

1.5, i). Let  $\alpha: F \Rightarrow G$  be a natural tf.

$$\Rightarrow \forall f: c \rightarrow c',$$

$$\begin{array}{ccc} Fc & \xrightarrow{\alpha_c} & Gc \\ Ff \downarrow & \curvearrowright & \downarrow Gf \\ Fc' & \xrightarrow{\alpha_{c'}} & Gc' \end{array} \quad \text{Now define } H \text{ as}$$

$$H: C \times \mathbb{2} \longrightarrow D.$$

$$\begin{array}{ccc} (c, 0) \longmapsto Fc & (c, 1) \longmapsto Gc \\ f \cdot 1_0 \downarrow \longmapsto \downarrow Ff & f \cdot 1_1 \downarrow \longmapsto \downarrow Gf & \text{and} \\ (c', 0) \longmapsto Fc' & (c', 1) \longmapsto Gc' \end{array}$$

$$\begin{array}{ccc} (c, 0) \longmapsto Fc & \text{Id} \\ f \cdot (0 \rightarrow 1) \downarrow \longmapsto \downarrow Gf \circ \alpha_c = \alpha_{c'} \circ Ff & \\ (c', 1) \longmapsto Gc' \end{array}$$

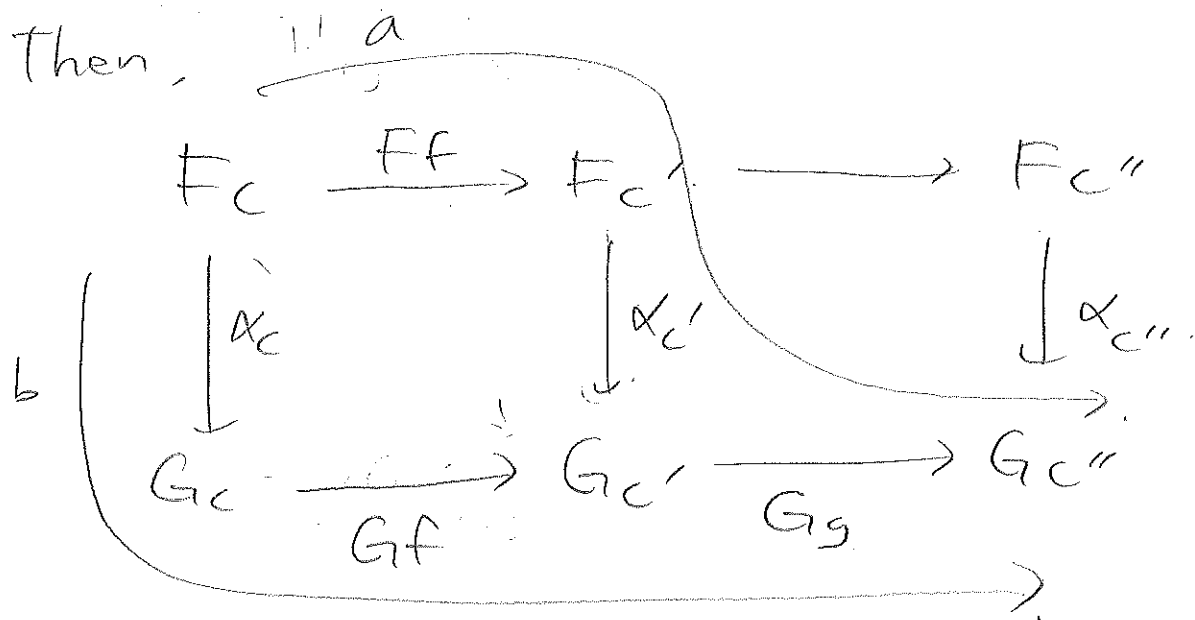
1) It satisfies functoriality.

