

# Natural Transformation

Ex)  $(-)^* \circ (-)^*$  where  $(-)^*: \text{Vect}_k^{\text{op}} \rightarrow \text{Vect}_k$   
 $V \mapsto V^*$

!Induces  $V' \cong V^{**}$

but this is cones from  $\text{ev}_V$   $\forall V \in \mathcal{V}$ .

No unnatural choice of basis is needed.

Def 1.4.1. (Natural Transformation)

$\mathcal{C}, \mathcal{D}$ : categories  $F, G: \mathcal{C} \Rightarrow \mathcal{D}$

$\alpha$ : Natural Transformation  $F \Rightarrow G = \{\alpha_c: c \rightarrow c\}$

st.  $\alpha_c: F_c \rightarrow G_c \in \text{Mor } \mathcal{D}$   $\forall c \in \mathcal{C}$

$\alpha$  component of  $\alpha$ .

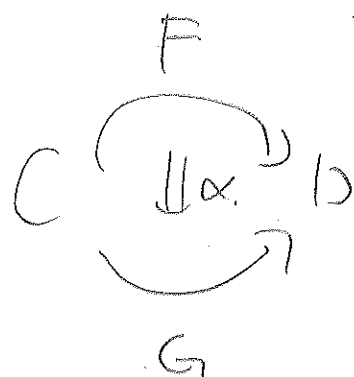
$$\begin{array}{ccc} F_c & \longrightarrow & G_c \\ Ff \downarrow & \circlearrowright & \downarrow Gf \\ F_{c'} & \longrightarrow & G_{c'} \end{array} \quad \begin{array}{l} \text{st. } \forall f: c \rightarrow c' \\ \in \text{Mor } \mathcal{C} \\ \text{commutes.} \end{array}$$

Natural isomorphism:  $\alpha$ : natural transformation

st.  $\alpha_c$  is iso  $\forall c \in \mathcal{C}$ .

Then  $\alpha: F \cong G$ .

"The arrows are natural"  $\Rightarrow$  collection of arrows define natural transf.



Ex 1.4.3.

$$(i) \quad \text{Id}_{\text{Vect}_k} \xrightarrow{\text{ev}} (-)^*, (-)^*$$

since  $\forall \phi: V \rightarrow W$  morphism in  $\text{Vect}_k$

$$\begin{array}{ccc}
 V & \xrightarrow{\text{ev}} & V^{**} \\
 \downarrow \phi & & \downarrow \phi^{**} \\
 W & \xrightarrow{\text{ev}} & W^{**}
 \end{array}$$

$f_v: g \in V^* \mapsto g(v)$   
 $\phi^{**}(f_v) = V \xrightarrow{\phi^*} W^* \xrightarrow{f_v} k$   
 $f_{\phi(v)}: g \in W^* \mapsto g(\phi(v))$

$$\begin{aligned}
 \phi^{**}(f_v) &:= W^* \xrightarrow{\phi^*} V^* \xrightarrow{f_v} k \\
 g &\mapsto g \circ \phi \mapsto g \circ \phi(v)
 \end{aligned}$$

$$\begin{aligned}
 f_{\phi(v)} &: \\
 g &\mapsto g(\phi(v))
 \end{aligned}$$

Same.

$$(ii) \quad \mathbb{I}_{\text{Vect}_K} \not\Rightarrow (-)^*$$

Since  $\mathbb{I}_{\text{Vect}_K}$  : Covariant

$(-)^*$  : Contravariant.

More significantly  $V \cong V^*$  cannot be defined without choice of basis (basis) which is not preserved by any nonidentity linear endomorphism.

$$(iii) \quad \mathbb{I}_{\text{Set}} \Rightarrow P. \quad \begin{array}{ccc} A & \xrightarrow{\quad} & P(A) \\ \mathbb{I}_{\text{Set}} \downarrow & & \downarrow \\ B & \xrightarrow{\quad} & P(B) \end{array}$$

$$\text{st.} \quad \begin{array}{ccc} A & \xrightarrow{\eta_A} & P(A) \\ f \downarrow & & \downarrow f \\ B & \xrightarrow{\eta_B} & P(B) \end{array} \quad \text{by } \eta_A: A \rightarrow P(A) \quad a \mapsto \{a\}$$

$$(iv) \quad \begin{array}{ccc} & e \mapsto x & \\ X, Y: BG & \rightrightarrows & C \\ & e \mapsto y & \end{array}$$

What is natural transf?

