

3.1. Limit and Colimit as a universal cone.

Recall: diagram (or shape) J : in a category C
 $\Rightarrow F: J \rightarrow C$ a functor.

Objective: Introduce \lim and colim of diagram
as universal ones over and under the
diagram.

Def 3.1.1 [Cone]

$\forall c \in C$, J : category,

$c: J \rightarrow C$ is called constant functor.

$$\begin{array}{ccc} x & \mapsto & c \\ f & \downarrow & \downarrow I_c \\ y & \mapsto & c \end{array}$$

$$\Rightarrow \Delta: C \longrightarrow C^J$$

$\begin{array}{ccc} c & \xrightarrow{\quad} & c : J \rightarrow C \\ f \downarrow & & \downarrow \alpha_f \\ d & \xrightarrow{\quad} & d : J \rightarrow C \end{array}$

α_f : natural transf
 $(\alpha_f)_c = f$

$\begin{array}{ccc} c & \xrightarrow{I_c} & c \\ \downarrow & & \downarrow \\ f & \xrightarrow{2. If} & d \end{array}$

is an embedding.

Def 3.1.2. A cone over $F: J \rightarrow C$
with summit (apex) $c \in C$.

= Natural transformation $\lambda: c \Rightarrow F$.

• Data $\lambda = \{\lambda_j : j \in J\}$

• From natural transf, for $f: j \rightarrow k \in J$

$$\begin{array}{ccc} c & \xrightarrow{\lambda_c} & c \\ \lambda_j \downarrow & \Downarrow & \downarrow \lambda_k \\ F_j & \xrightarrow[F_f]{} & F_k \end{array} \Leftrightarrow \begin{array}{ccc} c & & c \\ \lambda_j \swarrow & \Downarrow & \searrow \lambda_k \\ F_j & \xrightarrow[F_f]{} & F_k \end{array}$$

• Thus, if we have data satisfying the commutative diagram, then we can define cone.

A cone under F with nadir c

= Natural transf $\lambda: F \Rightarrow c$ s.t

$$\forall f: j \rightarrow k, \quad F_j \xrightarrow[F_f]{} F_k$$

$$\lambda_j \searrow \swarrow \lambda_k$$

Cone under F is called cocone
since it is dual of a cone.

(Cone over $F: J^{\text{op}} \rightarrow C^{\text{op}}$.)

$\lim F$: universal cone over F

Def) $F: J \rightarrow C$ diagram.

Cone $(-, F): C^{OP} \rightarrow \text{Set}$.

$$\begin{array}{ccc} C & \longmapsto & \{\lambda: \underset{J}{\underset{\downarrow}{C}} \Rightarrow F : \text{natural transformation}\} \\ f \downarrow & \longmapsto & \uparrow \alpha_f: C \rightarrow J \text{ as natural transformation.} \\ \downarrow & \longmapsto & \{\lambda: I \Rightarrow F, \text{ natural transformation}\} \end{array}$$

$\lim F$: representation of $\text{Cone}(-, F)$, i.e.

$\text{Cone}(-, F) \cong C(-, \lim F) \Leftrightarrow \text{terminal object}$
of $\int \text{Cone}(-, F)$

Similarly,

$$\begin{array}{ccc} \text{Cone}(F, -): C \rightarrow \text{Set} & & \text{limit cone:} \\ C & \longmapsto & \{\lambda: C \Rightarrow F\} \\ f \downarrow & & \downarrow \alpha_f \\ \downarrow & \longmapsto & \{\lambda: I \Rightarrow F\} \end{array}$$

Universal element
in this.

Colimit cone
(universal element)

$\text{Colim } F$: rep of $\text{Cone}(F, -)$. \Leftrightarrow initial object
of $\int \text{Cone}(F, -)$

Notes that if $\lambda_C: C \Rightarrow F$, $\lambda_d: d \Rightarrow F$, $f: C \rightarrow d$.
then α_f sends λ_d to λ_C via below diagram.

$$\begin{array}{ccc} \lambda_C & \xrightarrow{f} & \lambda_d \\ \downarrow \alpha_f & & \downarrow \alpha_d \\ F_j & \longrightarrow & F_k \end{array}$$

So: $\text{Cone}(-, F)$
is well-defined.

(Similar to $\text{Cone}(F, -)$).

Prop 3.1.7. $\lambda: l \Rightarrow F$, $\lambda': l' \Rightarrow F$ are the limit cones, $\Rightarrow l \stackrel{\exists!}{\cong} l'$. Unique iso. commutes with less of the limit cones.

pf) $(\text{lim } F, \lambda)$, $(\text{lim } F, \lambda')$ are terminal
 $\Rightarrow \text{``} \cong \text{''}$ by Corollary 2.3.2.
 And isomorphism in $\text{Scone}(-, F)$ gives commuting diagram:

Rank: There may be nontrivial automorphism between l and l' . but by Corollary 2.3.2 if it commutes with core, then it is the iso in prop 3.1.7.

Def 3.1.9. [Product]

Product = limit of a diagram indexed by
 J : a discrete category with only
Identity morphisms.

Thus, $F: J \rightarrow C$ is just collection
 $\{F_j\}_{j \in J}$. (Since J has no nontrivial morphism)

Thus, cone over J is $\lambda: c \Rightarrow F$
i.e. collection of $\{\lambda_j: c \rightarrow F_j\}_{j \in J}$.

Therefore its limit is $\prod_{j \in J} F_j$

with less $\pi_k: \prod_{j \in J} F_j \rightarrow F_k$.

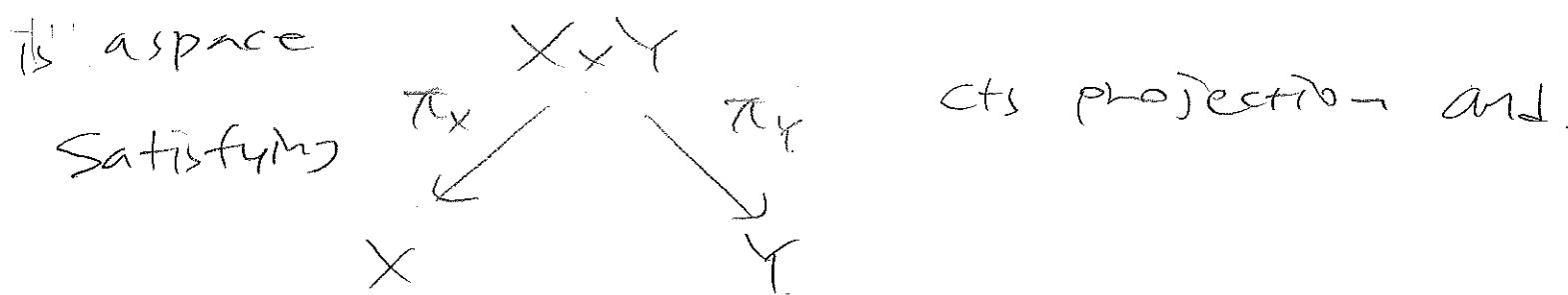
By the universal property of $\prod_{j \in J} F_j$ as a cone gives natural isomorphisms.

$$(C(c, \prod_{j \in J} F_j)) \xrightarrow[\cong]{(\pi_k)_*} \text{Cone}(c, F) \cong \prod_{k \in J} C(c, F_k)$$

(Yoneda lemma)

Just...

Ex 3.10. $X, Y \in \text{Top}$. Product of $X \times Y$



has universal property: For any $A \in \text{Top}$ with $f: A \rightarrow X, g: A \rightarrow Y, \exists h: A \rightarrow X \times Y$

St.

$$\begin{array}{ccc} f & A & \\ \swarrow \exists h & \searrow & \\ X & \xleftarrow{\pi_X} & Y \\ & \pi_X & \pi_Y \end{array}$$

commute

To construct $X \times Y$ as usual sense (Cartesian product)

By taking $A = \{*\}$, A as a constant functor

says that $X \times Y \supseteq X \times Y$ as a set.

(By similar argument, $X \times Y = X \times Y$ as a set)

This is because $\text{Top}(A, X \times Y) \cong \text{Top}(A, X) \times \text{Top}(A, Y)$
as a set.

To show $X \times Y = \underset{\text{Cartesian}}{X \times Y}$ as a top space

take $A = X \times Y$ as a set with
various topology.

By the universal property,

$X \times Y$ forced to defined as the coarsest
topology on cartesian product set.

π_X, π_Y are cts. (These are ^{case} when $J = \{\circ, -\}$)

(Do the saeths for any Mdex J .)

Def 3.1.11. Terminal Object = product
when the indexing category is empty.

i.e., if $J = \{\}$, Cone over J with sumpt C
is just $C \Rightarrow \int \text{Cone}(-, J \rightarrow C) \cong C$.

By def of limit, $\lim(J \rightarrow C)$ is terminal obj
of $\int \text{Cone}(-, J \rightarrow C) \cong C$.

Ex 3.1.12 $\{1 : \mathbb{Z} \cdot \mathbb{Z}^1\}$ called terminal category

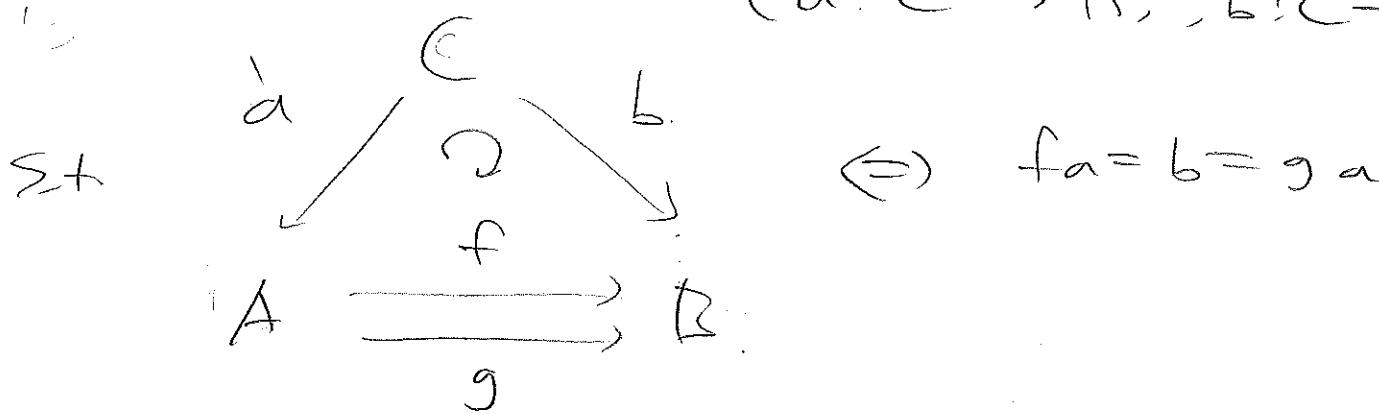
Since it is terminal obj of Cat or CAT

Def 3.1.13 [Equalizer].

Equalizer = Limit of $F: J \rightarrow \mathbf{C}$

Where $J = \{ \xrightarrow{\alpha}, \xrightarrow{\beta} \}$. parallel pair.

\Rightarrow Let $F(J) = A \xrightarrow{f} B$. A cone over F with sumit C is $(a: C \rightarrow A, b: C \rightarrow B)$



Thus cone over parallel pair $A \xrightarrow{f} B$ (with sumit C) is represented by a morphism $a: C \rightarrow A$
St $fa = ga$.

Hence Equalizer is the universal arrow with this property., $h: E \rightarrow A$. In particular.

Here $a: C \rightarrow A$, $\exists ! k: C \hookrightarrow E$ st



Ex 3.1.14. $\phi, \psi: G \rightarrow H$ \in Group

$$\Rightarrow \text{Eq}(\phi, \psi) = (\ker \phi \xrightarrow{\cong} G)$$

Since

$$\begin{array}{ccc} C & \xrightarrow{\alpha} & \\ \exists! k \downarrow & \searrow & \\ \ker \phi & \hookrightarrow G & \xrightarrow{\phi} H \\ & \downarrow \psi & \end{array}$$

$$\text{Eq}(\phi, \psi) = \left(\{ g \in G : \phi(g) = \psi(g) \} \hookrightarrow G \right)$$

If H is abelian, ($= \ker(\phi - \psi)$)

Actually, $\text{Eq}(A \xrightarrow{f} B) \hookrightarrow C$ is monomorphic
 $h: E \rightarrow A \dashrightarrow E$

Def 3.1.15 [Pullback]

Pullback = limit of $(f: \dots \rightarrow \circ \leftarrow \circ \dashrightarrow C)$

$$\text{Let } F(J) = B \xrightarrow{f} A \xleftarrow{g} C$$

Then cone over F with sumit B is

$$\begin{array}{ccc} B & \xrightarrow{c} & C & \text{triple morphism} \\ b \downarrow & \searrow a^2 & \downarrow g & \\ B & \xrightarrow{f} & A & \text{satisfies left} \\ & f & & \text{comm. diagram} \end{array}$$

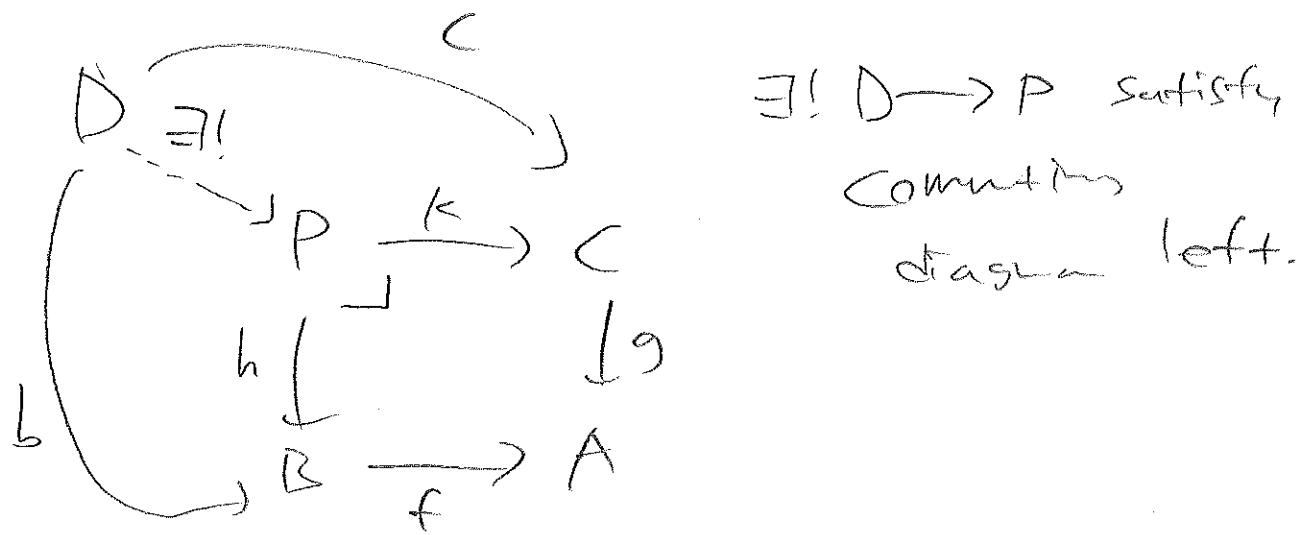
$$\Rightarrow gc = a = fb.$$

Thus we can represent this core by two morphism $B \xleftarrow{b} D \xrightarrow{c}$ satisfies rectangle.

Hence pull back : $B \xleftarrow{h} P \xrightarrow{k} C$

is universal one over $(B \xrightarrow{f} A \xleftarrow{g} C)$

i.e. for any core $B \xleftarrow{l} D \xrightarrow{c}$,



" \sqsubset " denote P is pullback, i.e. limit diagram
not just commutative square.

P is called "fiber product", denoted by $B \times_A C$.

If C is concrete category
thus $B = \{*\} \subseteq A$ as a set

$$\begin{array}{ccc} P & \xrightarrow{\quad} & C \\ \downarrow \perp & & \downarrow g \\ \{*\} & \xhookrightarrow{\quad} & A \end{array} \quad \text{implies} \quad P = \{c \in C : g(c) = *\}$$

$$\{*\} \xhookrightarrow{\quad} A = g^{-1}(*)$$

thus fiber of map g .

over $*$ (or $f(*)$)

Ex 3.1.18. $\rho: (\mathbb{R} \rightarrow S^1)$ cts.

$$t \mapsto e^{2\pi i t}.$$

What is P ?

$$P \dashrightarrow \mathbb{R}$$

$$\vdots \qquad \downarrow \rho \qquad = P^*(1)$$

$$1 \xleftarrow[\mathbb{Z}]{} S^1 \qquad = \{t \in \mathbb{R} : e^{2\pi i t} = 1\}$$

$$= \mathbb{Z}$$

Similar as equalizer.

Pullback $(B \xrightarrow{\quad} A \leftarrow C)$ defines $\underset{\text{mono.}}{(B \times_A C) \rightarrow C}$

Ex 3.1.19.