All we need are 4 cases. ①  $f \cdot 1_0 \xrightarrow{g \cdot (0 \rightarrow 1)}$

$$\text{② } f \cdot (0 \rightarrow 1) \xrightarrow{g \cdot 1_1}$$

$$\text{for } f: c \rightarrow c', g: c' \rightarrow c''$$

$$\text{① } \begin{array}{ccccc} (c, 0) & \xrightarrow{f \cdot 1_0} & (c', 0) & \xrightarrow{g \cdot (0 \rightarrow 1)} & (c'', 1) \\ \downarrow Ff & & \downarrow Fg & & \downarrow G(g \cdot (0 \rightarrow 1)) \\ Fc & \xrightarrow{Ff} & Fc' & \xrightarrow{Gg \circ \alpha_{c'}} & Gc'' \end{array} \quad \text{and} \quad \begin{array}{ccc} (c, 0) & \xrightarrow{f \cdot (0 \rightarrow 1)} & (c', 1) \\ \downarrow F(f \cdot (0 \rightarrow 1)) & & \downarrow G(f \cdot (0 \rightarrow 1)) \\ Fc & \xrightarrow{F(f \cdot (0 \rightarrow 1))} & Gc' \end{array}$$



$a$  is left diagram,  $b$ : right diagram

$\Rightarrow$  Commutes by natural transformation.

Other case is similar.

$$\begin{aligned} \text{And } H((C, \bar{x}) \xrightarrow{1_{(C, \bar{x})}} (C, \bar{x})) &= \begin{pmatrix} F_C \xrightarrow{F1_C} F_C & \bar{x}=0 \\ G_C \xrightarrow{G1_C} G_C & \bar{x}=1 \end{pmatrix} \\ &= 1_{F_C} \text{ or } 1_{G_C} \\ &= 1_{H(C, \bar{x})}. \end{aligned}$$

So  $H$  is a functor.

2)  $H$  satisfies  $\begin{array}{ccccc} C & \longrightarrow & K \times 2 & \longleftarrow & C \\ & \searrow \scriptstyle F & \downarrow \scriptstyle H & \swarrow \scriptstyle G & \\ & & D & & \end{array}$

$$\begin{array}{c} C \mapsto F_C \\ f \downarrow \quad \downarrow Ff \\ C' \mapsto F_{C'} \end{array} = \begin{array}{c} C \mapsto (C, 0) \mapsto F_C \\ f \downarrow \quad \downarrow f \cdot 1_0 \quad \downarrow Ff \\ C' \mapsto (C', -) \mapsto F_{C'} \end{array}$$

and the other way is similar.

1.5. ii.)

$T = \text{Obj}$  : finite set

Mor:  $S \longrightarrow T$

$\Rightarrow \theta: S \longrightarrow P(T) \quad S\text{-}t.$

①  $\theta(\alpha) \cap \theta(\beta) = \emptyset$  when  $\alpha \neq \beta$ .

$$S \xrightarrow{\theta} T \xrightarrow{\phi} U = S \xrightarrow{\psi} U$$

$$S\text{-}t. \quad \alpha \xrightarrow{\psi} \bigcup_{\beta \in \theta(\alpha)} \phi(\beta)$$

$$P(T) \longrightarrow (Fin_*^{op})^{op}$$

$$G: Fin_*^{op} \longrightarrow \Gamma$$

$$S \longrightarrow (S \cup \{s\}, s)$$

$$(S, s) \longrightarrow S \setminus s$$

$$f \downarrow \longmapsto \uparrow f^*$$

$$T \longrightarrow (T \cup \{t\}, t)$$

$$(T, t) \longmapsto T \setminus t$$

$$\theta^{-1}(\beta) = \begin{cases} \alpha & \text{if } \beta \in \theta(\alpha) \\ s & \text{o.w.} \end{cases} \quad f^*: \beta \longmapsto f^*(\beta)$$

First of all,  $\theta^{-1}(\beta)$  is well-defined since no element in  $T$  is contained in two preimage by ①.

Also,  $f^*$  is well-defined since it is a map from  $T \setminus t$  to  $P(S \setminus s)$  with distinct preimage.

