

0. Preface.

"Natural" \longleftrightarrow defined without arbitrary choice

ex) $V \xrightarrow{\cong} V^*$: need choice of the basis.

$V \approx V^{**}$! doesn't needed.

Eilenberg MacLane 45': Natural Transformation
Function (as source / target)
Category (")

Objectives

- ① Categorifying math. objects \Rightarrow isomorphism
- ② study itself \Rightarrow self-dual

In Chapter 1, ① Def ② Duality ③ Functor. Today!
④ Naturality ⑤ Equivalence (⑥ Diagram chase)
⑦ 2-categories

Def 1.1.1 C : Category consists of

- $\text{Obj}(C)$: Collection of objects, X, Y, Z, \dots
- $\text{Hom}(C)$ " morphisms f, g, h, \dots

(or $\text{Mor}(C)$)

Each morphism has domain and codomain.

1) Each morphism has domain and codomain.

2) Each object X has 1_X : identity morphism

3) Two morphism $f: X \rightarrow Y$ $g: Y \rightarrow Z$ s.t.
 $\text{cod}(f) = \text{dom}(g)$, $gf: X \rightarrow Z$ exists.

a) (unital) $\forall f: X \rightarrow Y \in \text{Hom}(C), f \cdot 1_X = f = 1_Y \cdot f$

b) (associates) for any composable triple,
 $f(gh) = (fg)h$.

In abstract category : Morphism = arrow or map.
(they're not functions)

Rem 1.1.5 Size issue.
Russell's Paradox \Rightarrow No collection contains itself.

To avoid,

Def 1.1.6. Small category : $\text{Hom}(C)$ is a set.
(Then, from $\text{Obj}(C) \xrightarrow{x} \text{Hom}(C) \xrightarrow{1_x} \text{Obj}(C)$ $\text{Obj}(C)$ is also a set.)
 $\Rightarrow \text{Hom}(C) \xrightleftharpoons[\text{Cod}]{\text{Dom}} \text{Obj}(C)$ are functions.

But none of examples of concrete category are small.

Def 1.1.7. Locally small category : $\forall X, Y \in \text{Obj}(C)$.
 $\text{Hom}(X, Y) = \text{Hom}(X, Y)$ is a set. (set of all morphisms ~~not~~ from X to Y)
 $= \text{Hom}(X, Y)$

Q: When is one thing the same as another thing?
 \Rightarrow iso. (actually we will define notion of equivalence)

Def: 1.1.9. Iso in category C is a morphism $f: X \rightarrow Y$
for which $\exists g: Y \rightarrow X$ st $fg = 1_Y$ $gf = 1_X$.

In this case write $X \cong Y$.
Endo: $f: X \rightarrow X$. Auto: endo + iso.

Ex 1.1.10.
(Set) bijection (Grp, Rng, Fied, Mod_p) bijective
(Top): Homeo (Htpy) homotopy equiv. (P, \leq) identity by antisymmetry

Q: In a concrete category, is even iso induced by bijection of underlying set?
A: Lem 5.6.1 (Yes!)

Def 1.1.11. Groupoid: a category in which every morphism is iso.

Cf) In algebra, groupoid: a group changing binary operation to partial function.
i.e. mult is not def on all ~~two~~ objects pair-set

Ex 1.1.12 $G(\text{Group})$: groupoid with 1 object.
(it is def of gr in category theory.)

$\pi_1(X)$: fundamental groupoid.
Obj: points in X Mor: Endpt preserving homotopy classes of paths.

Def) Subcategory D of C .

$\text{Obj } D \subseteq \text{Obj } C$, $\text{Mor } D \subseteq \text{Mor } C$ s.t.

- 1) $\text{Obj } D$ contains any domain or codomain of $f \in \text{Mor } D$
- 2) $\text{Mor } D$ " any identity morphism of $X \in \text{Obj } D$
- 3) Closed under composition.

Ex) $\text{CRing} \subseteq \text{Ring} \subseteq \text{Rng}$
(commutative) (unital) (maybe nonunital)

Lem 1.1.13 Any category C contains maximal groupoid.

Pf) Show collection of isomorphisms of C is subcategory

Ex). Fin : Obj: finite set, Mor: functions
 Fin_{iso} : " Mor: bijections

Fin_{iso} is maximal groupoid of Fin

(참고: Objects X can be identified with identity 1_X .

So cat def by morphism.

What we care: Morphism

Ex 1.1.3. (concrete category)

Name	Obj	Mom	Composite. is
Set	set X	function f	composition of f .
Top	top. space X	cts function f .	"
Set * , Top *	(X, x) \uparrow $x \in X$ base pt	base pt preserving \sim	"
Group	groups	gp homo	"
Ring.	rings	ring " "	"
Field:	fields	field " "	"
Mod $_R$	left R -module.	module " "	"
$_R$ Mod	right " "	Module π	G -structure. morphisms preserve structure.
Vect $_K =$ Mod $_K$	with $R=K$		
Ab = Mod $_Z$	with $R=Z$		
Graph	graphs	graph morphism (sending vertex to vertex edge to edge preserving incidence relation)	"
Digraph	dir " "		
Man	manifolds	smooth map.	
Meas	measurable space	measurable function	"
Poset	Posets	order preserving function	"
Ch $_R$	chain cpx of R -module	chain homo.	

Concrete Category (Precise: 1.6.17)

Obj has underlying set
Morph are functions between underlying set.

1.1.4 (Abstract Category)

Name	Obj	Morph	Composition
------	-----	-------	-------------

① Mat_R (R : unital ring) ↳ needed for identity morph.	$ P = \{1, 2, 3, \dots\}$	$A: n \rightarrow m$ is $(m \times n)$ R -valued matrix	matrix multiplication
--	---------------------------	---	-----------------------

② BG G : group.	$\{ \cdot \}$	$g \in G$	group product
------------------------	---------------	-----------	---------------

③ (P, \leq) P : poset	P	$x \leq y$ $\Leftrightarrow x \rightarrow y$ (identity from reflexivity)	clear by transitivity
------------------------------	-----	--	-----------------------

So we can do for pre order (transitivity and reflexivity)

④ Ordinals

①	\emptyset	\emptyset
$\mathbb{1}$	$\{\emptyset\}$	$\mathbb{1}_{\emptyset}$
$\mathbb{2}$	$\{\emptyset, \{\emptyset\}\}$	$\mathbb{1}_{\emptyset}, \mathbb{1}_{\{\emptyset\}},$ $\emptyset \rightarrow \{\emptyset\}$ ($0 \rightarrow 1$)

ω	$\{0, 1, 2, \dots\}$	$0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow \dots$
----------	----------------------	---

all morphisms
 $n \rightarrow m$
equals
 $n \rightarrow n + (1 \rightarrow n + 2) \dots$

⑤ A (set).	A	$\mathbb{1}_a \quad \forall a \in A$
-----------------	-----	--------------------------------------

"Discrete Category": Every morphism is an identity

⑥ Htop Htop^*) Measure	top space (X, \mathcal{X}) measure spaces	homotopy class base preserving equiv class of measurable function (= a.e.)
---	---	--

Exercise: 1.1 i) morphism can have at most 1 inverse.
 p.f) If $g, h: x \rightarrow y$ are inverse of $f: y \rightarrow x$

then, $g = g \circ I_x = g \circ f \circ h = I_y \circ h = h$.

1.1. ii). Maximal groupoid. \rightarrow done.

1.1. iii) Slice Category.

Let C : category $a \in \text{obj } C$.

$$C/C : \text{obj} := \bigcup_{x \in \text{obj } C} \text{Hom}(C, x)$$

under

$$\text{Mor} := \text{Mor}(f: C \rightarrow x, g: C \rightarrow y)$$

$$= \{ h: x \rightarrow y : \begin{array}{ccc} & C & \\ f \swarrow & \Delta & \searrow g \\ x & \xrightarrow{h} & y \end{array} \text{ commutes} \}$$

i.e. $g = hf$.

$$C/C : \text{obj} := \bigcup_{x \in \text{obj } C} \text{Hom}(x, C)$$

over

$$\text{Mor}(f: x \rightarrow C, g: y \rightarrow C)$$

$$= \{ h: x \rightarrow y : \begin{array}{ccc} & C & \\ f \nearrow & \Delta & \nwarrow g \\ x & \xrightarrow{h} & y \end{array} \text{ commutes} \}$$

i.e. $f = gh$.

Duality

Def 1.2.1. From C , C^{op} is opposite category

$$\bullet \text{Obj}(C^{op}) = \text{Obj}(C)$$

$$\bullet f^{op} \in \text{Hom}(C^{op}) \text{ for each } f \in C$$

$$\text{with } f^{op}: \text{cod}(f) \rightarrow \text{dom}(f)$$

Then, from structure of C ,

$$1) 1_X^{op} \text{ is id in } C^{op}$$

$$2) f^{op}: X \rightarrow Y \quad g^{op}: Y \rightarrow Z$$

$$\Rightarrow g^{op} f^{op}: X \rightarrow Z \quad \longleftrightarrow fg: Z \rightarrow X$$

Thus, C^{op} is category iff C is category

Ex 1.2.2. Mat_R^{op} : f^{op} is transpose of f

$$(P, \leq)^{op} : x \rightarrow y \iff y \leq x$$

$$\omega^{op} : \dots \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow 0$$

$$(BG)^{op} \cong B(G^{op}) \text{ where } G^{op} : \text{opposite gp.}$$

gp with right
multiplication

$$\text{i.e. } f \cdot g = gf \text{ in } G$$

Thus, any statement of proof about C also applies
to opposite: "dual theorem"

Ex).

Lemma 1.2.3 For any C , TFAE.

(1) $f: X \rightarrow Y$ is iso in C .

(2) $f_*: C(C, X) \rightarrow C(C, Y)$ is bijection $\forall C \in \text{Obj } C$
 $g \mapsto fg$.

(3) $f^*: C(Y, C) \rightarrow C(X, C)$ is bijection $\forall C \in \text{Obj } C$
 $g \mapsto gf$.

f_* : post-composition, f^* : pre-composition.

Pf)

(i) \Rightarrow (ii) f is iso, $\Rightarrow \exists g: Y \rightarrow X$ inverse of f .

Thus, $g_* f_*$ and $f_* g_*$ are identity $\forall C \in \text{Obj } C$.

(ii) \Rightarrow (i) Take $C = Y$. $\exists g \in C(Y, X)$. s.t.

$$f_*(g) = 1_Y \Rightarrow fg = 1_Y$$

Also,

$$f_*(gf) = fg = 1_Y f = f$$

$$f_*(1_X) = f = f$$

By bijectivity, $1_X = gf$.

(i) \Leftrightarrow (iii) Apply (i) \Leftrightarrow (ii) on C^{op} to get.

$f^{op}: Y \rightarrow X$ is iso $\Leftrightarrow f_*^{op}: C^{op}(C, Y) \rightarrow C^{op}(C, X)$
 $g^{op} \mapsto f^{op} g^{op}$
is bijection $\forall C \in C$.

Now, $C^{op}(C, Y) \leftrightarrow C(Y, C)$ $C^{op}(C, X) \leftrightarrow C(X, C)$

$g^{op} \mapsto g$, $f^{op} g^{op} \mapsto gf$.

Thus, f_*^{op} sends g to gf . $\Rightarrow f_*^{op} = f^*$.

① 150

$$a + 0b + 2c = 3$$

$$a + \mathbb{N} + \mathbb{O} = \mathbb{N}$$

2000 2001


$$\therefore (C = 6)$$


$$\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} \begin{bmatrix} a & b & c \\ d & e & f \\ g & h & i \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix}$$

$$2 \times 3 \quad 3 \times 3 \quad 3 \times 2$$

$$= \begin{bmatrix} a & b & c \\ d & e & f \end{bmatrix}$$

$$\begin{array}{ccc} \left[\right] & A & \left[\right] \\ a \times k & \cancel{k \times k} & n \times b \end{array}$$

~~~~

~~~~

$\Rightarrow f^{\text{op}}: Y \rightarrow X$ is iso $\Leftrightarrow f^*: (C(Y, C) \rightarrow (C(X, C)$
 bij. $\forall c \in C$.
 \Downarrow
 $f: X \rightarrow Y$ is iso. done.

Def: 1.2.7. $f: X \rightarrow Y$ morphism is

(1) monomorphism if $\forall h, k: W \Rightarrow X$ for some W ,
 mono (mon)
 monic (adj)
 $fh = fk \Rightarrow h = k$

(left cancellable)

(2) epimorphism if $\forall h, k: Y \Rightarrow W$ for any W ,
 epi
 epic
 $hf = kf \Rightarrow h = k$

(right cancellable)

f is mono. $\Leftrightarrow f_*: (C(C, X) \rightarrow (C(C, Y) \forall c \in C$.
 injective

" epi \Leftrightarrow

" surjective

Ex 1.2.8. $f: X \rightarrow Y$ mono in Set. take $x \in X$.

Let $\{x\} \xrightarrow[h]{g} X$, suppose $fg = fh \Rightarrow g = h$.

Thus f is injective.

Ex 1.2.9. (split epi / split mono)

s : section
 (right inverse)

$X \xrightarrow{s} Y \xrightarrow{r} X$

s.t. $rs = 1_X$. Then

r : retraction
 (retract)
 (left inverse)

In this case,

(s is mono)
 (r is epi.)

s : split mono

r : split epi.

Exercise 1.2.ii) (i). If f is split epimorphism, then
 $\exists s: y \rightarrow x$ s.t. $fs = 1_y$. Fix $c \in C$,
 let $g \in (C, y)$. $\Rightarrow sg \in (C, x)$
 $\Rightarrow f_*(sg) = fsg = 1_y g = g$. surjective.
 Conversely, take $C=y$. $\exists s \in (C, x)$ s.t.
 $fs = 1_y$. $\therefore f$ is split epi.
 1.2.ii) ii). In C^{op} , apply (i); we set
 $f^{op}: y^{op} \rightarrow x^{op}$ is split epi $\Leftrightarrow \forall c \in C^{op}$, $f_*^{op}: C^{op}(C, y) \rightarrow C^{op}(C, x)$
 is surj.
 \Uparrow \Downarrow
 $f: x \rightarrow y$ is split mono. $\forall c \in C$, $f^*: C(y, c) \rightarrow C(x, c)$
 is surj.
 1.2.iii) done.
 1.2.iv) What are monomorphisms of Field?
 ans) Every morphism is monomorphism.

$$\begin{array}{ccc} E & \xrightarrow{h} & F \\ & \searrow k & \downarrow g \\ & & G \end{array} \Rightarrow g \text{ is injective function, then}$$

$$g(h(x)) = g(k(x)) \Rightarrow h(x) = k(x)$$

But maybe not all epi is surjection. (e.g. $G \rightarrow E \Rightarrow E$ sending distinct into
 of f but not fixing G)

Ex 1.2.vi) Let $f: x \twoheadrightarrow y$. $\exists g: y \rightarrow x$
 s.t. $fg = 1_x$. By Lemma 1.2.ii) iii) f is epi.
 Apply it on C^{op} . Then $f^{op}: y^{op} \twoheadrightarrow x^{op}$, $g^{op}: x^{op} \rightarrow y^{op}$
 s.t. $f^{op}g^{op} = 1_{x^{op}}$. By above, f^{op} is iso.
 \Rightarrow in C , $f: x \twoheadrightarrow y$, $g: y \rightarrow x$ s.t. $gf = 1_y$
 then f is iso.

Ex 1.2.i).
By def, $C/C^{op} : Obj = \bigcup_{x \in Obj C} Hom_{op}(C, x)$

$$Mor_{op}(f^{op} : C \rightarrow x, g^{op} : C \rightarrow y) = \left\{ h^{op} : x \rightarrow y \text{ s.t. } \begin{array}{ccc} & C & \\ f^{op} \swarrow & \textcircled{2} & \searrow g^{op} \\ x & \xrightarrow{h^{op}} & y \end{array} \right\}$$

By duality, its opposite category is

$$(C/C^{op})^{op} : Obj = \bigcup_{x \in Obj C} Hom_{op}(C, x)$$

$$Mor_{op,op}(f^{op} : C \rightarrow x, g^{op} : C \rightarrow y)$$

$$= \left\{ (h^{op})^{op} : (h^{op}) \in Mor(g^{op} : C \rightarrow y, f^{op} : C \rightarrow x) \right\}$$

(Forget about commuting diagram: it is just abstract opposite.)

Thus, $(h^{op})^{op} \in Mor_{op,op}(f^{op} : C \rightarrow x, g^{op} : C \rightarrow y)$

$$\Leftrightarrow h^{op} \in Mor_{op}(g^{op} : C \rightarrow y, f^{op} : C \rightarrow x)$$

$$\Leftrightarrow \begin{array}{ccc} & C & \\ f^{op} \swarrow & \textcircled{2} & \searrow g^{op} \\ x & \xleftarrow{h^{op}} & y \end{array} \Leftrightarrow \begin{array}{ccc} & C & \\ f \swarrow & \textcircled{2} & \searrow g \\ x & \xrightarrow{h} & y \end{array}$$

Thus by sending $g \mapsto g^{op}$ for object,

$$h \mapsto (h^{op})^{op}$$

We can identify C/C by $(C/C^{op})^{op}$.

Ex 1.2.10 Not all epi or monic is surj or inj.

ex) $f: \mathbb{Z} \hookrightarrow \mathbb{Q}$ canonical inclusion.

1.2.v)
 f is monic since f is injective ($f(h(x)) = f(k(x)) \Rightarrow h(x) = k(x)$)
 f is epic since ($h(1) = k(1) \Rightarrow h(b) \cdot h(\frac{1}{b}) = k(b) \cdot k(\frac{1}{b})$
 $\wedge h(\frac{1}{b}) = h(b)^{-1}, k(\frac{1}{b}) = k(b)^{-1}$
 since $h(b) = k(b) \text{ by } f, h(b)^{-1} = k(b)^{-1}$
 $\Rightarrow h(\frac{1}{b}) = k(\frac{1}{b})$)

But f is not ring isomorphism! it is not bijective

Lemma 1.2.11. Use \twoheadrightarrow monic \rightarrow epic

(i) $f: x \twoheadrightarrow y, g: y \twoheadrightarrow z \Rightarrow g \circ f: x \twoheadrightarrow z$

(ii) $f: x \rightarrow y, g: y \rightarrow z$ st. $g \circ f: x \twoheadrightarrow z$
 $\Rightarrow f: x \twoheadrightarrow y$

pf) (i) : obvious (ii). suppose $w \xrightarrow[k]{h} x$ st. $fh = fk$

Then $gfh = gfk \Rightarrow h = k$

Dually,

(i) $f^{op}: y \twoheadrightarrow x, g^{op}: z \twoheadrightarrow y \Rightarrow f^{op} \circ g^{op}: z \twoheadrightarrow x$

(ii) $f^{op}: y \rightarrow x, g^{op}: z \rightarrow y$ and $f^{op} \circ g^{op}: z \twoheadrightarrow x$

$\Rightarrow f^{op}: y \twoheadrightarrow x$

Peleteer op.

Exercise 1.2.vii) Define sup, inf on poset (P, \leq)

Categorical sense, i.e. dual statement define inf.

Def: $x \in P$ is sup if $\forall y \in P, y \rightarrow x$ exists

(" inf " $x \rightarrow y$)

If x, y are sup, $x \rightleftharpoons y \Rightarrow$ by anti symmetry $x = y$

Functoriality.

Def. 1.3.1. $F: C \longrightarrow D$ functor (morphism of categories)

$$\textcircled{1} F_c \in D \quad \forall c \in C$$

$$\textcircled{2} Ff: F_c \longrightarrow F_{c'} \in D. \quad \forall f: c \longrightarrow c' \in C.$$

Satisfying "functoriality axioms"

$$a) f, g \text{ composable pair, } Fg \circ Ff = F(g \circ f)$$

$$b) \forall c \in C, F(1_c) = 1_{F_c}$$

Ex 1.3.2.

$$1) P: \text{Set} \longrightarrow \text{Set} \quad A \longmapsto P A \quad (f: A \longrightarrow B) \longmapsto f_*: P A \longrightarrow P B$$

2) Forgetful functor for (concrete) categories.

$$U: \text{Group} \longrightarrow \text{Set} \quad G \longmapsto G \quad f \longmapsto f \text{ as function.}$$

as grp as set as hom

Similar for Rng , Field , Top , ...

$$U, E: \text{Graph} \longrightarrow \text{Set} \quad G \longmapsto \begin{matrix} V(G) \\ E(G) \end{matrix} \quad f: G \longrightarrow H \longmapsto \begin{cases} V(G) \longrightarrow V(H) \\ E(G) \longrightarrow E(H) \end{cases}$$

$$V \sqcup E: \quad \quad \quad \begin{matrix} V(G) \sqcup E(G) \\ \uparrow \\ \text{disj} \end{matrix} \quad \quad \quad \begin{matrix} V(G) \sqcup E(G) \\ \longrightarrow V(H) \sqcup E(H) \end{matrix}$$

3) Another forgetful functor:

$$\text{Mod}_R \longrightarrow \text{Ab} \longleftarrow \text{Group}$$

$$\text{Field} \longleftarrow \text{Rng} \nearrow$$

$$4) \begin{matrix} \text{Rng} & \longrightarrow & \text{Set}^* \\ \text{Group} & \longrightarrow & \end{matrix}$$

$$R \longmapsto (R, e) \quad e: \text{identity}$$

functorial
because
homomorphism
preserves identity.

$$(5). \text{Top} \longrightarrow \text{Htpy} \quad (\text{Top}_* \longrightarrow \text{Htpy}_*)$$

$$X \longmapsto X \quad \text{base pt preserving} \dots$$

$$f \longmapsto \tilde{f} \text{ homotopy class of cts function.}$$

$$(6) \pi_1 : \text{Top}_* \longrightarrow \text{Group} \quad X \longmapsto \pi_1(X)$$

$$f : (Y, x) \longrightarrow (Y, y) \longmapsto f_* : \pi_1(X, x) \longrightarrow \pi_1(X, y)$$

"Fundamental group is homotopy invariants"

$$\Rightarrow \exists \text{ functor } \text{Htpy}_* \longrightarrow \text{Group} \quad \Sigma \dagger.$$

$$\pi_1 = \text{Top}_* \longrightarrow \text{Htpy}_* \longrightarrow \text{Group}.$$

$$(7) \pi_1 : \text{Top} \longrightarrow \text{Groupoid} \quad X \longmapsto \pi_1(X), \quad f \longmapsto f_*$$

since cts preserves path and path homotopy.

$$(8) \forall n \in \mathbb{Z}, \quad \begin{matrix} \mathbb{Z}_n \\ B_n \\ H_n \end{matrix} : C_n \begin{matrix} \longrightarrow \\ \longrightarrow \\ \longrightarrow \end{matrix} \text{Mod } R$$

$$n\text{-cycle } \mathbb{Z}_n C_n = \ker (d : C_n \longrightarrow C_{n-1})$$

$$n\text{-bdy } B_n C_n = \ker (d : C_{n+1} \longrightarrow C_n)$$

$$n\text{-homology } H_n C_n = \mathbb{Z}_n C_n / B_n C_n.$$

If you collect for all $n \in \mathbb{Z}$,

$$\begin{matrix} \mathbb{Z}_* \\ B_* \\ H_* \end{matrix} : C_* \begin{matrix} \longrightarrow \\ \longrightarrow \\ \longrightarrow \end{matrix} \text{Gr Mod } R \quad \text{graded } R\text{-module}$$

$$\text{ex) } \mathbb{Z}_* : C_* \longmapsto \bigoplus_{n \in \mathbb{Z}} \mathbb{Z}_n C_n$$

Singular homology of Top space

$$\text{Top} \longrightarrow C_n R \xrightarrow{H_*} \text{Gr Mod } R$$

(ix) $F: \text{Set} \rightarrow \text{Group}$.