What is P ? in $A\mathbb{Z}$.

$$P \xrightarrow{\alpha} \mathbb{Z}$$
$$b : \downarrow \quad \downarrow n$$

must $na = mb$.

$$\mathbb{Z} \xrightarrow{m} \mathbb{Z} \quad P : ab \text{ or } \text{s.t. } \forall p \in P$$

$$na(p) = mb(p)$$

Moreover, if it commutes with any other cores, satisfy $n(a(p)) = mb(p)$.

$\Rightarrow na(p) = mb(p)$ should be l.c.m
of $n, m,$

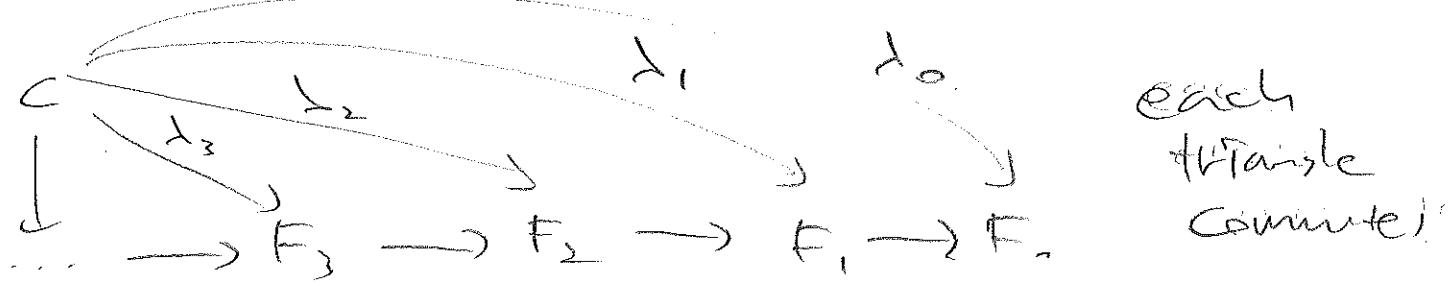
thus, we may take $P = \mathbb{Z}$ with

$$a(1) = \frac{\text{lcm}(n, m)}{n}, \quad b(1) = \frac{\text{lcm}(n, m)}{m}.$$

Def 3.1.2.1 [Inverse Limit].

Limit of $F: J \rightarrow C$ when $J = \omega^{\text{op}}$
 $\dots \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow 0, : \omega^{\text{op}}$

There's core over F with sumit C is



$\varprojlim F_n$: is the terminal cone.

(Thus direct limit is when $J = \omega$.)
and colimit of $F: J \rightarrow C$

Ex 3.1.22

$$\pi_p = \varprojlim \mathbb{Z}/p^n.$$

Their dual notions are

Coproduct : Colimit of discrete category

Initial object : " Empty "

Coequalizer : " $\xrightarrow{\quad} \xleftarrow{\quad}$ "

Pushout : " $\circ \leftarrow \circ \rightarrow \circ$ "

Direct limit : " ω .

(Sequential limit)

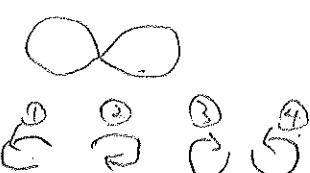
Ex 3.1.24 Pushout of $s' \leftarrow * \rightarrow s'$

$$is s' \vee s' \begin{array}{c} * \xrightarrow{+} s' \\ \downarrow \circ \quad \downarrow \circ \\ s' \rightarrow s' \vee s' \end{array}$$

Moreover,

$$s' \xrightarrow{aba^{-1}b^{-1}} s' \vee s' \begin{array}{c} \downarrow \circ \quad \downarrow \circ \\ D^2 \dashrightarrow T \end{array}$$

T : torus.
 $a, ba^{-1}b^{-1}$ loop



Ex). 1.25 Cokernel in Groups
 = Coequalizer ($\phi: G \rightarrow H$, $\alpha: G \rightarrow H$)

Ex). 1.26. (n -skeleta)

$$X_0 \hookrightarrow X_1 \hookrightarrow X_2 \hookrightarrow \dots$$

Colimit of sequence = $\bigcup_{n \geq 0} X_n$.

ex) CW complex = colimit of its n -skeleta.

$$\text{Zmk 3.1.27) } f: \coprod_{i \in I} A_i \rightarrow X \iff \{f_i: A_i \rightarrow X \mid i \in I\}$$

$$\text{Thus, } C(\coprod_{i \in I} A_i, X) \xrightarrow{\cong} \prod_{i \in I} C(A_i, X)$$

$$f \longmapsto (f_i)_{i \in I}.$$

by letting $\tilde{c}_i: A_i \rightarrow \coprod A_i$ and $f_i = f \circ \tilde{c}_i$.

$$\text{Dually, } g: X \rightarrow \prod_{j \in J} B_j \iff \{g_j: X \rightarrow B_j \mid j \in J\}$$

$$\Rightarrow C(X, \prod_{j \in J} B_j) \xrightarrow{\cong} \prod_{j \in J} C(X, B_j)$$

$$g \longmapsto g_i.$$

by letting $\pi_j: \prod_{i \in I} B_i \rightarrow B_j$ with $g_i = \pi_j \circ g$.

Then, if $\coprod_{i \in I} A_i \xrightarrow{f} \prod_{j \in J} B_j$, it shows

$$\begin{array}{ccc} \tilde{c}_i \uparrow & \square & \downarrow \pi_j \\ A_i & \dashrightarrow & B_j \end{array}$$

$$f(i, j) \dashrightarrow f(i', j')$$

I.e., a map from Coproduct to product is determined by a matrix of component maps.

In certain categories, Ab , Mod_R , Ch_R , have a notion of zero homomorphism between any pair of objects.

In any finite collection I ,

$$\coprod_{i \in I} A_i \xrightarrow{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}} \prod_{i \in I} A_i \quad \text{make sense;}$$

This induces isomorphism in "Abelian category".

$\Rightarrow \coprod_{i \in I} A_i \cong \prod_{i \in I} A_i$. So use notation $\bigoplus_{i \in I} A_i$, a direct sum.

Then any ~~map~~ composition of maps between direct sums is matrix multiplication.

Ex 3.1.i) Let $f: c \rightarrow d \in \text{Hom}(c)$. Then,

$\text{Cone}(c, F)$ $\alpha_f = \text{Gne}(f, F)$ st.

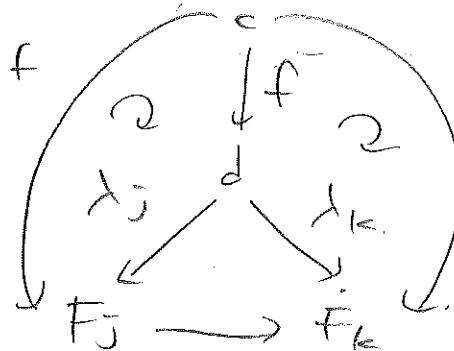
$$\alpha_f \uparrow$$

If $\lambda \in \text{Cone}(d, F)$ then

$\text{Cone}(d, F)$

$$\lambda_j \circ f$$

then,



$$\lambda_k \circ f$$

implies that

$$\alpha_f \circ \lambda \Rightarrow d \Rightarrow F \in \text{Cone}(c, F).$$

Dually,

$$\text{Cone}(F, f): \text{Cone}(F, c) \xrightarrow{\alpha_f} \text{Cone}(F, d)$$

st. if $\lambda \in \text{Cone}(F, c)$, $f \circ \lambda$ is a cone in $\text{Cone}(F, d)$

$$F_j \rightarrow F_k$$

in $\text{Cone}(F, d)$.



$$\therefore F \Rightarrow c \Rightarrow d$$

$$\in \text{Cone}(d, F)$$

Ex 3.1.ii) Notes that for any $\lambda \in \text{Gne}(c, F)$

$\lambda: c \Rightarrow F$ natural transformation, thus $\lambda \in \text{Hom}(\Delta(c), F)$

$$\Rightarrow \text{Cone}(c, F) \subseteq \text{Hom}(\Delta(c), F).$$

Conversely, let $\lambda \in \text{Hom}(\Delta(c), F)$. By definition, λ is a cone over F with summit $c \Rightarrow \text{Cone}(c, F) = \text{Hom}(\Delta(c), F)$.

Also, $\text{Hom}(\Delta(f), F) \cong \alpha_f = \text{Cone}(f, F)$

thus, $\text{Cone}(-, F) \cong \text{Hom}(\Delta(-), F)$.

Ex. 3.1. iii) WTS, $\int \text{Cone}(-, F) \cong \Delta \downarrow F$.

	$\int \text{Cone}(-, F)$	$\Delta \downarrow F$
Obj	$(c, \lambda), c \in C$ $\lambda \in \text{Cone}(c, F)$	$(\Delta(c), F, \lambda : \Delta(c) \Rightarrow F)$
Mor.	(c, λ) $f \downarrow$ (c', λ')	$\Delta(c), F, \lambda \vdash \alpha_f$ $(f, 1_F) \downarrow$ $\Delta(c) \rightarrow \Delta(c')$ $\lambda' \downarrow$ $(\Delta(c'), F, \lambda').$ F $\Rightarrow \alpha_f \circ \lambda' = \lambda$

Hence, there is bijection between them.

Now, for any $\top F \in (C^{\text{op}})^J$:

$$\int \text{Cone}^{\text{op}}(-, F) \cong \Delta^{\text{op}} \downarrow F^{\text{op}}$$

where $\Delta^{\text{op}} : C^{\text{op}} \rightarrow (C^{\text{op}})^J$, $F^{\text{op}} : 1 \rightarrow (C^{\text{op}})^J$.

And observe $\int \text{Cone}^{\text{op}}(-, F) = \int \text{cone}(F, -)$

$$\text{and } \Delta^{\text{op}} \downarrow F^{\text{op}} = F \downarrow \Delta^{\text{op}}$$

Ex 3.1. iv) From universal property,

$\begin{array}{ccc} x_j & \xrightarrow{\ell} & x_k \\ \downarrow \ell \circ b & & \downarrow x_j \\ x_j & \xrightarrow{\ell' \circ x_k} & x_j b = x_j \end{array}$ below diagram commutes for any
of triangle. Then, $x_j a = x_j$
 $x_j b = x_j \Rightarrow x_j b a = x_j \Rightarrow b a = 1_{\ell}$.
 $F_j \longrightarrow F_k$ Since $\text{Hom}(\ell, \lambda), (\ell, \lambda)) = \{1_{(\ell, \lambda)}\}$.

and \mathbf{f}_l induces $I_{(\ell, N)}$.

By the similar argument, $a b = I_{\ell'}$ \Rightarrow a, b are isomorphisms $\Rightarrow \ell \cong \ell'$.

Ex 3 (v) If $\lim F$ exists with limit cone λ , then for any cone $\phi: P \Rightarrow F$, $\exists! a: P \rightarrow \lim F$

$$\begin{array}{ccc} & P & \\ \phi_j \swarrow & \downarrow \exists! a & \searrow \phi_k \\ F_j & \xrightarrow{\quad} & F_k \end{array}$$

s.t. left diagram commutes
 \Rightarrow Let $\text{sumit}(P, \leq)$
be a subcategory
of (P, \leq) spanned by
elements in (P, \leq) which is
a sumit of a cone.

Then, $\lim F$ is "a" maximum
of $\text{sumit}(P, \leq)$

(Not "the" maximum since
preorder has two distinct
element on which both \geq and \leq
hold.)

Conversely, colimit is a minimum of subcategory
 $\text{nadir}(P, \leq)$, spanned by elements which is nadir
of some cocone over F .

EX 3.1, vi) Let $f', g': C \rightarrow E$ s.t. $hf' = hg'$.
 Let $a = hf' = hg'$. Then, $fa = fhf' = (gh)f' = g(a)$
 Thus universal property

i.e. $\exists! k: C \rightarrow E$ gives Unique $k: C \rightarrow E$
 $E \xrightarrow{h} A \xrightarrow{f} B$ s.t. $a = hk$.

This implies $f' = g'$.

Since uniqueness of k s.t.
 $a = hk$ implies $f' = k = g'$.

$\Rightarrow h$ is mono.

3.1, vii) Let $a, b: P \rightarrow P$ s.t. $ka = kb$.

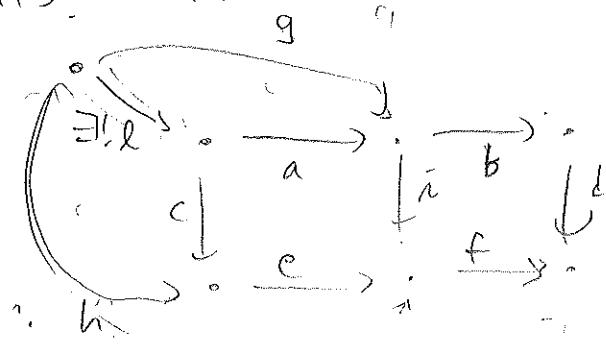
$P \xrightarrow{q} P \xrightarrow{k} C$ Then, $gka = gkb$
 $P \xrightarrow{b} P \xrightarrow{h} A$ $\Rightarrow fha = fhb$
 $\xrightarrow{d} B \xrightarrow{f} A$ $\Rightarrow ha = hb$

Now let $c = ka = kb$
 $d = ha = hb$

Then, $\exists! l: D \rightarrow P$ s.t. $c = fl$ $d = hl$.

This implies $a = l = b$ by uniqueness, $\Rightarrow k$ is
 mono,

3.1 viii) Suppose composite rectangle is pull back
 then, $\exists! l: D \rightarrow P$ for $b \circ g$ and h , the
 "big" pull back gives unique l



Set. $h = cl$, $bg = bal$.

Apply "small" (right) pull back on b_9 and e_h
 We get unique k s.t. $bk = b_9$ and
 $eh = ik$.

By uniqueness of k ,
 $ak = k = g$.

Thus, $h = cl$, $g = al$. This shows the
 existence of a map l to make $\begin{array}{c} \xrightarrow{a} \\ \downarrow \\ c \end{array} \xrightarrow{i} \begin{array}{c} \xrightarrow{e} \\ \downarrow \\ f \end{array} \xrightarrow{j} d$ (left)

pullback. To see uniqueness, suppose $\exists l' \in \mathcal{F}$,
 $g = al'$, $h = cl'$. Then, applying l 's pullback
 on b_9 and h gives $h = cl = cl'$.

Since l is unique, $l' = l$. \square

Conversely, both left and right one are pullbacks

then right pullback induce k ,

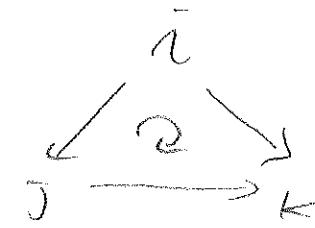
for g and e_h . s.t.

$$g = bk, eh = ik$$

$$\text{as } k = al \text{ and } h = cl$$

thus, $g = bal$, $h = cl$. Also if $\exists l' \in \mathcal{F}$ s.t. $g = bal'$, $h = cl'$
 then, by uniqueness of k and l , $al' = k = al \Rightarrow l = l'$.
 \therefore Composite rectangular is pullback

Ex 3.1-ix) Let i be initial in J . Then,
for any $j \rightarrow k \in J$.

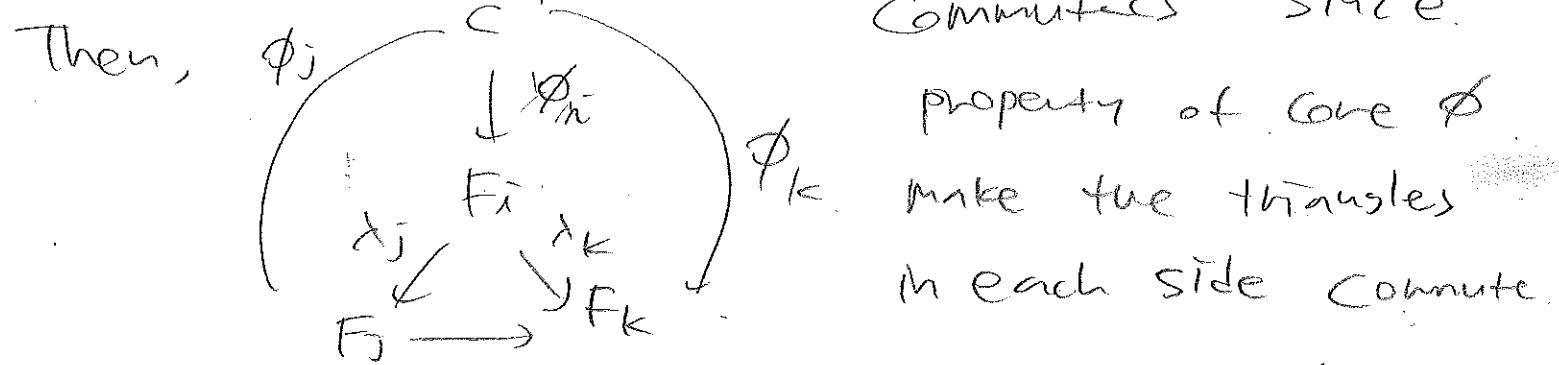


by uniqueness of $\text{Hom}(i, k)$,

Thus, let $\lambda_j = F(i \rightarrow j) \quad \forall j \in J$.