To see  $T \subseteq \text{Fin}_*^{\text{op}}$ .

Notes that  $FG: \text{Fin}_*^{\text{op}} \rightarrow T \rightarrow \text{Fin}_*^{\text{op}}$

$$(S, s) \mapsto S \setminus s \mapsto (S \setminus s \cup \{s\}, s)$$

$$f \downarrow \quad \mapsto \quad \uparrow f^{-1} \quad \mapsto \quad \downarrow h$$

$$(T, t) \mapsto T \setminus t \mapsto (T \setminus t \cup \{t\}, t)$$

To figure out  $h$ , notes that

If  $x \in S \setminus s$ , let  $y = f(x)$ .

$$\Rightarrow f^{-1}(y) = \{x \in S \setminus s : f(x) = y\}$$

Thus, for any  $x \in f^{-1}(y)$

$$h(x) = y.$$

Hence,  $f|_{S \setminus s} = h|_{S \setminus s}$  and  $h(S \setminus s) = T \setminus t$ .

Thus, define  $\alpha_{(S, s)}$

$$(S \setminus s \cup \{s\}, s) \xrightarrow{\alpha_{(S, s)}} (S, s)$$

$$\downarrow h \quad \quad \downarrow f$$

as usual incl with

$$(T \setminus t \cup \{t\}, t) \xrightarrow{\quad} (T, t)$$

$$S \setminus s \mapsto s$$

$$T \setminus t \mapsto t$$

Then, if  $x \in S \setminus s$  with  $f(x) = y$ .

$$\alpha_{(T,t)} \circ h(x) = \alpha_{(T,t)}(y) = y.$$

$$f \circ \alpha_{(S,s)}(x) = f(x) = y.$$

and for  $x = s$ ,

$$\alpha_{(T,t)} \circ h(s) = \alpha_{(T,t)}(T \setminus t) = t.$$

$$f \circ \alpha_{(S,s)}(s) = f(s) = t.$$

$\therefore$  diagram commutes.

Likewise

$$\begin{array}{ccccc} GF: T & \longrightarrow & \mathbf{Fin}_*^{op} & \longrightarrow & T \\ s & \longmapsto & (S \cup \{s\}, s) & \longmapsto & s, \\ \theta \downarrow & & \uparrow \theta^{-1} & & \downarrow h \\ T & \longmapsto & (T \cup \{t\}, t) & \longmapsto & T \end{array}$$

S.t. for  $x \in S$ ,  $\theta(x) \in T$ . Hence,  $\forall y \in \theta(x)$

$$\theta^{-1}(y) = x. \Rightarrow h(x) = \text{preimage of } x \text{ under } \theta^{-1} \\ = \theta(x). \quad \text{Monoid}$$

$$\Rightarrow GF = 1_T.$$

$$\text{Hence } T \cong \mathbf{Fin}_*^{op}.$$

In particular,  $\mathbf{Fin}_* \xrightarrow{M^-} \mathbf{Set}$  is a functor in 1.3.2(xi) and  $G$  induces a functor  $\hat{G}: T^{op} \rightarrow \mathbf{Fin}_*$   
 $\Rightarrow T^{op} \xrightarrow{\hat{G}} \mathbf{Fin}_* \xrightarrow{M^-} \mathbf{Set}$  are presheaves on  $T$ .

1.5. III) Any morphism  $f: a \rightarrow b$  and  $\left( \begin{array}{l} \text{fixed morphisms} \\ a \cong a', b \cong b' \end{array} \right)$

$\Rightarrow$  determine  $f': a' \rightarrow b'$  so that

$$\begin{array}{ccc} a \xleftarrow{\cong} a' & a \xrightarrow{\quad} a' & a \xleftarrow{\quad} a' & a \xrightarrow{\quad} a' \\ f \downarrow & f \downarrow & f \downarrow & f \downarrow \\ b \xrightarrow{\quad} b' & b \xrightarrow{\quad} b' & b \xleftarrow{\quad} b' & b \xleftarrow{\quad} b' \end{array}$$

$\cong$

pf) Define  $f' = B \circ f \circ A$

$\Rightarrow B \circ f = f' \circ A^{-1}, B^{-1} \circ f' = f \circ A, f' \circ B^{-1} \circ f' \circ A^{-1} = f$

1.5. IV)  $F: C \rightarrow D$  full and faithful

(i)  $f: c \rightarrow c'$  mon  $M \subset C$  s.t.  $Ff$  is iso in  $D$   
 $\Rightarrow f$  is iso.

