#### CSE599s, Spring 2014, Online Learning

Lecture 99 - 05/08/2014

## Adaptive Regret Bound

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### 1 Recap

A normal regret bound for fixed learning rate is as follows

$$Regret(u) \le \frac{1}{2\eta} \|u\|^2 + \eta \sum_{t=1}^{T} \|g_t\|_2^2$$
 (1)

We should note that the regret depends on choice of  $\eta$ . If we do not know T in advance, this bound could be very bad. Even if we know T and set  $\eta$  properly, we need to wait until we get T to get the regret bound. Ideally, we want our bound to hold for any T, this is where we need to introduce adaptive update.

Our goal is prove a bound in the following style.

$$Regret \le B\sqrt{\sum_{t=1}^{T} \|g_t\|^2} \ll GB\sqrt{T}$$
 (2)

We want to have a bound that does not need guess and double trick.

There are several class of related algorithms, that we will be discussed in a general framework OGD/Mirro Descent:

$$\hat{w}_{t+1} = \hat{w}_t - \eta_t g_t = \operatorname*{argmin}_{w} g_t w + \frac{1}{2\eta} \|w - w_t\|^2.$$
(3)

FTRL-Proximal

$$w_{t+1} = \underset{w}{\operatorname{argmin}} f_{1:t}(w) + \sum_{s=1}^{t} \frac{\sigma_s}{2} \|w - w_s\|^2.$$
 (4)

**Dual-Averaging** 

$$w_{t+1} = \underset{w}{\operatorname{argmin}} f_{1:t}(w) + \frac{\sigma_{1:t}}{2} ||w||^2.$$
 (5)

They are equivalent when we have no constraint. FTRL Proximal and dual averaging are equivalent when learning rate is constant.

# 2 General Framework for Adaptive Update

In this lecture, we will study the update rule in the following form:

$$w_{t+1} = \underset{w}{\operatorname{argmin}} f_{1:t}(w) + r_{0:t}(w) = \underset{w}{\operatorname{argmin}} h_{0:t}(w)$$
(6)

Note that we have  $h_0(w) = r_0(w)$ . Base on this update rule, we have a strong FTRL Lemma as follows

Lemma 1. Strong FTRL Lemma

$$Regret(u) \le r_{0:T}(u) + \sum_{t=1}^{T} [h_{0:t}(w_t) - h_{0:t}(w_{t+1}) - r_t(w_t)]$$
(7)

*Proof.* The bound can be proved by induction(see previous lecture note)

Before we prove the main theorem, we will need the following lemma

#### Lemma 2. Let

$$w_1 = \operatorname*{argmin}_w \phi_1,$$
 
$$w_2 = \operatorname*{argmin}_w \phi_2 = \operatorname*{argmin}_w [\phi_1(w) + \psi(w)],$$

where  $\phi_1$  is 1 strongly convex function with respect to norm  $\|.\|$ , and  $\psi(w)$  is convex function. Let  $b \in \partial \psi(w)$ , then we will have

$$\phi_2(w_1) - \phi_2(w_2) \le \frac{1}{2} ||b||_*,$$
  
 $||w_1 - w_2|| \le ||b||_*.$ 

We can verify that for a special case where  $\phi_1$  is quadratic function, and  $\psi$  is linear:  $\phi_1(w) = \frac{1}{2} ||w||^2$ ,  $\phi(w) = bw$ , this bound is tight.

**Theorem 3.** Assuming  $r_t(w) \ge 0$ ,  $r_t(w_t) = 0$ ,  $h_{0:t}(w)$  is 1 strongly convex with respect to  $||.||_t$ . Then the regret of general framework can be bounded by

$$Regret(u) \le r_{0:T}(u) + \frac{1}{2} \sum_{t=1}^{T} \|g_t\|_{(t,*)}^2.$$
 (8)

*Proof.* For fixed round t, Let

$$\phi_1(w) = f_{1:t-1}(w) + r_{1:t-1}(w) + r_t(w) = h_{0:t-1}(w) + r_t(w_t)$$

Note that  $w_t = \operatorname{argmin}_w r_t(w_t)$ , we have  $w_t = \operatorname{argmin}_w \phi_1(w)$ . Let  $\psi = f_t$ , and  $b = g_t, g_t \in \partial f_t(w)$ . The following inequality holds follows because of Lemma ??.

$$h_{0:t}(w_t) - h_{0:t}(w_{t+1}) = \phi_1(w_t) + f_t(w_t) - \phi_1(w_{t+1}) - f_t(w_{t+1})$$

$$= \phi_2(w_t) - \phi_2(w_{t+1}) \le \frac{1}{2} \|g_t\|_{(t,*)}^2.$$
(9)

Then the results follows by Lemma ??.