$\Rightarrow \alpha: X \rightarrow Y$  a single  $\alpha$  morphism in  $C$ .

$$\begin{array}{ccc} X & \xrightarrow{\alpha} & Y \\ g_* \downarrow & \circlearrowleft & \downarrow g_* \\ X & \xrightarrow{\alpha} & Y \end{array}$$

st. left diagram

Commutative.  
We say  $\alpha$  is "G-equivariant".

(V)  $\mathcal{O}: \text{Top}^{\text{op}} \Rightarrow \text{Set}$  by  $( )^c$ , taking complement.

i.e.,  $\mathcal{O} \xRightarrow{( )^c} C$

To see this, any cts  $f: Y \rightarrow X \in \text{Top}^{\text{op}}$

$$\begin{array}{ccc}
 Y & \xrightarrow{\mathcal{O}(Y) \xrightarrow{( )^c} C(Y)} & U^c \\
 \downarrow f^! & & \downarrow f^! \\
 \mathcal{O}(X) & \xrightarrow{( )^c} & C(X) \\
 f^!(U) & \xrightarrow{( )^c} & (f^!(U))^c = f^!(U^c)
 \end{array}$$

Then,  $(f^!(U))^c = X - f^!(U) = f^!(U^c)$

Also, it is natural iso since  $( )^c$  is bijection.

(VI)  $(-)^{\text{op}}: \text{Group} \rightarrow \text{Group}$

$$G \longmapsto G^{\text{op}}: \begin{cases} \text{obj: } \bullet \\ \text{Mor: } G \text{ s.t.} \end{cases}$$

$$g^{\text{op}} \cdot h^{\text{op}} = (hg)^{\text{op}}$$

Thus,  $\phi: G \rightarrow H$  induces  $\phi^{\text{op}}: G^{\text{op}} \rightarrow H^{\text{op}}$  by  $g \mapsto \phi(g)$

Let  $\tau_G: G \rightarrow G^{op}$   
 $g \mapsto g^{-1}$

it is not automorphism  $G \rightarrow G$   
 but it is homo  $G \rightarrow G^{op}$

since  $gh \mapsto (gh)^{-1} = (h^{-1} \cdot g^{-1})^{op} = \tau_G(g) \cdot \tau_G(h)$

(since multiple order is reversed)

Thus,

$$\begin{array}{ccc}
 g & \xrightarrow{\tau_G} & g^{-1} \\
 \downarrow \phi & & \downarrow \phi^{op} \\
 G & \xrightarrow{\tau_G} & G^{op} \\
 \downarrow \phi & & \downarrow \phi^{op} \\
 H & \xrightarrow{\tau_H} & H^{op} \\
 \phi(g) & \xrightarrow{\tau_H} & \phi(g)^{-1}
 \end{array}$$

and  $\phi(g^{-1}) = \phi(g)^{-1}$   
 in  $H$ .  
 $\Rightarrow \phi(g^{-1}) = \phi(g)^{-1}$   
 in  $H^{op}$ .

(VII).  $F: \text{Vect}_k \rightarrow \text{Vect}_k$   
 $V \mapsto V \otimes V$

Then,  $V \xrightarrow{\tau_V} V \otimes V$  commute,

$\phi \downarrow \quad \downarrow \phi \otimes \phi$  When  $\tau_V = 0$   
 $W \xrightarrow{\tau_W} W \otimes W$   $\forall V \in \text{Vect}_k$

However, there is no basis independent definable lin. map  $V \rightarrow V \otimes V$ .