(ii)  $x, y \in C$  s.t.  $F_x \cong F_y$  in  $D$

$\Rightarrow x \cong y$  in  $C$

pf) Since  $F$  is full and faithful,

$C(x, y)$  and  $D(F_x, F_y)$  are bijective.

So,  $\exists f, g \in D(F_y, F_x)$  s.t.  $Fg \circ Ff = 1_{F_x}$   
 $Ff \circ Fg = 1_{F_y}$

then  $F(gf) = 1_{F_x}$  Since  $C(x, x)$  and  $D(F_x, F_x)$   
 $F(fg) = 1_{F_y}$   $C(y, y)$  "  $D(F_y, F_y)$  are  
 bijective,  $gf = 1_x, fg = 1_y$ .

Similarly, if  $F_x \cong F_y$ ,  $\exists f \in D(F_x, F_y)$  which is isomorphism. By (i),  $f$  is iso in  $\mathcal{C}$ , thus  $x \cong y$ .

~~(Functor preserve iso)~~  
 Lem 1.3.8 is converse of this statement.

Ex 1.5.v.

$$\begin{array}{ccc} 2 & \longrightarrow & D \\ x \xrightarrow{f} y & & C_x \xrightleftharpoons[f^{-1}]{f} C_y \end{array}$$

It is faithful since  $2(x, y) \hookrightarrow D(C_x, C_y)$ .

but not reflect since  $x \rightarrow y$  in  $2$  is not iso, but in  $D$  it is.

Ex 1.5. vi)

(i) Composite of full, faithful, ess surj. again same.  $F: \mathcal{C} \rightarrow D$ ,  $G: D \rightarrow E$ .

$$\text{pf)} \quad \mathcal{C}(x, y) \xrightarrow{\hookrightarrow} D(C_x, C_y) \xrightarrow{\hookrightarrow} E(DC_x, DC_y)$$

$$\Rightarrow \mathcal{C}(x, y) \xrightarrow{\hookrightarrow} E(DC_x, DC_y)$$

And for essential surj,  $\forall e \in E$ ,  $\exists d \in D$  st.  $e \cong Gd$   
 and for  $d \in D \exists c \in C$  st.  $d \cong Fc$   
 $\Rightarrow e \cong Gd \cong GFc$  imple  $\forall e \in E \exists c$  st.  $GFc \cong e$ .

(ii). It is just from 1.5.9. Thm.

if we assume axiom of choice

Also,  $C \cong C$ ,  $C \cong D \Rightarrow D \cong C$ . trivially.

Ex 1.5.vii) Connected groupoid  $\Rightarrow$  all objects are isomorphic.  
if  $G = G(g, g)$ .

$$\begin{array}{ccc}
 G & \longrightarrow & BG \\
 h & \xrightarrow{\iota_h} & \bullet \\
 f \downarrow \cong & & \downarrow \alpha_h \\
 h' & \xrightarrow{\quad} & \bullet
 \end{array}
 \quad
 \begin{array}{ccc}
 & & \bullet \xrightarrow{\quad} g \\
 & & \downarrow \phi \\
 & & \bullet \xrightarrow{\quad} g
 \end{array}$$

$h \xrightarrow{f} h' \xrightarrow{(\alpha_{h'})^{-1}} g$

$$\phi \in G(g, g)$$

Then,

$$BG \longrightarrow G \longrightarrow BG = 1_{BG}$$

$$\begin{array}{ccc}
 \bullet & \xrightarrow{\quad} & g \xrightarrow{\quad} \bullet \\
 \phi \downarrow & \xrightarrow{\quad} & \downarrow \phi \\
 \bullet & \xrightarrow{\quad} & g \xrightarrow{\quad} \bullet
 \end{array}$$