$X \mapsto \text{Free gp gen by } X.$ $f \mapsto \text{induce } \downarrow \text{ sp home.}$

example of "free functor".

x) Euclid_* $\text{Obj}: (\mathbb{R}^n, a)$

$\text{Mor}: f: (\mathbb{R}^n, a) \rightarrow (\mathbb{R}^m, b)$ differentiable

consists of $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$ "

and $f(a) = b$.

Let $D: \text{Euclid}_* \rightarrow \text{Mat } \mathbb{R}$

$(\mathbb{R}^n, a) \mapsto n$

$f: (\mathbb{R}^n, a) \rightarrow (\mathbb{R}^m, b) \mapsto n \xrightarrow{df_a} m$

df_a : Jacobian
 $m \times n$ matrix

(Say correction on $\text{Mat}_\mathbb{R}$:)

It satisfies functoriality axiom due to chain rule.

xii). Fin_* : $\text{Obj}: (\{x_1, \dots, x_n, x_{n+1}\}, x_i)$

Mor : function preserving basept.

$M^+ : \text{Fin}_* \rightarrow \text{Set}$

Let $M^+ := ([n] \cup \{a\}, \{a\})$ (M : Commutative monoid)

$M^{n+} = M^{x^n}$ cartesian product (M^{0+} : Singleton)

$$\forall f: m_+ \longrightarrow n_+, \text{ def } M^f: M^m \longrightarrow M^n$$

$$(a_1, \dots, a_m) \longmapsto (b_1, \dots, b_n)$$

where
$$b_i = \begin{cases} \prod_{j \in f^{-1}(i)} a_j & \text{if } f^{-1}(i) \neq \emptyset \\ 1 & \text{if empty} \end{cases}$$

Then, M^f preserves unit. So,

$$\begin{array}{ccc} m_+ \rightarrow n_+ & n_+ & \text{Fin}_* \\ \downarrow & \downarrow & \downarrow \\ M^m & \xrightarrow{M^f} & M^n \\ \downarrow & & \downarrow \\ M^m & \xrightarrow{M^f} & M^n \\ (a_i) \mapsto (b_i) & & (M^n, (1, \dots, 1)) \end{array}$$

Segal 74

Cohomology on some suitable category

(Alg, k-theory from Quillen)

Ex 1.3.3 (Brouwer fixed pt Theorem)

Any cts endo $f: D^2 \longrightarrow D^2$ has a fixed pt.
 Pf) Let $r: D^2 \longrightarrow S^1$ by $(\cos x, \sin x)$

r is cts and $i: S^1 \longrightarrow D^2$ is inclusion

$\Rightarrow r i = 1_{S^1}$: r : split epi (retract) i : split mono (section)

Since π_1 is functor, $\text{Top}_* \longrightarrow \text{Group}$,

$$\pi_1(S_1, x) \xrightarrow{\pi_1(i)} \pi_1(D^2, x) \xrightarrow{\pi_1(r)} \pi_1(S^1, x)$$

$$\pi_1(r) \cdot \pi_1(i) = \pi_1(r.i) = \pi_1(1_{S'}) = 1_{\pi_1(S')}$$

But $\pi_1(S') = \mathbb{Z}$, $\pi_1(D^2) = 0$.

$\Rightarrow \mathbb{Z} \rightarrow 0 \rightarrow \mathbb{Z}$ is Identity, Contradiction

Def 1.3.5 (Contravariant functor) $F: C^{op} \rightarrow D$

① $F_c \in D$, $\forall c \in \text{obj } C$

② $Ff: F_{c'} \rightarrow F_c \quad \forall f: c \rightarrow c' \in \text{Hom } C$

St. 1) $f: x \rightarrow y, g: y \rightarrow z, \Rightarrow Ff \cdot Fg = F(gf)$
 $(F_y \rightarrow F_x, F_z \rightarrow F_y)$

2) $F1_c = 1_{F_c}$

Thus

$$\begin{array}{ccc} C^{op} & \xrightarrow{F} & D \\ c & \longmapsto & F_c \\ f \downarrow & \longmapsto & \uparrow Ff \\ c' & \longmapsto & F_{c'} \end{array}$$

(Contravariant Functor)

$$\begin{array}{ccc} C & \xrightarrow{F} & D \\ c & \longmapsto & F_c \\ f \downarrow & \longmapsto & \downarrow Ff \\ c' & \longmapsto & F_{c'} \end{array}$$

(Covariant) Functor

Ex 1.3.7

(1) $P: \text{Set}^{op} \rightarrow \text{Set}$

$$\begin{array}{ccc} A & \longmapsto & PA \\ f \downarrow & \longmapsto & \uparrow \\ B & \longmapsto & PB \end{array} \quad \begin{array}{c} f^{-1}(B')CA \\ \uparrow \\ B'CB \end{array}$$

$$(2) (-)^* : \text{Vect}_k^{\text{op}} \longrightarrow \text{Vect}_k$$

$$\begin{array}{ccc} V & \xrightarrow{\quad} & V^* \\ \emptyset \downarrow & \xrightarrow{\quad} & \uparrow \emptyset^* \\ W & \xrightarrow{\quad} & W^* \end{array} \quad \begin{array}{c} \text{w.o. } \emptyset \\ \uparrow \\ \bar{w} \end{array}$$

$$(3) \mathcal{O} : \text{Top}^{\text{op}} \longrightarrow \text{Poset}$$

$$\begin{array}{ccc} X & \xrightarrow{\quad} & \mathcal{O}(X) : \text{topology as poset by inclusion} \\ f \downarrow & \xrightarrow{\quad} & \uparrow f^{-1} \\ Y & \xrightarrow{\quad} & \mathcal{O}(Y) \end{array} \quad \begin{array}{c} f^{-1}(U_Y) \\ \uparrow \\ U_Y \end{array}$$

$$C : \text{Top}^{\text{op}} \longrightarrow \text{Poset} \text{ by collection of closed sets}$$

$$(4) \text{Spec} : \text{CRing}^{\text{op}} \longrightarrow \text{Top}$$

$$\begin{array}{ccc} R & \xrightarrow{\quad} & \text{Spec}(R) \text{ with Zariski top.} \\ f \downarrow & \xrightarrow{\quad} & \uparrow f_* \\ S & \xrightarrow{\quad} & \text{Spec}(S) \end{array} \quad \begin{array}{c} f^{-1}(P) \\ \uparrow \\ P \end{array}$$

(set-valued)

$$(5) \text{Presheaf: any functor } C \longrightarrow \text{Set} \text{ from small category } C.$$

$$\text{ex) } \mathcal{O}(X)^{\text{op}} \longrightarrow \text{Set}$$

$$\begin{array}{ccc} U & \xrightarrow{\quad} & C(U) = \{f : U \longrightarrow \mathbb{R} \text{ cts}\} \\ \text{incl} \downarrow & \xrightarrow{\quad} & \uparrow \text{res}(V, U) \\ V & \xrightarrow{\quad} & C(V) = \{f : V \longrightarrow \mathbb{R} \text{ cts}\} \end{array} \quad \begin{array}{c} f|_U \\ \uparrow \\ f \end{array}$$

(6) Presheaves on Δ : Simplex Category

Δ : Obj : $[n] = \{0, 1, \dots, n\}$, $n \in \mathbb{N}$ (finite nonempty ordinal)

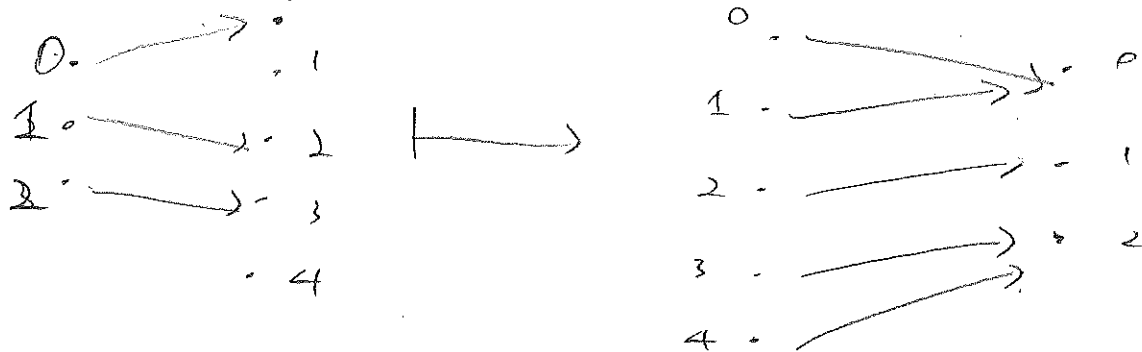
Mor : $f : [n] \rightarrow [m]$ order preserving function.

$$(i \leq j \mapsto f(i) \leq f(j))$$

A simplicial set : any presheaf on Δ , i.e.

$$X : \Delta^{op} \rightarrow \text{Set}$$

ex) X : sending $[n] \mapsto [n]$



Lem 1.3.8 Functor preserves iso.

pf) $F : C \rightarrow D \Rightarrow F(g)F(f) = F(gf) = F(1_x) = 1_{F(x)}$

$\left(\begin{array}{c} f : x \rightarrow y \\ g : y \rightarrow x \end{array} \right)$ iso. $F(f)F(g) = F(fg) = F(1_y) = 1_{F(y)}$

So does contravariant functor

Ex 1.3.9. G : gp. BG : group as a category.

Let $X : BG \rightarrow C$: any category

$\bullet \mapsto X$ some obj C

$$g \mapsto g_* : X \rightarrow X$$

By functoriality axiom,

$$\textcircled{1} h_* g_* = (hg)_*$$

$$\forall g, h \in G$$

$$\textcircled{2} e_* = 1_X$$

$$e : \text{id in } G$$

Thus $X: BG \rightarrow C$ defines an ^{left} action of the
 gp G on the object $X \in C$.

ex) If $C = \begin{pmatrix} \text{Set} \\ \text{Vect}_K \\ \text{Top} \end{pmatrix} \Rightarrow X: BG \rightarrow C$ is $\begin{pmatrix} G\text{-set} \\ G\text{-rep.} \\ G\text{-space} \end{pmatrix}$

And contravariant functor $BG^{op} \rightarrow C$ defines right
 action similar way.

Since Functor preserves iso, and all morphisms
 in BG are iso, so g_* is automorphism in
 that category.

Corollary 1.3.10. When G acts functorially on an
 object X in a category C (i.e. $\exists X: BG \rightarrow C$ exists)

g must act by automorphism $g_*: X \rightarrow X$.

and $(g_*)^{-1} = (g^{-1})_*$

Remark: Functor may not preserve mono or epic
 but preserves split mono / split epic.

pf) Almost same as 1.3.8.

Def 1.3.11. C locally small. $c \in \text{Obj } C$.
 Functor represented by $c: (C(c, -), C(-, c))$

St.

so contravariant functor defines right action

Also, since functor preserves isomorphism
in G -rep, any morphism mapped into
Automorphism
 $g_x: V \rightarrow V$

Corollary 1.3.10. G act functorially on an object
 X in Category C , $g \in G$ act by automorphisms
 $g_x: X \rightarrow X$ and $(g_x)^{-1} = (g^{-1})_x$.

Remark: Functor preserves split mono / split epi.

Def 1.3.11 C : locally small. $\forall C \in C$
define functor rep by C

$$\begin{array}{ccc}
 (C, -) & , & (-, C) : C \rightarrow \text{Set} \\
 \text{(covariant)} & & \text{(contravariant)} \\
 C \xrightarrow{\quad} \text{Set} & & C \xrightarrow{\quad} \text{Set} \\
 x \mapsto (C, x) & & x \mapsto (x, C) \\
 f \downarrow \quad \mapsto \quad \downarrow f_* & & f \downarrow \quad \mapsto \quad \uparrow f^* \\
 y \mapsto (C, y) & & y \mapsto (y, C)
 \end{array}$$

Thus, left action = covariant action
right " = contravariant "

$$C(C, -)$$

$$C(-, C)$$

$$C \longrightarrow \text{Set}$$

$$C^{\text{op}} \longrightarrow \text{Set}$$

$$x \longrightarrow C(C, x)$$

$$x \longrightarrow C(x, C)$$

$$\begin{array}{ccc} f \downarrow & \longrightarrow & \downarrow f_* \\ y & \longrightarrow & C(C, y) \end{array}$$

Post Composition

$$\begin{array}{ccc} f \downarrow & \longrightarrow & \uparrow f^* \\ y & \longrightarrow & C(y, C) \end{array}$$

Pre Composition

"post" Composition = Covariant action = left action

pre " = Contravariant " = right action

Bi-functor: ~~$C \times D$~~ or functor of two variables

Def 1.3.12 C, D : Category $C \times D$ Category

$$\text{Obj } C \times D : (c, d) \quad , \quad c \in \text{Obj } C, d \in \text{Obj } D$$

$$\text{Mor} : (f, g) : (c, d) \longrightarrow (c', d')$$

$$\begin{array}{l} \text{for } f : c \longrightarrow c' \in \text{Mor } C \\ g : d \longrightarrow d' \in \text{Mor } D. \end{array}$$

Def 1.3.13 , Two sided represented functor

$$C(-, -) : C^{\text{op}} \times C \longrightarrow \text{Set}$$

C : locally small

$$(x, y) \longrightarrow C(x, y)$$

$$\begin{array}{ccc} f^{\text{op}} \downarrow & \downarrow h & \downarrow (f^*, h_*) \\ (w, z) & \longrightarrow & C(w, z) \end{array}$$

$$(w, z) \longrightarrow C(w, z)$$

$$g \downarrow$$

$$hgf$$

(So it is covariant functor.)

In general $C^{\text{op}} \not\cong C$.

Ex 1, 3, 15. E/F : Galois ext. $\Leftrightarrow E/F$ finite ext.
 $|Aut(E/F)| = (E:F)$

Galois GP: $G = Aut(E/F)$

\mathcal{O}_G : Obj: G/H : ~~set~~ a left coset of H
(Set of all left cosets of H) by subgp H of G

Mor: G -function $G/H \rightarrow G/K$

Compute with left G -action
on the cosets.

Exercise: Every $G/H \rightarrow G/K$ has
the form $gH \mapsto gK$

Field E/F = obj: intermediate field
where $\exists \in G$ s.t. $\sigma H \sigma^{-1} \subset K$

morphism: $K \rightarrow L$ field hom fixing F .

Thus, $|Aut(E/F)| = \{ f \in \text{Hom}(E, E) : f \text{ is an automorphism} \}$

$\Phi: \mathcal{O}_G^{\text{op}} \rightarrow \text{Field } E/F$

$G/H \mapsto G \subseteq E$ subfield fixed by H
under the action of G

$G/H \rightarrow G/K \mapsto (x \mapsto \sigma x)$

(FTC of Galois Theory: Φ is bijective
in fact Φ is iso.

Thus we need a little bit relaxed concept
morphism of functor, i.e. Natural transformation.

Exercise 1.3.i) Functor between groups?

$$\begin{array}{ccc}
 G & \xrightarrow{F} & H \\
 \cdot \longmapsto \cdot & & \\
 g \downarrow & \longmapsto & h \downarrow \\
 \cdot \longmapsto \cdot & &
 \end{array}
 \quad
 \begin{array}{l}
 F(g) = h \text{ s.t.} \\
 a) F(gg') = F(g)F(g') \\
 b) F(e_g) = e_H
 \end{array}$$

\Rightarrow If F is functor, then
it is gp homo.

Conversely, gp homo: satisfy a) \Rightarrow b) \therefore by adding $\cdot \mapsto \cdot$.
it is functor.

1.3.ii) Functor between preorders?

$$(P, \leq) \quad (Q, \leq)$$

$$\begin{array}{ccc}
 x & \longmapsto & F(x) \\
 x \leq y \downarrow & \longmapsto & \downarrow \\
 y & \longmapsto & F(y)
 \end{array}
 \quad
 \begin{array}{l}
 F(x) \leq F(y) \quad \text{---} \Rightarrow a) \\
 b) \text{ : done}
 \end{array}$$

F is order preserving
function.

sketch

1.3.iii) Pr. Hasenmeyer Notation.

$$C: \begin{array}{ccc} \circ & \longrightarrow & \circ \\ a & & b \end{array} \quad \begin{array}{ccc} \circ & \longrightarrow & \circ \\ c & & d \end{array}$$

$$F: C \longrightarrow D$$

$$\begin{array}{ccc}
 a \longrightarrow b & \longmapsto & x \longmapsto y \\
 c \longrightarrow d & \longmapsto & y \longmapsto z
 \end{array}$$

$$D: \begin{array}{ccc} & y & \\ & \nearrow & \searrow \\ x & \longrightarrow & z \end{array}$$

$$F(C):$$

$$\begin{array}{ccc} & y & \\ & \nearrow & \searrow \\ x & \longrightarrow & z \end{array}$$

We have composable morphisms in $F(C)$
But composite doesn't exist.

Ex 1.3. iv) $f = 1_X \Rightarrow$

$$\begin{array}{ccccc}
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{\quad} & g \\
 f_x \downarrow & \xrightarrow{\quad} & \downarrow (f_x)_* & \downarrow & \text{Identity} \\
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{\quad} & (f_x g) = g
 \end{array}$$

$f: x \rightarrow y, \quad g: y \rightarrow z$

$$\Rightarrow \begin{array}{ccccc}
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{\quad} & h \\
 f \downarrow & & \downarrow f_* & & \downarrow \\
 y & \xrightarrow{\quad} & C(C, y) & \xrightarrow{\quad} & fh \\
 g \downarrow & & \downarrow g_* & & \downarrow \\
 z & \xrightarrow{\quad} & C(C, z) & \xrightarrow{\quad} & gfh
 \end{array}$$

Ex 1.3. v) Claim: $F: C^{op} \rightarrow D \iff G: C \rightarrow D^{op}$

Let $F: C^{op} \rightarrow D$. Construct $G: C \rightarrow D^{op}$ s.t.

