

Thus we need a little bit relaxed concept
morphism of functor, i.e. Natural transformation.

Exercise 1.3.i) Functor between groups?

$$\begin{array}{ccc}
 G & \xrightarrow{F} & H \\
 \cdot & \xrightarrow{\quad} & \cdot \\
 g \downarrow & \xrightarrow{\quad} & h \downarrow \\
 \cdot & \xrightarrow{\quad} & \cdot
 \end{array}
 \quad
 \begin{array}{l}
 F(g) = h \text{ s.t.} \\
 a) F(gg') = F(g)F(g') \\
 b) F(e_g) = e_H \\
 \Rightarrow \text{If } F \text{ is Functor, then} \\
 \text{it is gp homo.}
 \end{array}$$

Conversely, gp homo: satisfy a) \Rightarrow b) \therefore by adding $\cdot \rightarrow \cdot$
it is functor.

1.3.ii) Functor between preorders?

$$(P, \leq) \quad (Q, \leq)$$

$$\begin{array}{ccc}
 x & \xrightarrow{\quad} & F(x) \\
 x \leq y \downarrow & \xrightarrow{\quad} & \downarrow \\
 y & \xrightarrow{\quad} & F(y)
 \end{array}
 \quad
 \begin{array}{l}
 F(x) \leq F(y) \quad \dots a) \\
 b) : \text{done}
 \end{array}$$

F is order preserving function.

sketch

1.3.iii) Pr. Hasenmeyer Notation.

$$C: \begin{array}{ccc} \circ & \xrightarrow{\quad} & \circ \\ a & & b \end{array} \quad \begin{array}{ccc} \circ & \xrightarrow{\quad} & \circ \\ c & & d \end{array} \quad F: C \longrightarrow D$$

$$D: \begin{array}{ccc} & \circ & \\ \nearrow & & \searrow \\ \circ & \xrightarrow{\quad} & \circ \\ x & & z \end{array}$$

$$F(C):$$

$$\begin{array}{ccc} & \circ & \\ \nearrow & & \searrow \\ \circ & \xrightarrow{\quad} & \circ \\ x & & z \end{array}$$

We have composable morphisms in $F(C)$
But composite doesn't exist.

Ex 1.3. iv) $f = 1_x \Rightarrow$

$$\begin{array}{ccccc}
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{\quad} & g \\
 1_x \downarrow & \xrightarrow{\quad} & \downarrow (1_x)_* & \xrightarrow{\quad} & \downarrow \text{Identity} \\
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{\quad} & (1, g) = g
 \end{array}$$

$f: x \rightarrow y, \quad g: y \rightarrow z$

$$\Rightarrow \begin{array}{ccccc}
 x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{h} & x & \xrightarrow{\quad} & C(C, x) & \xrightarrow{h} \\
 f \downarrow & & \downarrow f_* & & g f \downarrow & & \downarrow & & \downarrow \\
 y & \xrightarrow{\quad} & C(C, y) & \xrightarrow{f h} & y & \xrightarrow{\quad} & C(C, y) & \xrightarrow{f h} \\
 g \downarrow & & \downarrow g_* & & & & \downarrow & & \downarrow \\
 z & \xrightarrow{\quad} & C(C, z) & \xrightarrow{g f h} & z & \xrightarrow{\quad} & C(C, z) & \xrightarrow{g f h}
 \end{array}$$

Ex 1.3. v) Claim: $F: C^{op} \rightarrow D \iff G: C \rightarrow D^{op}$

Let $F: C^{op} \rightarrow D$. Construct $G: C \rightarrow D^{op}$ s.t.