(or take  $\alpha \cdot 1_{BG} \Rightarrow BG \rightarrow G \rightarrow BG$  as identity)

$$\begin{array}{ccc}
 G & \longrightarrow & BG \longrightarrow G \\
 h & \xrightarrow{\quad} & \bullet \xrightarrow{\quad} g \\
 \psi \downarrow & \xrightarrow{\quad} & \downarrow \psi \\
 h' & \xrightarrow{\quad} & \bullet \xrightarrow{\quad} g
 \end{array}$$

$g \rightarrow h \xrightarrow{\psi} h' \rightarrow g$



we need to find iso. satisfying.

Thus,

$$\begin{array}{ccc}
 g & \xrightarrow{\alpha_h} & h \\
 \downarrow & \curvearrowright & \downarrow \psi \\
 g \xrightarrow{\psi} h \xrightarrow{\psi} g & & h' \\
 \downarrow & & \\
 g & \longrightarrow & h'
 \end{array}$$

And take

$$\alpha_h = g \rightarrow h$$

$$\alpha_{h'} = (h' \rightarrow g)^{-1}$$

Since such  $\alpha_h, \alpha_{h'}$  is fixed by construction of  $G \rightarrow BG$ , it is well-def and satisfy the commuting diagram

Ex 1.5 viii) Later

Ex 1.5 ix) Any category equiv to locally small category is locally small.

Pf) Let  $C \xrightleftharpoons[G]{F} D$ , equiv of category

By Thm 1.5.9,  $F, G$  : full, faithful.

$\Rightarrow$  If assume  $D$  locally small,

then  $(x, y)$  is a set since it is biject to a set  $D(x, y)$ .  $\Rightarrow C$  is locally small.

1.5. viii).

Affine : Obj = affine planes,  $A^2 K$   
where  $K$  is a field.

Mor = affine linear map.

$L \subset L$  :  $L$  : linear isomorphism

Proj : Obj = projective planes  $(\mathbb{P}^2 K, l)$   
with line of infinity.  $C$  : constant function

Mor = projective linear isomorphism

$F : \text{Proj} \longrightarrow \text{Affine}$

$(\mathbb{P}^2 K, l) \longmapsto A^2 K$  by deleting  $l$  in  $\mathbb{P}^2 K$

Notes that  $\mathbb{P}^2 K = K^2 \sqcup (K^1 \sqcup P^+)$

where  $K^2 = \{ [x:y:z] : z \neq 0 \} \longmapsto \{ (\frac{x}{z}, \frac{y}{z}) \in A^2 K \}$

$K = \{ [x:y:0] : y \neq 0 \}$

$P^+ = \{ [x:0:0] \}$

with  $l = K \sqcup P^+$ .

Thus,  $G : \text{Affine} \longrightarrow \text{Proj}$

is embedding  $A^2 K$  into  $(\mathbb{P}^2 K, l)$

by identifying  $A^2 K$  as  $K^2$  part.

$(x, y) \longmapsto [x:y:1]$

Claim 1:  $F, G$  maps pt to pt, line to line.  
(except  $\ell$ ).

Let  $L$  be a line in  $\mathbb{P}^2_K$  generated by

$$\vec{a} = [a_0 : a_1 : a_2], \quad \vec{b} = [b_0 : b_1 : b_2]$$

Then,  $L = \{ u\vec{a} + v\vec{b} : u, v \in K, (u, v) \neq (0, 0) \}$

Case 1: If  $\vec{a}, \vec{b} \in K^2$ ,  $F$  maps  $u\vec{a} + v\vec{b}$  into  $\left( \frac{u\frac{a_0}{a_2} + v\frac{b_0}{b_2}}{u+v}, \frac{u\frac{a_1}{a_2} + v\frac{b_1}{b_2}}{u+v} \right)$

$$\text{Thus, } F(L) = \left\{ t \left( \frac{a_0}{a_2}, \frac{a_1}{a_2} \right) + (1-t) \left( \frac{b_0}{b_2}, \frac{b_1}{b_2} \right) \right\}$$