$G(c) := F(c)$. From $\text{Hom}_D(Fb, Fa) \leftrightarrow \text{Hom}_{D^{op}}(Fa, Fb)$

So given $f \in \text{Hom}_C(a, b)$

define $G(f) = (ff)^{op}$ in $\text{Hom}_{D^{op}}(Fa, Fb)$

Thus, $F: C \rightarrow C^{op} \rightarrow D \quad G: C \xrightarrow{G} D^{op} \rightarrow D$

$$\begin{array}{ccccc}
 x & \xrightarrow{\quad} & x & \xrightarrow{\quad} & Fx \\
 f \downarrow & & f^{op} \uparrow & & \uparrow Ff \\
 y & \xrightarrow{\quad} & y & \xrightarrow{\quad} & Fy
 \end{array}$$

F, G are the same functor. Also, $F: C \rightarrow D \iff G: C^{op} \rightarrow D^{op}$

$$\begin{array}{ccccc}
 C & \rightarrow & D \\
 x & \downarrow & x \\
 y & \xrightarrow{\quad} & y
 \end{array}
 \quad
 \begin{array}{ccccc}
 C & \rightarrow & C^{op} & \rightarrow & D^{op} & \rightarrow & D \\
 x & \downarrow & x & \uparrow f^{op} & Fx & \downarrow Ff & Fx \\
 y & \downarrow & y & \uparrow (ff)^{op} & Fy & \downarrow Ff & Fy
 \end{array}$$

$$1.3.vi) \quad F: D \rightarrow C, \quad G: E \rightarrow C.$$

$$F \downarrow G : \text{Obj} := (d, e, f)$$

$$d \in D, e \in E, f: Fd \rightarrow Ge \in C$$

$$\text{Mor} : ((d, e, f), (d', e', f'))$$

$$= \{ (h, k) \in \text{Mor } D \times \text{Mor } E : \}$$

$$h: d \rightarrow d', k: e \rightarrow e' \text{ s.t.}$$

$$Fd \xrightarrow{f} Ge$$

$$Fh \downarrow \quad \curvearrowright \quad \downarrow Gk \quad \text{in } C, \text{ i.e.}$$

$$Fd' \xrightarrow{f'} Ge' \quad f' \cdot Fh = Gk \cdot f$$

It is category since,

$$\textcircled{1} \quad 1_{(d,e,f)} = (1_d, 1_e)$$

$$\textcircled{2} \quad (d, e, f) \xrightarrow{(h_1, k_1)} (d', e', f') \xrightarrow{(h_2, k_2)} (d'', e'', f'')$$

$$\Rightarrow \begin{array}{ccc} Fd & \xrightarrow{f} & Ge \\ Fh_2 \downarrow & \curvearrowright & \downarrow Fk_1 \\ Fd' & \xrightarrow{f'} & Ge' \end{array} \Rightarrow (Fk_2)(Fk_1) \cdot f = Fk_2(f') \cdot Fh_1 = f'' Fh_2 Fh_1$$

$$\begin{array}{ccc} Fd' & \xrightarrow{f'} & Ge' \\ Fh_2 \downarrow & \curvearrowright & \downarrow Fk_2 \\ Fd'' & \xrightarrow{f''} & Ge'' \end{array} \quad \begin{array}{ccc} Fd & \xrightarrow{f} & Ge \\ F(h_2 h_1) \downarrow & & \downarrow Fk_2 k_1 \\ Fd'' & \xrightarrow{f''} & Ge'' \end{array}$$

$$Fd'' \xrightarrow{f''} Ge'' \quad \Rightarrow f'' (F(h_2 h_1)) = F(k_2 k_1) f$$

Functor: $\text{dom}: F \downarrow G \longrightarrow D$, $\text{cod}: F \downarrow G \longrightarrow E$

$$(d, e, f) \longmapsto d \quad (d, e, f) \longmapsto e$$

$$(h, k) \downarrow \longmapsto h \quad (h, k) \downarrow \longmapsto k$$

$$(d', e', f') \longmapsto d' \quad (d', e', f') \longmapsto e'$$

Ex 1.3. vii) $D = (\{-3, 1\}, 1_0)$ $F: D \longrightarrow C$ $G = 1_C$

$E = C$ $\bullet \longmapsto c$

C/C
under

$\text{Obj} = (\bullet, x, f: C \longrightarrow x)$

$\text{Mor} : ((\bullet, x, f: C \longrightarrow x), (\bullet, y, g: C \longrightarrow y))$

$= \{ (1_0, h: x \longrightarrow y) \in \text{Mor } D \times \text{Mor } E$

s.t.

$$\begin{array}{ccc} C & \xrightarrow{f} & x \\ 1_C \downarrow & \circlearrowleft & \downarrow G \\ C & \xrightarrow[g]{} & y \end{array} \quad Gf = h \quad \text{i.e.} \quad \begin{array}{ccc} & C & \\ f \swarrow & & \searrow g \\ x & \xrightarrow[h]{} & y \end{array}$$

C/C
over

$D = C$ $F = 1_C$ $G = \bullet \longmapsto c$

$E = (\{-3, 1\}, 1_0)$

$\text{Obj} : (x, \bullet, f: x \longrightarrow c)$

$\text{Mor} ((x, \bullet, f: x \longrightarrow c), (y, \bullet, g: y \longrightarrow c))$

$= \{ (h: x \longrightarrow y, 1_0) \in \text{Mor } D \times \text{Mor } E$

s.t.

$$\begin{array}{ccc} x & \xrightarrow{f} & c \\ h \downarrow & \circlearrowleft & \downarrow 1_C \\ y & \xrightarrow[g]{} & c \end{array} \quad \text{i.e.} \quad \begin{array}{ccc} & c & \\ f \nearrow & & \nwarrow g \\ x & \xrightarrow[h]{} & y \end{array}$$

projection functor:

$$\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{D} = (\{ \cdot \}, 1)$$

$$\begin{aligned} (\cdot, x, f: \mathcal{C} \rightarrow x) &\longmapsto \cdot \\ (1, h) \downarrow &\longmapsto \int 1 \\ (\cdot, y, g: \mathcal{C} \rightarrow y) &\longmapsto \cdot \end{aligned}$$

$$\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{E} = \mathcal{C}$$

$$\begin{aligned} (\cdot, x, f: \mathcal{C} \rightarrow x) &\longmapsto x \\ (1, h) \downarrow &\longmapsto \int h \\ (\cdot, y, g: \mathcal{C} \rightarrow y) &\longmapsto y \end{aligned}$$

Thus: $\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{C}$

$$\begin{aligned} f: \mathcal{C} \rightarrow x &\longmapsto x \\ h \downarrow &\longmapsto \int h \\ g: \mathcal{C} \rightarrow y &\longmapsto y \end{aligned}$$

$\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{C}$

$$\begin{aligned} f: x \rightarrow \mathcal{C} &\longmapsto x \\ h \downarrow &\longmapsto \int h \\ g: y \rightarrow \mathcal{C} &\longmapsto y \end{aligned}$$

Similarly, $\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{D} = \mathcal{C}$

$$\begin{aligned} (x, \cdot, f: x \rightarrow \mathcal{C}) &\longmapsto x \\ (h, 1) \downarrow &\longmapsto \int h \\ (y, \cdot, g: y \rightarrow \mathcal{C}) &\longmapsto y \end{aligned}$$

$\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{E} = (\{ \cdot \}, 1)$

$$\begin{aligned} (x, \cdot, f: x \rightarrow \mathcal{C}) &\longmapsto \cdot \\ (h, 1) \downarrow &\longmapsto \int 1 \\ (y, \cdot, g: y \rightarrow \mathcal{C}) &\longmapsto \cdot \end{aligned}$$

Ex 1.3, viii) Ex: Functor need not

Find $f: \mathcal{C} \longrightarrow \mathcal{D}$, $f \in \text{Mor}(\mathcal{C})$

reflect isomorphism

st. f is iso in \mathcal{D}

f is not iso in \mathcal{C}

$$\begin{array}{c} \cdot \xrightarrow{f} \cdot \\ \mathcal{C} \end{array}$$

$$\begin{array}{c} \cdot \xrightleftharpoons[f^{-1}]{f} \cdot \\ \mathcal{D} \end{array}$$

More generically, $\mathcal{I} = (\{ \cdot \}, 1)$

$$\mathcal{C} \longrightarrow \mathcal{I} \quad \text{by} \quad c \longmapsto \cdot \\ f \longmapsto 1$$

ex) Homology
Homotopy
Functor:
(Quasi-iso
but not homeo.)
 $\forall f, \theta_c$

Ex 1.3. ix)

Source map	Group iso.	Group epi	Group
$Z(-)$	Yes (iso)	Yes	No ¹⁾
$C(-)$	Yes	Yes	Yes ²⁾
$\text{Aut}(-)$	Yes	Don't know.	No ³⁾

1). If $G \xrightarrow{f} H \Rightarrow Z(G) \cong Z(H)$

pf)

Claim 1: $G \xrightarrow{f} H$ surj, then $f(Z(G)) \leq Z(H)$

pf) $h \cdot f(g) = f(h'g) = f(gh') = f(g)h \quad \forall h \in H,$

$g \in Z(G), \quad h' : \text{preimage of } h.$

Claim 2: $G \xrightarrow{\text{inj } f} H$ then $f(Z(G)) = Z(H)$

pf)

If $h \in Z(H), \quad \exists g \in G \text{ st. } f(g) = h.$

$$g'g = f^{-1}(h'h) = f^{-1}(hh') = f^{-1}(h)f^{-1}(h') = gg'$$

$\Rightarrow g \in Z(G) \Rightarrow f(Z(G)) \supseteq Z(H). \quad \text{done.}$

Claim 3: $Z(-) : \begin{matrix} \text{Group iso} \\ \text{Group epi} \end{matrix} \longrightarrow \text{Group}$ is functorial.

pf) $G \xrightarrow{Z(-)} Z(G) \quad \text{Let } G \xrightarrow{\phi} H \xrightarrow{\psi} K$
 $\downarrow 1_G \quad \downarrow Z(1_G) \quad \text{Iso (or epi)}$
 $G \longrightarrow Z(G)$

Then from $\phi(Z(G)) \subseteq Z(H)$,
 $\psi(Z(H)) \subseteq Z(K)$

we have map

$$Z(G) \xrightarrow{\phi|_{Z(G)}} Z(H) \xrightarrow{\psi|_{Z(H)}} Z(K)$$

Thus, $\psi|_{Z(H)} \circ \phi|_{Z(G)} = (\psi \circ \phi)|_{Z(G)}$

implies ϕ first functoriality axiom.

Claim 4: $Z(-)$ Group \rightarrow Group is not a functor.

p.f) $I_{C_2} = C_2 \rightarrow S_3 \rightarrow S_3/A_3 \cong C_2$

$$1 \mapsto (12) \mapsto (12) + A_3 \mapsto 1$$

where C_2 : cyclic group of order 2,
 S_3 : Sym. " 3.

A_3 : alternating op of S_3 .

Then, If Z is a functor,

$$\begin{aligned} Z(C_2) &\rightarrow Z(S_3) \rightarrow Z(C_2) = Z(C_2) \xrightarrow{I_{Z(C_2)}} Z(C_2) \\ &= C_2 \rightarrow 0 \rightarrow C_2, \text{ Contradiction } \quad C_2 \xrightarrow{I_{C_2}} C_2 \end{aligned}$$

2). Claim 1: $G \xrightarrow{f} H \Rightarrow C(f(G)) = f(C(G)) \leq C(H)$

pf) $C(f(G)) = \{ f(x)f(y)f(x)^{-1}f(y)^{-1} : x, y \in G \}$
 $= \{ f(xy x^{-1}y^{-1}) : x, y \in G \}$
 $= f(C(G)) \leq \{ xyx^{-1}y^{-1} : x, y \in H \} = C(H)$

Claim 2: $C(-) : \begin{matrix} \text{Group}_{\text{iso}} & \longrightarrow & \text{Group} \\ \text{Group}_{\text{epi}} & \longrightarrow & \\ \text{Group} & \longrightarrow & \text{Group} \end{matrix}$ ~~is~~ are functors

pf). $G \xrightarrow{f} H \Rightarrow C(H) \cong f(C(G))$

$\Rightarrow C(G) \xrightarrow{f|_{C(G)}} C(H)$ is gp homo.

If $f = 1_G \Rightarrow C(f) = 1_{C(G)}$

$G \xrightarrow{\phi} H \xrightarrow{\psi} K$

$\Rightarrow \begin{matrix} \phi|_{C(G)} & \psi|_{C(H)} \\ C(G) \longrightarrow C(H) \longrightarrow C(K) \end{matrix} = (\psi \circ \phi)|_{C(G)}$

3)
 Claim 1. $\begin{matrix} \text{Aut} \\ \text{Group}_{\text{epi}} \longrightarrow \text{Group} \\ \text{Group} \longrightarrow \text{Group} \end{matrix}$ is not a functor

pf) $G = \mathbb{F}_n \rtimes \mathbb{F}_n^*$ by multiplication.

$$G = \{ (a, b) : a \in \mathbb{F}_n, b \in \mathbb{F}_n^\times \}$$

$$\text{and } (a, b) \cdot (c, d) = (a + bc, bd)$$

$$\text{Also, let } N = \{ (a, 1) : a \in \mathbb{F}_n \} \cong \mathbb{F}_n$$

$$K = \{ (0, b) : b \in \mathbb{F}_n^\times \} \cong \mathbb{F}_n^\times$$

N, K are subgp of G .

and by construction, N is normal;

$$\begin{aligned} (b, c) (a, 1) (b, c)^{-1} &= (b + ca, c) (b, c^{-1}) \\ &= (b + ca - bc, 1) \in N \end{aligned}$$

$$\Rightarrow NK = G \quad (\text{by construction})$$

By Schur-Zassenhaus lemma,

Any subgp of G with order $|K|$

conjugate to each other, i.e., if k_1, k_2 has order $|K|$, then $\exists g \in G$ s.t.

$$g k_1 g^{-1} = k_2$$

$$\begin{aligned} (a, 1) (c, d) &= (a, 1) \\ &= (a + bc, bd) \end{aligned}$$

Note, that $(0, 1)$ is identity in G .

$$\text{and } (a, b)^{-1} = \left(-\frac{a}{b}, b^{-1} \right)$$

Hence let $\alpha \in \text{Aut}(G)$

$\Rightarrow \alpha(K)$ is a subgp of G with order $|K|$

$\Rightarrow \exists g \in G$ s.t.

$$g\alpha(K)g^{-1} = K$$

\Rightarrow Let $\varphi_g \in \text{Inn}(G)$ s.t. $\varphi_g(g') = gg's^{-1}$

$\Rightarrow \varphi_g \circ \alpha$ fix K .

Also, notes that $\varphi_g \circ \alpha$ fix N since N is unique 11-sylow subgp of G .

Thus, if $\varphi_g \circ \alpha(1,1) = (a,1)$ then

then $a \neq 0$ since $\varphi_g \circ \alpha^{-1}(0,1) = (0,1)$
(identity.)

thus $\exists b \in \mathbb{F}_n^\times$ s.t. $b^{-1}a = 1$

$$\therefore (0,b)(a,1)(0,b^{-1}) = (ab,b)(0,b^{-1}) = (ab,1) = (1,1)$$

$$\therefore \varphi_b \varphi_g \circ \alpha(1,1) = (1,1)$$

And since $(0,b)(0,1)(0,b^{-1}) = (0,1)$,

φ_b still fix K .

Let $\beta = \varphi_b \circ \varphi_g \circ \alpha$.

Then, $\beta(k) = k$, $\beta(1,1) = (1,1)$.

and β is auto. ($\beta(N) = N$)

If $a \neq 0$ in F_{11}

$$\Rightarrow \beta((a,1)) = \beta((0,a)(1,1)(0,a^{-1})) \\ = (\beta(0,a))(1,1)(\beta(0,a))^{-1}$$

Since β fix N , $\beta(a,1) = (b,1)$

And β fix $K \Rightarrow \beta(0,a) = (0,c)$

for some
 $b \in F_{11}^*$
 $c \in F_{11}^*$

$$\text{Thus, } (b,1) = (0,c)(1,1)(0,c^{-1}) \\ = (c,c)(0,c^{-1}) = (c,1)$$

$\Rightarrow b=c$. And

$$\beta(0,a) = (0,c) = (0,b)$$

Thus, $\gamma: F_{11} \rightarrow N \xrightarrow{\beta} N \rightarrow F_{11}$

$$a \mapsto (a,1) \mapsto (b,1) \mapsto b \\ 0 \mapsto (0,1) \mapsto (0,1) \mapsto 0$$

induces a map on F_{11} (we didn't show it's homo)

similarly, $F_{11}^* \rightarrow K \xrightarrow{\beta} K \rightarrow F_{11}^*$

$$a \mapsto (0,a) \mapsto (0,b) \mapsto b$$

we claim that γ is field homo.

$$\sigma(a)\sigma(b) = (\sigma(a), 1) \cdot (\sigma(b), 1) \quad \text{in } G.$$

$$\begin{aligned} (c, 1) &= (0, \sigma(a)) (\sigma(b), 1) (0, \sigma(a)^{-1}) \\ &= \beta(0, a) \beta(b, 1) \beta(0, a^{-1}). \end{aligned}$$

$$\begin{aligned} \text{Since } \sigma(a) &=: a \mapsto (0, a) \mapsto \beta(0, a) \\ &= (0, b) \\ &\text{for some } c \\ &\mapsto c. \end{aligned}$$

$$\text{thus } \beta(0, a) = (0, c) = (0, \sigma(a)).$$

$$\text{and } b \mapsto (b, 1) \mapsto \beta(b, 1) = (c, 1) \text{ for some } c.$$

$$\Rightarrow \sigma(b) = c$$

$$\Rightarrow \beta(b, 1) = (c, 1) = (\sigma(b), 1)$$

$$= \beta(ab, 1) = (c, 1) \text{ for some } c \in F.$$

$$\text{and } \Rightarrow \sigma(a)\sigma(b) = c.$$

$$ab \mapsto (ab, 1) \mapsto \beta(ab, 1) = (c, 1)$$

$$\Rightarrow \sigma(ab) = c.$$

Hence $\sigma(ab) = \sigma(a)\sigma(b)$. Thus β induces a field homomorphism on F .

Thus $\text{Aut}(G) \cong \text{Inn}(G) \cdots \text{Aut}(F_{11})$.

$$\cong \text{Inn}(G)$$

$$\cong G$$

Since any automorphism of G is field homomorphism times inner automorphism.

and $\text{Aut}(F_{11}) = 1$ since there are only F_{11}^\times gp homomorphisms and only 1 homomorphism fix $1 \mapsto 1$.

Now think about a map.

$$\phi: F_{11}^\times \longrightarrow G \longrightarrow G/N \cong F_{11}^\times: \text{isomorphism}$$

$$\text{Aut}(-) \text{ induces } \text{Aut}(F_{11}^\times) \longrightarrow \text{Aut}(G) \longrightarrow \text{Aut}(F_{11}^\times)$$

$$\text{But } |\text{Aut}(F_{11}^\times)| = (\mathbb{Z}/4\mathbb{Z}) = 4 \quad \text{and } (4|110) = 1.$$

$$|\text{Aut}(G)| = |G| = 110.$$

$$\Rightarrow \text{Aut}(F_{11}^\times) \longrightarrow \text{Aut}(G) \text{ is zero.}$$

$$\text{But } \phi = 1_{F_{11}^\times} \text{, thus}$$

$$\text{Aut}(\phi) = 1_{\text{Aut}(F_{11}^\times)} \neq 0. \text{ So } \text{Aut} \text{ doesn't satisfy functoriality axiom.}$$

1.3.X.) Let G, H gr, $f: G \rightarrow H$ homo,
 X_G, X_H , set of conjugacy classes, of G, H .

$g: G \rightarrow X_G, h: H \rightarrow X_H$ class functions

i.e. $g(a) = g(bab^{-1}) \quad \forall b \in G.$

Let $\text{Conj}: \text{Group} \xrightarrow{\text{class functions}} \text{Set}$

$$\begin{array}{ccc} G & \xrightarrow{\quad} & X_G \\ \phi \downarrow & \xrightarrow{\quad} & \downarrow \text{Conj}(\phi) \\ H & \xrightarrow{\quad} & X_H \end{array} \quad \begin{array}{c} \overline{g} \\ \downarrow \\ \overline{h} \end{array}$$

① $\text{Conj}(\phi)$ is well-defined;

if $g' \in \overline{g}$, then, $g' = bab^{-1}$ for some $b \in G$.

Thus, $\phi(g') = \phi(b)\phi(g)\phi(b)^{-1} \Rightarrow \phi(g') \in \overline{\phi(g)}.$

② It is functorial; If $\psi: H \rightarrow J$ morphism,

$$\text{Conj}(\psi) \cdot \text{Conj}(\phi) = \text{Conj}(\psi \circ \phi)$$

$$\text{and } \text{Conj}\left(\begin{smallmatrix} G \\ 1_G \end{smallmatrix}\right) = 1_{X_G}.$$

Thus, if $|X_G| \neq |X_H|$, then $\text{Conj}(\phi)$

is not isomorphism, since functor preserves

iso, $\phi: G \rightarrow H$ is not iso.

Natural Transformation

Ex) $(-)^* \circ (-)^*$ where $(-)^*: \text{Vect}_k^{\text{op}} \rightarrow \text{Vect}_k$
 $V \mapsto V^*$

!Induces $V' \cong V^{**}$

but this iso comes from $\text{ev}_V \quad \forall V \in \mathcal{V}$.

No unnatural choice of basis is needed.

Def 1.4.1. (Natural Transformation)

C, D : categories $F, G: C \Rightarrow D$

α : Natural Transformation $F \Rightarrow G = \{\alpha_c: c \in C\}$

s.t. $\alpha_c: F_c \rightarrow G_c \in \text{Mor } D \quad \forall c \in C$

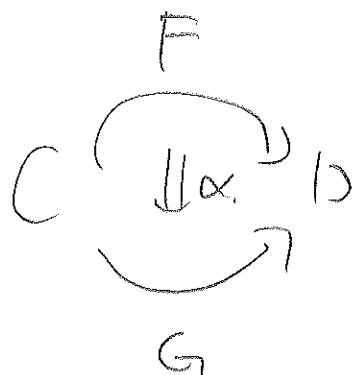
\uparrow Component of α .

$$\begin{array}{ccc} F_c & \longrightarrow & G_c \\ Ff \downarrow & \curvearrowright & \downarrow Gf \\ F_{c'} & \longrightarrow & G_{c'} \end{array} \quad \begin{array}{l} \text{s.t. } \forall f: c \rightarrow c' \\ \in \text{Mor } C, \\ \text{commutes.} \end{array}$$

Natural Isomorphism: α : natural transformation
s.t. α_c is iso $\forall c \in C$.

Then $\alpha: F \cong G$.

"The arrows are natural" \Rightarrow collection of arrows define natural transf.



Ex 1.4.3.

$$(i) \quad \text{Vect}_k \xrightarrow{\text{ev}} (-)^*, (-)^*$$

since $\forall \phi: V \rightarrow W$

morphism in Vect_k

$$f_v: g \in V^* \mapsto g(v)$$

$$\begin{array}{ccc} V & \xrightarrow{\text{ev}} & V^{**} \\ \downarrow \phi & & \downarrow \phi^{**} \\ W & \xrightarrow{\text{ev}} & W^{**} \end{array}$$

$$\begin{array}{c} \downarrow \phi^{**}(W^* \rightarrow k) = V^{**\phi^*} \\ \phi^{**}(f_v) \end{array}$$

$$f_{\phi(v)}: g \in W^* \mapsto g(\phi(v))$$

$$\phi^{**}(f_v) := W^* \xrightarrow{\phi^*} V^* \xrightarrow{f_v} k$$

$$g \mapsto g \circ \phi \mapsto g \circ \phi(v)$$

$$f_{\phi(v)}:$$

$$g \mapsto g(\phi(v))$$

Same.

$$(ii) \quad \mathbb{1}_{\text{Vect}_K} \not\Rightarrow (-)^*$$

Since $\mathbb{1}_{\text{Vect}_K}$: Covariant

$(-)^*$: Contravariant.