$G(c) := F(c)$. From $\text{Hom}_D(Fb, Fa) \leftrightarrow \text{Hom}_{D^{op}}(Fa, Fb)$

So given $f \in \text{Hom}_C(a, b)$

define $G(f) = (f^*)^{op}$ in $\text{Hom}_{D^{op}}(Fa, Fb)$

Thus, $F: C^{op} \rightarrow D \xrightarrow{F} D$ $G: C \xrightarrow{G} D^{op} \xrightarrow{G} D$

$$\begin{array}{ccccc}
 x & \xrightarrow{\quad} & x & \xrightarrow{\quad} & Fx \\
 f \downarrow & & f^{op} \uparrow & & \uparrow Ff \\
 y & \xrightarrow{\quad} & y & \xrightarrow{\quad} & Fy
 \end{array}
 \qquad
 \begin{array}{ccccc}
 x & \xrightarrow{\quad} & Fx & \xrightarrow{\quad} & Fx \\
 f \downarrow & & \downarrow (f^*)^{op} & & \uparrow Ff \\
 y & \xrightarrow{\quad} & Fy & \xrightarrow{\quad} & Fy
 \end{array}$$

F, G are the same functor. Also, $F: C \rightarrow D \iff G: C^{op} \rightarrow D^{op}$

$$\begin{array}{ccccc}
 C & \xrightarrow{\quad} & D \\
 x \downarrow & & \downarrow Ff \\
 y & \xrightarrow{\quad} & y
 \end{array}
 \qquad
 \begin{array}{ccccc}
 C & \xrightarrow{\quad} & C^{op} & \xrightarrow{\quad} & D^{op} & \xrightarrow{\quad} & D \\
 x \downarrow & & \uparrow f^{op} & & \uparrow (f^*)^{op} & & \downarrow Ff \\
 y & \xrightarrow{\quad} & y & \xrightarrow{\quad} & Fy & \xrightarrow{\quad} & Fy
 \end{array}$$

$$1.3.vi) F: D \rightarrow C, G: E \rightarrow C.$$

$$F \downarrow G : \text{Obj} := (d, e, f)$$

$$d \in D, e \in E, f: Fd \rightarrow Ge \in C$$

$$\text{Mor} : ((d, e, f), (d', e', f'))$$

$$= \{ (h, k) \in \text{Mor } D \times \text{Mor } E : \}$$

$$h: d \rightarrow d', k: e \rightarrow e' \text{ s.t.}$$

$$Fd \xrightarrow{f} Ge$$

$$Fh \downarrow \quad \quad \downarrow Gk \quad \text{in } C, \text{ i.e.}$$

$$Fd' \xrightarrow{f'} Ge'$$

$$f' \cdot Fh = Gk \cdot f$$

It is category since,

$$\textcircled{1} 1_{(d,e,f)} = (1_d, 1_e)$$

$$\textcircled{2} (d, e, f) \xrightarrow{(h_1, k_1)} (d', e', f') \xrightarrow{(h_2, k_2)} (d'', e'', f'')$$

$$\Rightarrow \begin{array}{ccc} Fd & \xrightarrow{f} & Ge \\ Fh_2 \downarrow & & \downarrow Fk_1 \\ Fd' & \xrightarrow{f'} & Ge' \end{array} \Rightarrow (Fk_2)(Fk_1) \cdot f = Fk_2(f') \cdot Fh_1 = f'' Fh_2 Fh_1$$

$$\begin{array}{ccc} Fd' & \xrightarrow{f'} & Ge' \\ Fh_2 \downarrow & & \downarrow Fk_2 \\ Fd'' & \xrightarrow{f''} & Ge'' \end{array} \quad \begin{array}{ccc} Fd & \xrightarrow{f} & Ge \\ F(h_2 h_1) \downarrow & & \downarrow Fk_2 k_1 \\ Fd'' & \xrightarrow{f''} & Ge'' \end{array}$$

$$Fd'' \xrightarrow{f''} Ge'' \quad \Rightarrow f'' F(h_2 h_1) = F(k_2 k_1) f$$

Functor: $\text{dom}: F \downarrow G \longrightarrow D$, $\text{cod}: F \downarrow G \longrightarrow E$

$$(d, e, f) \longmapsto d \quad (d, e, f) \longmapsto e$$

$$(h, k) \downarrow \longmapsto h \quad (h, k) \downarrow \longmapsto k$$

$$(d', e', f') \longmapsto d' \quad (d', e', f') \longmapsto e'$$

Ex 1.3. vii) $D = (\{-3, 1\})$ $F: D \longrightarrow C$ $G = 1_C$
 $E = C$ $\bullet \longmapsto c$