Then, $\{\lambda_j\}$ is a cone over F with summit F_i .

Let $\phi: C \Rightarrow F$ a cone with summit C .



Now suppose $f: C \rightarrow F_i$ s.t. $\alpha_f(\lambda) = \phi$.

Then, ϕ_j

$\Rightarrow f = \phi_i$. Hence in $\text{Cone}(-, F)$

(F_i, λ) is terminal element.
 $\therefore F_i = \text{lim } F$.

Now suppose J is indexed by W , i.e.

$\exists G: W \rightarrow J$ exists.

If J has terminal i , then $F(J \rightarrow C)$ has limit

F_i .

Then, if ϕ is a cocone with nadir c and
 λ is colimit core,

The diagram shows a cocone ϕ with vertex c . The objects are $F_0 \rightarrow F_1 \rightarrow \dots$. The core λ consists of maps $\lambda_i : F_i \rightarrow F_{\bar{i}}$ and $\lambda_k : F_{\bar{i}} \rightarrow F_k$. The nadir c is at the bottom. A vertical arrow labeled $\phi_{\bar{i}}$ points from c to $F_{\bar{i}}$. A curved arrow labeled ϕ^0 points from c to the top of the cone.

implies that

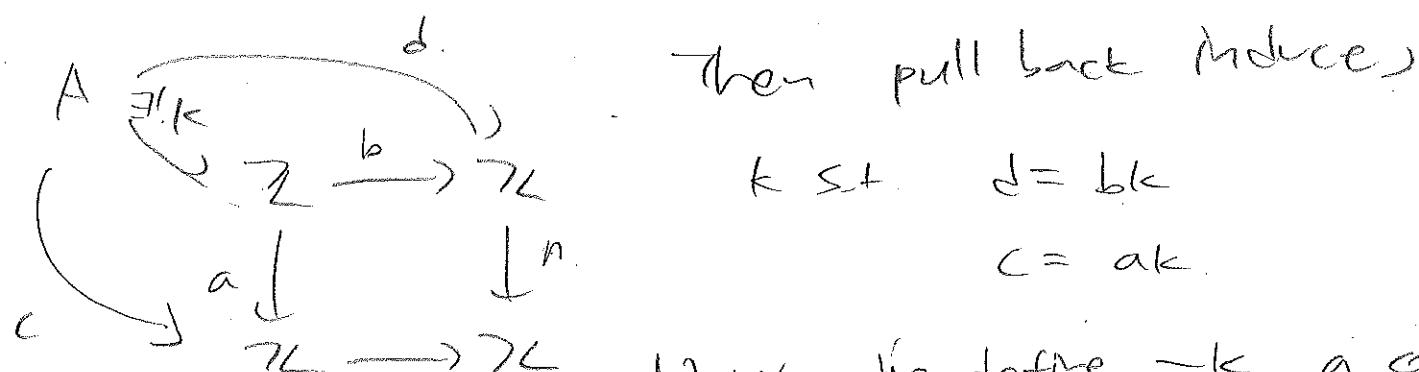
$$\phi_{\bar{i}} \circ \lambda_k = \phi_k.$$

Since J is indexed by W , this implies
 that for $k > \bar{i}$, $\lambda_k \circ (F_{\bar{i}} \rightarrow \dots \rightarrow F_k) = 1_{F_{\bar{i}}}$.

$\Rightarrow F_k = F_{\bar{i}}$. Hence $J = i$, a subcategory
 of W spanned by $G(0), G(1), \dots, G(\bar{i})$,

And colimit of F is $F_{\bar{i}}$.

Ex 3.1.x). Let $A \xrightarrow{\quad f \quad} \mathcal{U}$ be a map s.t.



then pull back induces

$$k \text{ s.t. } d = bk$$

$$c = ak.$$

Now, we define $-k$, a composition
 of k with inversion. Then
 $(-k) \circ (b) = k \circ b$ and $-a \circ k = ak$.

Thus, $(-a, -b)$ also has a
 universal property.

However,

$$\begin{array}{ccc}
 \mathbb{Z} & \xrightarrow{\quad b \quad} & \mathbb{Z} \\
 \downarrow -1 & \nearrow -b & \downarrow n \\
 \mathbb{Z} & \xrightarrow{-a} & \mathbb{Z} \\
 \downarrow a & & \downarrow n \\
 \mathbb{Z} & \xrightarrow{m} & \mathbb{Z}
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathbb{Z} & \xrightarrow{\quad -b \quad} & \mathbb{Z} \\
 \downarrow -1 & \nearrow b & \downarrow n \\
 \mathbb{Z} & \xrightarrow{-a} & \mathbb{Z} \\
 \downarrow a & & \downarrow n \\
 \mathbb{Z} & \xrightarrow{m} & \mathbb{Z}
 \end{array}$$

Shows that $(\mathbb{Z}, (a, b)) \cong (\mathbb{Z}, (-a, -b))$

in the category of cores. So, this $(-a, -b)$ is just another limit one naturally isomorphic to (a, b) .

Ex 3.1.xi) Note, that each leg is denoted by

$$\bar{c}_j : A_j \rightarrow \coprod_{j \in I} A_j$$

Since this category has a morphism for each pair of objects, let $f_{ij} : A_i \rightarrow A_j \quad \forall i \in I, f_{ii} = 1_{A_i}$ we can pick such f_{ij} by given assumption.

Then, A_j is coproduct over I , so universal property gives that

$$A_{\bar{j}} \xrightarrow{\bar{c}_{\bar{j}}} \coprod_{j \in I} A_j$$

$$\begin{matrix}
 \forall i \in I, & \swarrow & \downarrow \exists k_{ji} \\
 f_{ij} & & A_j
 \end{matrix}$$

Now let $\bar{i} = j$ Then

$$\begin{array}{ccc}
 A_j & \xrightarrow{\bar{c}_j} & \coprod_{j \in I} A_j \\
 \downarrow & \swarrow & \downarrow \exists k_{jj} \\
 \exists k_{jj} & & A_j
 \end{array}
 \Rightarrow \bar{i} \text{ is split mono.}$$

3. (xii) Let $\lambda: C \Rightarrow F$ is J -shaped cone.
 $\eta: C \Rightarrow FE$ is I -shaped cone.

Then, for $f: i \rightarrow j' \in J$ essential surjectivity
 gives $E(i) \xrightarrow{\cong} i'$, $E(j) \xrightarrow{\cong} j'$ for some $i, j \in I$.

Hence, $J(E(i), E(j)) \cong J(i', j')$ since
 composition of isomorphisms induce bijection.

on $J(-, -)$. Thus, $\exists f: i \rightarrow j$ s.t.

$$E(f) = [E(i) \xrightarrow{\cong} i' \xrightarrow{f'} j' \xrightarrow{\cong} E(j)].$$

Hence

$$\begin{array}{ccc} & C & \\ i & \swarrow \alpha & \searrow \beta \\ FE(i) & \xrightarrow{FE(f)} & FE(j) \\ \cong \{fa_i\} & \xrightarrow{\cong} & \cong \{f\alpha_j\} \end{array}$$

Thus, sending

$$\lambda: fa_i \mapsto \{F\alpha_j \circ \lambda_j\}$$

$$\text{Sends } (c, \lambda) \in \text{Cone}(-, F) \text{ to } (c, \lambda) \in \text{Cone}(-, FE)$$

$$F(i') \xrightarrow{Ff'} F(j')$$

$$\text{to } (c, \lambda) \in \text{Cone}(-, F)$$

$$\begin{array}{ccc} & C & \\ i' & \swarrow \alpha' & \searrow \beta' \\ \lambda & \xrightarrow{\cong} & \lambda' \end{array}$$

Conversely, sending

$$\lambda \mapsto \lambda a_i^{-1} := \{F(a_i^{-1}) \circ \lambda_i\}$$

$$\text{Sends } (c, \lambda) \in \text{Cone}(-, F) \text{ to } (c, \lambda a_i^{-1}) \in \text{Cone}(-, FE)$$

$$\text{to } (c, \lambda) \in \text{Cone}(-, FE)$$

And this map is inverse to each other.

So if we show these are functor, then $\text{Cone}(-, F)$
 and $\text{Cone}(-, FE)$ are isomorphism of category.

If $(C, \alpha) \xrightarrow{g} (C', \alpha') \in \text{Scone}(-, \text{FE})$
 then, $g: C \rightarrow C'$ and $\text{Scone}(g, \text{FE})(\alpha') = \alpha$

And, $\text{Scone}(g, \text{FE}) = \alpha_g: C \Rightarrow C'$.

$$\text{So, } \alpha' \circ \alpha_g = \alpha.$$

$$\Rightarrow \begin{array}{c} n'_i \xrightarrow{\beta} n_j \\ \downarrow f \downarrow \quad \downarrow \alpha \\ n'_i \xrightarrow{\text{FE}(\beta)} n_j \\ \text{FE}(\alpha) \xrightarrow{\text{FE}(f)} \text{FE}(\alpha) \end{array}$$

$$F(a_i) \supseteq F(a_j)$$

$$F(\alpha'_i) \xrightarrow{Ff'} F(\alpha'_j)$$

Now sending $\alpha \mapsto \alpha a$.
 and $\alpha' \mapsto \alpha' a$.

has

$$\alpha' a \circ \alpha_g = \alpha a.$$

$$\text{So } \alpha \mapsto \alpha a.$$

$$\alpha_g \mapsto \alpha_g$$

$$\alpha' \mapsto \alpha' a.$$

is function.

In the similar way, we can conclude that

$$\text{Scone}(-, \text{FE}) \cong \text{Scone}(-, \text{F}).$$

Hence If \lim_{FE} exists, $(\lim_{\text{FE}}, \lambda)$ is initial
 in $\text{Scone}(-, \text{FE})$. thus, $(\lim_{\text{FE}}, \lambda a)$ is initial
 since they are isomorphic category.
 $\Rightarrow \lim_{\text{F}} = \lim_{\text{FE}}.$

Ex 3, 1, (iii). Coproduct of CRngs.

Let $\coprod_{j \in J} R_j$ be a coproduct of CRngs.

Then, $i_j: R_j \rightarrow \coprod_{j \in J} R_j$ exists and

for any $f_j: R_j \rightarrow S$, $\exists! k: \coprod_{j \in J} R_j \rightarrow S$

st. $R_j \xrightarrow{i_j} \coprod_{j \in J} R_j$

$$\begin{array}{ccc} f_j & \swarrow & \downarrow k \\ & & S \end{array}$$

When $J=2$, then

Right diagram holds.

$$R_1 \xrightarrow{i_1} R_1 \amalg R_2 \xleftarrow{i_2} R_2$$

$$\begin{array}{ccc} f_1 & \swarrow & \downarrow k \quad ? \\ & & S \swarrow f_2 \end{array}$$

Notes that every comm. ring is \mathbb{Z} -module.

So, $R_1 \times R_2 \xrightarrow{\otimes} R_1 \otimes_{\mathbb{Z}} R_2$ from the universal

$$\begin{array}{ccc} f_1(r_1) \cdot f_2(s_2) & \swarrow & \downarrow \exists! f \\ & & S \end{array}$$

property of tensor product

We claim that $R_1 \otimes_{\mathbb{Z}} R_2 \cong R_1 \amalg R_2$.

$R_1 \xrightarrow{\bar{i}_1} R_1 \amalg R_2$ induces induced in $R_1 \times R_2 \rightarrow R_1 \amalg R_2$

So universal property of tensor product gives unique $\bar{i}: R_1 \otimes_{\mathbb{Z}} R_2 \rightarrow R_1 \amalg R_2$. Also, universal property of coproduct on $R_1 \xrightarrow{- \otimes 1} R_1 \otimes R_2, R_2 \xrightarrow{1 \otimes -} R_1 \otimes R_2$

Induces unique $\kappa: R_1 \amalg R_2 \rightarrow R_1 \otimes_{\mathbb{Z}} R_2$

\Rightarrow

$$\begin{array}{ccccc}
 & \bar{i}_1 & \nearrow & \bar{i}_2 & \\
 R_1 & \swarrow & \downarrow \kappa & \searrow & R_2 \\
 & \xrightarrow{- \otimes 1} & R_1 \otimes_{\mathbb{Z}} R_2 & \xleftarrow{1 \otimes -} &
 \end{array}$$

$$\text{Now } \bar{i} \circ \kappa \circ \bar{i}_1 = \bar{i}(- \otimes 1) = \bar{i}_1.$$

\Rightarrow Thus $R_1 \xrightarrow{\bar{i} \circ \kappa \circ \bar{i}_1} R_1 \amalg R_2$ with universal property
 $R_2 \xrightarrow{\bar{i} \circ \kappa \circ \bar{i}_2} R_1 \amalg R_2$

of Coproduct Induce $\exists! R_1 \amalg R_2 \rightarrow R \amalg R_2$

which is identity. $\Rightarrow \bar{i} \circ \kappa = 1_{R_1 \amalg R_2}$.

Conversely, $\kappa \circ \bar{i} \circ (1 \otimes -) = \kappa \circ \bar{i}_2 = 1 \otimes -$.

$$\begin{array}{ccc}
 \text{Hence, } R_1 \times R_2 & \xrightarrow{\otimes} & R_1 \otimes R_2 \\
 & \downarrow & \downarrow \exists! \overline{\kappa \circ \bar{i} \circ (1 \otimes -)} \\
 & \xrightarrow{\kappa \circ \bar{i} \circ (1 \otimes -)} & R_1 \otimes R_2 \\
 & \xrightarrow{\otimes} &
 \end{array}
 \Rightarrow \begin{aligned}
 \kappa \circ \bar{i} \circ (1 \otimes -) \\ = 1_{R \otimes R}
 \end{aligned}$$

Hence, κ and \bar{i} are isomorphisms.

3.2 Limits in the category of sets.

Def 3.2.1. $F: J \rightarrow C$ diagram is small if J is small category.

C is complete $\Leftrightarrow \forall F: J \rightarrow C$ small diagram $\exists \lim F$

" Cocomplete $\Leftrightarrow \forall F: J \rightarrow C$ " $\exists \text{colim } F$.

Claim: $F: J \rightarrow \text{Set}$ small diagram. If $\lim F$ exists, then $\lim F \cong \text{Cone}(1, F)$.
 pf). $\text{Set}(X, \lim F) \cong \text{Cone}(X, F) \quad \forall X \in \text{Set}$.
 Take $X = 1$, singleton. Then.

$$\lim F \cong \text{Set}(1, \lim F) \cong \text{Cone}(1, F).$$

Def 3.2.3. $\forall F: J \rightarrow \text{Set}$ small diagram,
 $\lim F := \text{Cone}(1, F)$. and define $\lambda: \lim F \Rightarrow F$
 by $\lambda_i, \lambda_j: \lim F \rightarrow F_j$
 $\lambda_i: \lim F \xrightarrow{\cong} \mu_i: 1 \rightarrow F_j$.

(Notes that μ_j can be identified as element of F_j)

Thm 3.2.6. 3.2.3 is well-defined i.e., set is complete
 pf) We need to show that λ is a cone and limit cone.

Let $f: J \rightarrow K \in \mathcal{J} \Rightarrow$

induces $ff(\lambda_j(u)) = \lambda_k(Ff(u))$

Now, $u: I \Rightarrow F$. Hence $\mu_j \xrightarrow{1} u_k \Rightarrow Ff(\mu_j) = Ff(u_k)$

$$FJ \xrightarrow{2} FK = u_k$$

$$\Rightarrow ff(\lambda_j(u)) = ff(u_j) = u_k = \lambda_k(u).$$

So the triangle commute. $\Rightarrow \lambda$ is a cone.

To see λ is universal cone, let $\{\mathbb{x}: X \Rightarrow F$ be a cone. For each $x \in X$, $\mathbb{x}_x: I \Rightarrow F$ in the sense that $(\mathbb{x}_x)_j := \mathbb{x}|_{\{x\}}$.

Hence let $r: X \rightarrow \lim F$. This induces

$$x \mapsto \mathbb{x}_x.$$

s.t. $\lambda_j(r(x)) = \lambda_j(\mathbb{x}_x)$
 $= (\mathbb{x}_x)_j = \mathbb{x}_x|_{\{x\}} = \mathbb{x}(x)$

Moreover, any map $X \rightarrow \lim F$ should send x to \mathbb{x}_x , otherwise the above diagram doesn't commute.

$\Rightarrow r$ is unique. $\Rightarrow \lim F$ is well-defined.

Ex 3.2.7). Let $F: \mathcal{J} \rightarrow \mathcal{C}$ with \mathcal{J} is discrete.
 Then $\text{Im } F = \prod_{j \in \mathcal{J}} A_j$. By thm 3.2.6, $\text{Im } F \cong \text{Cone}(1, F)$

Notes that $\lambda: 1 \Rightarrow F$ consists of $\lambda_j: 1 \rightarrow A_j$, i.e.,
 elements of A_j , for each $j \in \mathcal{J}$.