More significantly $V \cong V^*$ cannot be defined without choice of basis (basis) which is not preserved by any nonidentity linear endomorphism.

$$(iii) \quad \mathbb{1}_{\text{Set}} \Rightarrow P. \quad \begin{array}{ccc} A & \xrightarrow{\quad} & P(A) \\ \mathbb{1}_{\text{Set}} : \text{Set} & \xRightarrow{\quad} & \text{Set} \\ A & \xrightarrow{\quad} & A \end{array}$$

$$\text{s.t.} \quad \begin{array}{ccc} A & \xrightarrow{\eta_A} & P(A) \\ f \downarrow & & \downarrow f \\ B & \xrightarrow{\eta_B} & P(B) \end{array} \quad \begin{array}{l} \text{by } \eta_A : A \rightarrow P(A) \\ a \mapsto \{a\} \end{array}$$

$$(iv) \quad \begin{array}{ccc} & e \mapsto x & \\ X, Y : BG & \xRightarrow{\quad} & C \\ & e \mapsto y & \end{array}$$

What is natural transf?

$$\Rightarrow \alpha : X \rightarrow Y \quad \begin{array}{l} \text{a single } \alpha. \\ \text{homomorphism in } C. \end{array}$$

$$\begin{array}{ccc} X & \xrightarrow{\alpha} & Y \\ g_* \downarrow \circlearrowleft & & \downarrow \circlearrowright g_* \\ X & \xrightarrow{\alpha} & Y \end{array}$$

s.t. left diagram

Commutative.
We say α is "G-equivariant".

(V) $\mathcal{O}: \text{Top}^{\text{op}} \Rightarrow \text{Set}$ by $()^c$, taking
Complement.

i.e., $\mathcal{O} \xRightarrow{()^c} C$.

To see this, any cts $f: Y \rightarrow X \in \text{Top}^{\text{op}}$.

$$\begin{array}{ccc}
 Y & \xrightarrow{\mathcal{O}(Y) \xrightarrow{()^c} C(Y)} & U^c \\
 \downarrow f^! & & \downarrow f^! \\
 \mathcal{O}(X) & \xrightarrow{()^c} & C(X) \\
 f^!(U) & \xrightarrow{()^c} & (f^!(U))^c = f^!(U^c)
 \end{array}$$

Then, $(f^!(U))^c = X - f^!(U) = f^!(U^c)$.

Also, it is natural iso since $()^c$ is
bijection.

(VI) $(-)^{\text{op}}: \text{Group} \rightarrow \text{Group}$.

$$G \longmapsto G^{\text{op}}: \begin{cases} \text{obj: } \bullet \\ \text{mon: } G \text{ s.t.} \end{cases}$$

$$g^{\text{op}} \cdot h^{\text{op}} = (hg)^{\text{op}}.$$

Thus, $\phi: G \rightarrow H$ induces $\phi^{\text{op}}: G^{\text{op}} \rightarrow H^{\text{op}} \hookrightarrow g \mapsto \phi(g)$.

Let $\tau_G: G \rightarrow G^{op}$
 $g \mapsto g^{-1}$.

it is not automorphism $G \rightarrow G$
 but it is homo $G \rightarrow G^{op}$.

since $gh \mapsto (gh)^{-1} = (h^{-1} \cdot g^{-1})^{op} = \tau_G(g) \cdot \tau_G(h)$

(since multiple order is reversed.)

Thus

$$\begin{array}{ccc}
 & g & \\
 & \downarrow & \\
 G & \xrightarrow{\tau_G} & G^{op} \\
 \downarrow \phi & & \downarrow \phi^{op} \\
 H & \xrightarrow{\tau_H} & H^{op} \\
 \phi(g) & \xrightarrow{\quad} & \phi(g)^{-1}
 \end{array}$$

and $\phi(g^{-1}) = \phi(g)^{-1}$
 in H .
 $\Rightarrow \phi(g^{-1}) = \phi(g)^{-1}$
 in H^{op} .

(VII). $F: Vect_K \rightarrow Vect_K$

$V \mapsto V \otimes V$.

Then, $V \xrightarrow{\tau_V} V \otimes V$ commute,

$\phi \downarrow \quad \downarrow \phi \otimes \phi$ when $\tau_V = 0$
 $W \xrightarrow{\tau_W} W \otimes W$ $\forall V \in Vect_K$

However, there is no basis independent definable lin. map. $V \rightarrow V \otimes V$.

Ex) Ab_{fg} : Category of fin. gen. ab gp.

For given ab gp A , let TA : torsion subgp of A .

By classification Thm of ab gp,

$$A \cong TA \oplus (A/TA)$$

Prop 1.4.4. $A \cong TA \oplus (A/TA)$ are not natural.
 $\forall A \in \text{Ab}_{fg}$.

pf) If it were natural,

$$\alpha: A \twoheadrightarrow A/TA \hookrightarrow TA \oplus (A/TA) \cong A$$

acting as a natural endomorphism of
 the $\mathbb{I}_{\text{Ab}_{fg}}$, i.e. $\mathbb{I}_{\text{Ab}_{fg}} \xrightarrow{\alpha} \mathbb{I}_{\text{Ab}_{fg}}$.

Claim 1: If η : natural endomorphism of $\mathbb{I}_{\text{Ab}_{fg}}$

$$\Rightarrow \eta: x \mapsto nx \quad \text{for some } n \in \mathbb{Z}$$

pf). Since $\eta: \mathbb{Z} \rightarrow \mathbb{Z}$ is multiple by $n \in \mathbb{Z}$.

and $\forall A \in \text{Ab}_{gp}$ with map $\mathbb{Z} \xrightarrow{\alpha} A$,
 $1 \mapsto a$.

By choosing $x=1$,
 $\eta_A(a) = na$.

$$\begin{array}{ccccc} x & \xrightarrow{\eta_{\mathbb{Z}}} & nx & & \\ \downarrow & \eta_{\mathbb{Z}} & \downarrow & & \\ \mathbb{Z} & \xrightarrow{\quad} & \mathbb{Z} & & \\ \downarrow & \eta & \downarrow & & \\ A & \xrightarrow{\eta_A} & A & \xrightarrow{ax} & \\ \downarrow & \eta_A & \downarrow & & \\ ax & \xrightarrow{\quad} & \eta_A(ax) = x\eta_A(a) & & \end{array}$$

Thus, $\eta_A: A \rightarrow A$
 $a \mapsto na.$ \square

Now, if $A = \mathbb{Z}$, $\alpha: a \mapsto na$ by \mathbb{Z} .
 claim. Then $\alpha: A \rightarrow A/\eta_A = A \rightarrow A \xrightarrow{\cong} A$.
 is Iso, thus, $n \neq 0$.

If $A = \mathbb{Z}/2n\mathbb{Z}$, $\eta_A = A$, thus

$$\alpha_{\mathbb{Z}/2n\mathbb{Z}}: A \rightarrow 0 \rightarrow A \rightarrow A$$

\therefore zero map.

But $n \neq 0$ in $\mathbb{Z}/2n\mathbb{Z}$ thus.

$$\begin{array}{ccc} \mathbb{Z} & \xrightarrow{\alpha_{\mathbb{Z}}} & \mathbb{Z} \\ \pi \downarrow & & \downarrow \pi \\ \mathbb{Z}/2n\mathbb{Z} & \xrightarrow{\alpha_{\mathbb{Z}/2n\mathbb{Z}}} & \mathbb{Z}/2n\mathbb{Z} \end{array}$$

$\alpha_{\mathbb{Z}/2n\mathbb{Z}}$

doesn't
commute,
contradiction

$n \neq 0$,

0 .

Ex 1.4.6. (Riesz Representation Thm)

$$\begin{array}{ccc} \mathcal{I} : \text{cHaus} & \longrightarrow & \text{Ban.} \\ \text{cpt} & & \text{Banach space.} \\ \text{Hausdoff} & & \end{array}$$

$$\begin{array}{ccccc} X & \longmapsto & \mathcal{I}(X) := \{ \text{signed Baire measure} \} \\ f \downarrow & \longmapsto & \downarrow f^* & \downarrow \mu \\ Y & \longmapsto & \mathcal{I}(Y) & \mu \circ f^* \end{array}$$

is a functor.

$$C^* : \text{cHaus} \longrightarrow \text{Ban.}$$

$$X \longmapsto C(X)^* : \text{dual of the Banach space } C(X) \text{ of cts real valued functions on } X.$$

$$\text{For each } \mu \in \mathcal{I}(X), \quad \phi_\mu : C(X) \rightarrow \mathbb{R} \in C(X)^* \\ g \mapsto \int_X g d\mu$$

$$\text{Let } \eta : \mathcal{I} \Rightarrow C^* \quad \eta,$$

$$\begin{array}{ccccc}
 \mu & \xrightarrow{\quad} & & & \phi_\mu \\
 \downarrow & \Sigma(X) \xrightarrow{\eta_X} C(X)^* & & & \downarrow \\
 & \downarrow (-\circ f^{-1}) & & & \downarrow (-\circ f) \\
 & \Sigma(Y) \xrightarrow{\eta_Y} C(Y)^* & & & \downarrow \\
 \mu \circ f^{-1} & & & & \phi_{\mu \circ f^{-1}} \\
 & \xrightarrow{\quad} & \phi_{\mu \circ f^{-1}} & &
 \end{array}$$

st. $\phi_{\mu \circ f^{-1}}(g) = \int_Y g \, d(\mu \circ f^{-1})$, $\phi_{\mu \circ f} = \int_X g \circ f \, d\mu$.

But from real analysis, they are equal.

thus $\eta: \mu \mapsto \phi_\mu$ is natural transf of $\Sigma \Rightarrow C^*$.

RRT: This is natural iso.

Ex 1.4.7. \mathcal{L} : locally shall

$$f: w \rightarrow z, \quad h: y \rightarrow z$$

$$\begin{array}{ccc}
 C(x, y) & \xrightarrow{h \circ -} & C(y, z) \\
 \downarrow - \circ f & \wr & \downarrow - \circ f \\
 C(w, y) & \xrightarrow{h \circ -} & C(w, z)
 \end{array}$$

where $h \circ - = h_x$
 $- \circ f = f^*$

(post and pre composition)

Since $C(-, y), C(-, z)$ are functors,

$$h_x : C(-, y) \Rightarrow C(-, z)$$

is natural transformation. (n.t)

Similarly,

$$f^* : C(x, -) \Rightarrow C(w, -) \text{ n.t.}$$

Ex 1.4.9. For $A, B \in \text{Sets}$. $A+B := \text{disjoint union of } A, B$.

$$A^B := \{B \rightarrow A\}$$

$$\Rightarrow A \times (B+C) \cong (A \times B) + (A \times C), (A \times B)^C \cong A^C \times B^C$$

$$A^{B+C} \cong A^B \times A^C$$

$$(A^B)^C \cong A^{B \times C}$$

Actually, this is natural iso between

$$\text{Set} \times \text{Set} \times \text{Set} \rightarrow \text{Set} \text{ functors.}$$

or (contravariant)

Now, restrict this natural iso in Fin Iso ,

[Obj: finite set Mor: bijection]

Then let $| - | : \text{Fin}_{\text{iso}} \rightarrow \mathbb{N}$ be functor.

(\mathbb{N} : discrete category of natural #)

$$\begin{array}{ccc} A & \xrightarrow{\text{iso}} & |A| \\ \downarrow & & \downarrow \\ B & \xrightarrow{\text{iso}} & |B| = |A| \end{array} \quad \begin{array}{c} \text{1}_{|A|} \\ \text{1}_{|B|} \end{array}$$

Then this induce:

$$a(b+c) = ab + ac, \quad (ab)^c = a^c b^c.$$

$$a^{b+c} = a^b a^c, \quad (a^b)^c = a^{bc}.$$

a laws of mult and add in \mathbb{N} .

deategorification of $\mathbf{Fin}_{iso} \Rightarrow \mathbb{N}$ and
these law.

Categorification of $\mathbb{N} = \mathbf{Fin}_{iso}$

Exercise 1.4. i.

Since α is natural iso, $\forall C \in \mathcal{C}$,

$$\exists \alpha_C^{-1} \text{ s.t. } \alpha_C \circ \alpha_C^{-1} = \text{id}_{G_C}$$

$$\alpha_C^{-1} \circ \alpha_C = \text{id}_{F_C}$$

Thus, let α^{-1} : collection of α_C^{-1} . Then,

$$\forall f: C \rightarrow C'$$

WTS

$$G_C \xrightarrow{\alpha_C^{-1}} F_C$$

$$Ff \circ \alpha_C^{-1} = \alpha_{C'}^{-1} \circ Gf$$

$$Gf \downarrow$$

$$\downarrow Ff$$

Since $\alpha_{C'} \circ Gf = Ff \circ \alpha_C$,

$$G_{C'} \xrightarrow{\alpha_{C'}^{-1}} F_{C'}$$

$$\Rightarrow \alpha_{C'} \circ Gf \circ \alpha_C^{-1} = Ff \circ \text{id}_{G_C}$$

$$\Rightarrow \text{id}_{F_{C'}} \circ Ff \circ \alpha_C^{-1} = \alpha_{C'}^{-1} \circ Gf$$

1.4. ii. If $\phi: G \rightarrow H$ be a function,

$$\text{then, } \phi(e_G) = e_H \quad \phi(g) \cdot \phi(g') = \phi(gg')$$

Thus, ϕ is a homomorphism

$$\text{If } \phi: G \rightarrow H$$

then, if $\phi \stackrel{\alpha}{\Rightarrow} \psi$ exists,

$$\textcircled{1} \alpha \in H$$

$$\forall g \in G,$$

$$\begin{array}{ccc} & \xrightarrow{\alpha} & \\ \phi(g) \downarrow & & \downarrow \psi(g) \\ & \longrightarrow & \end{array}$$

$$\textcircled{2} \psi(g) \alpha = \alpha \phi(g)$$

$$\Rightarrow \alpha^{-1} \psi(g) \alpha = \phi(g)$$

Hence, $\alpha \in \text{Inn}(M)$ s.t. $\alpha \circ \psi = \phi$.

Ex 1.4. iii) $(P, \leq) \xrightleftharpoons[F]{F}(Q, \leq)$,

then F, G are order preserving function.

If $F \xrightarrow{\alpha} G$, then $\forall f: P \rightarrow P'$
 $(P \leq P')$

$$F_P \xrightarrow{\alpha_P} G_P \quad \text{i.e. } G_{P'} \geq F_{P'} \geq F_P.$$

$$\downarrow \quad \downarrow \quad \text{if } G_{P'} \geq G_P \geq F_P.$$

$$F_{P'} \xrightarrow{\alpha_{P'}} G_{P'} \Rightarrow G_P \geq F_P \quad \forall P.$$

Thus, natural transformation of F, G
 order preserving function is $F \leq G$.

i.e. $\forall p \in P, F_p \leq G_p$.

(Notes that $F \leq G \Leftrightarrow$ Natural transformation hold.)

(1.4. iv)

Ex 1.4 iv). Each f_*, g_* , f^*, g^* defines
 n.t. by Ex 1.4.7.

If $f_* = g_*$ \Rightarrow $\forall h: Y \rightarrow X$

$$\begin{array}{ccc} C(x, c) & \xrightarrow{f_*} & C(x, d) \\ \downarrow \cdot h & \wr & \downarrow \cdot h \\ C(y, c) & \xrightarrow{f_*} & C(y, d) \end{array}$$

$\Rightarrow \forall l \in C(x, c). \quad l \mapsto flh$

$flh = glh$. Now pick $h = 1_x$

$\Rightarrow \therefore fl = gl$, Pick $x = c, l = 1_c$.

$\Rightarrow f = g$, contradiction. $\therefore f_* \neq g_*$ as
 natural transformation. (f^*, g^* : similar manner)

Ex 1.4 v.

$$\begin{array}{ccccc} F \downarrow G & \xrightarrow{\text{dom}} & D & \xrightarrow{F} & C \\ (d, e, f) & \longmapsto & d & \longrightarrow & Fd \\ (h, k) \downarrow & \longmapsto & \downarrow h & & \downarrow Fh \\ (d', e', f') & \longmapsto & d' & \longrightarrow & Fd' \end{array}$$

$$\begin{array}{ccccc}
 F \downarrow G & \xrightarrow{\text{cod}} & E & \xrightarrow{G} & C \\
 (d, e, f) & \longmapsto & e & \longmapsto & Ge \\
 (h, k) \downarrow & & \downarrow k & & \downarrow Gk \\
 (d', e', f') & \longmapsto & e' & \longmapsto & Ge'
 \end{array}$$

Thus, For each $(d, e, f) \xrightarrow{(h, k)} (d', e', f')$
 we need to find $\alpha_{(d, e, f)}$, $\alpha_{(d', e', f')}$
 s.t.

$$\begin{array}{ccc}
 F_d & \xrightarrow{\alpha_{(d, e, f)}} & Ge \\
 Fh \downarrow & \curvearrowright & \downarrow Gk \\
 F_{d'} & \xrightarrow{\alpha_{(d', e', f')}} & Ge'
 \end{array}
 \quad \text{Commutative}$$

Set $\alpha_{(d, e, f)} = f$. By construction of $F \downarrow G$, it holds.

Ex 1.4 vi) If F, G has different target category,
 i.e., $F: A \times B \times B^{\text{op}} \rightarrow D$ with $D \neq D'$.
 $G: A \times C \times C^{\text{op}} \rightarrow D'$

then, we ~~don't know how to~~ cannot have a
 morphism $F(a, b, b) \rightarrow G(a, c, c)$ in general

1.5. Equivalence of Categories.

$$2: \begin{array}{ccc} \bullet & \longrightarrow & \bullet \\ & & \uparrow \\ & & \bullet \end{array}$$

$$\begin{array}{ccccc} C & \xrightarrow{1_0} & (C, 0) & \xleftarrow{1_1} & (C, 1) & \xleftarrow{1_C} & C \\ & & \downarrow & & \downarrow & & \\ C & \xrightarrow{1_0} & C \times 2 & \xleftarrow{1_1} & C & & \end{array}$$

$$\textcircled{1}: \begin{array}{ccc} & & \\ & \searrow & \downarrow H \\ F & & D \end{array} \quad \begin{array}{ccc} & & \\ & \swarrow & \uparrow G \\ & & D \end{array}$$

Lemma 1.5.1. $F, G: C \Rightarrow D$ Then,

$\{ \alpha: F \Rightarrow G \mid \text{natural transformation} \}$

$\{ H: \text{satisfying diagram} \}$

bijection

2f) Construction of H is in Ex 1.5.1)

To see bijective corresp, notes that given

$$H, \text{ with } (C, 0) \xrightarrow{f \cdot (0 \rightarrow 1)} (C', 1)$$

$H(C, 0) = F_C$, $H(C', 1) = G_{C'}$ by commutative diagram. Now define

$$\alpha_C = H \left((C, 0) \xrightarrow{1_C \cdot (0 \rightarrow 1)} (C, 1) \right).$$

Then, $F_C \xrightarrow{\alpha_C} G_C$. If we show this diagram commutes, done.

$$\textcircled{2}: \begin{array}{ccc} F_C & \xrightarrow{\alpha_C} & G_C \\ \downarrow Ff & & \downarrow Gf \\ F_{C'} & \xrightarrow{\alpha_{C'}} & G_{C'} \end{array}$$

Actually it is taking H on

$$\begin{array}{ccc}
 (C, 0) & \xrightarrow{1_C \cdot (0 \rightarrow 1)} & (C, 1) \\
 f \cdot 1_0 \downarrow & \curvearrowright & \downarrow f \cdot 1_1 \\
 (C', 0) & \xrightarrow{1_{C'} \cdot (0 \rightarrow 1)} & (C', 1)
 \end{array} \Rightarrow \text{Commutative!}$$

Thus, H induces a natural transformation α .

Hence the bijection occurs \square

If $C = 2$, 2×2 is depicted as

$$\begin{array}{ccc}
 (0, 0) & \longrightarrow & (1, 0) \\
 \downarrow & \searrow & \downarrow \\
 (0, 1) & \longrightarrow & (1, 1)
 \end{array}$$

then H sends γ to

$$\begin{array}{ccc}
 F_0 & \xrightarrow{F(\gamma)} & F_1 \\
 \alpha_0 \downarrow & \searrow \gamma & \downarrow \alpha_1 \\
 G_0 & \xrightarrow{G(\gamma)} & G_1
 \end{array}$$

If we change 2 to

$$\text{II: } \begin{array}{ccc} \circ & \xrightarrow{f} & \circ \\ \circ & \xleftarrow{f \cdot 1} & \circ \end{array} \text{ in the}$$

lemma, then

Function satisfying Comm. diagram

\updownarrow Bijection
Natural Isomorphism.

Def 1.5.4 Equivalence of Categories, consists of.

$$\textcircled{1} \quad F: C \rightleftarrows D: G$$

$$\textcircled{2} \quad \eta: 1_C \cong GF, \quad \varepsilon FG \cong 1_D$$

natural isomorphism.

In this case write $C \cong D$.

(cf. $C \cong D$: isomorphism of category)

Lem 1.5.5. $C \cong D$ is equivalence relation

pf) Ex 1.5.6.