C/C
under

Obj = $(\bullet, x, f: c \longrightarrow x)$

Mor: $((\bullet, x, f: c \longrightarrow x), (\bullet, y, g: c \longrightarrow y))$
 $= \{ (1_\bullet, h: x \longrightarrow y) \in \text{Mor } D \times \text{Mor } E$

\leq

$$\begin{array}{ccc} c & \xrightarrow{f} & x \\ 1_c \downarrow & \circlearrowright & \downarrow Gh=h \\ c & \xrightarrow{g} & y \end{array} \quad \text{i.e.} \quad \begin{array}{ccc} & c & \\ f \swarrow & & \searrow g \\ x & \xrightarrow{h} & y \end{array}$$

C/C
over

$D = C$ $F = 1_C$ $G = \bullet \longmapsto c$
 $E = (\{-3, 1\})$

Obj: $(x, \bullet, f: x \longrightarrow c)$

Mor $((x, \bullet, f: x \longrightarrow c), (y, \bullet, g: y \longrightarrow c))$
 $= \{ (h: x \longrightarrow y, 1_\bullet) \in \text{Mor } D \times \text{Mor } E$

\leq

$$\begin{array}{ccc} x & \xrightarrow{f} & c \\ h \downarrow & \circlearrowright & \downarrow 1_c \\ y & \xrightarrow{g} & c \end{array} \quad \text{i.e.} \quad \begin{array}{ccc} & c & \\ f \nearrow & & \nwarrow g \\ x & \xrightarrow{h} & y \end{array}$$

projection functor:

$$\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{D} = (\{ \cdot \}, 1)$$

$$\begin{aligned} (\cdot, x, f: \mathcal{C} \rightarrow x) &\longmapsto \cdot \\ (1, h) \downarrow &\longmapsto \downarrow 1. \\ (\cdot, y, g: \mathcal{C} \rightarrow y) &\longmapsto \cdot \end{aligned}$$

$$\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{E} = \mathcal{C}$$

$$\begin{aligned} (\cdot, x, f: \mathcal{C} \rightarrow x) &\longmapsto x \\ (1, h) \downarrow &\longmapsto \downarrow h. \\ (\cdot, y, g: \mathcal{C} \rightarrow y) &\longmapsto y \end{aligned}$$

Thus: $\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{C}$

$$\begin{aligned} f: \mathcal{C} \rightarrow x &\longmapsto x \\ h \downarrow &\longmapsto \downarrow h. \\ g: \mathcal{C} \rightarrow y &\longmapsto y \end{aligned}$$

$$\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{C}$$

$$\begin{aligned} f: x \rightarrow \mathcal{C} &\longmapsto x \\ h \downarrow &\longmapsto \downarrow h. \\ g: y \rightarrow \mathcal{C} &\longmapsto y \end{aligned}$$

Similarly, $\text{dom} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{D} = \mathcal{C}$

$$\begin{aligned} (x, \cdot, f: x \rightarrow \mathcal{C}) &\longmapsto x \\ (h, 1) \downarrow &\longmapsto \downarrow h. \\ (y, \cdot, g: y \rightarrow \mathcal{C}) &\longmapsto y \end{aligned}$$

$$\text{cod} : \mathcal{C}/\mathcal{C} \longrightarrow \mathcal{E} = (\{ \cdot \}, 1)$$

$$\begin{aligned} (x, \cdot, f: x \rightarrow \mathcal{C}) &\longmapsto \cdot \\ (h, 1) \downarrow &\longmapsto \downarrow 1. \\ (y, \cdot, g: y \rightarrow \mathcal{C}) &\longmapsto \cdot \end{aligned}$$

Ex 1.3, viii) Ex: Functor need not

Find $f: \mathcal{C} \longrightarrow \mathcal{D}$, $f \in \text{Mor}(\mathcal{C})$

reflect isomorphism

s.t. f is iso in \mathcal{D}

f is not iso in \mathcal{C} .