Since  $\frac{u}{u+v}, \frac{v}{u+v}$  can be mapped into  $t$  and

$1-t$ , for any  $t \in K$ ,

$$= \left\{ \left( \frac{b_0}{b_2}, \frac{b_1}{b_2} \right) + t \left( \frac{a_0}{a_2} - \frac{b_0}{b_2}, \frac{a_1}{a_2} - \frac{b_1}{b_2} \right) \right\}$$

Hence  $F(L)$  is a line in  $\mathbb{A}^2_K$ .

Case 2: If one of  $\vec{a}, \vec{b}$  is in  $\ell$ .

Then, by adding suitable  $u\vec{a}$  to  $\vec{b}$ ,

We can change this as a case 1.

Case 3: Both  $\vec{a}, \vec{b}$  are in  $\ell$

Then,  $L = \ell$  and  $F$  drops  $\ell$ .

Now, let  $L$  be a line in  $\mathbb{A}^2_K$ .

$$\Rightarrow L = \{tA + (1-t)B : t \in K\}$$

for some  $A = (a_0, a_1)$ ,  $B = (b_0, b_1)$

$$\Rightarrow G(L) = \left\{ t [a_0 : a_1 : 1] + (1-t) [b_0 : b_1 : 1] \right. \\ \left. t \in K \right\}.$$

To see that  $G(L)$  is contained in a line in  $\mathbb{P}^2_K$ , let  $L'$  be a line gen by

$$[a_0 : a_1 : 1], [b_0 : b_1 : 1]. \text{ Then } L' \supseteq G(L)$$

So  $G(L)$  matches with only one line  $L'$ .

$$\text{Thus, } FG(L) \subseteq F(L') = L$$

$$\text{and } GF(L) = G\left(t\left(\frac{a_0}{a_2}, \frac{a_1}{a_2}\right) + (1-t)\left(\frac{b_0}{b_2}, \frac{b_1}{b_2}\right)\right)$$

$$\subseteq L' \text{ where } L' \text{ is gen by}$$

$$\left[\frac{a_0}{a_2}, \frac{a_1}{a_2}, 1\right], \left[\frac{b_0}{b_2}, \frac{b_1}{b_2}, 1\right]$$

$$= [a_0 : a_1 : a_2], [b_0 : b_1 : b_2]$$

$$\Rightarrow = L.$$

Hence,  $FG(\mathbb{A}^2_K) = \mathbb{A}^2_K$  since they

$GF(\mathbb{P}^2_K) = \mathbb{P}^2_K$  preserves pt and lines.

Thus, equivalence is induced by  $\alpha$  as identity for each object.

1.5. xi) See Hame, P. "Examples of some properties of functors"

	Full	Faithful	Essentially Surj.
$Ab \rightarrow Gp$	✓	✓	✗
$Rng \rightarrow Ab$	✗	✓	✗
$Rng \xrightarrow{(-)^*} Gp$	✗	✗	✗
$Rng \rightarrow Rng$	✗	✓	✗
$Field \rightarrow Rng$	✓	✓	✗
$M_n(R) \rightarrow Ab$ dep on R		✓	✓

(b)  $Ring \rightarrow Ab$ : ① Not full since. No ring homo  $\mathbb{Z}/n\mathbb{Z} \rightarrow \mathbb{Z}$ .

Since  $f\left(\underbrace{1+\dots+1}_n\right) = 0$  but  $f(1)+\dots+f(1)=n$   
Contradiction.

(However,  $Rng$  on  $Ab$   $\exists$  trivial homo.)

② Not essentially Surj:

Let  $\mathbb{Z}(p^\infty) := \{e^{2\pi i m/p^n} : m, n \in \mathbb{N}\}$

Prüfer  $p$ -group.