Ex 1.5.6.

$$(i) \quad (-)_+ : \text{Set}^2 \longrightarrow \text{Set}_*$$

in 1.3.

$$U : \text{Set}_* \longrightarrow \text{Set}^2$$

are actually equiv of category by

$$1_{\text{Set}^2} = U(-)_+ \quad \text{and}$$

$$\eta: 1_{\text{Set}_*} \cong (U-)_+ \quad \text{with}$$

$$\eta_{(X, X)}: (X, X) \longrightarrow (X \setminus \{x\} \cup \{X \setminus \{x\}\}, X \setminus \{x\})$$

$$y \longmapsto y \quad \text{if } y \neq x.$$

$$x \longmapsto X \setminus \{x\}.$$

$$(2) \text{Mat}_{\mathbb{K}} \xrightleftharpoons[\text{H.}]{\mathbb{K}^{(-)}} \text{Vect}_{\mathbb{K}}^{\text{basis}} \xrightleftharpoons[\text{C.}]{\text{U}} \text{Vect}_{\mathbb{K}}^{\text{fd.}}$$

Where $\mathbb{K}^{(-)} (n \xrightarrow{M} m) \Rightarrow \mathbb{K}^n \xrightarrow{LM} \mathbb{K}^m.$

U : forgetful functor.

C : Sending V-space by choosing a basis.

H : Sending V.s to dim

and linear transformation to matrix over given bases.

Aim : WTS

$$\text{Mat}_{\mathbb{K}} \simeq \text{Vect}_{\mathbb{K}}^{\text{basis}} \simeq \text{Vect}_{\mathbb{K}}^{\text{fd.}}$$

where $\text{Vect}_{\mathbb{K}}^{\text{basis}}$: Category of f.d. v.s with a chosen basis.

Def: 1.5.7. $F: C \rightarrow D$ a functor is

① full if $\forall x, y \in C, C(x, y) \rightarrow D(Fx, Fy)$ is surjective

② faithful " " is injective.

③ Essentially surjective if $\forall d \in D, \exists c \in C$ s.t. $Fc \simeq d.$

Rem 1.5.8

- (1) Full, faithful : local conditions
- (2) Faithful and injective on object = embedding
- (3) full + faithful = fully faithful
- fully faithful + injective on object = full embedding

In case of full embeddings, image of domain
= full subcategory of the codomain

Thm 1.5.9. $f: C \cong D: G \iff f, G$ are full + faithful + ess. surj.

(under axiom of choice)

Lem 1.5.10. For $f: a \rightarrow b$ with $a \cong a', b \cong b'$

$\exists f': a' \rightarrow b'$ st. any arrow in

$$\begin{array}{ccc} a & & a' \\ f \downarrow & & \downarrow f' \\ b & & b' \end{array}$$

the left square makes it commutes.

pf) Ex 1.5.iii)

pf of Thm) Let $f, g: C \Rightarrow C'$ with $ff = fg$.

$$\Rightarrow \begin{array}{ccccc} C & \xrightarrow{\eta_C} & GF C & \xleftarrow{\eta_C} & C \\ f \downarrow & \curvearrowright & GF f \downarrow & GF g \curvearrowright & \downarrow g \\ C' & \xrightarrow{\eta_{C'}} & GF C' & \xleftarrow{\eta_{C'}} & C' \end{array} \quad \begin{array}{l} \text{with } GF f = GF g \\ (\text{since } ff = fg) \end{array}$$

By 1.5.10,
$$\begin{array}{ccccc} C & \xrightarrow[\cong]{\eta_C} & GF_C & \xrightarrow[\cong]{\eta_C^{-1}} & C \\ f \downarrow & & & & \downarrow g \\ C' & \xrightarrow[\eta_{C'}]{\cong} & GF_{C'} & \xrightarrow[\eta_{C'}^{-1}]{\cong} & C' \end{array}$$
 commutes

$\Rightarrow g = f. \Rightarrow F$ is faithful.

(So is G by applying same argument on FG
 \Rightarrow "by symmetry")

Let $g \in D(F_C, F_{C'})$

$$f := \eta_{C'}^{-1}(Gg) \eta_C \in (C_C, C')$$

$$\Rightarrow \begin{array}{ccccc} GF_C & \xrightarrow[\eta_C^{-1}]{\eta_C} & C & \xrightarrow{\eta_C} & GF_C \\ GFf \downarrow & \cong & f \downarrow & \cong & \downarrow Gg \\ GF_{C'} & \xrightarrow[\eta_{C'}^{-1}]{\eta_{C'}} & C' & \xrightarrow{\eta_{C'}} & GF_{C'} \end{array}$$

commutes
from
equivalence
and
def of f and g .

$\Rightarrow Gg = GFf$. From G is faithful,

$g = Ff \Rightarrow F$ is full.

Now, for $d \in D$, $\varepsilon_d : FGd \xrightarrow{\cong} d$

$\Rightarrow F$ is essentially surjective.

By symmetry, G is also full and ess. sur.

Conversely let $F: C \rightarrow D$; full, faithful, and ess. sur. Want to construct $G: D \rightarrow C$ equiv.

By axiom of choice and essential surj.

$$\forall d \in D, \exists c \in C \text{ st. } d \cong Fc$$

$$\text{Let } Gd := c \Rightarrow \varepsilon_d: d \cong FGd$$

Then, given $f \in D(d, d')$

$$\begin{array}{ccc} \text{We have } FGd & \xrightarrow{\varepsilon_d} & d \\ \downarrow g & & \downarrow f \\ FGd' & \xrightarrow{\varepsilon_{d'}} & d' \end{array} \Rightarrow \text{Let } g = \varepsilon_{d'}^{-1} f \varepsilon_d \Rightarrow \text{diagram commutes}$$

Then, from F is fully faithful, $\exists h \in C(FGd, FGd')$ st. $Fh = g$. Let $Gf := h$

Hence, by this definition

$$\begin{array}{ccc} FGd & \xrightarrow{\varepsilon_d} & d \\ FGf \downarrow & \cap & \downarrow f \\ FGd' & \xrightarrow{\varepsilon_{d'}} & d' \end{array} \Rightarrow \varepsilon_{d'} \circ FG \Rightarrow 1_D$$

Claim 1: G is functorial

For first condition,

$$\begin{array}{ccccc}
 FGd & \xrightarrow{\varepsilon_d} & d & \xrightarrow{\varepsilon_d^{-1}} & FGd \\
 \downarrow FG I_d & \circlearrowleft \textcircled{1} & \downarrow I_d & \circlearrowright \textcircled{2} & \downarrow F(I_{Gd}) \\
 FGd & \xrightarrow{\varepsilon_d} & d & \xrightarrow{\varepsilon_d^{-1}} & FGd
 \end{array}$$

① commutes by definition of $G I_d$.

② " functoriality of F .

$$\Rightarrow FG I_d = F(I_{Gd}) \Rightarrow G(I_d) = I_{Gd} \text{ by faithfulness of } F.$$

For second condition,

$$\text{let } f: d \rightarrow d', \quad f': d' \rightarrow d''$$

$$\begin{array}{ccccc}
 \Rightarrow FGd & \xrightarrow{\varepsilon_d} & d & \xrightarrow{\varepsilon_d^{-1}} & FGd \\
 \downarrow GF & & \downarrow f & & \downarrow F(Gf) \\
 FG(f'f) & \xrightarrow{\varepsilon_{f'f}} & f'f & \xrightarrow{\varepsilon_{f'}^{-1}} & FGd' \\
 \downarrow & & \downarrow f' & & \downarrow F(Gf') \\
 FGd'' & \xrightarrow{\varepsilon_{d''}} & d'' & \xrightarrow{(\varepsilon_{d''})^{-1}} & FGd''
 \end{array}$$

① commutes by definition of $G(f')$, $G(f)$, and $G(f'f)$.
 (We define these to commute such diagram)

$$\Rightarrow FG(f'f) = F(G(f')) \cdot F(G(f)) \\ = F(G(f')G(f))$$

$$\Rightarrow G(f'f) = G(f')G(f) \\ \text{by faithfulness of } F.$$

$\Rightarrow G$ is functor.

Claim 2: $\eta: 1_C \Rightarrow GF$ is natural iso.

Define η_c is preimage of $\epsilon_{F_c}^{-1}$ by F .

(It is well-def since F is fully faithful.)

i.e. $F\eta_c = \epsilon_{F_c}^{-1}$, for any $f: C \rightarrow C'$,

$$\begin{array}{ccc} F_c & \xrightarrow{F\eta_c} FG F_c & \xrightarrow{\epsilon_{F_c}} F_c \\ \downarrow Ff & \textcircled{1} \quad \downarrow FG Ff & \textcircled{2} \quad \downarrow Ff \\ F_{c'} & \xrightarrow{\epsilon_{\eta_{c'}}} FG F_{c'} & \xrightarrow{\epsilon_{F_{c'}}} F_{c'} \end{array}$$

② Commutes by naturality of ϵ .

Since ϵ is natural iso, ① also commutes.

$$\Rightarrow FGFf \circ F\eta_c = \epsilon_{\eta_{c'}} \circ FGF = F\eta_{c'} \circ FGF$$

(Inverse - ?)

$$\Rightarrow F(GFf \circ \eta_c) = F(\eta_{c'} \circ f)$$

$\Rightarrow GFf \circ \eta_c = \eta_{c'} \circ f$ by faithfulness of F .
and by Ex 1.5 IV, $\eta_c, \eta_{c'}$ are iso $\Rightarrow \eta$ is natural \square

Corollary 1.5.11. $\text{Mat}_K \cong \text{Vect}_K^{\text{fd}}$ for any field K .

Pf.) $\text{Mat}_K \xrightleftharpoons{\text{basis}} \text{Vect}_K^{\text{fd}} \xrightleftharpoons{\text{fd}} \text{Vect}_K^{\text{fd}}$

Those are ess sur, full, and faithful.

(full: set of matrix $\xleftrightarrow{\text{bijection}}$ all lin trans of other basis
faithful $\xleftrightarrow{\text{bijection}}$ lin trans.)

(ess sur: objects are 1-1 corresp.)

Def) Category is connected if $\forall x, y \in C$.
 x, y connected by a finite zig-zag morphisms.

Prop 1.5.12. Any connected groupoid is equiv. to automorphism gp of any of its object as a category.

2-f) G : groupoid. Fix $g \in \text{Obj } G$.

Let $H = G(g, g)$.

$\Rightarrow BG \hookrightarrow G$ is fully faithful since
 $\bullet \mapsto g$ $BG(\bullet, \bullet) \cong G(g, g)$ by def.
 $\emptyset \begin{bmatrix} \mapsto g \\ \mapsto d \\ \mapsto g \end{bmatrix}$ and ess, surj.
 (Groupoid is connected, (all mor = iso)
 \Rightarrow every pair of objects is isomorphic.)

Cor 1.5.13. X : path-connected.

x : base pt of X .

$$\Rightarrow \pi_1(X, x)$$

Pf) Let $\pi_1(X)$: fundamental groupoid

By Prop 1.5.12, for any $x, x' \in X$

$$\pi_1(X, x) \xrightarrow{\sim} \pi_1(X) \xleftarrow{\sim} \pi_1(X, x')$$

This gives $\pi_1(X, x) \cong \pi_1(X, x')$

Claim: $C \cong D$ with C, D : 1 object category
 $\Rightarrow C \cong D$ as a group. i.e. C, D are group.

Pf) Let $F: C \rightarrow D$ be equivalence

$\Rightarrow C(\cdot, \cdot) \cong D(\cdot, \cdot)$ as bijection.

By 1.3.1) Exercise, F is gp homo.
on $C(\cdot, \cdot)$ and $D(\cdot, \cdot)$

From bijection, F is bijective homo.

\Rightarrow GP Isomorphism

$\Rightarrow C \cong D$ as a group.

(Moreover $C \cong D$ as category since

gp is o F^T induces a functor st. $I_C = FF^T$
 $I_D = F^T F$)

Zen 1.5.14 Top_*^{pc} : path cony spaces

$$\pi_1: Top_*^{pc} \xrightarrow{\pi_1} Group \longleftrightarrow Cat$$

$$\Pi_1: Top_*^{pc} \xrightarrow{U} Top \xrightarrow{\Pi_1} Groupoid \longleftrightarrow Cat.$$

Inclusion of $\pi_1(x, x) \leq \pi_1(y)$ gives

Natural transformation

$$\pi_1 \Rightarrow \Pi_1$$

$$\left(\begin{array}{ccc} \pi_1(x, x) \longleftrightarrow \pi_1(x) & & \\ \pi_1(f) \downarrow & \curvearrowright & \downarrow \Pi_1(f) \\ \pi_1(y, y) \longleftrightarrow \pi_1(y) & & \end{array} \right)$$

And this inclusion is a functor.

Moreover, " is equivalence of categories

since fully faithful (as 1 obj subcategory)

and ess. surj. (from connected groupoid)

However, $\Pi_1(y) \rightarrow \pi_1(x, x)$ inverse equivalence requires axiom of choice for its construction.

(In this case, $\forall p \in X$, choose a path $p \rightarrow x$)

And these chosen paths ($p \rightarrow x$) need not be preserved by morphism in Top_*^{pc} .

Def 1.5.15 C : Category is skeletal

If it contains 1-object in each isomorphism class.
class. $sk C$: a skeletal category equiv to C
(Unique up to iso)

Rem 1.5.16

$sk C$ construction: Choose 1 object in each iso class and $sk C$: full subcategory of C having these objects.

By thm 1.5.9, $sk C \hookrightarrow C$ is full (by def) and faithful (by construction) and ess surj. by choice of representative of iso class.
 $\Rightarrow sk C \cong C$.

But $sk(-) : CAT \rightarrow CAT$ is not a functor.
since $sk(F)$ may not be a functor.

$$(ex) \quad C: \begin{array}{ccc} 0 & \longrightarrow & 1 \cong 2 \\ 0 & & 1 \end{array} \quad F: \begin{array}{ccc} 0 & \longrightarrow & 0 \\ 1 & \longrightarrow & 2 \\ 2 & \longrightarrow & 1 \end{array}$$

$$sk C: \begin{array}{ccc} 0 & & 1 \end{array}$$

$\Rightarrow sk F := F|_{sk C}$ sends 1 to 2 but 2 is not in $sk C$.

Ex 15.17)

(i) G : Connected Groupoid

$$\Rightarrow \text{sk } G = G(g, g) \quad (\text{by construction})$$

(ii) $\text{sk}(P, \geq) = \text{poset!}$ (since we delete any isomorphic but not equal one)
 \uparrow preorder

(iii) $\text{sk}(\text{Vect}_{\mathbb{K}}^{\text{fd}}) \cong \text{Mat}_{\mathbb{K}}$

(every v.s with same dim is isomorphic)

(iv) sk Fin iso (Fin iso: obj: finite set)
 Mor: bijection

= obj: positive integer

Mor: $\text{Hom}(n, n) = S_n$ permutation of n

$\text{Hom}(n, m) = \emptyset$ if $m \neq n$

Ex 15.18

$X: BG \rightarrow \text{Set}$: a left G -set.

Translation Groupoid: $T_G X$: obj = $X(\cdot) \in \text{Set}^G$

Mor: $g: x \rightarrow y$
 if $\exists g \in G$ s.t.

$\Rightarrow \text{sk } T_G X$: obj: connected components in $T_G X$
 i.e. orbit of G -action.

Mor: for two distinct orbit, \emptyset .

Let $x \in X$, O_x : orbit of x

$$\text{Hom}_{\text{sk } T_G X} (O_x, O_x) \cong \text{Hom}_{T_G X} (x, x) =: G_x$$

\uparrow

from $\text{sk } T_G X \cong T_G X$. equivalence.
 \Rightarrow fully faithful

i.e, $\text{Hom}_{\text{sk } T_G X} (O_x, O_x)$ is stabilizer G_x of x .

\Rightarrow Any pair in the same orbit should have
Isomorphic stabilizer.

Also, for any fixed $x \in X$, it has disjoint union

$$\bigcup_{y \in O_x} \text{Hom}_{T_G X} (x, y) = G$$

Since $|\text{Hom}_{T_G X} (x, y)| = |G_x| \quad \forall y$,

$$|O_x| \cdot |G_x| = |G|.$$

\therefore orbit-stabilizer Thm.

Def) C : essentially small $\Rightarrow C \cong D$, D is
small category
 C : " discrete $\Rightarrow C \cong D$, D is
discrete category

C : locally small, $D \simeq C \Rightarrow D$ is locally small
 groupoid \Rightarrow " groupoid.
 $\Rightarrow D^{op} \simeq C^{op}$

$$C \simeq D, \quad C' \simeq D' \Rightarrow C \times C' \simeq D \times D'$$

$$\left(\begin{array}{l} f: x \rightarrow y \in C \text{ iso.} \\ F: C \simeq D \end{array} \right) \Leftrightarrow Ff \text{ is iso.}$$

$F: C \rightarrow D$ fully faithful.

Then "Essential image" of F \leq
 = full subcategory of objects isomorphic to
 some Fc for $c \in C$.

1.5, i). Let $\alpha: F \Rightarrow G$ be a natural tf.

$$\Rightarrow \forall f: c \rightarrow c',$$

$$\begin{array}{ccc} Fc & \xrightarrow{\alpha_c} & Gc \\ Ff \downarrow & \curvearrowright & \downarrow Gf \\ Fc' & \xrightarrow{\alpha_{c'}} & Gc' \end{array} \quad \text{Now define } H \text{ as}$$

$$H: C \times \mathbb{2} \longrightarrow D.$$

$$\begin{array}{ccc} (c, 0) \longmapsto Fc & (c, 1) \longmapsto Gc \\ f \cdot 1_0 \downarrow \longmapsto \downarrow Ff & f \cdot 1_1 \downarrow \longmapsto \downarrow Gf & \text{and} \\ (c', 0) \longmapsto Fc' & (c', 1) \longmapsto Gc' \end{array}$$

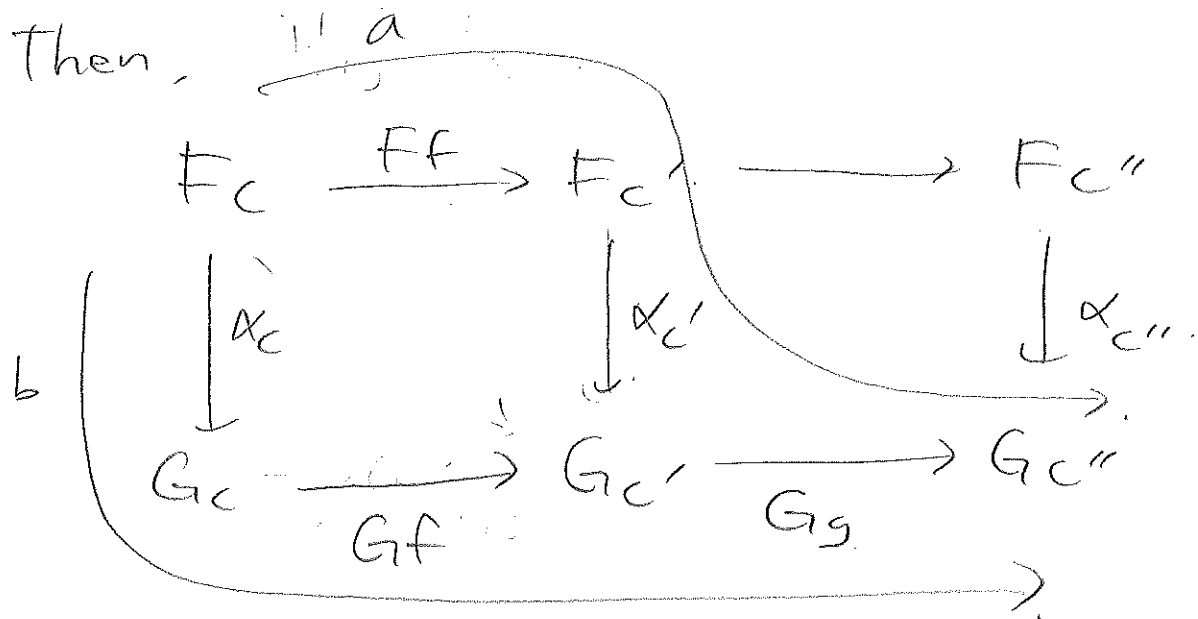
$$\begin{array}{ccc} (c, 0) \longmapsto Fc & \downarrow Ff & \\ f \cdot (0 \rightarrow 1) \downarrow \longmapsto \downarrow & Gf \circ \alpha_c = \alpha_{c'} \circ Ff & \\ (c', 1) \longmapsto Gc' & & \end{array}$$

1) It satisfies functoriality.

All we need are 4 cases. ① $f \cdot 1_0$, $g \cdot (0 \rightarrow 1)$,
② $f \cdot (0 \rightarrow 1)$, $g \cdot 1_1$,

for $f: c \rightarrow c'$, $g: c' \rightarrow c''$

$$\begin{array}{ccc} \textcircled{1} & (c, 0) \xrightarrow{f \cdot 1_0} (c', 0) \xrightarrow{g \cdot (0 \rightarrow 1)} (c'', 1) & (c, 0) \xrightarrow{g \cdot (0 \rightarrow 1)} (c'', 1) \\ \downarrow Ff & \downarrow Gg \circ \alpha_{c'} & \downarrow G(g \cdot 0) = \alpha_{c''} \\ & Fc \xrightarrow{Ff} Fc' \xrightarrow{Gg \circ \alpha_{c'}} Gc'' & Fc \xrightarrow{G(g \cdot 0) = \alpha_{c''}} Gc'' \end{array}$$



a is left diagram, b : right diagram
 \Rightarrow Commutes by natural transformation.

Other case is similar.

$$\begin{aligned} \text{And } H((c, \bar{\alpha}) \xrightarrow{1_{(c, \bar{\alpha})}} (c, \bar{\alpha})) &= \begin{pmatrix} F_c \xrightarrow{F1_c} F_c & \bar{\alpha}=0 \\ G_c \xrightarrow{G1_c} G_c & \bar{\alpha}=1 \end{pmatrix} \\ &= 1_{F_c} \cup 1_{G_c} \\ &= 1_{H(c, \bar{\alpha})}. \end{aligned}$$

So H is a functor.