$\Rightarrow \prod_{j \in \mathcal{J}} A_j := \{(\alpha_j \in A_j)_{j \in \mathcal{J}}\}$, a Cartesian product.

Ex 3.2.8). Terminal object of Set = Cone with sum 1
 over empty diagram.

$= 1$.

Ex 3.2.9.) $F: (\cdot \xrightarrow{\exists} \cdot)$ $\rightarrow \mathcal{C}$: equalizer.

$\text{Cone}(1, F) := \{ \lambda: 1 \Rightarrow F \}$.

$\begin{array}{ccc} \lambda_j & \xrightarrow{1} & \\ \downarrow & \nearrow \lambda_k & \\ F_j & \xrightarrow{f} & F_k \\ \downarrow g & & \end{array} \Rightarrow$ So identify λ_j as
 an element of F_j ,

$\text{Cone}(1, F) = \{ x \in F_j : f_j x = g x \}$

Ex 3.2.10. $F: \mathbf{W}^{\text{op}} \rightarrow \text{Set}$. If $\lambda \in \text{Cone}(1, F)$,

$\Rightarrow \begin{array}{ccccc} 1 & \xrightarrow{1} & x_1 & \xrightarrow{x_0} & \\ \downarrow & \nearrow \lambda_i & & & \\ \dots & \rightarrow & F_3 & \xrightarrow{f_{23}} & F_2 \xrightarrow{f_{12}} F_1 \xrightarrow{f_{10}} F_0 \end{array}$

By identifying each x_i as element of F_i ,

$\text{Cone}(1, F) = \{ (x_n) \in \prod_{n \in \mathbb{N}} F_n \mid f_{n,m}(x_n) = x_{n-1} \}$

Ex. 3.2.11. For pull back, $\lambda \in \text{Cone}(I; F)$

implies $\begin{array}{ccc} I & \xrightarrow{\lambda_1} & C \\ \downarrow \lambda_2 & \swarrow g & \downarrow g \\ B & \xrightarrow{f} & A \end{array}$ so by identifying λ_1 and λ_2 as elements of C and B resp.,

$$\text{Cone}(I, F) =: B \underset{A}{\times} C := \{(b, c) \in B \times C : f(b) = g(c)\}$$

Ex 3.2.12. $X: BG \rightarrow \text{Set}$. Then $\lambda \in \text{Cone}(I, X)$

satisfy

$$\begin{array}{ccc} I & & \\ \lambda \swarrow & \searrow \lambda & \text{By identifying } \lambda \in X, \\ \downarrow & \downarrow & \\ X & \xrightarrow{g_X} & X \end{array} \quad g\lambda = \lambda \quad \therefore \lambda \text{ is fixed by } g$$

for all $\lambda \in \text{Cone}(I, X)$. $\Rightarrow \text{Im } F = \text{Cone}(I, X) = X^G$
 Set of G -fixed pt of X .

Rmk.) Let $\exists \longrightarrow B \times C \xrightarrow{\begin{array}{c} (b, c) \mapsto f(b) \\ (b, c) \mapsto g(c) \end{array}} A$. Then, its equalizer E

$$E = \{(x, y) \in B \times C \mid f(x) = g(y)\} = B \underset{A}{\times} C.$$

This is not a coincidence.

Thm. 3.2.13. If $F: J \rightarrow \text{Set}$ small diagram,

$$\lim_J F \longrightarrow \prod_{J \in \text{obj}} F_J \xrightarrow{\begin{array}{c} \hookrightarrow \\ \downarrow \text{forget } J \end{array}} \prod_{J \in \text{obj}} F(\text{cod } f)$$

is an equalizer diagram, i.e. $\lim_J F$ is an equilizer of c and d .

pf) $\lim_{\leftarrow} F = \text{Cone}(I, F)$. Pick $\lambda : I \rightarrow F$.

Then for any $f \in \text{Mor } I$,
 $\lambda_{\text{dom } f} / \lambda_{\text{codom } f}$
 $F(\lambda_{\text{dom } f}) \xrightarrow{Ff} F(\lambda_{\text{codom } f})$

Hence we may identify λ_j be element of F_j .

$\Rightarrow \lambda$ is identified by $(\lambda_j) \in \prod_{j \in \text{Ob } I} F_j$

Now, define.

$c: (\lambda_j)_{j \in \text{Ob } I} \longmapsto (\lambda_{\text{codom } f})_f \in \prod_{f \in \text{Mor } I} F(\text{codom } f)$

by copying existing components of $(\lambda_j)_{j \in \text{Ob } I}$.

$d: (\lambda_j)_{j \in \text{Ob } I} \longmapsto (Ff(\lambda_{\text{dom } f}))_f \in \prod_{f \in \text{Mor } I} F(\text{dom } f)$

Then, If $c(\lambda) = d(\lambda)$, then for each $f \in \text{Mor } I$,

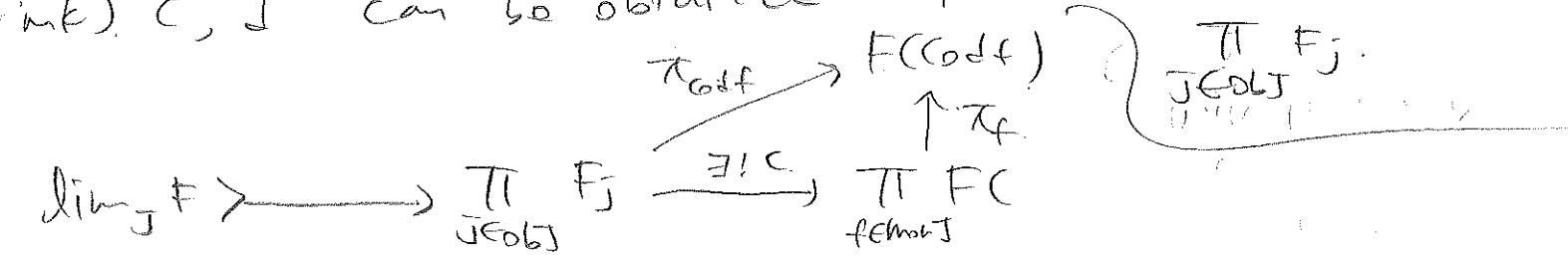
$Ff(\lambda_{\text{dom } f}) = \lambda_{\text{codom } f} \Rightarrow \lambda \in \text{Cone}(I, F)$.

\Rightarrow Equalizer $\subseteq \lim_{\leftarrow} F$.

And every element in $\lim_{\leftarrow} F$ satisfy $c(\lambda) = d(\lambda)$.

$\Rightarrow \lim_{\leftarrow} F$ is Equalizer of c and d . \square

Q.E.D. c, d can be obtained by universal property of \lim_{\leftarrow} .



$$\lim_{\leftarrow} F \longrightarrow \prod_{j \in \text{Ob}} T_j \xrightarrow{\exists! \downarrow} F(\text{Codf})$$

\prod_{Ob}

$\int \pi_f.$

$$F(\text{domf}) \xrightarrow{\text{ff}} F(\text{codf})$$

Then, we can use the same argument showing that \amalg_F is equalizer:

Ex. 3.2.16 (Idempotent) $e: A \rightarrow A$ is idempotent if $e \cdot e = e$. So $F = (G \circ Qe) \rightarrow \text{Set}$, s.t. $e \cdot e = e$.

If $\lambda: 1 \Rightarrow F$, then,

$$A \xrightarrow{e} A \xrightarrow{e} A$$

$$e\lambda = \lambda \text{ and } e \cdot e\lambda = e\lambda.$$

Thus $\amalg_F = \{a \in A : ea = a\}$
 $= A^e : \text{fixed pt of idempotent.}$

Is the Exercise 3.2.ii., $A^e \xrightarrow{s} A \xrightarrow[e]{1} A$.

as an equalizer. Then, $\exists!$ r.s.t below diagram commutes.

$$\begin{array}{ccc} A & & \\ r \downarrow \swarrow e & & \\ A & \xrightarrow[e]{1} & A \end{array} \Rightarrow sr = e \text{ i.e. } sr \text{ splits } e.$$

Now $srs = es = s$. Thus

$$\begin{array}{ccc}
 A^e & \xrightarrow{s} & \\
 \downarrow r_s & \downarrow \text{id}_{A^e} & \\
 A^e & \xrightarrow{s} & A \xrightarrow{e} A
 \end{array}
 \Rightarrow rs = \text{id}_{A^e}$$

by uniqueness

Also notes that $r|_B$ surjective: $\forall a \in A^e$,
 $\exists a' \in A$ s.t. $s(a') = a$. and s is canonical.

Injection, thus $r(a) = a$.

Conversely, $B \xrightarrow{\cong} A \xrightarrow{r} B$ with $rs = \text{id}_B$
 gives an idempotent sr. on A .

($sr \circ sr = s \circ r = sr$) which is split by B .

Ex. 32) Take a pullback P of below diagram.

$$\begin{array}{ccc}
 P & \dashrightarrow & \text{mor}(C) \\
 \downarrow & \downarrow \text{dom} & \text{Then } P := \underset{\text{Ob}(C)}{\text{mor}(C) \times \text{mor}(C)} \\
 \text{mor}(C) & \longrightarrow & \text{Ob}(C)
 \end{array}$$

$\doteq \{(g, f) \in \text{mor}(C) \times \text{mor}(C) \mid \text{dom}g = \text{cod}(f)\}$

Thus, If $(g, f) \in P$, then $g \circ f$ is possible composition.

And we have $m: P \rightarrow \text{mor}(C)$ it is well-defined
 $(g, f) \mapsto gf$. function.

Now we have

$$\begin{array}{ccc}
 P & \xrightarrow{m} & \text{mor}(C) \xleftarrow{\text{dom}} \text{Ob}(C) \\
 & & \xrightarrow{\text{cod}}
 \end{array}$$

Then $m(gf)$ has dom don't
 cod cods.

We claim that $B \times_A B \xrightarrow{m} B \xrightleftharpoons[f]{h} A$.

with ① $f(m(b_1, b_2)) = f(b_2)$
 i.e. ② $h(m(b_1, b_2)) = h(b_1)$ has a category structure
 ③ $m(m(b_3, m(b_2, b_1)), b_1) = m(m(b_3, b_2), b_1)$ ④ $m(g(a), b) = b = m(a, g(b))$
 ⑤ $m(b, g(a)) = b$

⑥ Object = A
 ⑦ Morphism = B i.e. by definition $b \in B$
 as $f(b) \rightarrow h(b)$.

⑧ Composition, m. Suppose (b_1, b_2, b_3) are
 composable, i.e., $(b_2, b_1), (b_3, b_2) \in P$.

Then, $f \circ m(b_2, b_1) = f(b_1)$ $f \circ m(b_3, b_2) = f(b_2)$
 $h \circ m(b_2, b_1) = h(b_2)$ $h \circ m(b_3, b_2) = h(b_3)$

Hence, $f(b_3) = h(b_2) = h \circ m(b_2, b_1)$

implies $(b_3, m(b_2, b_1)) \in B \times_A B$.

Also, $f \circ m(b_3, b_2) = f(b_2) = h(b_1) \Rightarrow (m(b_3, b_2), b_1) \in B \times_A B$.

Then,

$$m(m(b_3, m(b_2, b_1)), b_1) = m(m(b_3, b_2), b_1)$$

shows associativity.

Also, if $g(a)$ has a role of identity morphism

thus, the diagram with ①, ②, ③, ④, ⑤
 gives a small category structure \mathcal{J} .

with Obj $\mathcal{J} = A$

Mor $\mathcal{J} = B$.

Ex 3.2(iii) Change $\prod F(\text{obj } f)$, with $\prod_{\text{fHom}(J)}$, with $F(\text{obj } f)$
 $\prod_{\text{fHom}(J)}$

in the diagram p. 88. In this case, $\prod_{\text{fHom}(J)}$ is
 not a discrete category then we still have
 C and D which is the same as above.
 So equalizer must be same.

In case of J is discrete, then $C = D$.
 thus, $\prod_{\text{fHom}(J)} F = \prod_{j \in J} F_j$, a product, as
 be expected.

Ex 3.2(iii) If $\begin{array}{ccc} a & \xrightarrow{\quad} & C \\ f \downarrow & \Downarrow g & \\ b & \xrightarrow{\quad} & d \end{array}$, then, $gf = gh$.

$\Rightarrow P \rightarrow \text{Hom}(a, c)$ let P be a pull back,
 $\downarrow \quad \downarrow g^*$ then,
 $\text{Hom}(b, d) \xrightarrow[\text{p}_x]{\quad} \text{Hom}(a, d)$ $P = \frac{\text{Hom}(a, c) \times \text{Hom}(b, d)}{\text{Hom}(a, d)}$
 set $(h, l) \in \text{Hom}(a, c) \times \text{Hom}(b, d)$
 $g^*(h) = f_x l$.
 i.e. $gh = lf$ }
 thus, $Sq(f, g) := \frac{\text{Hom}(a, c) \times \text{Hom}(b, d)}{\text{Hom}(a, d)}$.

Ex 32 ir Let J^S be a category set.

$\text{Ob } J^S = \text{Mor } J$. $\text{Mor } J^S$: for each $f: x \rightarrow y \in J$

$$\begin{aligned} I_x &\rightarrow f & \in \text{Mor } J^S \\ I_y &\rightarrow f \end{aligned}$$

Let $\alpha \in \text{Hom}(F, G)$. Then, for any $f: x \rightarrow y \in J$.

$$\begin{array}{ccc} Fx & \xrightarrow{\alpha_x} & Gy \\ Ff \downarrow & \Downarrow & \downarrow Gf \\ Fy & \xrightarrow{\alpha_y} & Gy \end{array}$$

Now construct $H: J^S \rightarrow \text{Set}$ i.e.,

$$\begin{array}{ccc} I_x & \longmapsto & \text{Hom}(Fx, Gy) \\ \downarrow & \longmapsto & \downarrow (Gf)^* \\ f: x \rightarrow y & \longmapsto & \text{Hom}(Fx, Gy) \\ \uparrow & & \uparrow (Ff)^* \\ I_y & \longmapsto & \text{Hom}(Fy, Gy) \end{array}$$

If $\lambda \in \text{Core}(A, H)$, then,

$$\begin{array}{ccc} 1 & & \lambda_x \\ & \swarrow \lambda_x & \searrow \lambda_y \\ 2 & \xrightarrow{\lambda_f} & 2 \end{array}$$

$$\begin{array}{ccc} \text{Hom}(Fx, Gy) & \longrightarrow & \text{Hom}(Fx, Gy) \leftarrow \text{Hom}(Fy, Gy) \\ (\text{Gf})^* & & (\text{Ff})^* \end{array}$$

By identification, this shows that

$$(\text{Gf})(\lambda_x) : Fx \rightarrow Gy = \lambda_f = \text{Ff}^*(\lambda_y)$$

i.e.

$$\begin{array}{ccc} x & \xrightarrow{\text{Ff}} & y \\ \lambda_x \downarrow 1 & \Downarrow & \downarrow \lambda_y \\ 2 & \xrightarrow{\text{Gf}} & 2 \end{array}$$

Hence we identify λ as
 $\lambda \in \text{Hom}(F, G)$

Conversely, if $\lambda \in \text{Hom}(F, G)$

then define $\lambda_f = Gf \circ \lambda_x = \lambda_y \circ Ff$

for any $f: x \rightarrow y \in J$ nontrivial morphism

then, by identities λ_x and λ_y as a map

$1 \rightarrow \text{Hom}(Fx, Gx)$ and $1 \rightarrow \text{Hom}(Fy, Gy)$

we have a diagram

$$\begin{array}{ccccc} & & 1 & & \\ & \swarrow \lambda_x & \downarrow f & \searrow \lambda_y & \\ \sim & \xrightarrow{(Gf)^*} & \sim & \xleftarrow{(FF)_x} & \sim \end{array}$$

which commutes every triangle.

Hence such λ is a cone.

So, $\text{Hom}(F, G) \cong \text{Lim} H$ as a set.

Ex. 3, 2 U) It suffices to show that equalizer
is $\text{Hom}(F, G)$. Let E be equalizer of
diagram. Then,

$$E = \left\{ (\lambda_j)_{j \in \text{Ob}J} \mid \forall i \in \text{Ob}I, \quad \forall j \in \text{Ob}J, \quad \lambda_i = \lambda_j \right\}.$$

$$\Leftarrow. (Ff^*) \circ \pi_j, (\lambda_j)_{j \in \text{Ob}J}$$

$$= Gf_x \circ \pi_j (\lambda_j)_{j \in \text{Ob}J}$$

i.e., $\lambda_j \circ Ff = Gf \circ \lambda_j$ i.e. $\begin{array}{ccc} Fj & \xrightarrow{\lambda_j} & Gj \\ \downarrow f & \lrcorner & \downarrow \\ Fj' & \xrightarrow{\lambda'_j} & Gj' \end{array}$

Thus $E \subseteq \text{Hom}(F, G)$. Conversely, any $\lambda \in \text{Hom}(F, G)$, as identified by $(\lambda_j)_{j \in \text{Ob}C} \in \prod_{j \in \text{Ob}C} ((F_j, G_j))$ satisfy the diagram $\Rightarrow \lambda \in E \Rightarrow E = \text{Hom}(F, G)$. Ex 32. vi).