Ex)  $\text{Ab}_{fg}$ : Category of fin. gen. ab gp.

for given ab gp  $A$ , let  $TA$ : torsion subgp of  $A$ .

By classification Thm of ab gp,

$$A \cong TA \oplus (A/TA)$$

Prop 1.44.  $A \cong TA \oplus (A/TA)$  are not natural.  
 $\forall A \in \text{Ab}_{fg}$ .

pf) If it were natural,

$$\alpha: A \longrightarrow A/TA \longrightarrow TA \oplus (A/TA) \cong A$$

acting as a natural endomorphism of  
 the  $\mathbb{I}_{\text{Ab}_{fg}}$ , i.e.  $\mathbb{I}_{\text{Ab}_{fg}} \xrightarrow{\alpha} \mathbb{I}_{\text{Ab}_{fg}}$ .

Claim 1: If  $\eta$ : natural endomorphism of  $\mathbb{I}_{\text{Ab}_{fg}}$

$$\Rightarrow \eta: x \mapsto nx \quad \text{for some } n \in \mathbb{Z}$$

pf). Since  $\eta: \mathbb{Z} \rightarrow \mathbb{Z}$  is multiple by  $n \in \mathbb{Z}$ .

and  $\forall A \in \text{Ab}_{gp}$  with map  $\mathbb{Z} \xrightarrow{\alpha} A$   
 $1 \mapsto a$ .

By choosing  $x=1$ ,

$$\eta_A(a) = na.$$

$$\begin{array}{ccc} x & \xrightarrow{\eta_{\mathbb{Z}}} & nx \\ \downarrow & \eta_{\mathbb{Z}} & \downarrow \\ \mathbb{Z} & \xrightarrow{\quad} & \mathbb{Z} \\ \downarrow & \eta & \downarrow \\ A & \xrightarrow{\eta_A} & A \\ ax. & \eta_A(ax) = x\eta_A(a) & \end{array}$$

Thus,  $\eta_A: A \rightarrow A$   
 $a \mapsto na. \quad \square$

Now, if  $A = \mathbb{Z}$ ,  $\alpha: a \mapsto na$  by  
 claim. Then  $\alpha: A \rightarrow A/\eta_A = A \rightarrow A \xrightarrow{\cong} A$   
 is Iso, thus,  $n \neq 0$ .

If  $A = \mathbb{Z}/2n\mathbb{Z}$ ,  $\eta_A = A$ , thus

$$\alpha_{\mathbb{Z}/2n\mathbb{Z}}: A \rightarrow 0 \rightarrow A \rightarrow A$$

$\therefore$  zero map.

But  $n \neq 0$  in  $\mathbb{Z}/2n\mathbb{Z}$  thus.

$$\begin{array}{ccc}
 \mathbb{Z} & \xrightarrow{\alpha_{\mathbb{Z}}} & \mathbb{Z} \\
 \downarrow \pi & & \downarrow \pi \\
 \mathbb{Z}/2n\mathbb{Z} & \xrightarrow{\alpha_{\mathbb{Z}/2n\mathbb{Z}}} & \mathbb{Z}/2n\mathbb{Z} \\
 \downarrow \eta & & \downarrow \eta \\
 \mathbb{Z} & \xrightarrow{\alpha_{\mathbb{Z}}} & \mathbb{Z}
 \end{array}$$

doesn't  
 commute,  
 contradiction

$n \neq 0$ ,

# Ex 1.4.6. (Riesz Representation Thm)

$$\begin{array}{ccc} \mathcal{I} : \text{cHaus} & \longrightarrow & \text{Ban.} \\ \text{cpt} & & \text{Banach space} \\ \text{Hausdoff} & & \end{array}$$

$$\begin{array}{ccccc} X & \longmapsto & \mathcal{I}(X) := \{ \text{signed Baire measure} \} \\ f \downarrow & \longmapsto & \downarrow f^{-1} & \downarrow \mu & \\ Y & \longmapsto & \mathcal{I}(Y) & \mu \circ f^{-1} & \end{array}$$

Is a functor.