$$\begin{array}{c} \cdot \xrightarrow{f} \cdot \\ \mathcal{C} \end{array}$$

$$\begin{array}{c} \cdot \xrightleftharpoons[f^{-1}]{f} \cdot \\ \mathcal{D} \end{array}$$

More generically, $\mathcal{I} = (\{ \cdot \}, 1)$

$$\mathcal{C} \longrightarrow \mathcal{I} \quad \text{by} \quad \begin{aligned} c &\longmapsto \cdot \\ f &\longmapsto 1. \end{aligned}$$

ex) Homology

Homotopy

Functor:

(Quasi-iso but not homeo.)

$\forall f, \forall c.$

Ex 1.3. ix)

Source map	Group iso.	Group epi	Group
$Z(-)$	Yes (iso)	Yes	No ¹⁾
$C(-)$	Yes	Yes	Yes ²⁾
$\text{Aut}(-)$	Yes	Don't know.	No ³⁾

1). If $G \xrightarrow{f} H \Rightarrow Z(G) \cong Z(H)$

pf)

Claim 1: $G \xrightarrow{f} H$ surj, then $f(Z(G)) \leq Z(H)$

pf) $h \cdot f(g) = f(h'g) = f(gh') = f(g)h \quad \forall h \in H,$

$g \in Z(G), \quad h': \text{preimage of } h.$

Claim 2: $G \xrightarrow{\text{surj}} H$ then $f(Z(G)) = Z(H)$

pf)

If $h \in Z(H), \quad \exists g \in G \text{ st. } f(g) = h$

$$g'g = f^{-1}(h'h) = f^{-1}(hh') = f^{-1}(h)f^{-1}(h') = gg'$$

$\Rightarrow g \in Z(G) \Rightarrow f(Z(G)) \supseteq Z(H). \quad \text{done.}$

Claim 3: $Z(-): \begin{matrix} \text{Group iso} \\ \text{Group epi} \end{matrix} \longrightarrow \text{Group is functorial.}$

pf) $G \xrightarrow{Z(-)} Z(G) \quad \text{Let } G \xrightarrow{\phi} H \xrightarrow{\psi} K$
 $\downarrow 1_G \quad \downarrow Z(1_G) \quad \text{Iso (or epi)}$
 $G \longrightarrow Z(G)$

Then, from $\phi(Z(G)) \subseteq Z(H)$,
 $\psi(Z(H)) \subseteq Z(K)$

We have map

$$Z(G) \xrightarrow{\phi|_{Z(G)}} Z(H) \xrightarrow{\psi|_{Z(H)}} Z(K)$$

Thus, $\psi|_{Z(H)} \circ \phi|_{Z(G)} = (\psi \circ \phi)|_{Z(G)}$
implies ϕ first functoriality axiom.

Claim 4: $Z(-)$ Group \rightarrow Group is not a functor.

Pf) $1_{C_2} = C_2 \rightarrow S_3 \rightarrow S_3/A_3 \cong C_2$

$$1 \mapsto (12) \mapsto (12) + A_3 \mapsto 1$$

where C_2 : cyclic group of order 2,
 S_3 : Sym. " 3.

A_3 : alternating op of S_3 .

Then, If Z is a functor,

$$Z(C_2) \rightarrow Z(S_3) \rightarrow Z(C_2) = Z(C_2) \xrightarrow{1_{Z(C_2)}} Z(C_2)$$

$$= C_2 \rightarrow 0 \rightarrow C_2, \text{ contradiction } C_2 \xrightarrow{f_{C_2}} C_2$$

2). Claim 1: $G \xrightarrow{f} H \Rightarrow C(f(G)) = f(C(G)) \leq C(H)$

pf) $C(f(G)) = \{ f(x)f(y)f(x)^{-1}f(y)^{-1} : x, y \in G \}$
 $= \{ f(xy x^{-1}y^{-1}) : x, y \in G \}$
 $= f(C(G)) \leq \{ xyx^{-1}y^{-1} : x, y \in H \} = C(H)$

Claim 2: $C(-) : \begin{matrix} \text{Group}_{\text{iso}} & \longrightarrow & \text{Group} \\ \text{Group}_{\text{epi}} & \longrightarrow & \\ \text{Group} & \longrightarrow & \text{Group} \end{matrix}$ ~~are~~ ^{are} functors

pf). $G \xrightarrow{f} H \Rightarrow C(H) \cong f(C(G))$

$\Rightarrow C(G) \xrightarrow{f|_{C(G)}} C(H)$ is gp homo.