Now we need to make use Theorem ?? to analyze FTRL-Proximal algorithm. A first simple fact is that if  $r_t$  is  $\sigma_t$  strongly convex with respect to ||.||, then  $r_{0:t}$  is 1 strongly convex with respect to  $||u||_t = \sqrt{\sigma_{1:t}}||.||$ . For FTRL-Proximal, we have

- $r_0(w) = I_W(w)$
- $r_t = \frac{\sigma_t}{2} ||w w_t||^2$ , note  $\eta_t = \frac{1}{\sigma_{1:t}}$
- $||g||_{t,*} = \frac{1}{\sqrt{\sigma_{1:t}}} ||g||_2$

Applying Theorem ??, we can get the following bound for adaptive learning rate.

$$Regret(u) \le \frac{(2B)^2}{2\eta_T} + \frac{1}{2} \sum_{t=1}^T \eta_t ||g_t||^2$$
 (10)

We still need to decide how we can choose  $\eta_t$ , an important bound that we will use, is stated by following Lemma

**Lemma 4.** For sequence  $a_1, a_2, \dots, a_n$ ,  $a_i \geq 0$  the following inequality holds.

$$\sum_{i=1}^{n} \frac{a_i}{\sqrt{\sum_{j=1}^{i} a_j}} \le 2\sqrt{\sum_{i=1}^{n} a_i} \tag{11}$$

*Proof.* Let  $x_i = \sum_{j=1}^i a_j$ ,  $x_0 = 0$ , first note the integral equality

$$\int_0^{x_n} \frac{1}{\sqrt{z}} dz = 2\sqrt{x_n} - 2\sqrt{0} \tag{12}$$

This is because  $2\partial_x \sqrt{x} = \frac{1}{\sqrt{x}}$ . Then we can think how we can "compute" the integral in the left side numerically. We can first discretize the interval into small pieces of length  $a_1, a_2, a_3 \cdots$ , then take the right end of the function to approximate the function value in that interval. Note that the right end of function  $\frac{1}{\sqrt{z}}$  is smaller than the functions in the interval, we can get a lower bound of integral:

$$\int_0^{x_n} \frac{1}{\sqrt{z}} dz = \sum_{i=0}^{n-1} \int_{x_{i+1}}^{x_i} \frac{1}{\sqrt{z}} dz \ge \sum_{i=0}^{n-1} \frac{x_{i+1} - x_i}{\sqrt{x_{i+1}}} = \sum_{i=1}^n \frac{a_i}{\sqrt{\sum_{j=1}^i a_j}}$$
(13)

As a special case (take  $a_i = 1$ ), we have  $\sum_{t=1}^T \frac{1}{\sqrt{t}} \leq 2\sqrt{T}$ . If we choose  $\eta_t = \frac{\sqrt{2}B}{G\sqrt{t}}$ , we have

$$Regret(u) \le \frac{(2B)^2}{2\eta_T} + \frac{1}{2} \sum_{t=1}^{T} \frac{\sqrt{2}B}{G\sqrt{t}} G^2 \le 2\sqrt{2}GB\sqrt{T}$$
 (14)

We can also let  $a_i = \|g_t\|^2$ ,  $\eta_t = \frac{\alpha}{\sqrt{\sum_{s=1}^t \|g_s\|^2}}$ 

$$\frac{1}{2} \sum_{t=1}^{T} \eta_t \|g_t\|^2 \le \alpha \sqrt{\sum_{t=1}^{T} \|g_t\|^2}$$
 (15)

The adaptive regret bound is given by

$$Regret(u) \le \frac{(2B)^2}{2\alpha} \sqrt{\sum_{t=1}^{T} \|g_t\|^2} + \alpha \sqrt{\sum_{t=1}^{T} \|g_t\|^2} = \left(\frac{2B^2}{\alpha} + \alpha\right) \sqrt{\sum_{t=1}^{T} \|g_t\|^2}$$
 (16)