If  $U(R) \cong \mathbb{Z}(p^\infty)$  then  $1 \mapsto a \in \mathbb{Z}(p^\infty)$

$\Rightarrow |1| = n$  in  $R \Rightarrow \forall r \in R, nr = 0$

Since  $r(1+\dots+1) = r \cdot 0 = 0$

Thus every element in  $U(R)$  has order at most  $n$ . But  $\mathbb{Z}(p^\infty)$  has element with greater order.  $\square$

$$(c) (-)^{\times} \text{Rings} \rightarrow \text{Grp.}$$

① Not full.

Since  $\mathbb{Z}$  is initial in  $\text{Rings}$ ,  $\forall R \in \text{Rings}$   
 $\exists! f: \mathbb{Z} \rightarrow R \Rightarrow |\text{Rings}(\mathbb{Z}, R)| = 1$

But  $\mathbb{Z}^{\times} \cong \mathbb{Z}/2$ . Take  $R = \mathbb{F}_p$   
 finite  $p$ -field.

$$\Rightarrow R^{\times} \cong \mathbb{Z}/(p-1)$$

And  $|\text{Grp}(\mathbb{Z}^{\times}, R^{\times})| = 2$  since 0

and  $\mathbb{Z}/2 \rightarrow \mathbb{Z}/(p-1)$  by  $1 \mapsto \frac{p-1}{2}$   
 are the homomorphisms.

② Not faithful.

Claim 1:  $\phi: K[t] \rightarrow K[t]$  automorphism  
 fixing  $K \Rightarrow \phi(t) = at + b$

pf)  $\phi(t)$  cannot be degree more than one  
 otherwise it is not automorphism.

Also,  $\phi(t)$  cannot have degree 0 by sneaky  $\frac{1}{a}$

And all  $\phi(t) = at + b$  is auto since

$$\phi\left(\frac{1}{a}t - b\right) = t.$$

□

Thus by Claim 1,

$\text{Ring}(K[t], K[t])$  is determined by

$$\{at + b : a \in K^{\times}, b \in K\}$$

And  $K[t]^{\times} = K^{\times}$ .

But  $\text{Grp}(K^{\times}, K^{\times})$  is determined by

generators of  $K^{\times}$  (since  $K^{\times}$  is cyclic)

and any map in  $\text{Ring}(K[t], K[t])$

induces map  $K^{\times} \rightarrow K^{\times}$  fixing all  $K^{\times}$ .

Thus it is not faithful.

③ Not ess. surj.

[Pearson, Schneider, Lam]

Not every cyclic gp is isomorphic to  
the gp of units of some ring.

$$(\mathbb{Z}/5 \not\cong R^{\times} \quad \forall R \in \text{Rings})$$

(d)  $R_m \rightarrow R_n$ .

① Not full. : Zero homomorphism  
is not a homomorphism  
in  $R_m$

② Faithful : Any unital ring homo  
is also ring homo  
and distinct unital ring homos  
are " ring homos.

③ Not ess. surj. : Ring without mult.  
identity is not  
iso to unital ring.

(e)  $\text{Field} \leftrightarrow R_m$

① fully faithful.

Since every field homo is  
ring homo.

② Not ess. surj. : Not every ring  
is a field.



(7)  $U: \text{Mod}_R \rightarrow \text{Ab}$ .

① faithful: Any distinct  $R$ -homo  
is also distinct  $\text{gp}$  homo.

(Since distinction determined  
by value, not property.)

② Ess surj:  $\forall A \in \text{Ab}$ , make  
trivial  $R$ -module structure  
s.t.  $\forall a \in A, \forall r \in R$   
 $ra = a$ .

③ Not necessarily full.

$\text{End}(R) \cong R$  in  $\text{Mod}_R$

but  $\text{End}(R) \not\cong R$  in  $\text{Ab}$  in general.

(If  $R \neq \mathbb{Z}$ , then it is full.)