If $f = I_G \Rightarrow C(f) = I_{C(G)}$

$G \xrightarrow{\phi} H \xrightarrow{\psi} K$

$\Rightarrow \begin{matrix} \phi|_{C(G)} & \psi|_{C(H)} & (\psi \circ \phi)|_{C(G)} \\ C(G) & \longrightarrow C(H) & \longrightarrow C(K) \end{matrix} =$

3) Claim 1. $\begin{matrix} \text{Aut} \\ \text{Group}_{\text{epi}} & \longrightarrow & \text{Group} \\ \text{Group} & \longrightarrow & \text{Group} \end{matrix}$ is not a functor

pf) $G = \mathbb{F}_n \rtimes \mathbb{F}_n^{\times}$ by multiplication

$$G = \{ (a, b) : a \in F, b \in F^\times \}$$

$$\text{and } (a, b) \cdot (c, d) = (a + bc, bd)$$

$$\text{Also, let } N = \{ (a, 1) : a \in F \} \cong F$$

$$K = \{ (0, b) : b \in F^\times \} \cong F^\times$$

N, K are subgp of G .

and by construction, N is normal;

$$\begin{aligned} (b, c) (a, 1) (b, c)^{-1} &= (b + ca, c) (b, c^{-1}) \\ &= (b + ca - bc, 1) \in N \end{aligned}$$

$$\Rightarrow NK = G \quad (\text{by construction})$$

By Schur-Zassenhaus lemma,

Any subgp of G with order $|K|$

conjugate to each other, i.e., if k_1, k_2 has order $|K|$, then $\exists g \in G$ s.t.

$$g k_1 g^{-1} = k_2$$

$$\begin{aligned} (a, 1) (c, d) &= (a, 1) \\ &= (a + bc, bd) \end{aligned}$$

Note, that $(0, 1)$ is identity in G .

$$\text{and } (a, b)^{-1} = \left(-\frac{a}{b}, b^{-1} \right)$$

Hence let $\alpha \in \text{Aut}(G)$

$\Rightarrow \alpha(K)$ is a subgp of G with order $|K|$

$\Rightarrow \exists g \in G$ s.t.

$$g\alpha(K)g^{-1} = K$$

\Rightarrow Let $\varphi_g \in \text{Inn}(G)$ s.t. $\varphi_g(g') = gg's^{-1}$

$\Rightarrow \varphi_g \circ \alpha$ fix K .

Also, notes that $\varphi_g \circ \alpha$ fix N since N is unique 11-Sylow subgp of G .

Thus, if $\varphi_g \circ \alpha(1,1) = (a,1)$

then, $a \neq 0$ since $\varphi_g \circ \alpha^{-1}(0,1) = (0,1)$
(identity)

thus $\exists b \in \mathbb{F}_n^\times$ s.t. $ba = 1$

$$\therefore (0,b)(a,1)(0,b^{-1}) = (ab,b)(0,b^{-1}) = (ab,1) = (1,1)$$

$$\therefore \varphi_b \varphi_g \circ \alpha(1,1) = (1,1)$$

And since $(0,b)(0,1)(0,b^{-1}) = (0,1)$,

φ_b still fix K .

Let $\beta = \varphi_b \circ \varphi_g \circ \alpha$.