$\lim F \cong \text{Cone}(A, F)$. Let $\lambda \in \lim F$. Then, $\lambda_c: A \rightarrow F_c$ can be identified as an element of F_c .

Thus, for each $\lambda \in \lim F$,

let $G: \text{Cone}(A, F) \rightarrow \text{Section}(\Pi)$.

$$\begin{aligned} \lambda &\longmapsto F_\lambda: C \rightarrow SF \\ &\quad c \mapsto (c, \lambda_c). \end{aligned}$$

Conversely, let $H \in \text{Section}(\Pi)$.

Then, $\forall c \in \text{Ob}C, H_c = (c, \eta_c)$

for some $\eta_c \in F_c$. And for all $\exists c \xrightarrow{f} d \in C$,

$$Hf: (c, \eta_c) \rightarrow (d, \eta_d)$$

st. $f: c \rightarrow d$ and $Ff(\eta_c) = \eta_d$.

Hence,

$$\begin{array}{ccc} \eta_c & \xrightarrow{f} & \eta_d \\ \downarrow \eta_c & & \downarrow \eta_d \\ F_c & \xrightarrow{Ff} & F_d \end{array} \Rightarrow \begin{cases} \eta_c : (c \in \text{Ob}C) \\ \text{can be identified as} \end{cases}$$

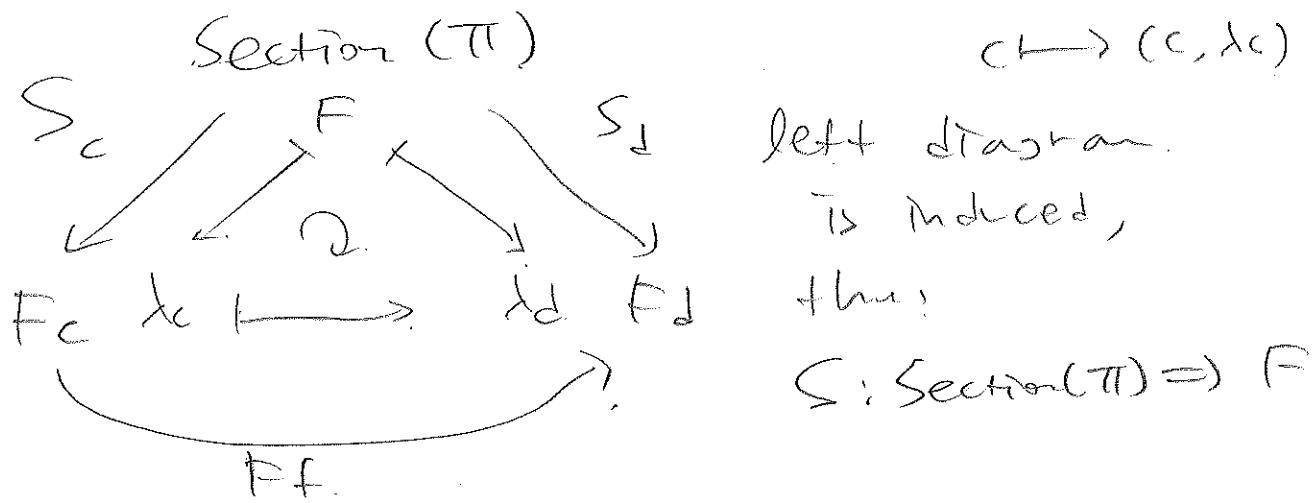
a cone $A \Rightarrow F$.

And this is inverse of G .

Recall limit cone λ is $\lambda : \lim_{\leftarrow} F$ s.t.
 $\forall M : X \Rightarrow F$ "a cone," $\exists! f : X \rightarrow \lim_{\leftarrow} F$ ESet.
 s.t. $\lambda \circ f = M$,
 Then, $\forall x \in X$, $M(x) : I \Rightarrow F$. Thus,
 $\lambda \circ f(x) : I \Rightarrow F$
 Since $f : f|_{\{x\}}$ is an iso, we know that
 $M(x) \cong \lambda(f(x))$ as a natural transformation.

Thus, every cone $M : X \Rightarrow F$ is just a set of
 $M(x) : I \Rightarrow F$ s.t. $M(x)$ is a section.

Thus limit cone is just a cone S induced by
 Section (II) s.t. for $f : c \rightarrow d \in \text{Mor } C$,
 when $F : C \rightarrow SF$
 $c \mapsto (c, \lambda_c)$



Can be regarded as a limit. For any $M : X \Rightarrow F$,
 $M(x) : I \Rightarrow F$ so $M(x) \in \text{Section}(II)$. Thus, the map
 $X \rightarrow \text{Section}(II)$ by $x \mapsto M(x)$ induces desired diagram.

3.3. Preservation, reflection, creation.

Def 3.3.1. For any class \mathcal{L} of diagrams $K: J \rightarrow C$ and $F: C \rightarrow D$

① F preserves those \mathcal{L} -limits:

$\forall K: J \rightarrow C$ in \mathcal{L} , and $\ell: \text{lim} K \Rightarrow K$ is a limit cone $\Rightarrow FK$ is a limit cone of FK , i.e. $F(\text{lim } K) = \text{lim } FK$ and $FK: \text{lim } K \Rightarrow FK$.

② F reflects "

$\forall K: J \rightarrow C$ in C , and FK is a limit cone of $FK \Rightarrow K$ is a limit cone, i.e. $X = \text{lim } K$ and K has universal property

③ F creates "

If $\text{lim } FK$ exists, $\exists K: \text{lim } FK \Rightarrow FK$ a limit cone s.t. it can be lifted to a limit cone over K and F reflects the limits.

(For colimit, dualize these.)

In other words, if $K: \text{lim } K \Rightarrow K$ is a limit cone, $K': X \Rightarrow K$ is a cone (thus $F(\text{lim } K) = \text{lim } FK$)

① preserve: " $F(K)$ is a limit cone $\Rightarrow K$ is a limit cone"

② reflects: " $F(K')$ is a limit cone $\Rightarrow K'$ is a limit cone"

③ creates: FK has a limit cone $\Rightarrow K$ has a limit cone and F reflects.

Prop 3.3.3. (1) $F: \mathcal{D} \rightarrow \mathcal{C}$ creates \mathcal{D} 's limit. (2) \mathcal{D} has a limit
 $\Rightarrow \mathcal{C}$ has a limit and F preserves it.

(1) Let $k: J \rightarrow C$ be in the class.

$\mu: J \rightarrow FK$ is a limit cone.

Since F creates these limits, $\exists \lambda: C \rightarrow k$
a limit cone s.t. $F\lambda$ is a limit cone.
By Ex. 3.1.N, $F\lambda \cong \mu$.

To see F preserves them, let $\lambda': C' \rightarrow k$ be
another limit cone. By prop 3.1.7, $\lambda' \cong \lambda$
 $\Rightarrow F\lambda' \cong F\lambda \cong \mu$. $\Rightarrow F\lambda'$ is a limit cone.

Def 3.3.4 (Sheaf) $\mathcal{O}(X)$: poset of open subsets of X
ordered by inclusion. Let I be a discrete category.

$\{U_i\}_{i \in I}$ cover U . If U_i 's are mage of I .

and $\text{colim}_I (\{U_i\}_{i \in I} \xleftarrow{U_i \hookrightarrow U_i \cap U_j \rightarrow U_j}) = U$.

Notes that colimit in this diagram is an open set
 U' with $U_i \rightarrow U'$ and $U_i \cap U_j \rightarrow U'$.

$\Rightarrow U' = \bigcup_{i \in I} U_i$. So this definition is consistent with
usual definition of open coverings.

Let $F: \mathcal{O}(X)^{\text{op}} \rightarrow \text{Set}$. By def in 1.3.7(U)

since $\mathcal{O}(X)$ is small). F is presheaf. Sheaf is a presheaf
preserving these colimits \wedge pairwise and $H^{\text{index } I}$
covering U .

Now let I' be a small category s.t. $I' \cong I$. s.t.

① Indexes $\{U_i\}_{i \in I}$ a cover of U .

② I' has objects (i, j) corresponds to

$U_i \cap U_j$, $i, j \in I$.

Let $k: (I')^{\text{op}} \rightarrow \mathcal{O}(X)^{\text{op}}$ and $F: \mathcal{O}(X)^{\text{op}} \rightarrow \text{Set}$

as sheaf. Notes that any morphism $U_i \rightarrow U_j$ actually induces $U_i \cap U_j \hookrightarrow U_j$.

Then, Thm 3.213 yields

$$\lim_{\substack{\longrightarrow \\ \text{op}}} Fk \longrightarrow \prod_{(i,j) \in I'} Fk(i,j) \xrightarrow{c} \prod_{\substack{\text{op.} \\ f \in \text{mor}(I')}} F(\text{cod } f)$$

Let f be an arbitrary morphism $\text{mor}(I')^{\text{op}}$.

We may regards it as $U_i \cap U_j \hookrightarrow U_i$.

by the above argument,

Then, c is just sends $F(U_i \cap U_j)$ component

the middle to the same component in the last,

by seeing p. 88. diagram

Now, d is a little different. $\prod_{\text{op.}} = \prod_{f \in \text{mor}(I')}$.

so if $f: U_i \cap U_j \hookrightarrow U_i$ then d sends

$F(U_i)$ in the middle to $F(U_i \cap U_j)$ in the last

i.e. $F(U_i \cap U_j) \hookrightarrow U_i$.

Now Let $\lambda \in \prod_{(i,j) \in I'} Fk(i,j)$ If $c(\lambda) = d(\lambda)$, then,

$Ff(\lambda|_{U_i}) = \lambda|_{U_i \cap U_j} = Ff(\lambda|_{U_j})$ by applying f twice and care.

This shows that if $\text{FF}(\lambda|_{U_i}) = \text{FF}(\lambda|_{U_j})$ for all i, j , then $\lambda \in \lim_{I'} F_k$.

Thus, instead of c , we may use this condition.

$$\lim_{I'} F_k \rightarrow \prod_{(i,j) \in I'} F_k(i,j) \xrightarrow{\quad} \prod_{i \in I'} \text{F}(\text{cod } f^{\circ i})^{\text{op}} \\ \text{FF}(\lambda|_{U_i})_{\text{on } i} \xrightarrow{\quad} \text{FF}(\lambda|_{U_j})_{\text{on } j}^{\text{op}} \text{ for } (i,j) \in I'$$

Since those two maps are only depends on $F_k(i,j)$ (with $i=j$), we may delete component with $i \neq j$, so.

$$\lim_{I'} F_k \rightarrow \prod_{i \in I'} F_k(i) \xrightarrow{\quad} \mathbb{P} //$$

Also, since all those two map only come of $F_k(i,j)$ with $i \neq j$. So we may delete $F(U_i)$ parts in the last

$$\rightarrow \lim_{I'} F_k \rightarrow \prod_{i \in I} F_k(i) \xrightarrow{\quad} \prod_{i \in I} \text{F}(U_i \cap U_j)^{\text{op}} \\ f: U_i \cap U_j \rightarrow U_i \\ f: U_i \cap U_j \rightarrow U_j$$

then, by change of notation, we may get.

$$\lim_{I'} F \rightarrow \prod_{i \in I} \text{F}(U_i) \xrightarrow{\quad} \prod_{i, j \in I} \text{F}(U_i \cap U_j)$$

as desired.

Lemma 3.3.5 $F: C \rightarrow D$ fully faithful.

$\Rightarrow F$ reflects limits and colimits which are in C domain.

Pf) Let $\mu: c \Rightarrow k$ be a core of $k \rightarrow c$.

S.t. $F\mu: Fc \Rightarrow Fk$ is a limit one.

Let $\mu': c' \Rightarrow k$, WTS. $\exists! f: c' \rightarrow c$ s.t.

$$\begin{array}{ccc} & c' & \\ \mu'_i \swarrow & \downarrow \exists! f & \searrow \mu_j \\ & c & \\ \mu_i \swarrow & \downarrow \exists! f & \searrow \mu_j \\ k_i & \longrightarrow & k_j \end{array}$$

Left diagram commutes.

Now note that $F\mu'$ is still a core; this comes from functoriality of F .

Hence $\exists! \bar{f}: Fc' \rightarrow Fc$ s.t. $F\mu' = \bar{f} \circ F\mu$.

By fully faithfulness, $\exists f: c' \rightarrow c$ s.t. $f = \bar{f}$.

Then, $F\mu'_i = f(\mu_i \circ f) \Rightarrow \mu'_i = \mu_i \circ f$

Since fully faithful implies $(C(c, k_i)) \xrightarrow{F(-)} D(Fc, Fk_i)$

is bijection.

Thus, the above diagram commutes. So F reflects limit.

For the colimit, apply this on contravariant functor and get reflection of colimit by duality.

Lemma 3.3.6. Equivalence of Categories
 preserves, reflects, and creates any limit
 and colimit that are present in either this
 domain or codomain.

(f) Let $F: C \rightarrow D$ $G: D \rightarrow C$ be equivalence
 of categories. Then F, G are fully faithful and
 essentially surjective.
 It suffices to show that each F and G creates
 limit in D or C . Suppose we showed it
 ① reflectivity: due by Lemma 3.3.5

② preserves limit/colimit: If λ is a limit cone in C .

then, G creates a limit cone λ' in D .

St. $G\lambda' \cong \lambda$. Thus, $G\lambda'$ is a limit cone.

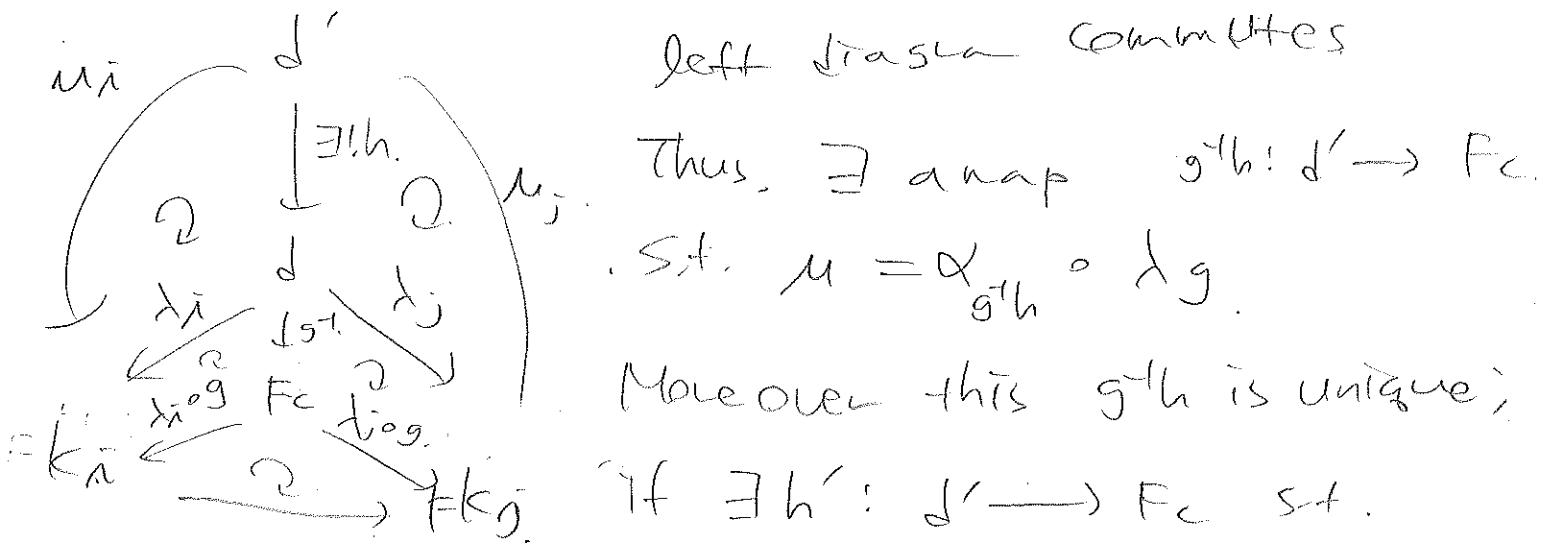
Hence $F(\lambda) \cong FG(\lambda') \cong \lambda'$ implies F
 preserves a limit.

Thus, all we have to do is show that F creates
 limit/colimit.

Let λ be a limit cone in D , i.e. $\lambda: d \Rightarrow Fk$.

By essential surjectivity, $\exists c \in C$ s.t. $Fc \xrightarrow{g} d$.

