is true. So $\Box A$ is true at x.

These two claims are also equivalent, though I will not prove this.

- 2. ¬□¬A

Big Picture

These claims are both logically true.

- 1. $\Box \neg A \leftrightarrow \neg \diamondsuit A$
- $2. \ \diamondsuit \neg A \leftrightarrow \neg \Box A$

To move a negation sign outside of a modal operator, either \square or \diamondsuit , you have to rotate the operator by 45 degrees.

Normality

This sentence is also true no matter what the model looks like, and no matter what sentence A is.

$$\Box(A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$$

- · Assume it is false at w.
- So $\Box(A \to B)$ is true at w and $(\Box A \to \Box B)$ is false at w.
- So □A is true at w and □B is false at w.
- So at every where y such that wRy, A must be true (since □A is true at w), and A → B must be true (since □(A → B) is true at w).
- If A and A → B are true at y, B must be true at y as well.
- But this implies that B is true all y such that wRy, contradicting the assumption that □B is false at w.

Normality

This principle has a very important role in the history of modal logics.

$$\Box(A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$$

Having this be an axiom is one of two conditions on what have come to be called **normal** modal logics.