Then, $\beta(k) = k$, $\beta(1,1) = (1,1)$.

and β is auto. ($\beta(N) = N$)

If $a \neq 0$ in F_{11}

$$\Rightarrow \beta((a,1)) = \beta((0,a)(1,1)(0,a^{-1})) \\ = (\beta(0,a))(1,1)(\beta(0,a))^{-1}$$

Since β fix N , $\beta(a,1) = (b,1)$

And β fix $K \Rightarrow \beta(0,a) = (0,c)$

for some
 $b \in F_{11}^*$
 $c \in F_{11}^*$

$$\text{Thus, } (b,1) = (0,c)(1,1)(0,c^{-1}) \\ = (c,c)(0,c^{-1}) = (c,1)$$

$\Rightarrow b=c$. And

$$\beta(0,a) = (0,c) = (0,b)$$

Thus, $\gamma: F_{11} \rightarrow N \xrightarrow{\beta} N \rightarrow F_{11}$

$$a \mapsto (a,1) \mapsto (b,1) \mapsto b$$

$$0 \mapsto (0,1) \mapsto (0,1) \mapsto 0$$

induces a map on F_{11} (we didn't show it's homo)

similarly, $F_{11}^* \rightarrow K \xrightarrow{\beta} K \rightarrow F_{11}^*$

$$a \mapsto (0,a) \mapsto (0,b) \mapsto b$$

we claim that γ is field homo.

$$\sigma(a)\sigma(b) = (\sigma(a), 1)(\sigma(b), 1) \quad \text{in } G$$

$$(\sigma(a), 1) = (0, \sigma(a))(\sigma(b), 1)(0, \sigma(a)^{-1})$$

$$= \beta(0, a)\beta(b, 1)\beta(0, a^{-1})$$

$$\text{Since } \sigma(a) =: a \mapsto (0, a) \mapsto \beta(0, a)$$

$$= (0, b)$$

for some c

$$\mapsto c$$

$$\text{thus } \beta(0, a) =: (\sigma, c) = (0, \sigma(a))$$

$$\text{and } b \mapsto (b, 1) \mapsto \beta(b, 1) = (c, 1)$$

for some c

$$\Rightarrow \sigma(b) = c$$

$$\Rightarrow \beta(b, 1) = (c, 1) = (\sigma(b), 1)$$

$$= \beta(ab, 1) = (c, 1) \quad \text{for some } c \in \mathbb{F}_n$$

$$\text{and } \Rightarrow \sigma(a)\sigma(b) = c$$

$$ab \mapsto (ab, 1) \mapsto \beta(ab, 1) = (c, 1)$$

$$\Rightarrow \sigma(ab) = c$$

Hence $\sigma(ab) = \sigma(a)\sigma(b)$. Thus β induces a field homomorphism on \mathbb{F}_n .

Thus $\text{Aut}(G) \cong \text{Inn}(G) \cdot \text{Aut}(F_{11})$.

$$\cong \text{Inn}(G)$$

$$\cong G$$

Since any automorphism of G is field homomorphism times inner automorphism.

and $\text{Aut}(F_{11}) = 1$ since there are only F_{11}^\times gp homomorphisms and only 1 homomorphism fix 1 to 1.

Now Think about a map.

$$\phi: F_{11}^\times \longrightarrow G \longrightarrow G/N \cong F_{11}^\times: \text{isomorphism}$$

$$\text{Aut}(-) \text{ induce, } \text{Aut}(F_{11}^\times) \rightarrow \text{Aut}(G) \rightarrow \text{Aut}(F_{11}^\times)$$

$$\text{But } |\text{Aut}(F_{11}^\times)| = (\mathbb{Z}/47\mathbb{Z}) = 4 \quad \text{and } (4|110) = 1.$$

$$|\text{Aut}(G)| = |G| = 110.$$

$$\Rightarrow \text{Aut}(F_{11}^\times) \rightarrow \text{Aut}(G) \text{ is zero.}$$

$$\text{But } \phi = 1_{F_{11}^\times} \text{, thus}$$

$$\text{Aut}(\phi) = 1_{\text{Aut}(F_{11}^\times)} \neq 0. \text{ So Aut doesn't satisfy functoriality axiom.}$$

1.3.X.) Let G, H gr, $f: G \rightarrow H$ homo,
 X_G, X_H set of conjugacy classes of G, H .