Now we claim $\lambda_g: Fc \Rightarrow Fk$ is also a limit cone to
 see this, let $\mu: d' \Rightarrow Fk$ be a cone. Then,



Left diagram commutes

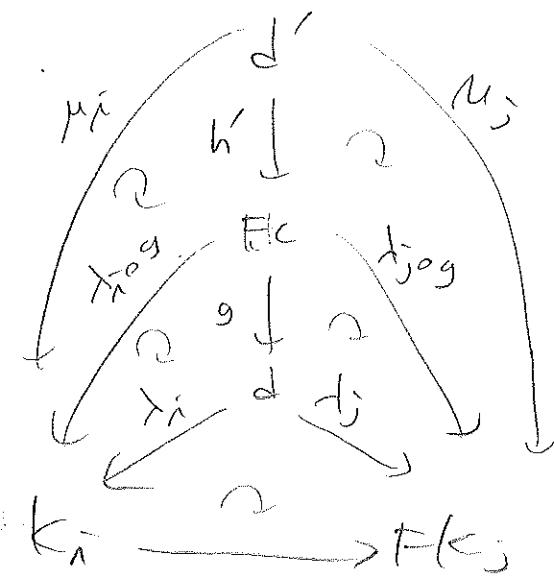
Thus, \exists a map $g'h: d' \rightarrow F_k$.

$$\text{S.t. } \mu = \alpha_{g'h} \circ \lambda_g.$$

Moreover this $g'h$ is unique;

If $\exists h': d' \rightarrow F_k$ s.t.

the above one commutes, then left diagram



commutes, thus, universal property of $d: d \Rightarrow F_k$ shows that

$$gh' = h \quad \dots \quad .$$

$$\Rightarrow h' = g^{-1}h. \quad \dots \quad .$$

Now, full-faithfulness of F shows that

any $\lambda_i^o g$ comes from $\overline{\lambda_i^o g} \in C(C, k_i)$

via F . Thus, $\exists \bar{x}: C \Rightarrow L$ s.t. $F\bar{x} = \lambda_i^o g$.

Since F reflects a limit, \bar{x} is a limit cone.

Hence, (since F is already reflects the limit)

F creates a limit. Since F was chosen

arbitrarily, both F and G creates limits

By duality, \Rightarrow colimits.

$K \in \pi(K)$ and $\lambda: K \Rightarrow K$ is a cone induced by a diagram $K: J \rightarrow \mathcal{Y}_C$, with sumit C .

WTS if $K: J \xrightarrow{(K, k)} \mathcal{Y}_C \xrightarrow{\pi} C$

admits a limit cone in C , then this cone lifts to define a unique limit cone over (K, k) in \mathcal{Y}_C .

Let $\lambda: l \Rightarrow K$ be a limit cone in C . Also we have $k: c \Rightarrow K$. Hence $\exists! t: c \rightarrow l$ st. $K = \lambda \circ \alpha_t$.

Now, we can identify $\lambda_i: l \rightarrow k_i$ as a morphism from $t: c \rightarrow l$ to $(K, k)_i: c \rightarrow k_i$ in \mathcal{Y}_C . Thus, below left, Δ commutes since right triangle commutes by G-1 ii .

$$\begin{array}{ccc} & (l, t: c \rightarrow l) & \\ \lambda_i & \swarrow \quad \searrow & \\ (K, k)_i & \xrightarrow{\quad} & (K, k)_j \\ & (K, k)(\lambda \rightarrow j) & \end{array}$$

$$\begin{array}{ccc} & (c & (K, k)_j) \\ & \downarrow t & \\ (k_i & \xrightarrow{\quad} & k_j \\ & (K, k)(\lambda \rightarrow j) & \end{array}$$

So, λ is a cone $t \Rightarrow (K, k)$.

To see it is a limit cone, suppose that

$\exists M: r \Rightarrow (K, k)$ in \mathcal{Y}_C . Then, by letting $r: c \rightarrow r$,

Def) 3.3.7. $F! : \mathcal{C} \rightarrow \mathcal{D}$ strictly creates \lim_{inj} .

If $\forall K : J \rightarrow C$ in the class and λ limit core over FK ,

① $\exists !$ lift $^{\lambda}$ of λ over K in C as a core

② $\mathcal{A} \cdot \mathcal{D}$ limit core.

Def) J : category is connected if $J \neq \emptyset$

and in the underlying directed graph, every pair of objects is connected by a finite zig-zag of morphism.

A diagram $K : J \rightarrow C$ is connected $\Leftrightarrow J$ is connected.

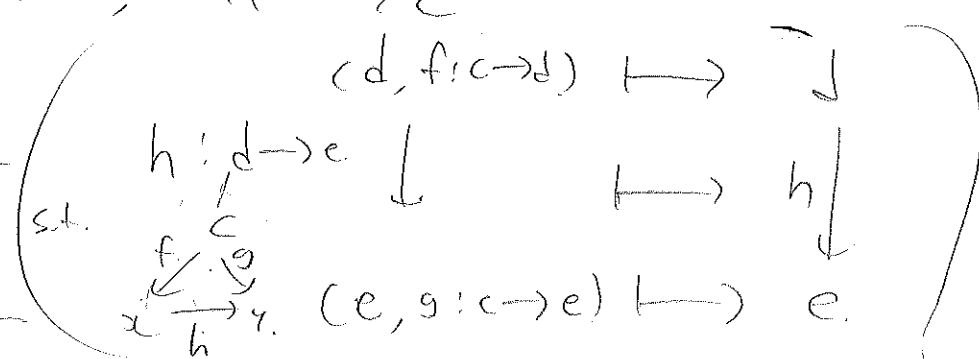
Prop 3.3.8 $\forall c \in \text{Ob } C$, $\pi_c : \mathcal{G}_c \rightarrow C$

strictly creates

(1) all limits

(2) all connected

limits



(3) Let $K : J \rightarrow \mathcal{G}_c$ be a diagram. Then

$\forall i \rightarrow j \in \text{Mor } J$, $K(i \rightarrow j)$ is a morphism

$h : \text{cod}(K_i) \rightarrow \text{cod}(K_j)$ s.t.

always



Hence K induces a cone $K : c \Rightarrow (J \rightarrow \mathcal{G}_c \Rightarrow c)$ in C .

thus, we use notation $K = (K, h) : J \rightarrow \mathcal{G}_c$ s.t.

we have

$$\begin{array}{ccc}
 (r', r: C \rightarrow r') & & \\
 \downarrow m_i \quad \downarrow n_j & \Rightarrow & \\
 (k, k)_i \xrightarrow{(k, k)(i \rightarrow j)} (k, k)_j & & \\
 & &
 \end{array}$$

Since λ is a limit core, $\exists t: r' \rightarrow l$ is s.t.

$$\begin{array}{ccc}
 \begin{array}{c} \text{left diagram commutes, hence} \\ \text{rot}' = t, \text{ by uniqueness.} \\ \text{of } t. \end{array} & & \\
 \begin{array}{c} \text{Figure 1} \\ \text{Thus, } (t, t: a, !\text{morphism in } \mathcal{C}) \end{array} & &
 \end{array}$$

Hence, $(r', r: C \rightarrow r')$

left diagram

$$\begin{array}{ccc}
 \begin{array}{c} \text{commutes} \\ \text{by Figure 1.} \end{array} & & \\
 \begin{array}{c} \text{Figure 1} \\ \text{Thus, } (t, t: C \rightarrow l, \\ \text{!morphism in } \mathcal{C}) \end{array} & &
 \end{array}$$

In case of colimit, let $u: k \Rightarrow p$ is a colimit core and J is connected. Still we have $k: C \Rightarrow p$, a cone.

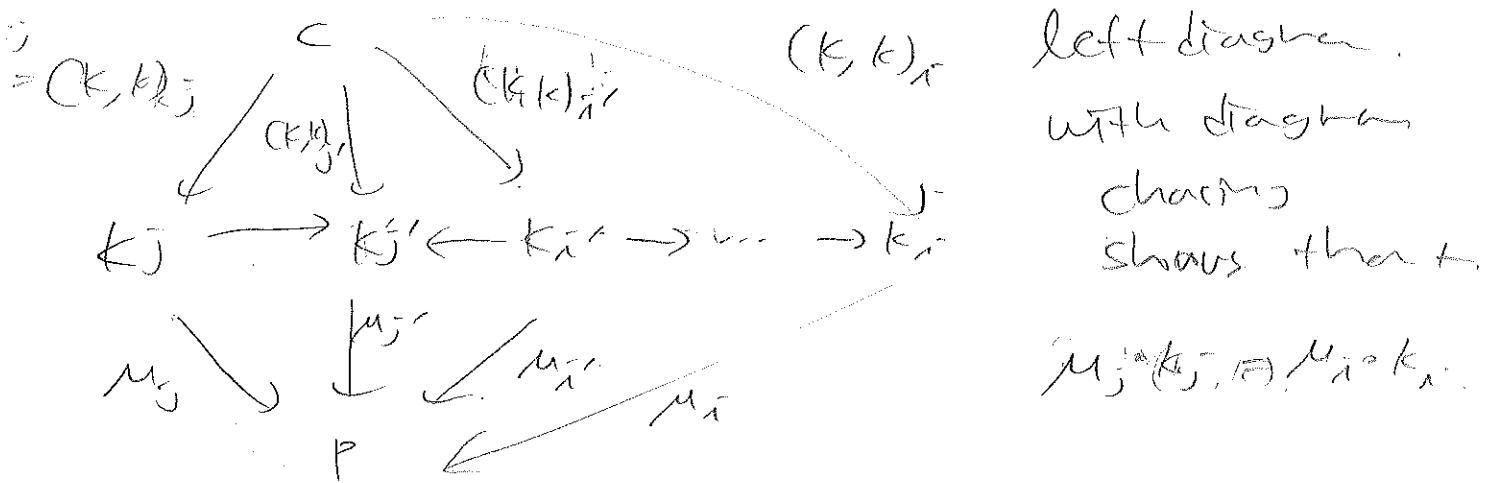
Then, get $n: b \Rightarrow (a, !\text{morphism in } \mathcal{C})$.

$$n_j: C \xrightarrow{k_j} k_j \xrightarrow{m_j} p$$

We claim that $\mathcal{C} \rightarrow \mathcal{P}$ is an object of \mathcal{C} .

It is independent of choice of $j \in J$. To see this

for any $i, i \in \text{ob } J$, \exists zig-zag morphism. Thus



left diagram.

With diagram

charts

shows that +

$\mu_j^i : (K_j, \alpha) \rightarrow (K_i, \alpha)$.

Now let $\phi : (J \rightarrow \mathcal{C}) \Rightarrow f$ be a cone in \mathcal{C} under $(l, f : C \rightarrow l)$. Then,

$$(l, f : C \rightarrow l)$$

$$\phi_i \nearrow \downarrow \nwarrow \phi_j \Rightarrow$$

$$(k_i : C \rightarrow K_i) \rightarrow (k_j : C \rightarrow K_j)$$

$$\begin{array}{ccc} \phi_i & \nearrow & l \\ & \uparrow f & \\ k_i & \nearrow & k_j \\ & \searrow & \\ & l & \end{array}$$

Now, $(l, f : C \rightarrow l)$

$$\begin{array}{ccc} \phi_i & \nearrow & l \\ & \uparrow u & \\ k_i & \nearrow & \nwarrow \phi_j \\ & \uparrow & \\ (P, h : C \rightarrow P) & & \\ \mu_j \nearrow & \nwarrow \mu_i & \\ & \searrow & \\ & P & \end{array}$$

$$(k_i : C \rightarrow K_i) \rightarrow (k_j : C \rightarrow K_j)$$

$$\begin{array}{ccc} \mu_j & \nearrow & l \\ & \uparrow & \\ P & \nearrow & \mu_i \\ & \searrow & \\ & l & \end{array}$$

Since u is a colimit cone,
 $\exists! u : P \rightarrow l$. So

$$u \circ \mu_i = \phi_i$$

So it suffices to show that

u is a morphism in \mathcal{C} .

Since $f = \phi_i \circ k_i$, $\begin{array}{ccc} p & \xrightarrow{u} & l \\ \downarrow & \nearrow f & \downarrow \\ n & \xrightarrow{\mu_i \circ k_i} & c \end{array}$

Thus $u \circ n = u \circ \mu_i \circ k_i = \phi_i \circ k_i = f$.

$\Rightarrow u$ is a morphism in \mathcal{C} .

Thus $\mu: \eta \Rightarrow (J \rightarrow \mathcal{C})$ is a colimit cone.

These colimit and limit construction is unique since it depends only on \otimes cone (a core) structure of C . $\Rightarrow \Pi$ strictly creates concrete colimits or limits.

Rank: If $k: J \rightarrow C^A$ is a diagram, then $\lim_J K$ is a functor from $A \rightarrow C$ s.t whose evaluation at a is a limit of $J \rightarrow C$ by $j \mapsto F_j(a)$. (This is prop 3.3.9)
 Prop 3.3.9.) Let A : small category. $\Pi: C^A \rightarrow C^{ob A}$ be a forgetful functor where $ob A$ is a discrete subcategory of A . Then Π strictly creates all limit and colimit that exist in C .

These limits are defined objective, meaning that $\forall a \in A$, $ev_a: C^A \rightarrow C$ preserves all limit and colimits existing in C .

pf) $C^{\text{ob}A}$ is isomorphic to $\prod_{\text{ob}A} C$; since any functor in $C^{\text{ob}A}$ is just choosing $\text{ob}A$ -indexed elements in $\text{Ob } C$.

By applying universal property of product in CAT, any diagram $k: J \rightarrow C^{\text{ob}A} \cong \prod_{\text{ob}A} C$ is just $\text{ob}A$ -indexed family of diagram $J \rightarrow C$.

Let $\lim k$ be a limit in $\prod_{\text{ob}A} C$, λ is a limit one.

Then, for any $\mu: (C_i)_{i \in \text{ob}A} \Rightarrow k$, with $i \rightarrow j \in J$

$$\begin{array}{ccc} (C_i)_{i \in \text{ob}A} & & \\ \downarrow \mu_i & \searrow \lambda_i & \\ & \lambda_j & \\ & \downarrow \mu_j & \\ (k_i) & \xrightarrow{k(i \rightarrow j)} & k_j \end{array}$$

$$\begin{array}{ccc} M_{it} & & M_{jt} \\ \downarrow & \downarrow \text{id}_t & \downarrow \\ (k_i)_t & \xrightarrow{(k(i \rightarrow j))_t} & (k_j)_t \end{array}$$

but in the product of category, every homomorphism can be decomposed into components indexed by $\text{ob}A$.

So for any index $t \in \text{ob}A$,

Hence, $\lim k$ is assertible of limit in each C .

also is its limit cone.

Thus, C admit $\lim \prod_{\text{ob}A} C$ (dually for colimit).

Also, $e_{\lambda}: C^{\text{ob}A} \rightarrow C$ can be viewed as a projection, thus it preserves limit. (dually colimit).

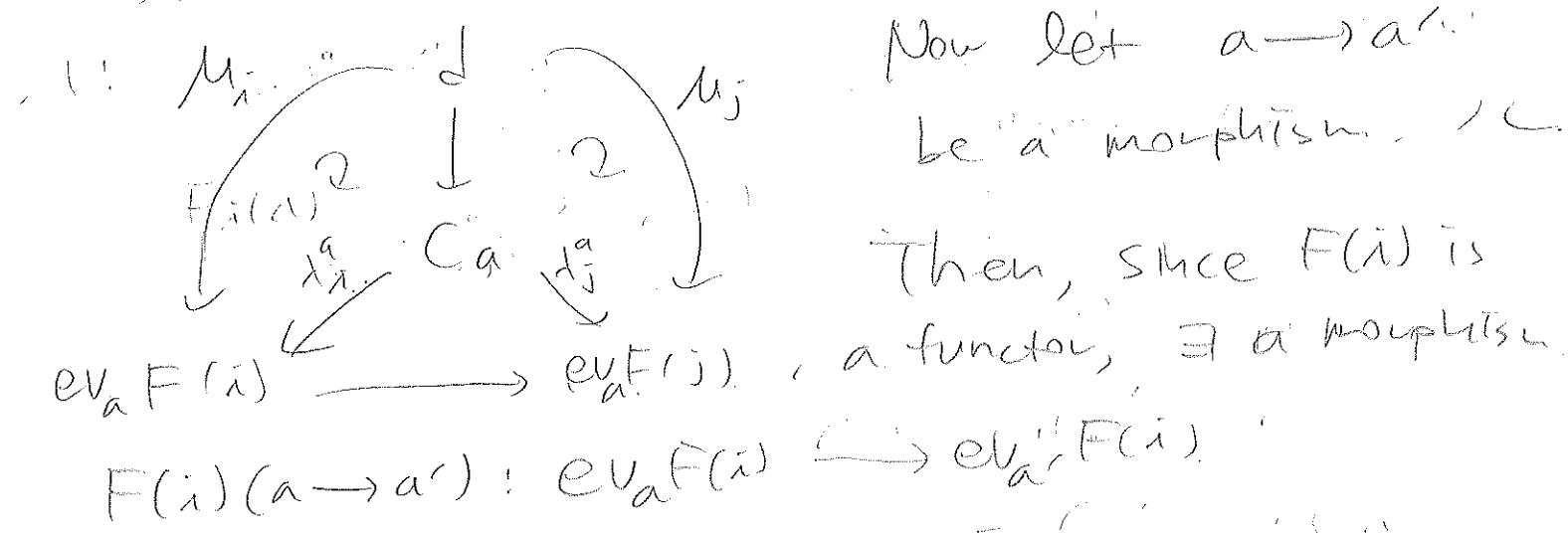
To show $\pi: C^A \rightarrow C^{ob A}$ strictly creates limits,

Let $F: J \rightarrow C$ a diagram s.t. $J(E)$ has a limit
and π limit cones. $\lambda: (C_i)_{\text{neq} A} \rightarrow \pi F$.