2) H satisfies $\begin{array}{ccccc} C & \longrightarrow & K \times L & \longleftarrow & C \\ & \searrow F & \downarrow H & \swarrow G & \\ & & D & & \end{array}$

$$\begin{array}{c} C \mapsto F_c \\ f \downarrow \quad \downarrow Ff \\ C' \mapsto F_{c'} \end{array} = \begin{array}{c} C \mapsto (C, 0) \mapsto F_c \\ f \downarrow \quad \downarrow f \cdot 1_0 \quad \downarrow Ff \\ C' \mapsto (C', 0) \mapsto F_{c'} \end{array}$$

and the other way is similar.

1.5. ii.)

$T = \text{Obj}$: finite set.

Mor: $S \longrightarrow T$

$$\Rightarrow \theta: S \longrightarrow P(T) \quad S.t.$$

$$\textcircled{1} \quad \theta(\alpha) \cap \theta(\beta) = \emptyset \text{ when } \alpha \neq \beta.$$

$$S \xrightarrow{\theta} T \xrightarrow{\phi} U = S \xrightarrow{\psi} U$$

$$S.t., \quad \alpha \mapsto \bigcup_{\beta \in \theta(\alpha)} \phi(\beta)$$

$$\text{Fin} T \longrightarrow (\text{Fin}_*^{\text{op}})$$

$$G: \text{Fin}_*^{\text{op}} \longrightarrow \Gamma$$

$$S \longrightarrow (S \cup \{s\}, s)$$

$$(S, s) \longmapsto S \setminus s$$

$$f \downarrow \longmapsto \uparrow f^T$$

$$T \longrightarrow (T \cup \{t\}, t)$$

$$(T, t) \longmapsto T \setminus t$$

where

$$\theta^{-1}(\beta) = \begin{cases} \alpha & \text{if } \beta \in \theta(\alpha) \\ s & \text{o.w.} \end{cases} \quad f^T: \beta \mapsto f^T(\beta)$$

First of all, $\theta^{-1}(\beta)$ is well-defined since no element in T is contained in two preimage by $\textcircled{1}$.

Also, f^T is well-defined since it is a map from $T \setminus t$ to $P(S \setminus s)$ with distinct preimage.

To see $T \subseteq \text{Fin}_*^{\text{op}}$.

Notes that $FG: \text{Fin}_*^{\text{op}} \rightarrow T \rightarrow \text{Fin}_*^{\text{op}}$

$$(S, s) \mapsto S \setminus s \mapsto (S \setminus s \cup \{s\}, s)$$

$$f \downarrow \quad \mapsto \quad \uparrow f^{-1} \quad \mapsto \quad \downarrow h$$

$$(T, t) \mapsto T \setminus t \mapsto (T \setminus t \cup \{t\}, t)$$

To figure out h , notes that

If $x \in S \setminus s$, let $y = f(x)$.

$$\Rightarrow f^{-1}(y) = \{x \in S \setminus s : f(x) = y\}$$

Thus, for any $x \in f^{-1}(y)$

$$h(x) = y.$$

Hence, $f|_{S \setminus s} = h|_{S \setminus s}$ and $h(S \setminus s) = T \setminus t$.

Thus, $(S \setminus s \cup \{s\}, s) \xrightarrow{\alpha_{(S, s)}} (S, s)$

define

$$\alpha_{(S, s)}$$

$$h \downarrow$$

$$\downarrow f$$

as usual incl with

$$(T \setminus t \cup \{t\}, t) \longrightarrow (T, t)$$

$$S \setminus s \mapsto s$$

$$T \setminus t \mapsto t$$

Then, if $x \in S \setminus s$ with $f(x) = y$.

$$\alpha_{(t,t)} \circ h(x) = \alpha_{(t,t)}(y) = y.$$

$$f \circ \alpha_{(s,s)}(x) = f(x) = y.$$

and for $x = s$,

$$\alpha_{(t,t)} \circ h(s) = \alpha_{(t,t)}(T \setminus t) = t.$$

$$f \circ \alpha_{(s,s)}(s) = f(s) = t.$$

\therefore diagram commutes.

Likewise

$$\begin{array}{ccccc} GF: T & \longrightarrow & Fm_*^{op} & \longrightarrow & T \\ s & \longmapsto & (S \cup \{s\}, s) & \longmapsto & s, \\ \theta \downarrow & & \uparrow \theta^{-1} & & \downarrow h \\ T & \longmapsto & (T \cup \{t\}, t) & \longmapsto & T \end{array}$$

S.t. for $x \in S$, $\theta(x) \in T$. Hence, $\forall y \in \theta(x)$

$$\theta^{-1}(y) = x. \Rightarrow h(x) = \text{preimage of } x \text{ under } \theta^{-1} = \theta(x). \quad \text{Monoid}$$

$$\Rightarrow GF = 1_T.$$

$$\text{Hence } T \cong Fm_*^{op}.$$

In particular, $Fm_* \xrightarrow{M^-} \text{Set}$ is a functor in 1.3.2(xi) and G induce a functor $\hat{G}: T^{op} \rightarrow Fm_*$
 $\Rightarrow T^{op} \xrightarrow{\hat{G}} Fm_* \xrightarrow{M^-} \text{Set}$ are presheaves on T .

1.5. iii) Any morphism $f: a \rightarrow b$ and $\left(\begin{array}{l} \text{fixed morphisms} \\ a \cong a', b \cong b' \end{array} \right)$

\Rightarrow determine $f': a' \rightarrow b'$ so that

$$\begin{array}{ccccc} a \xleftarrow{\cong} a' & a \rightarrow a' & a \xleftarrow{\cong} a' & a \rightarrow a' & a \xrightarrow{\cong} a' \\ f \downarrow & f \downarrow & f \downarrow & f \downarrow & f \downarrow \\ b \xrightarrow{\cong} b' & b \rightarrow b' & b \xleftarrow{\cong} b' & b \rightarrow b' & b \xrightarrow{\cong} b' \end{array}$$

pf) Define $f' = B \circ f \circ A$

$$\Rightarrow B \circ f = f' \circ A^{-1}, \quad B^{-1} \circ f' = f \circ A, \quad f' \circ B^{-1} \circ f' \circ A^{-1} = f.$$

1.5. iv) $F: C \rightarrow D$ full and faithful

(i) $f: c \rightarrow c'$ mon $M \in C$ s.t. Ff is iso in D

$\Rightarrow f$ is iso.

(ii) $x, y \in C$ s.t. $F_x \cong F_y$ in D

$\Rightarrow x \cong y$ in C

*f) Since F is full and faithful,

$C(x, y)$ and $D(F_x, F_y)$ are bijective.

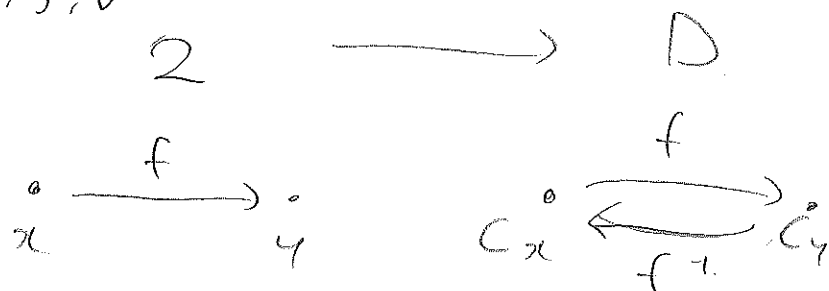
$$\text{So, if } Fg \in D(F_y, F_x) \text{ s.t. } Fg \circ Ff = 1_{F_x} \\ Fg \circ Fg = 1_{F_y}.$$

Then $F(gf) = 1_{F_x}$ Since $C(x, x)$ and $D(F_x, F_x)$
 $F(fg) = 1_{F_y}$ $C(y, y)$ " $D(F_y, F_y)$ are
 bijective, $gf = 1_x, fg = 1_y$.

Similarly, if $F_x \cong F_y$, $\exists f \in D(F_x, F_y)$ which is isomorphism. By (i), f is iso in \mathcal{C} , thus $x \cong y$.

~~(Function preserve iso)~~
 Lem 1.3.8 is converse of this statement.

Ex 1.5.v.



It is faithful since $2(x, y) \hookrightarrow D(C_x, C_y)$.

but not reflect since $x \rightarrow y$ in 2 is not iso, but in D it is.

Ex 1.5. vi)

(i) Composite of full, faithful, ess surj. again same. $F: \mathcal{C} \rightarrow D$, $G: D \rightarrow E$.

$$\text{pf)} \quad \mathcal{C}(x, y) \xrightarrow{\hookrightarrow} D(C_x, C_y) \xrightarrow{\hookrightarrow} E(DC_x, DC_y)$$

$$\Rightarrow \mathcal{C}(x, y) \xrightarrow{\hookrightarrow} E(DC_x, DC_y)$$

And for essential surj, $\forall e \in E$, $\exists d \in D$ st. $e \cong Gd$.
 and for $d \in D \exists c \in \mathcal{C}$ st. $d \cong Fc$.
 $\Rightarrow e \cong Gd \cong GFc$ in E . $\forall e \in E \exists c$ st. $GFc \cong e$.

(ii). It is just from 1.5.9. Thm.

if we assume axiom of choice

Also, $C \cong C$, $C \cong D \Rightarrow D \cong C$. trivially.

Ex 1.5.vii) Connected groupoid \Rightarrow all objects are isomorphic.
if $G = G(g, g)$.

$$\begin{array}{ccc} G & \longrightarrow & BG \\ h & \xrightarrow{\alpha_h} & \bullet \\ f \downarrow \cong & & \downarrow \alpha_h \\ h' & \xrightarrow{\quad} & \bullet \end{array} \quad \begin{array}{ccc} & \xrightarrow{f} & h \xrightarrow{(\alpha_{h'})^{-1}} g \\ & & \downarrow \alpha_h \end{array} \quad \begin{array}{ccc} BG & \longrightarrow & G \\ \bullet & \xrightarrow{\quad} & g \\ \emptyset \downarrow & \xrightarrow{\quad} & \downarrow \emptyset \\ \bullet & \xrightarrow{\quad} & g \end{array}$$

$$\emptyset \in G(g, g)$$

Then,

$$BG \longrightarrow G \longrightarrow BG = 1_{BG}$$

$$\begin{array}{ccc} \bullet & \xrightarrow{\quad} & g \xrightarrow{\quad} \bullet \\ \emptyset \downarrow & \xrightarrow{\quad} & \downarrow \emptyset \\ \bullet & \xrightarrow{\quad} & g \xrightarrow{\quad} \bullet \end{array}$$

(or take $\alpha: 1_{BG} \Rightarrow BG \rightarrow G \rightarrow BG$
as identity.)

$$\begin{array}{ccc} G & \longrightarrow & BG \longrightarrow G \\ h & \xrightarrow{\quad} & \bullet \xrightarrow{\quad} g \\ \psi \downarrow & \xrightarrow{g \rightarrow h \xrightarrow{\psi} h' \rightarrow g} & \downarrow g \rightarrow h \xrightarrow{\psi} h' \rightarrow g \\ h' & \xrightarrow{\quad} & \bullet \xrightarrow{\quad} g \end{array}$$

we need to find Iso. satisfy,

Thus,

$$\begin{array}{ccc}
 g & \xrightarrow{\alpha_h} & h \\
 \downarrow & \curvearrowright & \downarrow \psi \\
 g \xrightarrow{\psi} h \xrightarrow{\psi} g & & h' \\
 \downarrow & & \\
 g & \longrightarrow & h'
 \end{array}$$

And take

$$\alpha_h = g \rightarrow h$$

$$\alpha_{h'} = (h' \rightarrow g)^{-1}$$

Since such $\alpha_h, \alpha_{h'}$ is fixed by construction of $G \rightarrow BG$, it is well-def and satisfy the commuting diagram

Ex 1.5 viii) Later

Ex 1.5 ix) Any category equiv to locally small category is locally small.

Pf) Let $C \xrightleftharpoons[G]{F} D$, equiv of category

By Thm 1.5.9, F, G : full, faithful.

\Rightarrow If assume D locally small,

then $C(x, y)$ is a set since it is biject to a set $D(x, y)$. $\Rightarrow C$ is locally small

1.5. viii).

Affine : Obj = affine planes, $A^2 K$
where K is a field.

Mor = affine linear map.

$L \rightarrow C$: L : linear isomorphism

Proj¹ : Obj = projective planes $(\mathbb{P}^2 K, \ell)$
w/ ℓ : line of infinity.
 C : Constant function

Mor = projective linear isomorphism

$F : \text{Proj}^1 \longrightarrow \text{Affine}$

$(\mathbb{P}^2 K, \ell) \longmapsto A^2 K$ by deleting ℓ in $\mathbb{P}^2 K$.

Notes that $\mathbb{P}^2 K = K^2 \sqcup (K^1 \sqcup P^+)$

where $K^2 = \{ [x:y:z] : z \neq 0 \} \longmapsto \{ (\frac{x}{z}, \frac{y}{z}) \in A^2 K \}$

$K = \{ [x:y:0] : y \neq 0 \}$

$P^+ = \{ [x:0:0] \}$

with $\ell = K \sqcup P^+$.

Thus, $G : \text{Affine} \longrightarrow \text{Proj}^1$

is embedding $A^2 K$ into $(\mathbb{P}^2 K, \ell)$

by identifying $A^2 K$ as K^2 part.

$(x, y) \longmapsto [x:y:1]$

Claim 1: F, G maps pt to pt, line to line.
(except ℓ).

Let L be a line in \mathbb{P}^2_K generated by

$$\vec{a} = [a_0 : a_1 : a_2], \quad \vec{b} = [b_0 : b_1 : b_2]$$

Then, $L = \{ u\vec{a} + v\vec{b} : u, v \in K, (u, v) \neq (0, 0) \}$

Case 1: If $\vec{a}, \vec{b} \in K^2$, F maps $u\vec{a} + v\vec{b}$ into $\left(\frac{u\frac{a_0}{a_2} + v\frac{b_0}{b_2}}{u+v}, \frac{u\frac{a_1}{a_2} + v\frac{b_1}{b_2}}{u+v} \right)$

$$\text{Thus, } F(L) = \left\{ t \left(\frac{a_0}{a_2}, \frac{a_1}{a_2} \right) + (1-t) \left(\frac{b_0}{b_2}, \frac{b_1}{b_2} \right) \right\}$$

Since $\frac{u}{u+v}, \frac{v}{u+v}$ can be mapped into t and

$1-t$, for any $t \in K$,

$$= \left\{ \left(\frac{b_0}{b_2}, \frac{b_1}{b_2} \right) + t \left(\frac{a_0}{a_2} - \frac{b_0}{b_2}, \frac{a_1}{a_2} - \frac{b_1}{b_2} \right) \right\}$$

Hence $F(L)$ is a line in \mathbb{A}^2_K .

Case 2: If one of \vec{a}, \vec{b} is in ℓ .

Then, by adding suitable $u\vec{a}$ to \vec{b} ,

we can change this as a case 1.

Case 3: Both \vec{a}, \vec{b} are in ℓ

Then, $L = \ell$ and F drops ℓ .

Now, let L be a line in A^2_K .

$$\Rightarrow L = \{tA + (1-t)B : t \in K\}$$

for some $A = [a_0; a_1]$, $B = [b_0; b_1]$

$$\Rightarrow G(L) = \left\{ t [a_0 : a_1 : 1] + (1-t) [b_0 : b_1 : 1] \right. \\ \left. t \in K \right\}.$$

To see that $G(L)$ is contained in a line in \mathbb{P}^2_K , let L' be a line gen by

$$[a_0 : a_1 : 1], [b_0 : b_1 : 1]. \text{ Then } L' \supseteq G(L)$$

So $G(L)$ matches with only one line L' .

$$\text{Thus, } FG(L) \subseteq F(L') = L.$$

$$\text{and } GF(L) = G\left(t\left(\frac{a_0}{a_2}, \frac{a_1}{a_2}\right) + (1-t)\left(\frac{b_0}{b_2}, \frac{b_1}{b_2}\right)\right)$$

$$\subseteq L' \text{ where } L' \text{ is gen by}$$

$$\left[\frac{a_0}{a_2}, \frac{a_1}{a_2}, 1\right], \left[\frac{b_0}{b_2}, \frac{b_1}{b_2}, 1\right]$$

$$= [a_0 : a_1 : a_2], [b_0 : b_1 : b_2]$$

$$\Rightarrow = L.$$

Hence, $FG(A^2_K) = A^2_K$ since they
 $GF(\mathbb{P}^2_K) = \mathbb{P}^2_K$ preserves pt and lines.
Thus, equivalence is induced by α as identity for each object.

1.5. xi) See Haine, P. "Examples of some properties of functors"

	Full	Faithful	Essentially Surj.
$Ab \rightarrow Gr$	✓	✓	✗
$Rng \rightarrow Ab$	✗	✓	✗
$Rng \xrightarrow{(-)^*} Gr$	✗	✗	✗
$Ans \rightarrow Png$	✗	✓	✗
$Field \rightarrow Rng$	✓	✓	✗
$Mod_R \rightarrow Ab$	dep on R	✓	✓

(b) $Ring \rightarrow Ab$: ① Not full since. No ring homo $\mathbb{Z}/n\mathbb{Z} \rightarrow \mathbb{Z}$.

Since $f(\underbrace{1+\dots+1}_n) = 0$ but $f(1)+\dots+f(1)=n$

Contradiction.

(However, Rng on Ab \exists trivial homo.)

② Not essentially Surj:

Let $\mathbb{Z}(p^\infty) := \{ e^{2\pi i m/p^n} : m, n \in \mathbb{N} \}$

Prüfer p -group.

If $U(R) \cong \mathbb{Z}(p^\infty)$ then $1 \mapsto a \in \mathbb{Z}(p^\infty)$

$\Rightarrow |1| = n$ in $R \Rightarrow \forall r \in R, nr = 0$

Since $r(1+\dots+1) = r \cdot 0 = 0$

Thus every element in $U(R)$ has order at most n . But $\mathbb{Z}(p^\infty)$ has element with greater order. \square

(C) $(-)^{\times} \text{Rings} \rightarrow \text{Grp}$.

① Not full.

Since \mathbb{Z} is initial in Rings , $\forall R \in \text{Rings}$
 $\exists! f: \mathbb{Z} \rightarrow R \Rightarrow |\text{Rings}(\mathbb{Z}, R)| = 1$

But $\mathbb{Z}^{\times} \cong \mathbb{Z}/2$. Take $R = \mathbb{F}_p$.
finite p -field.

$$\Rightarrow R^{\times} \cong \mathbb{Z}/(p-1)$$

And $|\text{Grp}(\mathbb{Z}^{\times}, R^{\times})| = 2$ since 0

and $\mathbb{Z}/2 \rightarrow \mathbb{Z}/(p-1)$ by $1 \mapsto \frac{p-1}{2}$
are the homomorphisms.

② Not faithful.

Claim 1: $\phi: K[t] \rightarrow K[t]$ automorphisms
fixing $K \Rightarrow \phi(t) = at + b$.

pf) $\phi(t)$ cannot be degree more than one
otherwise it is not automorphism.

Also, $\phi(t)$ cannot have degree 0 by sneaky ϕ .

And all $\phi(t) = at + b$ is auto since

$$\phi\left(\frac{1}{a}t - b\right) = t.$$

□

Thus by claim 1,

$\text{Ring}(K[t], K[t])$ is determined by
 $\{at + b : a \in K^x, b \in K\}$

And $K[t]^x = K^x$.

But $\text{Grp}(K^x, K^x)$ is determined by

generators of K^x . (since K^x is cyclic)

and any map in $\text{Ring}(K[t], K[t])$

induces map K^x, K^x fixing all K^x .

Thus it is not faithful.

③ Not ess. surj.

[Pearson, Schneider, 1970]

Not every cyclic gp is isomorphic to
the gp of units of some ring.

$(\mathbb{Z}/5 \not\cong R^x \quad \forall R \in \text{Ring})$

(d) $\text{Ring} \rightarrow \text{Ring}$.

① Not full. : Zero homomorphism
is not a homomorphism
in Ring

② Faithful : Any unital ring homo
is also ring homo
and distinct unital ring homos
are " ring homos.

③ Not ess. surj. : Rings without mult.
identity is not
iso to unital rings.

(e) $\text{Field} \leftrightarrow \text{Ring}$.

① fully faithful.

Since every field homo is
ring homo.

② Not ess. surj. : Not every ring
is a field.

(7) $U: \text{Mod}_R \rightarrow \text{Ab}$.

① faithful: Any distinct R -homo
is also distinct \mathbb{Q} -homo.

(Since distinction determined
by value, not property.)

② Ess surj: $\forall A \in \text{Ab}$, make
trivial R -module structure
st. $\forall a \in A, \forall r \in R$
 $ra = a$.

③ Not necessarily full.

$$\text{End}(R) \cong R \quad \text{in } \text{Mod}_R$$

but $\text{End}(R) \not\cong R$ in Ab . in general.

(If $R \neq \mathbb{Z}$, then it is full.)

1.6. Art of the diagram chase.

Diagram (informal): directed graph.

Commutative (") : any two composable arrows are the same.

Def 1.6.2 (Monoid) $M \in \text{Set}$ with

$\mu: M \times M \rightarrow M$, $\eta: I \rightarrow M$ s.t.

$$M \times M \times M \xrightarrow{I \times \mu} M \times M \xrightarrow{\mu \times I} M \xrightarrow{\eta \times I} M \times M \xleftarrow{I \times \eta} M$$

$$\begin{array}{ccc} \mu \times I_M \downarrow & \circlearrowright & \downarrow \mu \\ M \times M & \xrightarrow{\mu} & M \end{array} \quad , \quad \begin{array}{ccc} & \circlearrowright & \\ I_M \searrow & & \downarrow \\ & M & \swarrow I_M \end{array}$$

μ : multiplication.