$g: G \rightarrow X_G, h: H \rightarrow X_H$ class functions

$$i.e. \quad g(a) = g(bab^{-1}) \quad \forall b \in G.$$

Let $\text{Conj}: \text{Group} \rightarrow \text{Set}$

$$\begin{array}{ccc} G & \xrightarrow{\quad} & X_G \\ \phi \downarrow & \xrightarrow{\quad} & \downarrow \text{Conj}(\phi) \\ H & \xrightarrow{\quad} & X_H \end{array} \quad \begin{array}{c} \overline{g} \\ \downarrow \\ \overline{h} \end{array}$$

① $\text{Conj}(\phi)$ is well-defined;

if $g' \in \overline{g}$, then, $g' = bab^{-1}$ for some $b \in G$.

Thus, $\phi(g') = \phi(b) \phi(g) \phi(b)^{-1} \Rightarrow \phi(g') \in \overline{\phi(g)}$.

② It is functorial; If $\psi: H \rightarrow J$ morphism,

$$\text{Conj}(\psi) \cdot \text{Conj}(\phi) = \text{Conj}(\psi \circ \phi)$$

$$\text{and } \text{Conj}\left(\begin{smallmatrix} G \\ 1_G \end{smallmatrix}\right) = 1_{X_G}.$$

Thus, if $|X_G| \neq |X_H|$, then $\text{Conj}(\phi)$
 is not isomorphism, since functor preserves
 iso, $\phi: G \rightarrow H$ is not iso.

Exercise 1.4. i.

Since α is natural iso, $\forall C \in \mathcal{C}$,

$$\exists \alpha_C^{-1} \text{ s.t. } \alpha_C \circ \alpha_C^{-1} = \text{id}_{G_C}$$

$$\alpha_C^{-1} \circ \alpha_C = \text{id}_{F_C}$$

Thus, let α^{-1} : collection of α_C^{-1} . Then,

$$\forall f: C \rightarrow C'$$

WTS

$$G_C \xrightarrow{\alpha_C^{-1}} F_C$$

$$Ff \circ \alpha_C^{-1} = \alpha_{C'}^{-1} \circ Gf$$

$$Gf \downarrow \quad \downarrow Ff$$

Since $\alpha_{C'} \circ f = Gf \circ \alpha_C$,

$$G_{C'} \xrightarrow{\alpha_{C'}^{-1}} F_{C'}$$

$$\Rightarrow \alpha_{C'} \circ f \circ \alpha_C^{-1} = Gf \circ \text{id}_{G_C}$$

$$\Rightarrow \text{id}_{F_{C'}} \circ Ff \circ \alpha_C^{-1} = \alpha_{C'}^{-1} \circ Gf$$

1.4. ii. If $\phi: G \rightarrow H$ be a functor,

$$\text{then, } \phi(e_G) = e_H \quad \phi(g) \cdot \phi(g') = \phi(gg')$$

Thus, ϕ is an α -homomorphism

If $\phi: G \xrightarrow{\alpha} H$, then, if $\phi \xrightarrow{\alpha} \psi$ exists,

$$\textcircled{1} \alpha \in H$$

$$\forall g \in G, \quad \begin{array}{ccc} & \xrightarrow{\alpha} & \\ \phi(g) \downarrow & & \downarrow \psi(g) \\ & \xrightarrow{\quad} & \end{array} \quad \textcircled{2} \psi(g) \alpha = \alpha \phi(g)$$

$$\Rightarrow \alpha^{-1} \psi(g) \alpha = \phi(g)$$

Hence, $\alpha \in \text{Inn}(H)$ s.t. $\alpha \circ \psi = \phi$.

Ex 1.4. iii) $(P, \leq) \xrightleftharpoons[F]{F} (Q, \leq)$,

then F, G are order preserving function.

If $F \xrightarrow{\alpha} G$, then $\forall f: P \rightarrow P'$
($P \leq P'$)

$$F_P \xrightarrow{\alpha_P} G_P \quad \text{i.e. } G_{P'} \geq F_{P'} \geq F_P$$

$$\downarrow \quad \downarrow \quad \text{H} \quad G_{P'} \geq G_P \geq F_P$$

$$F_{P'} \xrightarrow{\alpha_{P'}} G_{P'} \Rightarrow G_P \geq F_P \quad \forall P.$$

Thus, natural transformation of F, G
order preserving function is $F \leq G$.

i.e. $\forall p \in P, F_p \leq G_p$.