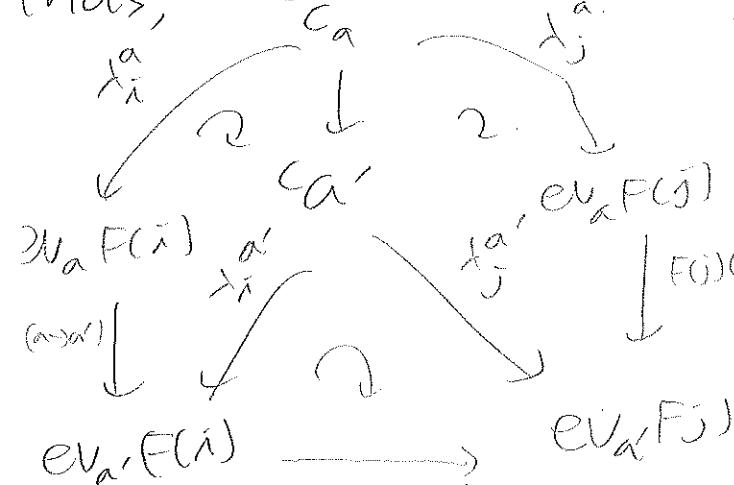
This equivalent to say that C admits a limit.

$\therefore \exists J \in \text{Ob } C^A \xrightarrow{\text{ev}_a} C$ admits a limit for any $a \in \text{Ob } A$. Let $\lambda^a: C_a \rightarrow \text{ev}_a F$ be a limit cone.

Then, for any $M: d \rightarrow \text{ev}_a F$, $i \mapsto j, i \in J$,



Thus, we have the unique map $G_a: C_a \rightarrow C_{a'}$ s.t. the diagram commutes.



Now define $G: A \rightarrow C$ s.t. $G(a) = C_a$.

$G(a \rightarrow a') = (a \rightarrow a')$ induced by the universal property of limit cone $\lambda^{a'}$.

Then functionality is satisfied by the uniqueness of universal property) I_a goes to I'_{ca} by uniqueness and $a \rightarrow a' \rightarrow a''$ goes to $a \rightarrow a''$ by uniqueness too.

Thus G is a functor. Now we claim

G is $\lim(\mathcal{I} \rightarrow \mathcal{C})^A$

Let $n: H \Rightarrow F$ be a cone. Then,

$$\begin{array}{ccc} n_i & H & n_j \\ \downarrow & & \downarrow \\ F(i) & \longrightarrow & F(j) \\ & & F(i \rightarrow j) \end{array} \quad \text{Thus, by } ev_a, \text{ we set}$$

$$\begin{array}{ccc} & H(a) & \\ (n_i)_a & \swarrow h_a & \searrow (n_j)_a \\ & ca & \\ \downarrow & & \downarrow \\ {}_a F(i)(a) & \xrightarrow{\quad ev_a F(i) \quad} & ev_a F(j) \\ & & ev_a F(i \rightarrow j) \end{array}$$

Then, the universal property gives left diagram
with $h: H(a) \rightarrow G(a) = G(a)$

Now we claim that $h = \{h_a : a \in \mathcal{I}\}$ is a natural transformation because of uniqueness;

$$\begin{array}{ccccc} F(i)(a) & \xrightarrow{\quad (n_i)_a \quad} & F(j)(a) & \xrightarrow{\quad (n_j)_a \quad} & H(a) \rightarrow G(a) \rightarrow G(a) \\ \downarrow & \swarrow h_a & \downarrow & \searrow h_a & \downarrow \\ H(a) & \xrightarrow{\quad h_a \quad} & G(a) & & \\ \downarrow & & \downarrow & & \\ H(a') & \xrightarrow{\quad h_{a'} \quad} & G(a') & & \\ \downarrow & & \downarrow & & \\ (n_i)_{a'} & \swarrow h_{a'} & \searrow (n_j)_{a'} & & \\ & & & & \\ F(i)(a') & \xrightarrow{\quad (n_i)_{a'} \quad} & F(j)(a') & \xrightarrow{\quad (n_j)_{a'} \quad} & H(a') \rightarrow G(a') \rightarrow G(a') \end{array}$$

and $H(a) \rightarrow H(a') \rightarrow G(a')$
gives the commutativity
cone structure \Rightarrow uniqueness
of limit cone give the
same map.

Finally, this natural transformation make cones commute
via left diagram
 $\Rightarrow G$ is $\lim_{\leftarrow} F$.

Ex. 3.3.i) Let $\tilde{\eta} \rightarrow \tilde{\zeta} \in \mathcal{J}$. $\text{Colim } F\mathcal{K}$ be a colimit core.

(i)

$F\text{Colim } \mathcal{K}$

$\exists! t: \text{Colim } \mathcal{K} \rightarrow F\text{Colim } \mathcal{K}$

$F K i \xrightarrow{2} F K j$

$F K i \rightarrow F K j \Rightarrow F d$ is a cone in D .

then left diagram commutes by functoriality

Thus, the universal property of $\text{Colim } \mathcal{K}$ gives

$\exists! t: \text{Colim } \mathcal{K} \rightarrow F\text{Colim } \mathcal{K}$

$F K i \rightarrow F K j$

$F K i \rightarrow F K j$

$M_j > \text{Colim } \mathcal{K}$

$\exists! t: \text{Colim } \mathcal{K} \rightarrow F\text{Colim } \mathcal{K}$

$F K i \rightarrow F K j$

st. left diagram commutes
($M: F K i \rightarrow \text{Colim } \mathcal{K}$
is a colimit cone)

(ii) If t is an isomorphism, then $F d$ is also a colimit core since $\mu_i = t^{-1} f d_i \circ t_i$. So if $\eta: F \mathcal{K} \Rightarrow X$ is a cocone then, $\exists! t': \eta \rightarrow \text{Colim } \mathcal{K}$ st. $\eta_i = t' \circ \mu_i \circ t_i$. $\Rightarrow \eta_i = t' \circ t^{-1} \circ f d_i \circ t_i$. And $t' \circ t$ uniquely determined by the universal property of $\text{Colim } \mathcal{K}$.

Thus, F preserves colimit.

3.3 ii), 3.3 iii) already done in lecture note.

Ex 3.3iv). Let $K: J \rightarrow C$ be in class. and

$\exists \mu: J \rightarrow K$ s.t. $\mu_i: J_i \rightarrow K_i$ is a limit cone.

By (i), $\exists \mu: (\Rightarrow)K$ is a limit cone and
 $F\mu$ is a limit cone in D . Thus FK has
a limit cone over $FK: (F\mu)$ which can
be lifted to a limit cone (μ) over K .

Thus, we need to show that F reflects limits.

Let $\lambda: C \Rightarrow K$ st. $F\lambda$ is a limit cone.

Then, $F\lambda \cong F\mu$ by Prop 3.1.7 (Uniqueness
of limit cone)

$\Rightarrow Fc \cong Fc' \Rightarrow c \stackrel{\cong}{\sim} c'$ by (ii) (F reflects iso.)

Then, it suffices to show that λ is a limit cone

Notes that for any $\gamma: C'' \Rightarrow K$,

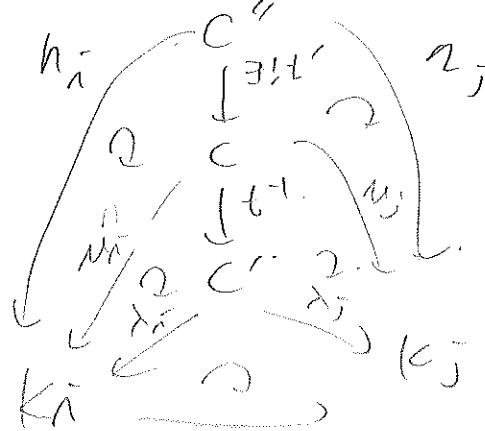
then left diagram commutes

$\Rightarrow \exists !: t^t: C'' \rightarrow C$

st. cones @ commutes,

$\Rightarrow \lambda$ is a limit cone.

□



Ex 3.3, v) Let $(A, *) \in (B, \#) \in \text{Set}_{\text{f.g.}}$
Then their coproduct D disjoint union of A, B
but picking apt from A or B .

Thus, if point picked from B , then

$(A, *) \rightarrow (A, *) \amalg (B, \#)$ is not a function.

$A \rightarrow A \amalg B$ in Set. Hence it is not a
coproduct, in Set, since it is possible that
in Set, $*$ and $\#$ can be sent to different
objects.

(This happen in Top $_X$ too.)

If, J was connected, then \exists morphism S -t.
sending $x \mapsto *$. or vice versa in limit.

Ex 3.3 vi) done. In lecture note.

3.4. Representable nature of \$\mathrm{lim}\$ and \$\mathrm{colim}\$.

Thm 3.4.2 For any \$F: \mathcal{J} \rightarrow \mathcal{C}\$ s.t. \$\mathrm{lim} F\$ exists.
and \$\mathcal{C}(-, \mathrm{lim}_\mathcal{J} F) \cong \mathrm{lim}_\mathcal{J} \mathcal{C}(-, F-)\$.

pf) First of all, Let \$X \in \text{Set}\$, Then, we claim

$$\mathrm{lim}_\mathcal{J} \mathcal{C}(X, F-) \cong \text{Cone}(X, F). \quad (\cong \mathcal{C}(-, \mathrm{lim}_\mathcal{J} F)$$

by definition of \$\mathrm{lim}\$.

If. \$\lambda \in \text{Cone}(X, F)\$ then \$\lambda\$ satisfies (Regularity).

$$\text{In Thm 3.2.3. } \Rightarrow \lambda \in \mathrm{lim}_\mathcal{J} (X, F-).$$

Conversely, any. \$\lambda \in \mathrm{lim}_\mathcal{J} (X, F-)\$ is an element of \$\prod_{j \in \mathcal{J}} \mathcal{C}(X, F_j)\$, and \$\forall i \xrightarrow{j} j \in \mathcal{J}

$$\begin{array}{c} \lambda_i \xrightarrow{X} \lambda_j \\ \downarrow \quad \downarrow \\ F_i \xrightarrow{Ff} F_j \end{array} \Rightarrow \lambda \text{ is a cone over } F \text{ with sumit } X$$

\$\therefore \exists \text{ isomorphism.}\$

To see this is natural, think when \$X \rightarrow X'\$
is a set map. Then, for any \$\lambda \in \mathrm{lim}_\mathcal{J} (X', F)\$

$$\begin{array}{c} \begin{array}{ccc} g & X' & y \\ \downarrow & \nearrow & \downarrow \\ x & & \end{array} & \lambda \in \prod_{j \in \mathcal{J}} \mathcal{C}(X', F_j) \\ \begin{array}{ccc} \lambda_i & \xrightarrow{X'} & \lambda_j \\ \downarrow & \nearrow & \downarrow \\ F_i & \xrightarrow{Ff} & F_j \end{array} & \text{if } g \text{ is left adj.} \\ \text{thus, } \lambda g \in \prod_{j \in \mathcal{J}} \mathcal{C}(X, F_j) & \end{array}$$

But as a 'cone', \$\lambda g\$ is left adj. which is the same one.

$\cong \mathrm{lim}_\mathcal{J} \mathcal{C}(X, F-) \xrightarrow{\cong} \text{Cone}(X, F)$

$\downarrow g_X \quad \downarrow g_X$

$\mathrm{lim}_\mathcal{J} \mathcal{C}(X', F-) \xrightarrow{\cong} \text{Cone}(X', F)$

(Thm 3.4.3') natural is o expressive?

"representable universal property of the limit"

Im_F is defined representantly as $\text{Im}(X_F)$ for chosen X .

Ex. 3, 4, 4.) $F: \mathbb{I} \rightarrow \mathcal{C}$ is an iterated product,

i.e. $\text{Im } F = A^T := \bigcap_{\mathbb{I}} A$, then the property gives $C(X, A^T) \cong \bigcap_{\mathbb{I}} C(X, A)$.

i.e. any morphism $h: X \rightarrow A^I$ is determined by $h_i: X \rightarrow A$ for all $i \in I$ with each product projection $\pi_i: A^I \rightarrow A$.

In a category of set, this is already known.

Prop 3.4.5: $C(X, -) : C \rightarrow \text{Set}$ preserves limit.

If C is locally small.

If C is locally small.
 $\text{pf})$ Thm 3.4.2 shows $(X, \Omega^F) \cong \Omega^{\infty}(\mathcal{T} \rightarrow C \rightarrow \text{Set})$

pf) Thm 3.4.2 shows (X, \mathcal{F}_X) is a
Moreover, remember that the natural isomorphism
is constructed via $\text{JET}(X, \mathcal{F}_X)$. Hence,
JET

$$\lim_{\leftarrow}^{\vee} C(X; F \rightarrow) \cong \text{Cone}(X; F) \cong C(X, \lim_{\leftarrow} F)$$

And if $x_j : \lim_{j \in J} F_j \rightarrow F_j$ is a limit of limit one,
 $\pi_{\{j\}}(x_j)$ is a limit one of $\pi_{\{j\}}(F_j)$.

$C(X, \lambda) = \prod_{j \in J} C(X, \lambda_j)$ is RHS diagonal map, and LHS is product of less of limit cone.

More precisely,

$\mu \searrow$ is constructed by Thm 3.2.13;

Actually Thm 3.2.13 construction shows that
 λ is a limit core.

Then $\text{Core}(X, F) \rightarrow C(X, \text{lim}_j F)$

$\downarrow \quad ? \quad \swarrow (\alpha_j)$ is just application
 $\prod_{j \in J} (X, F_j)$ of the universal property
of limit core λ ;

any cone $\lambda': X \Rightarrow F$ has a map $\exists! k: X \rightarrow \text{lim}_j F$.

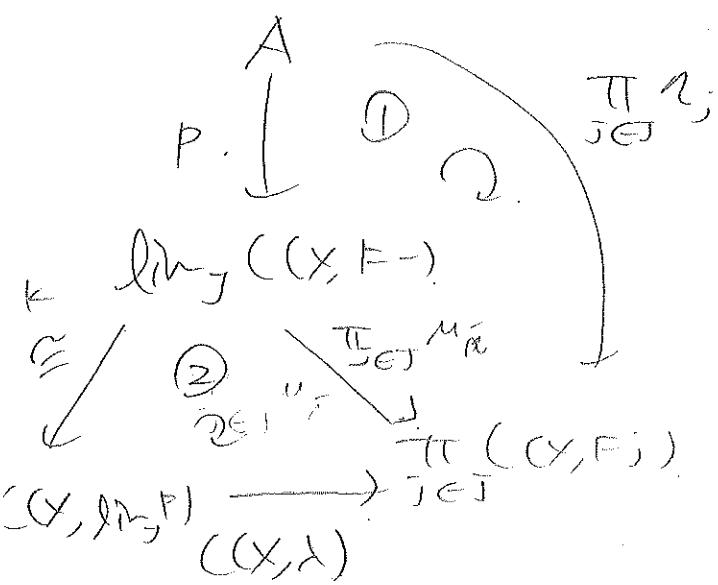
Ext. $\begin{array}{ccc} X & & \\ \downarrow \alpha_1 \quad \downarrow \alpha_2 \quad \downarrow \alpha_j & \nearrow \text{lim}_j F & \\ f_1 & \longrightarrow & f_j \end{array}$ commutes
So, $X' \in \text{Core}(X, F)$
 f_i can be sent to k ,
the unique map.

and $k \circ \alpha_i = \alpha'_i \quad \forall i \in J$ shows that
the triangle commutes via $C(X, \lambda)$.

Hence, by identifying $C(X, \lambda)$ as a cone $\lambda: X \Rightarrow$
the diagram shows that it acts

as a limit core; i.e. for any
cone $\eta: A \Rightarrow C(X, F-)$ in Set

$\exists!$ unique map $p: A \rightarrow \lim_j ((X, F) \rightarrow)$ s.t.



① commutes

Moreover, ② also commutes by above

\triangleleft Tagra

\Rightarrow Unique map k_p .

s.t. $A \xrightarrow{k_p} ((X, \lim_j F) \xrightarrow{(X, \lambda)} \lim_{j \in J} ((X, F)))$

(λ must equal to $\lim_{j \in J} h_j$;

\square

This implies $\gamma_j = ((X, \lambda)) \circ k_p \Rightarrow ((X, \lambda))$ is a limit core.