$\eta: I \rightarrow M$ I : singleton.

thus $\eta(I)$ is multiplicative identity.

Def 1.6.3 (Topological monoid) $M \in \text{Top}$ with the same commutative diagram (so M is cts)

(Unital rings) $R \in \text{Ab}$ with $R \otimes_R R$

instead of $R \times R$ (Monoidal structure)

(K -algebras) $R \in \text{Vect}_K$ with $R \otimes_K R$

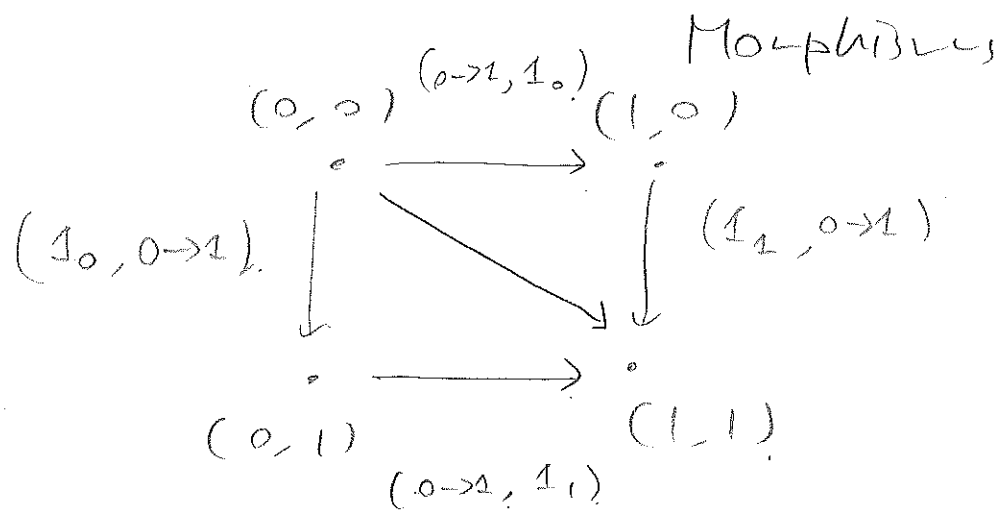
instead of $R \times R$ (")

Def 1.6.4 Diagram of $C = F: J \rightarrow C$
 where J : indexing category is small.

Diagram is commutative.

\Rightarrow Any composite relation in J
 must hold at C via F .

Ex 1.6.6 2×2 : Obj $(0,0), (0,1), (1,0), (1,1)$



Notes that $(1_1, 0 \rightarrow 1) \circ (0 \rightarrow 1, 1_0)$
 $= (0 \rightarrow 1, 0 \rightarrow 1)$
 $= (0 \rightarrow 1, 1_1) \circ (1_0, 0 \rightarrow 1)$

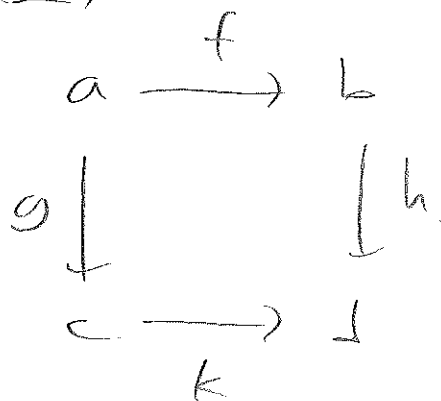
So diagonal arrow is unique.

"Commutative Square!"

Def 1.6.7: "Shape" as indexing category

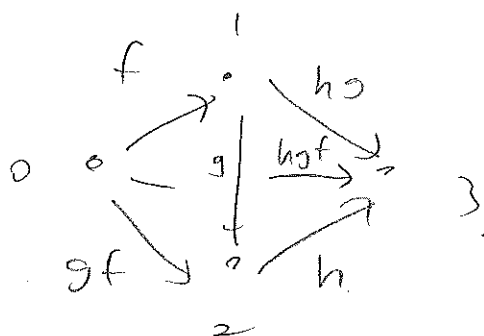
Shape = directed graph with
 specified commutativity relation.

ex) 2×2

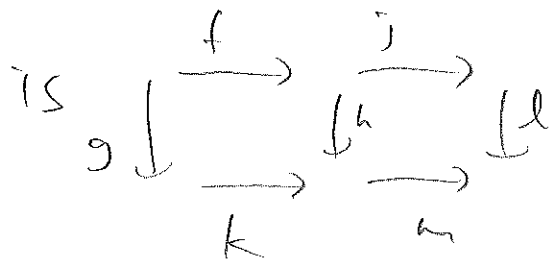
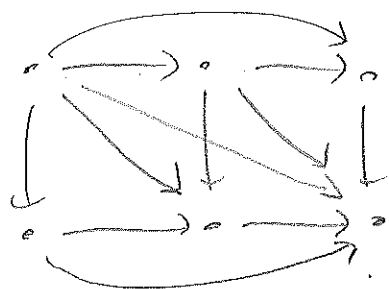


with $hf = kg$

4



2×3



with $hf = ks$, $lj = mh$

Lemma 1.6.11. $f_1 \dashv\dashv f_n$: composable path.

$$f_{k+1} \dashv\dashv f_n = g_m \dashv\dashv g_i$$

$$\Rightarrow f_n \dashv\dashv f_1 = f_n \dashv\dashv f_{k+1} g_m \dashv\dashv g_i$$

pf) $g=h \Rightarrow fg=fh$ for any composable f .

Diagram chasing : showing two paths are equal.

Lemma 1.6.12

$$\begin{array}{ccc} & & h \\ f \downarrow & \searrow & \downarrow \\ & a & \\ & \nearrow & \downarrow \\ & & g \end{array}, f \text{ iso} \Rightarrow \begin{array}{ccc} & & h \\ f^{-1} \uparrow & \searrow & \downarrow \\ & a & \\ & \nearrow & \downarrow \\ & & g \end{array}$$

(Dually,

$$\begin{array}{ccc} & i & \\ k \uparrow & \nwarrow & \nearrow \\ & a & \\ & \nwarrow & \nearrow \\ & j & \end{array} \text{ and } k \text{ iso} \Rightarrow \begin{array}{ccc} & i & \\ k^{-1} \uparrow & \nwarrow & \nearrow \\ & a & \\ & \nwarrow & \nearrow \\ & j & \end{array})$$

pf). $gf = h \Rightarrow gff^{-1} = hf^{-1} \Rightarrow g = hf^{-1}$

Lemma 1.6.13.

$$\begin{array}{ccc} & \alpha & \\ \gamma \downarrow & \xrightarrow{\alpha} & \downarrow \beta \\ & \delta & \end{array} \Rightarrow \alpha^{-1}\beta^{-1} = \gamma^{-1}\delta^{-1}$$

pf) $\beta\alpha = \delta\gamma \Rightarrow \alpha^{-1}\beta^{-1} \cdot (\beta\alpha) \gamma^{-1}\delta^{-1} = \alpha^{-1}\beta^{-1}(\delta\gamma)\gamma^{-1}\delta^{-1}$
 $\Rightarrow \gamma^{-1}\delta^{-1} = \alpha^{-1}\beta^{-1}$

Def 1.6.14. $i \in C$ is initial if $\forall c \in C$

$$\exists ! i \rightarrow c$$

$t \in C$ is terminal if $\forall c \in C$

$$\exists ! c \rightarrow t$$

Ex 1.6.15.

Category	initial	terminal
Set	\emptyset	singleton

Category	Initial	Terminal
Set	\emptyset	Singleton
Top	"	"
Set*	Singleton	
Mod _R	0	
Group	0	
Rng	\mathbb{Z}	0
Rng (non unital)	0	0
Field	Do not exist. (diff characteristic \Rightarrow No homo)	
Cat	\emptyset	\mathbb{I}
(P, \leq)	global minimum	global maximum
	(if exist)	

Lem 1.6.16. $f_1 \dots f_n$ composable seq.
 $g_1 \dots g_m$

$$\text{s.t. } \text{dom}(f_1) = \text{dom}(g_1), \quad \text{cod}(f_n) = \text{cod}(g_m)$$

If either $\text{dom}(f_1) = \text{initial}$ or $\text{cod}(f_n) = \text{terminal}$

$$\Rightarrow f_n \dots f_1 = g_m \dots g_1 \quad \text{pf) Uniqueness of morphism from / to initial / terminal.}$$

Def 1.6.17. \mathcal{C} : Concrete Category

If $U: \mathcal{C} \rightarrow \text{Set}$ a faithful functor exists.

Ex 1.6.18: (\mathcal{C} Concrete Category) = Ex 1.1.3.

Graph: $U \sqcup E: \text{Graph} \rightarrow \text{Set}$ is faithful

Lemma 1.6.19: $U: \mathcal{C} \rightarrow \mathcal{D}$ faithful then

for any diagram in \mathcal{C} whose image

commutes in \mathcal{D} also commutes in \mathcal{C} .

pf) Let $f_1, \dots, f_n, g_1, \dots, g_m$ parallel seq
of composable morphisms s.t.

$$Uf_1, \dots, Uf_n = Ug_1, \dots, Ug_m$$

$$\Rightarrow U(f_1, \dots, f_n) = U(g_1, \dots, g_m) \text{ by Functoriality}$$

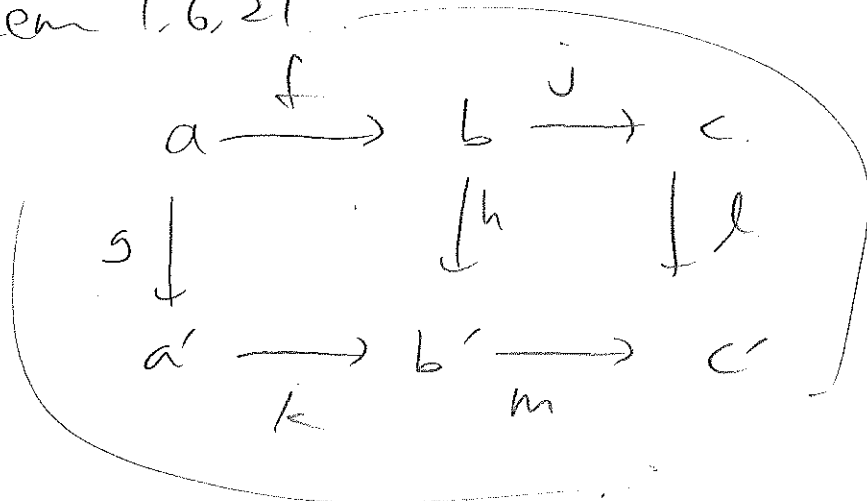
$$\Rightarrow f_1, \dots, f_n = g_1, \dots, g_m \text{ by faithfulness.}$$

Rem 1.6.20. Even outer rectangular commutes
inner rectangular may not commute

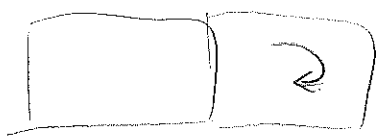

ex)

$$\begin{array}{ccccc} \mathbb{Z} & \xrightarrow{1_{\mathbb{Z}}} & \mathbb{Z} & \longrightarrow & 0 \\ \downarrow & & \downarrow 1_{\mathbb{Z}} & & \downarrow \\ 0 & \longrightarrow & \mathbb{Z} & \xrightarrow{1_{\mathbb{Z}}} & \mathbb{Z} \end{array}$$

Lemma 1.6.21



and $ljf = mkg$.

If either ①  or ②  $f: \text{epi}$

$m: \text{mono}$

then the diagram commutes

pf) Assume ①: $ljf = mkg \Rightarrow hf = kg$
 \parallel
 mkg by mono.

② is dual case of ①.

Ex 1.6.1) Let \bar{t} , initial, t : terminal

$\Rightarrow \exists ! g: \bar{t} \rightarrow t$ and ~~$\exists ! f: t \rightarrow \bar{t}$~~

If $\exists f: t \rightarrow \bar{t} \Rightarrow g \circ f: t \rightarrow t$

Since $t \rightarrow t$ is unique, $gf = 1_t$.

Also, $fg: \bar{t} \rightarrow \bar{t}$ is unique $\Rightarrow fg = 1_{\bar{t}}$.

Ex 1.6.ii) Let t_1, t_2 are the terminal object. $\Rightarrow \exists ! f: t_1 \rightarrow t_2$ and $\exists ! g: t_2 \rightarrow t_1$

Thus $fg: t_2 \rightarrow t_2$ unique $\Rightarrow fg = 1_{t_2}$
 $gf: t_1 \rightarrow t_1$ " $\Rightarrow gf = 1_{t_1}$

Ex 1.6.iii) Let $f: C \rightarrow C'$ Set.

$Ff: C \rightarrow C'$ is mono in \mathcal{P} .

Let $g_1, g_2 \in (b, c)$ Set.

$$fg_1 = fg_2$$

$$\Rightarrow F(fg_1) = F(fg_2) \Rightarrow Ff Fg_1 = Ff Fg_2$$

$$\Rightarrow Fg_1 = Fg_2 \quad \text{by } Ff \text{ is mono}$$

$$\Rightarrow g_1 = g_2 \quad \text{by faithfulness}$$

$$\Rightarrow f \text{ is mono in } C.$$

Thus, if C is concrete category, then faithful

$U: C \rightarrow \text{Set}$ exists. thus if $f \in \text{Mor } C$

st. $Uf = \text{injection}$, then f is mono.

By duality, faithful functor reflects epi.

Ex 1.6. IV) C : category $f: c \rightarrow c'$.
 not epi or mono.
 $2: 0 \rightarrow 1$: iso.

$$F: 2 \rightarrow C \quad F(0 \rightarrow 1) = f.$$

$$\begin{array}{ccc} 0 & & c \\ \downarrow & \mapsto & \downarrow f \\ 1 & & c' \end{array} \Rightarrow \text{neither epi or mono.}$$

Ex 1.6. V) DN : Category of divisible group.
 $(G, +)$ is divisible if $\forall n \in \mathbb{N}, g \in G$,
 $\exists y \in G$ s.t. $ny = g$.

Let: $\pi: \mathbb{Q} \rightarrow \mathbb{Q}/\mathbb{Z}$ and $f, g: G \rightarrow \mathbb{Q}$ s.t.
 $\pi \circ f = \pi \circ g$. Let $x \in G$.

$$\Rightarrow \pi \circ f(x) = \pi \circ g(x) \Rightarrow f(x) - g(x) = n \in \mathbb{Z}.$$

$\nexists n \neq 0$.
 By divisibility, $\exists y \in G$ s.t. $2ny = x$.

$$\Rightarrow f(2ny) = g(2ny) \Rightarrow \frac{1}{2n} f(x) = \frac{1}{2n} g(x)$$

$$\Rightarrow f(y) - g(y) = \frac{1}{2n} (f(x) - g(x)) = \frac{1}{2},$$

contradiction.

$$\Rightarrow n=0, \therefore f(x) = g(x), \forall x.$$

So π is mono. but not injective.

Also, $\pi: \mathbb{Z} \rightarrow \mathbb{Q}$ in \mathbf{Rho} is epi
but not surjective.

$\text{tsm} = \text{tsm}$

Suppose $f, g: \mathbb{Q} \rightarrow R$ s.t. $f \circ \pi = g \circ \pi$.

$$\begin{aligned} \text{Then, } f\left(\frac{a}{b}\right) &= f\left(\frac{1}{b}\right) f(a) = f\left(\frac{1}{b}\right) \cdot g(a) \\ &= g\left(b \frac{a}{b}\right) f\left(\frac{1}{b}\right) = g\left(\frac{a}{b}\right) g(b) f\left(\frac{1}{b}\right) \\ &= g\left(\frac{a}{b}\right) f(b) f\left(\frac{1}{b}\right) \\ &= g\left(\frac{a}{b}\right) \end{aligned}$$

(Assume comm ring; but it holds for any associative unital ring)

Ex. $\therefore f = g$.

1.6. VI) Let (C, γ) be a terminal.

Then for any (d, ϕ) algebra.

$$\exists! f: (d, \phi) \xrightarrow{f} (C, \gamma) \text{ s.t.}$$

$$\begin{array}{ccc} d & \xrightarrow{\quad} & C \\ \phi \downarrow & \curvearrowright & \downarrow \gamma \\ Td & \xrightarrow{\quad} & TC \\ & T_f & \end{array}$$

Thus $C \xrightarrow{\gamma} TC \Rightarrow TC \xrightarrow{T\gamma} T^2C$
 has a coalgebra map (unique)

$$f' : (TC, T\gamma) \longrightarrow (C, \gamma)$$

s.t.

$$\begin{array}{ccc} TC & \xrightarrow{f} & C \\ T\gamma \downarrow & \curvearrowright & \downarrow \gamma \\ T^2C & \xrightarrow{Tf} & TC \end{array}$$

$$\text{Thus, } Tf \circ T\gamma_i = \gamma \circ f$$

And $C \xrightarrow{\gamma} TC$ by functoriality

$$\begin{array}{ccc} C & \xrightarrow{\gamma} & TC \\ \gamma \downarrow & \curvearrowright & \downarrow T(\gamma) \\ TC & \xrightarrow{T(\gamma)} & T^2C \end{array}$$

$$\Rightarrow T\gamma : (C, \gamma) \longrightarrow (TC, T\gamma) \text{ is a morphism}$$

$$\Rightarrow f \circ \gamma : C \longrightarrow C \text{ s.t.}$$

$$\begin{array}{ccc} C & \xrightarrow{f \circ \gamma} & C \\ \gamma \downarrow & \curvearrowright & \downarrow T\gamma \\ TC & \xrightarrow{T(f \circ \gamma)} & TC \end{array}$$

Since (C, γ) is terminal,
 $f \circ \gamma = 1_C$

Hence, $\sigma \circ f = T \circ f \circ T_2 = \frac{T(f \circ 2)}{\text{(first square)}}$
 $= T(1_c)$
 $= 1_{Tc}$

$\Rightarrow \sigma$ is iso.

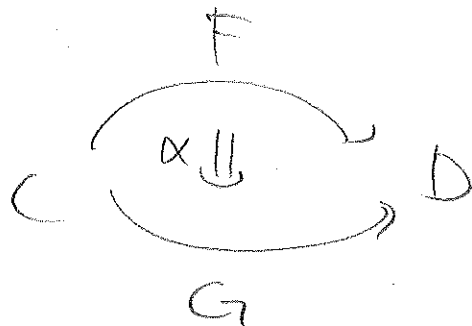
1.7. 2-categories of categories

C, D : categories

D^C : functor category

Obj : $F : C \rightarrow D$ functor.

Mon : α : natural iso.



$\Rightarrow \quad 1_F : F \Rightarrow F \quad \text{by } (1_F)_c := 1_{Fc}$

is identity morphism.

Lemma 1.7.1 (Vertical Composition)

$\alpha : F \Rightarrow G, \quad \beta : G \Rightarrow H$

$\Rightarrow \exists \beta \cdot \alpha : F \Rightarrow H \quad \text{s.t. } (\beta \cdot \alpha)_c := \beta_c \cdot \alpha_c$

pf)

$$\begin{array}{ccccc} F_c & \xrightarrow{\alpha_c} & G_c & \xrightarrow{\beta_c} & H_c \\ \downarrow f & \circlearrowleft & \downarrow g & \circlearrowleft & \downarrow \\ F_{c'} & \xrightarrow{\alpha_{c'}} & G_{c'} & \xrightarrow{\beta_{c'}} & H_{c'} \end{array}$$

from def of natural transt β, α .

$\Rightarrow (\beta \cdot \alpha)$ is natural transt.

Cor 1.7.2. D^c is well-def.

Req 1.7.3.	Size of C	Size of D	Size of D^c
	small	small	small
	(small)	locally small	locally small
	locally small	"	(?)

$$\Rightarrow \text{Cat}^{\text{op}} \times \text{Cat} \rightarrow \text{Cat}.$$

$$\text{Cat}^{\text{op}} \times \text{CAT} \rightarrow \text{CAT}.$$

Vertical Composition: Composition in Lem 1.7.1.

i.e.

$$C \begin{array}{c} \xrightarrow{\alpha \Downarrow} \\ \xrightarrow{G} \\ \xrightarrow{\beta \Downarrow} \end{array} D = C \begin{array}{c} \xrightarrow{\Downarrow \beta \alpha} \\ \xrightarrow{H} \end{array} D$$

Horizontal Composition:

$$C \begin{array}{c} \xrightarrow{F} \\ \xrightarrow{\Downarrow \alpha} \\ \xrightarrow{G} \end{array} D \begin{array}{c} \xrightarrow{H} \\ \xrightarrow{\Downarrow \beta} \\ \xrightarrow{K} \end{array} E = C \begin{array}{c} \xrightarrow{HF} \\ \xrightarrow{\Downarrow \beta * \alpha} \\ \xrightarrow{KG} \end{array} E$$

def by $HF_c \xrightarrow{F_{G_c}} KG_c$

$$\begin{array}{ccc} H\alpha_c \downarrow & \xrightarrow{(\beta * \alpha)_c} & \downarrow K\alpha_c \\ & \searrow & \\ HG_c & \xrightarrow{\beta_{G_c}} & KG_c \end{array}$$

Lemma 1.7.7. (Middle four interchange)

$$\begin{array}{ccc} & F & J \\ \begin{array}{c} \text{---} \curvearrowright \\ \Downarrow \alpha \\ \text{---} \\ G \\ \Downarrow \beta \\ \text{---} \end{array} & D & \begin{array}{c} \text{---} \curvearrowright \\ \Downarrow \gamma \\ \text{---} \\ K \\ \Downarrow \delta \\ \text{---} \end{array} E \\ H & & L \end{array}$$

$$\Rightarrow \begin{array}{ccc} & F & J \\ \begin{array}{c} \text{---} \curvearrowright \\ \Downarrow \beta \cdot \alpha \\ \text{---} \end{array} D & \begin{array}{c} \text{---} \curvearrowright \\ \Downarrow \delta \cdot \gamma \\ \text{---} \end{array} E & = \begin{array}{c} JF \\ \text{---} \curvearrowright \\ \Downarrow \delta \cdot \gamma \cdot \alpha \\ \text{---} \\ KG \\ \Downarrow \delta \cdot \gamma \cdot \beta \\ \text{---} \\ LH \end{array} E \\ H & & L \end{array}$$

Def 1.7.8 A 2-category

Obj: (ex) Categories,
 1-Morphism (ex) Functors,
 2- " (ex) Natural transf.