(Notes that $F \leq G \Leftrightarrow$ Natural transf. hold.)

1.4. iv)

Ex 1.4 iv). Each f_*, g_*, f^*, g^* defines
n.t. by Ex 1.4.7. $\forall h: Y \rightarrow X$.

$$\text{If } f_* = g_* \Rightarrow \begin{array}{ccc} C(x, c) & \xrightarrow{f_*} & C(x, d) \\ \downarrow \cdot h & \wr & \downarrow \cdot h \\ C(y, c) & \xrightarrow{f_*} & C(y, d) \end{array}$$

$$\Rightarrow \forall l \in C(x, c). \quad l \mapsto f l h$$

$$f l h = g l h. \quad \text{Now pick } h = 1_x$$

$$\Rightarrow f l = g l, \quad \text{Pick } x = c, \quad l = 1_c.$$

$\Rightarrow f = g$, contradiction. $\therefore f_* \neq g_*$ as
natural transformation. (f^*, g^* : similar manner)

Ex 1.4 v.

$$\begin{array}{ccccc} F \downarrow G & \xrightarrow{\text{dom}} & D & \xrightarrow{F} & C \\ (d, e, f) & \longmapsto & d & \longrightarrow & Fd \\ (h, k) \downarrow & \longmapsto & \downarrow h & & \downarrow Fh \\ (d', e', f') & \longmapsto & d' & \longrightarrow & Fd' \end{array}$$

$$\begin{array}{ccccc}
 F \downarrow G & \xrightarrow{\text{cod}} & E & \xrightarrow{G} & C \\
 (d, e, f) & \longmapsto & e & \longmapsto & Ge \\
 \downarrow (h, k) & & \downarrow k & & \downarrow Gk \\
 (d', e', f') & \longmapsto & e' & \longmapsto & Ge'
 \end{array}$$

Thus, For each $(d, e, f) \xrightarrow{(h, k)} (d', e', f')$
 we need to find $\alpha_{(d, e, f)}$, $\alpha_{(d', e', f')}$
 ≤ 1 .

$$\begin{array}{ccc}
 F_d & \xrightarrow{\alpha_{(d, e, f)}} & Ge \\
 \downarrow Fh & \curvearrowright & \downarrow Gk \\
 F_{d'} & \xrightarrow{\alpha_{(d', e', f')}} & Ge'
 \end{array}$$

Commutative,

Set $\alpha_{(d, e, f)} = f$. By construction of $F \downarrow G$, it holds.

Ex 1.4 vi) If F, G has different target category,
 i.e., $F: A \times B \times B^{\text{op}} \rightarrow D$
 $G: A \times C \times C^{\text{op}} \rightarrow D'$ with $D \neq D'$.

then, we ~~don't know how to~~ cannot have a
 morphism $F(a, b, c) \rightarrow G(a, c, c)$ in general

$$\begin{array}{ccccc}
 F \downarrow G & \xrightarrow{\text{cod}} & E & \xrightarrow{G} & C \\
 (d, e, f) & \longmapsto & e & \longmapsto & Ge \\
 \downarrow (h, k) & & \downarrow k & & \downarrow Gk \\
 (d', e', f') & \longmapsto & e' & \longmapsto & Ge'
 \end{array}$$

Thus, For each $(d, e, f) \xrightarrow{(h, k)} (d', e', f')$
 we need to find $\alpha_{(d, e, f)}$, $\alpha_{(d', e', f')}$
 s.t.

$$\begin{array}{ccc}
 Fd & \xrightarrow{\alpha_{(d, e, f)}} & Ge \\
 \downarrow Fh & \searrow & \downarrow Gk \\
 Fd' & \xrightarrow{\alpha_{(d', e', f')}} & Ge'
 \end{array}
 \quad \text{Commutative}$$

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