$$C^* : \text{cHaus} \longrightarrow \text{Ban.}$$

$$X \longmapsto C(X)^* : \text{dual of the Banach space } C(X) \text{ of cts real valued functions on } X.$$

$$\text{For each } \mu \in \mathcal{I}(X), \quad \phi_\mu : C(X) \rightarrow \mathbb{R} \in C(X)^* \\ g \mapsto \int_X g d\mu$$

$$\text{Let } \eta : \mathcal{I} \Rightarrow C^* \quad \eta_X$$



$$\begin{array}{ccccc}
 \mu & \xrightarrow{\quad} & & & \phi_\mu \\
 \downarrow & \Sigma(X) \xrightarrow{\eta_X} C(X)^* & & & \downarrow \\
 & \downarrow (-\circ f^{-1}) & & & \downarrow (-\circ f) \\
 & \Sigma(Y) \xrightarrow{\eta_Y} C(Y)^* & & & \downarrow \\
 \mu \circ f^{-1} & & & & \phi_{\mu \circ f} \\
 & \xrightarrow{\quad} & \phi_{\mu \circ f^{-1}} & & 
 \end{array}$$

st.  $\phi_{\mu \circ f^{-1}}(g) = \int_Y g \, d(\mu \circ f^{-1})$ ,  $\phi_{\mu \circ f} = \int_X g \circ f \, d\mu$ .

But from real analysis, they are equal.

thus  $\eta: \mu \mapsto \phi_\mu$  is natural transf of  $\Sigma \Rightarrow C^*$ .

RRT: This is natural iso.

Ex 1.4.7.  $\mathcal{L}$ : locally shall

$$f: W \rightarrow Z, \quad h: Y \rightarrow Z$$

$$\begin{array}{ccc}
 C(x, y) & \xrightarrow{h \cdot -} & C(x, z) \\
 \downarrow -\circ f & \eta & \downarrow -\circ f \\
 C(w, y) & \xrightarrow{h \cdot -} & C(w, z)
 \end{array}$$

Where  $h \cdot - = h_x$   
 $-\circ f = f^*$

(post and pre composition)

Since  $C(-, y), C(-, z)$  are functors,

$$h_x : C(-, y) \Rightarrow C(-, z)$$

is natural transformation. (n.t.)

Similarly,

$$f^* : C(x, -) \Rightarrow C(w, -) \quad \text{n.t.}$$

Ex 1.4.9. For  $A, B \in \text{Sets}$ .  $A+B := \text{disjoint union of } A, B$ .

$$A^B := \{B \rightarrow A\}$$

$$\Rightarrow A \times (B+C) \cong (A \times B) + (A \times C), \quad (A \times B)^C \cong A^C \times B^C$$

$$A^{B+C} \cong A^B \times A^C$$

$$(A^B)^C \cong A^{B \times C}$$

Actually, this is natural iso between

$$\text{Set} \times \text{Set} \times \text{Set} \rightarrow \text{Set} \text{ functors}$$

or (contravariant)

Now, restrict this natural iso in  $\text{Fin Iso}$ ,

Obj: finite set Mor: bijection

Then let  $| - | : \text{Fin Iso} \rightarrow \mathbb{N}$  be functor.

( $\mathbb{N}$ : discrete category of natural #)

$$\begin{array}{ccc} A & \xrightarrow{\quad} & |A| \\ \downarrow & & \downarrow \\ B & \xrightarrow{\quad} & |B| = |A| \end{array} \quad \begin{array}{c} \text{f} \\ \text{I}_{|A|} \end{array}$$

Then this induces

$$a(b+c) = ab+ac, \quad (ab)^c = a^c b^c.$$

$$a^{b+c} = a^b a^c, \quad (a^b)^c = a^{bc}.$$

a laws of mult and add in  $\mathbb{N}$ .

deategorification of  $\mathbf{Fin}_{\text{iso}} \Rightarrow \mathbb{N}$  and  
these laws.

Categorification of  $\mathbb{N} = \mathbf{Fin}_{\text{iso}}$

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