4-Mon: morphism between pair of objects
 2 " " " " functors

St. 1) Obj + 4Mon : category

2) For any $C, D \in \text{Obj}$

D^C with 2-mon between the ele in D^C
 form a category.

pf) $HF_c \xrightarrow{\beta_{Fc}} KF_c$ Commutes by $\beta: H \Rightarrow K$.

$$\begin{array}{ccc} H\alpha_c \downarrow & & \downarrow K\alpha_c \\ HG_c \xrightarrow{\beta_{Gc}} & & KG_c \end{array} \Rightarrow (\beta * \alpha)_c \text{ is well-def.}$$

To show $\beta * \alpha$ is natural transf.

Let $f: c \rightarrow c' \in \text{Hom } C$. Then

$$\begin{array}{ccccc} HF_c & \xrightarrow{H\alpha_c} & HG_c & \xrightarrow{\beta_{Gc}} & KG_c \\ Hff \downarrow & & \downarrow Hgf & & \downarrow Kgf \\ HF_{c'} & \xrightarrow{H\alpha_{c'}} & HG_{c'} & \xrightarrow{\beta_{G_{c'}}} & KG_{c'} \end{array} \quad \text{Commutes,}$$

Com by
naturality
of α .

Comm by $\beta: H \Rightarrow K$

and H preserves
comm diagram.

$$\Rightarrow \text{Also note that } \frac{H\alpha_c}{\xrightarrow{\quad}} \frac{\beta_{Gc}}{\xrightarrow{\quad}} = (\beta * \alpha)_c$$

$$\frac{\quad}{\xrightarrow{H\alpha_{c'}}} \frac{\quad}{\xrightarrow{\beta_{G_{c'}}}} = (\beta * \alpha)_{c'}$$

$$\Rightarrow KGf (\beta * \alpha)_c = (\beta * \alpha)_{c'} \cdot Hff. \quad \text{done.}$$

3) $Obj = Obj \quad \mathbb{F}$
 $Mor := C \begin{pmatrix} \Downarrow \\ \Downarrow \end{pmatrix} D$

Composition: G horizontal composition.

For a category

4) Law of middle four interchange holds.

Ex 1.7.2).

Ex 1.7.2) Let $F, G: C \Rightarrow D$

Then $F(\text{ob } C)$, $G(\text{ob } C)$,
 $F(\text{Mor } C)$, $G(\text{Mor } C)$ is a set.

Since $|^*| \leq |C|$

Thus, $\bigoplus_{c \in C} D(F_c, G_c)$ is a set

Since C is set and $D(F_c, G_c)$ is a set.

Now take a map

$$D^c(F, G) \longrightarrow \bigoplus_{c \in C} D(F_c, G_c)$$

$$\alpha \longmapsto \{(\alpha_c)_{c \in C}\}$$

It is mono, since $(\alpha_c)_{c \in C} = (\beta_c)_{c \in C}$

$$\Rightarrow \alpha = \beta$$

Thus $D^c(F, G)$ is a set

$\Rightarrow D^c$ is locally small.

1.7.ii Let $f: C \rightarrow C'$

$$\begin{array}{ccc}
 LHF_C & \xrightarrow{L\beta_{F_C}} & LKF_C \\
 LHFf \downarrow & & \downarrow LKff \\
 LHF_{C'} & \xrightarrow{L\beta_{F_{C'}}} & LKF_{C'}
 \end{array}$$

$$\text{WTS } LKFf \circ L\beta_{F_C} = L\beta_{F_{C'}} \circ LHFf$$

$$L(Kff \circ \beta_{F_C}) = L(\beta_{F_{C'}} \circ H(ff))$$

Since $ff: F_C \rightarrow F_{C'}$, from β natural,

$$\begin{array}{ccc}
 HF_C & \xrightarrow{\beta_{F_C}} & KF_C \\
 HFf \downarrow & \curvearrowright & \downarrow Kff \\
 HF_{C'} & \xrightarrow{\beta_{F_{C'}}} & KF_{C'}
 \end{array}
 \Rightarrow
 \begin{aligned}
 & Kff \circ \beta_{F_C} \\
 &= \beta_{F_{C'}} \circ HFf.
 \end{aligned}$$

done.

Thus $L\beta_F$: Natural transf.

1.7.iii)

$$C \xrightarrow{1_C} C \begin{array}{c} \xrightarrow{F} D \\ \Downarrow \alpha \\ \xrightarrow{G} D \end{array} \begin{array}{c} \xrightarrow{H} E \\ \Downarrow \beta \\ \xrightarrow{K} E \end{array} \xrightarrow{1_E} E$$

By whiskering, we have

$$H\alpha : HF \Rightarrow HG$$

$$K\alpha : KF \Rightarrow KG$$

$$\beta F : HF \Rightarrow KF$$

$$\beta G : HG \Rightarrow KG$$

Then,

$$C \begin{array}{c} \xrightarrow{HF} \\ \Downarrow H\alpha \\ \xrightarrow{HG} \\ \Downarrow \beta G \\ \xrightarrow{KG} \end{array} E \quad \text{and} \quad C \begin{array}{c} \xrightarrow{HF} \\ \Downarrow K\alpha \\ \xrightarrow{KF} \\ \Downarrow \beta F \\ \xrightarrow{KG} \end{array} E$$

By vertical composition, we have

$$\beta G \cdot H\alpha : HF \Rightarrow KG, \quad K\alpha \cdot \beta F : HF \Rightarrow KG$$

s.t.

$$\begin{array}{ccc} HF_c & \xrightarrow{H\alpha_c} & HG_c \\ \downarrow \beta F_c & & \downarrow \beta G_c \\ KF_c & \xrightarrow{K\alpha_c} & KG_c \end{array} \quad \text{with } ((\beta G H\alpha)_c)$$

And we already showed in the proof of Lem 1.7.4

that this diagram commutes

$$\Rightarrow \beta G H\alpha = K\alpha \beta F$$

$$\Rightarrow (\beta * \alpha)_c = " = "$$

1.7. iv) Given

$$\begin{array}{ccc}
 & F & J \\
 & \downarrow \alpha & \downarrow \beta \\
 C & \xrightarrow{G} D & \xrightarrow{K} E \\
 & \uparrow \beta & \uparrow \alpha \\
 & H & L
 \end{array}$$

RHS)

From 1.7. iii) we know that

$$\partial * \alpha = \partial G \cdot J\alpha = K\alpha \partial F$$

$$\delta * \beta = \delta H \cdot K\beta = L\beta \delta G$$

$$\begin{aligned}
 \Rightarrow (\delta * \beta) \cdot (\partial * \alpha) &= L\beta \delta G \cdot K\alpha \partial F \\
 &= \quad \quad \cdot \partial G \cdot J\alpha \\
 &= \delta H \cdot K\beta \cdot K\alpha \partial F \\
 &= \quad \quad \cdot \partial G \cdot J\alpha
 \end{aligned}$$

LHS)

$$\begin{aligned}
 (\beta \cdot \alpha) * (\delta \cdot \gamma) &= (\delta \cdot \gamma H) \cdot (J\beta \cdot \alpha) \\
 &= (L\beta \cdot \alpha) \cdot (\delta \cdot \gamma F)
 \end{aligned}$$

Notes that $\delta \cdot \gamma H = (\delta \cdot \gamma)_{Hc} = \delta_{Hc} \cdot \gamma_{Hc}$

$$J\beta \cdot \alpha = J(\beta \cdot \alpha)_c = J(\beta_c \cdot \alpha_c)$$

$$\Rightarrow \text{one expression of LHS} = \delta_{Hc} \gamma_{Hc} J\beta_c J\alpha_c$$

and

$$(f * \beta)(\partial * \alpha) = f_H \cdot k_\beta \partial G \cdot J_\alpha$$

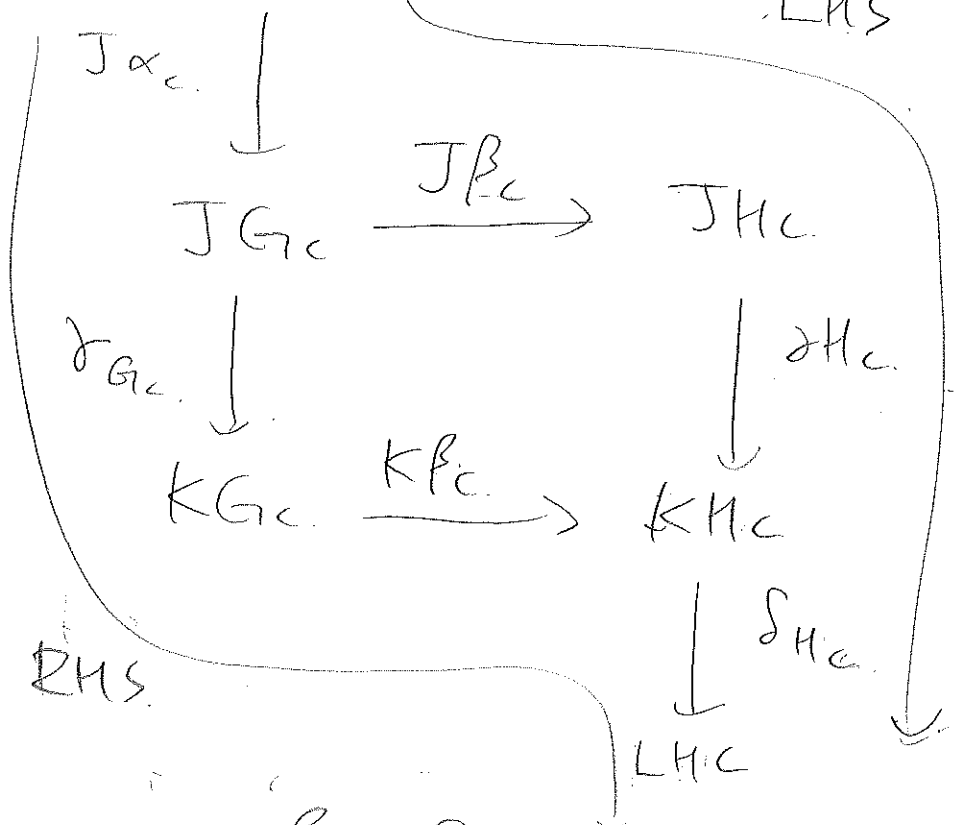
$$= (\delta H)_c \cdot (k\beta)_c \cdot (\partial G)_c \cdot (J_\alpha)_c$$

$$= \delta H_c \cdot K\beta_c \cdot \partial G_c \cdot J\alpha_c$$

LHS

$$\Rightarrow J F_c$$

LHS



It suffices to show that middle rectangle commutes.

Since $\beta_c : G_c \rightarrow H_c$

and $\partial : J \Rightarrow K$

we have

$$JG_c \xrightarrow{J(\beta_c)} JH_c$$

$$\partial G_c \downarrow \quad \quad \quad \downarrow \partial H_c$$

$$KG_c \xrightarrow{K(\beta_c)} KH_c$$

Thus, $RHS = LHS$.

□

Ex 1.7 v.

$$\text{Let } Z(C) := \hat{C}(1_C, 1_C)$$

$$\text{Then } \forall \alpha, \beta \in Z(C)$$

$$\textcircled{1} \quad \alpha \cdot \beta = \beta \cdot \alpha$$

Note that $\forall \alpha \in Z(C)$, take $f: c \rightarrow c'$ then

$$\begin{array}{ccc} c & \xrightarrow{\alpha_c} & c \\ f \downarrow & \curvearrowright & \downarrow f \\ c' & \xrightarrow{\alpha_{c'}} & c' \end{array}$$

Thus, choose f as

$$f = \beta_c : c \rightarrow c$$

Then,

$$\begin{array}{ccc} c & \xrightarrow{\alpha_c} & c \\ \beta_c \downarrow & & \downarrow \beta_c \\ c & \xrightarrow{\alpha_c} & c \end{array}$$

$$\Rightarrow \alpha_c \cdot \beta_c = \beta_c \cdot \alpha_c \quad \forall c$$

$$\Rightarrow \alpha \cdot \beta = \beta \cdot \alpha$$

$$\textcircled{2} \quad 1 \in Z(C)$$

1 : natural iso s.t.
 $(1)_c = 1_c$

$$\text{Then } \alpha \cdot 1 = 1 \cdot \alpha = \alpha$$

$\textcircled{3}$ closed under vertical composition

Remark: actually $\textcircled{1}$ follows from Horizontal composition.

(Π, \bar{V}) First we can do is use:

$$D \begin{array}{c} \xrightarrow{FG} \\ \Downarrow \varepsilon \\ \xrightarrow{1_D} \\ \Downarrow \eta' \\ \xrightarrow{G'F'} \end{array} D \quad \text{to get another natural transf.} \\ \varepsilon * \eta'. \quad \text{Then,}$$

$$D \begin{array}{c} \xrightarrow{FG} \\ \Downarrow \varepsilon * \eta' \\ \xrightarrow{G'F'} \end{array} D \quad \text{exists and natural equivalence} \\ \text{since both } \varepsilon, \eta' \text{ are natural} \\ \text{equivalence.}$$

$$\text{Then, } GG'F'F \stackrel{\textcircled{1}}{\cong} GF GF \stackrel{\textcircled{2}}{\cong} 1_C \cdot 1_C \cong 1_C$$

$$F'FGG' \stackrel{\textcircled{1}}{\cong} F'G'F'G' \stackrel{\textcircled{2}}{\cong} 1_E \cdot 1_E \cong 1_E$$

For example.

$$C \xrightarrow{G} D \begin{array}{c} \xrightarrow{G'F'} \\ \Downarrow FG \end{array} D \xrightarrow{F} C \quad \text{is whiskered}$$

Composite, thus Ex 1.7.ii gives $\textcircled{1}$. And

$$C \begin{array}{c} \xrightarrow{GF} \\ \Downarrow \eta \\ \xrightarrow{1_C} \end{array} C \begin{array}{c} \xrightarrow{GF} \\ \Downarrow \eta \\ \xrightarrow{1_C} \end{array} C = C \begin{array}{c} \xrightarrow{GF GF} \\ \Downarrow \eta * \eta \\ \xrightarrow{1_C \cdot 1_C} \end{array} C \quad \text{is natural equn} \\ \text{by Horizontal} \\ \text{Composition.} \\ \Rightarrow \textcircled{2} \text{ holds.}$$

By $\textcircled{1}$, $\textcircled{2}$, we get desired natural equivalence.

1.7 v (i).

Let $F: C \times D \rightarrow E$ be bifunctor. Then let

$$F': C \rightarrow E^D.$$

$$c \mapsto F(c, -).$$

$$f \downarrow \mapsto F(f, -)$$

$$c' \mapsto F(c', -)$$

where $F(f, -)$ is
natural transformation
between $F(c, -)$ and
 $F(c', -)$.

To define $F(f, -)$, notes that for any

$$d \xrightarrow{g} d',$$

$$F(c, d) \xrightarrow{F(f, d)} F(c', d)$$

$$F(c, g) \downarrow \qquad \qquad \downarrow F(c', g)$$

$$F(c, d') \xrightarrow{F(f, d')} F(c', d')$$

$$\text{with } F(c, g) = F(1_c, g).$$

$$F(c', g) = F(1_{c'}, g)$$

$$F(f, d) = F(f, 1_d)$$

$$F(f, d') = F(f, 1_{d'})$$

Then, since F is bifunctor,

$$F(c', g) \circ F(f, d) = F(1_{c'}, g) \circ F(f, 1_d) = F(f, g)$$

$$F(f, d') \circ F(c, g) = F(f, 1_{d'}) \circ F(1_c, g) = F(f, g)$$

Thus $F(f, -)$ is well-def natural transformation.

Thus $(\)_{\text{obj}} E^{\text{op}} \times D \xrightarrow{\text{obj}} (E^D)^C$ is
 well-def map.

Conversely, for $F' : C \xrightarrow{\quad} E^D$
 $f' : C \xrightarrow{\quad} F'(C)$
 $f' \downarrow \xrightarrow{\quad} F'(f)$ = natural trans
 $C' \xrightarrow{\quad} F'(C')$

$$\begin{aligned} \text{let } UF : C \times D &\longrightarrow E \\ (c, d) &\longmapsto F'(c)(d) \\ (f, g) \downarrow &\longmapsto F'(f)(g) \\ (c', d') &\longmapsto F'(c')(d') \end{aligned}$$

where $F'(f)(g)$ is defined as follow;

$$\begin{array}{ccc} F'(c)(d) & \xrightarrow{F'(f)_d} & F'(c')(d) \\ F'(c)(g) \downarrow & & \downarrow F'(c')(g) \\ F'(c)(d') & \xrightarrow{F'(f)_{d'}} & F'(c')(d') \end{array}$$

Since $F'(f)$ is natural transformation, let.

$$\begin{aligned} F'(f)(g) &:= F'(c')(g) \circ F'(f)_g \quad \text{It is} \\ &= F'(f)_{g'} \circ F'(c)(g) \end{aligned}$$

is well-defined morphism.

Thus $U(\) : \text{Obj}(E^D)^C \longrightarrow \text{Obj}(E^{C \times D})$
 is well-defined map.

(Actually we need to check compositions;
 Identity is clear. For f' , apply it on

$$c \xrightarrow{f} c' \xrightarrow{f'} c'' \quad \text{give}$$

$$F(c, -) \xrightarrow{F(f, -)} F(c', -) \xrightarrow{F(f', -)} F(c'', -)$$

and $F(f'f, -) = F(f', -) \circ F(f, -)$ since

$$\text{for any } d \xrightarrow{g} d'$$

$$F(c, d) \xrightarrow{F(f, 1_d)} F(c', d) \xrightarrow{F(f', 1_d)} F(c'', d)$$

$$\begin{array}{ccc} F(c, g) \downarrow & F(c', g) \downarrow & F(c'', g) \downarrow \\ F(c, d') & \xrightarrow{F(f, 1_{d'})} & F(c', d') \xrightarrow{F(f', 1_{d'})} F(c'', d') \end{array}$$

$$F(c, d') \xrightarrow{F(f, 1_{d'})} F(c', d') \xrightarrow{F(f', 1_{d'})} F(c'', d')$$

and functoriality of F gives $F(f', 1_d) \circ F(f, 1_d) = F(f'f, 1_d)$

Thus $F(f'f, -)_d = F(f', -)_d \circ F(f, -)_d \quad \forall d \in D$

hence $F(f'f, -) = F(f', -) \cdot F(f, -)$ as
 natural iso.

In case of U_F apply it on

$$(c, d) \xrightarrow{(f, g)} (c', d') \xrightarrow{(f', g')} (c'', d'')$$

we get

$$F'(c)(d) \xrightarrow{F'(f)(g)} F'(c')(d') \xrightarrow{F'(f')(g')} F'(c'')(d'')$$

WTS $F'(f')(g') \circ F'(f)(g) = F(f'f)(g'g)$

$$F'(c)(d) \xrightarrow{F(f'f)(g'g)} F'(c'')(d'')$$

$$\begin{array}{ccc} \downarrow & \begin{array}{c} \text{ } \\ \text{ } \end{array} & \downarrow \\ \downarrow & \begin{array}{c} \text{ } \\ \text{ } \end{array} & \downarrow \end{array}$$

$$F'(c)(d) \xrightarrow{F'(f)(g)} F'(c')(d') \xrightarrow{F'(f')(g')} F'(c'')(d'')$$

$$\begin{array}{ccc} \downarrow & \begin{array}{c} \text{ } \\ \text{ } \end{array} & \downarrow \\ \downarrow & \begin{array}{c} \text{ } \\ \text{ } \end{array} & \downarrow \end{array}$$

$$F'(c)(d'') \xrightarrow{F(f'f)(g'g)} F'(c'')(d'')$$

✓

Notes that each rectangular commutes by def of $F'(f)(g)$, $F'(f')(g)$, $F'(f')(g')$, $F'(f)(g)$, $F(1_{c''}, g'g)$

Hence outer rectangular commutes.

So it suffices to show that each outer edge is $F(f'f)(1_d)$ and $F(1_{c''})(g', g)$

But

$$\begin{array}{ccc}
 F'(c)(1_d) & & \\
 F'(f)(1_d) \downarrow & \searrow & \\
 F'(c')(1_d) & & F(f', 1_d) \\
 F'(f')(1_d) \downarrow & & = F'(f'f)_d \circ F'(c)(1_d) \\
 F'(c'')(1_d) & \swarrow &
 \end{array}$$

Given

$$\begin{aligned}
 & F'(f')(1_d) \cdot F'(f)(1_d) \\
 &= F'(f')_d \circ F'(c')(1_d) \cdot F'(c')(1_d) \circ F'(f)_d \\
 &= \underline{F'(f')_d \cdot F'(c')(1_d)} \cdot F'(f)_d \\
 &= \underline{F'(f')_d \cdot 1_{F'(c')(1_d)}} \cdot F'(f)_d \quad \text{by Functionality of } F'(c') \\
 &= F'(f')_d \cdot F'(f)_d \quad \downarrow \text{by functionality of } F'(f) \\
 &= F'(f'f)_d
 \end{aligned}$$

done. (Vertical one is similar)

Hence given map is well-def, thus bijection.