Thm 3.4.6 (Second interpretation of Th 3.4.2) \square

C: locally small

(i) $C(X, -)$ preserves all limits that exist in S

(ii) The covariant Yoneda embedding

$y: C \hookrightarrow \text{Set}^{C^{\text{op}}}$ preserves and reflects limits.

pf) (i) is just result of Prop 3.4.5.

(ii) Yoneda embedding is fully faithful \Rightarrow reflects limit.

Also, $y(\lim_j F) = ((-, \lim_j F))$

And for any $x \rightarrow x' \in C$, $((x', \lim_j F)) \cong \lim_j ((x, F) \rightarrow)$

By Thm 3.4.2

$((x, \lim_j F)) \cong \lim_j ((x, F) \rightarrow)$

This shows $\gamma(\lim_j F) = ((-\lim_j F) \cong \lim_j (-, F))$
 as we've constructed in Prop 3.3.9. \square

Finally, for any $F: J \rightarrow C$, $X \in C$, let

$$((F-, X)) := J^{\text{op}} \xrightarrow{F} C^{\text{op}} \xrightarrow{(-, X)} \text{Set}.$$

Thm 3.2.6 shows that a limit exists.

Thm 3.2.73 shows construction of limit:

By the very similar argument of pf of

Thm 3.4.2, $\lim_{J^{\text{op}}} ((F-, X)) \cong \text{Gae}(F, X)$.

Thm 3.4.7 For any $F: J \rightarrow C$ "a",

$$\exists \text{ natural } i_0 \quad ((\text{colim}_J F, X)) \cong \lim_{J^{\text{op}}} ((F-, X))$$

"Representable uni pro. of colimit".

(Ex. 3.4.8) $Y, Y, Z \in \text{ob } C$. Thm 3.4.7 implies

$$((X \sqcup Y, Z)) \cong ((X, Z) \times (Y, Z)).$$

Ex 3.4.9. Iterated Coproduct of $A \in C$: "copowers"
 or "tensors".

$$\Rightarrow ((\coprod_I A, X)) \cong ((A, X))^I.$$

i.e. $h: \coprod_I A \rightarrow X$ is determined by $h_i: A \rightarrow X$.

with coproduct inclusion $i_i: A \rightarrow \coprod_I A$.

Ex 3.4.10. By rep universal
 property.

$$\begin{array}{ccc} U & \xrightarrow{f} & V \\ \downarrow g & \lrcorner & \downarrow h \\ W & \xrightarrow{k} & P \end{array} \Rightarrow \begin{array}{ccc} ((P, X)) & \xrightarrow{((h, X))} & ((V, X)) \\ \downarrow ((k, X)) & & \downarrow ((f, X)) \\ ((W, X)) & \xrightarrow{((g, X))} & ((U, X)) \end{array}$$

pull back.

Thm 3.4.11. C locally small

(i) $\mathbf{C}(-, X)$ carry colimits to limit

(ii) $\gamma: \mathbf{C}^{\text{op}} \hookrightarrow \text{Set}^C$ Yoneda embedding preserves and reflects limits in \mathbf{C}^{op}

(pf) Dualize the argument. for covariant case.

Thm 3.4.12 (Generalization of 3.2.13)

$F: \mathcal{I} \rightarrow C$ small diagram. Suppose C has coproduct and coequalizers.

$$\coprod_{f \in \text{Mor}(\mathcal{I})} F(\text{dom } f) \xrightarrow{d} \coprod_{j \in \text{Ob}(\mathcal{I})} F_j \longrightarrow \text{Colim}_j F.$$

In particular, if C admits coproduct and coequalizers then \mathcal{I} is complete. Dually,

$$\text{Colim}_j F \longrightarrow \prod_{j \in \text{Ob}(\mathcal{I})} F_j \xrightarrow{c} \prod_{f \in \text{Mor}(\mathcal{I})} F(\text{cod } f)$$

and if C admits product and equalizers, the C is complete.

(pf) Using Remark 3.2.15, we have below diagram

$$\begin{array}{ccc} & \text{F(dom } f) & \\ & \swarrow \bar{c}_{\text{dom } f} & \\ \mathcal{C} & \leftarrow \coprod_{j \in \text{Ob}(\mathcal{I})} F_j & \xleftarrow{d} \coprod_{f \in \text{Mor}(\mathcal{I})} F(\text{dom } f) \\ & \uparrow \bar{c}_{\text{cod } f} & \uparrow \bar{c}_f \\ & \text{F(cod } f) & \end{array}$$

where $(d)_f = \bar{c}_{\text{dom } f}$
 $(c)_f = \bar{c}_{\text{cod } f} \circ F_f$.

By hypothesis C (a coequalizer of c and d) exists.

By Thm. 3.4.11(ii), Yoneda embedding carries coequalizer diagram to an equalizer diagram. Thus (for any $X \in \mathcal{C}$),

$$\begin{array}{ccccc}
 & & C(F(\text{dout}), X) & & \\
 & C(\bar{i}_{\text{dout}}, X) \nearrow & & & \uparrow C(i_f, X) \\
 (C) & \rightarrow C(\coprod_{j \in \text{Ob}^{\text{op}}} F_j, X) & \xrightarrow{C(c, X)} & C(\coprod_{f \in \text{Mor}^{\text{op}}} F(\text{dout}^f), X) & \\
 & \downarrow C(\bar{i}_{\text{cod}}^f, X) & & & \downarrow C(i_f, X) \\
 C(F(\text{cod}^f), X) & \longrightarrow & C(F(\text{dout}^f), X)
 \end{array}$$

Applying 3.4.11(i) carries colimit to a limit, we have

$$\begin{array}{ccc}
 \pi_{\text{cod}^f} & \longrightarrow & C(F(\text{cod}^f), X) \\
 \uparrow \pi_c & & \\
 C(C, X) \longrightarrow \prod_{j \in \text{Ob}^{\text{op}}} (F_j, X) & \xrightarrow{c} & \prod_{f \in \text{Mor}^{\text{op}}} C(F(\text{cod}^f), X) \\
 \pi_{\text{dout}^f} \downarrow & & \downarrow \pi_f \\
 C(F(\text{dout}^f), X) & \xrightarrow{C(F, X)} & C(F(\text{cod}^f), X)
 \end{array}$$

(choose from $F(\text{dout}^f)$ + $F(\text{cod}^f)$ occur by indexation with \mathcal{I}^{op}). where $c := C(c, X)$

If X is fixed, by Prop 3.3.9, with $\text{ev}_X : \text{Set}^{\mathcal{I}^{\text{op}}} \rightarrow \text{Set}$, the above one is an equalizer diagram in Set .

By Th 3.2.13., The drazen's equalizer

is equal to $\lim_{\mathcal{J}^{\text{op}}} ((F_{-}, \times))$.

$\Rightarrow \lim_{\mathcal{J}^{\text{op}}} ((F_{-}, \times)) \cong (\mathbf{C}(C, \times)) \quad \forall X \in C.$

Now these are assembled into a function

in Set^{obc} . Using forgetful function $\text{Set}^{obc} \rightarrow \text{Set}^{ob}$
with prop 3.3.9, $(\mathbf{C}(C, \times)) \cong \lim_{\mathcal{J}^{\text{op}}} ((F_{-}, \times))$

By Thm 3.4.7, $(\mathbf{C}(C, \times)) \cong \varprojlim_{\mathcal{J}^{\text{op}}} ((F_{-}, \times)) \cong ((\text{colim}_{\mathcal{J}} F_{-}) \times)$

$\forall X \in C. \Rightarrow C \cong \text{colim}_{\mathcal{J}} F_{-}$.

□

Thm 3.4.15. (Generalized element)

checks

Direct pf of above Thm is just C and it
has the universal property that defines $\text{colim}_{\mathcal{J}} F_{-}$.
This can be viewed as the same thing of what we actually
did for getting proof, via 'philosophy of generalized
elements' =

Ex) $\forall a \in A^{\text{Set}}$, we can identify a as $a : \mathbb{1} \rightarrow A$.

In this case, $\forall f : A \rightarrow B$, $f(a)$ is identified as fa .

In any category C, a "generalized element" of A
is a morphism $a : X \rightarrow A$ with codomain A.
Then, $f : A \rightarrow B$ act on X-shaped element by composition
 $\Rightarrow \text{Hom}(-, A) : C^{\text{op}} \rightarrow \text{Set}$ send object X to the
set of X-shaped generalized element of A.

Since $C \hookrightarrow \text{Set}^{\text{op}}$ Yoneda embedding
 is fully faithful, no information about A
 is lost by considering not A but $\text{Hom}(A)$.

For example, if we want to show $f_*: A \rightarrow B$
 are equal in this philosophy, it suffices to show that

$$f_* = g_* : \text{Hom}(X, A) \rightarrow \text{Hom}(X, B)$$

If we want to show a cone is limit cone,
 then show that induced function in Set is a
 cone in Set , etc...

Lemma 3.4.16 Let C : category with terminal

$$\text{object } 1. \Rightarrow A \times B \xrightarrow{\pi_B} B \quad \begin{matrix} \text{pull back} \\ \text{defines} \\ \text{product } A \times B, \end{matrix}$$

$$\pi_A \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \text{of } A, B \in C.$$

Pf). Let $A \leftarrow \begin{smallmatrix} \pi_A' \\ \sqsubset \\ \pi_B' \end{smallmatrix} \rightarrow B$ a diagram.

Then, universal property of pull back gives
 a map $\exists ! k: C \rightarrow A \times B$ st.

$$\begin{smallmatrix} \pi_A' & \sqsubset & \pi_B' \\ \swarrow & \downarrow k & \searrow \\ A & \leftarrow \begin{smallmatrix} \sqsubset \\ A \times B \end{smallmatrix} \rightarrow B \end{smallmatrix} \quad \text{left diagram commutes.}$$

$$A \leftarrow \begin{smallmatrix} \sqsubset \\ A \times B \end{smallmatrix} \rightarrow B \quad \text{By deleting } \sqsubset \text{ and } \sqsupset,$$

$\begin{smallmatrix} \pi_A & \sqsubset \\ \swarrow & \downarrow \\ I & \leftarrow \begin{smallmatrix} \sqsubset \\ I \end{smallmatrix} \rightarrow I \end{smallmatrix}$ this is just universal property
 of product of two elements.

Lem 3.4 (7) C : category with binary products.

\Rightarrow pullback defines an equalizer, i.e.

$$\begin{array}{ccc} E & \xrightarrow{e} & A \\ l \downarrow \lrcorner & \downarrow (f, g) & \rightsquigarrow \\ B & \xrightarrow{(1_B, 1_B)} & B \times B \end{array}$$

pf) The representable universal property of the product

$$\Rightarrow \text{Hom}(X, B \times B) \cong \text{Hom}(X, B) \times \text{Hom}(X, B)$$

Also, for pull back we have (by Thm 3.4.6 (1))

$$\begin{array}{ccc} \text{Hom}(X, E) & \xrightarrow{e_k} & \text{Hom}(X, A) \\ I_X \downarrow & \downarrow & \downarrow (f_X, g_X) \\ \text{Hom}(X, B) & \xrightarrow{(I_k, I_k)} & \text{Hom}(X, X \times B) \times \text{Hom}(X, B) \end{array}$$

By concrete construction of pullback in Set
as we've seen in Ex 3.2.11,

$$\text{Hom}(X, E) \cong \{ (h_1, h_2) \in \text{Hom}(X, A) \times \text{Hom}(X, B) \mid$$

$$\text{s.t. } (I_X, I_X)(h_2) = (f_X, g_X)(h_1) \}$$

$$\Leftrightarrow f h_1 = h_2 = g h_1$$

Thus, $\text{Hom}(X, E) \cong \{ h \in \text{Hom}(X, A) : f h = g h \} \subseteq \text{Hom}(X, A)$,
which is equalizer of $(f_X \text{ and } g_X)$.

Thus, $\text{Hom}(X, E) \xrightarrow{e_k} \text{Hom}(X, A) \xrightarrow{f_X} \text{Hom}(X, B)$
is an equalizer dia sham. ~~No other~~ ~~the~~ ~~3.4.6 (1)~~ shows

Now Thm 3.4.6.(ii) shows that E is an equalizer of f, g . \square

Ex 3.4.7) Done in lecture note.

Ex 3.4.ii) Yoneda embedding preserves limit via Thm 3.4.2

" reflects limit since it is fully faithful.

But it does not create limit.

since limit may be not representable;

Ex 3.4.iii) $F: \mathcal{C} \rightarrow \text{Set}$, $\Pi: \int F \rightarrow \mathcal{C}$.

(i) Π strictly creates all limits that C admits and F preserves.

Let $k': J \rightarrow \int F$ a small diagram, s.t $k'_i = (k_i, x_i)$

By this notation, we may define $k: \text{Ob } J \rightarrow \mathcal{C}$ by $i \mapsto k'_i$.

Also, if $i \rightarrow j \in \text{Mor } J$, $k'(i \rightarrow j): (k_i, x_i) \xrightarrow{f} (k_j, x_j)$

and f comes from a morphism $k_i \rightarrow k_j$ s.t

$Ff(x_i) = x_j$. Hence by defining $k(i \rightarrow j) = f$, we may define $k: J \rightarrow \mathcal{C}$ a function.

Now let $\delta: (\mathcal{C}, x) \Rightarrow k'$ be a limit cone.

Then $(\frac{x_i}{2} \xrightarrow{c, x_i} x_j) \Rightarrow (\frac{x_i}{2} \xrightarrow{k_i} x_j)$ thus this λ

Induces a cone λ in C , which is a limit iff
 since other cone in C must have at least
 one distinct λ_i , which induces distinct legs M_i
 ST.

Now it suffices to show that λ is a limit cone in C .
 Let $M: d \Rightarrow k$ be a cone. Then,

$$\begin{array}{ccc} M_i \downarrow_{C} & \xrightarrow{d} & M_j \\ \downarrow_{k_i \longrightarrow k_j} & & \downarrow_{M_i(y)} \end{array} \Rightarrow \begin{array}{ccc} M_i \downarrow_{\lambda} & \xrightarrow{(d,y)} & M_j \\ \downarrow_{(k_i, M_i(y))} & & \downarrow_{(k_j, M_j(y))} \end{array} \text{ by } \text{YCF}_d.$$

By universal property of λ , $\exists! f: (d, y) \rightarrow (\lambda, x)$

st. left diagram commutes.

$$\begin{array}{ccc} M_i & \xrightarrow{(d, y)} & M_j \\ \downarrow_{k_i \longrightarrow k_j} & \downarrow f & \downarrow_{M_i(y)} \\ \downarrow_{(k_i, M_i(y))} & & \downarrow_{(k_j, M_j(y))} \end{array} \Rightarrow \begin{array}{ccc} M_i & \xrightarrow{\lambda} & M_j \\ \downarrow_{k_i \longrightarrow k_j} & \downarrow f & \downarrow_{M_i(y)} \\ \downarrow_{(k_i, M_i(y))} & & \downarrow_{(k_j, M_j(y))} \end{array}$$

$\therefore \lambda$ is a limit cone in C .

ii) π strictly creates connected colimit.

Let λ be a colimit cone. Then

$$(k_i, \lambda_i) \rightarrow (k_j, \lambda_j) \Rightarrow \begin{array}{c} k_i \rightarrow k_j \\ \downarrow_{\lambda_i \longrightarrow \lambda_j} \end{array}$$

So λ induces a cocore in C .

To see λ is a colimit in C , let $M: K \Rightarrow d$ a cocone. Then, from $i \rightarrow j \in J$,

$$\begin{array}{ccc} k_i & \longrightarrow & k_j \\ \downarrow \mu_i & \nearrow \mu_j & = \\ M_i & & M_j \\ & & \downarrow \mu_i \quad \downarrow \mu_j \\ & & (d, M_i(x_i)) \quad \text{and} \\ & & \qquad \qquad \qquad M_j(x_j) \end{array}$$

Notes that this μ is unique core in \mathcal{SF} since connectedness fixes y regardless of choice of i, j . precisely, if we fix x_i for some $i \in J$, x_j is fixed via zig zags up over the core $M_i \Rightarrow y$ is fixed.

Now universal property of d shows that,

$$\begin{array}{ccc} (k_i, x_i) & \longrightarrow & (k_j, x_j) \\ \downarrow \gamma_i & \nearrow \gamma_j & \\ M_i & \longrightarrow & M_j \\ & & \downarrow \mu_i \quad \downarrow \mu_j \\ & & (d, M_i(x_i)) \quad M_j(x_j) \end{array} \Rightarrow \begin{pmatrix} k_i & \longrightarrow & k_j \\ 2 & \nearrow & \downarrow \\ 2 & 1 & 2 \end{pmatrix}_d$$

Thus λ is a colimit core in C .

Ex. 3.4.iv) done in lecture note.

3.5. Completes and CoCompletes Categories.

A function is continuous / cocontinuous if it preserves all small limits or colimits.

Prop 3.5.1. Set isocomplete.

pf) By thm 3.4.12, it suffices to show that Set has coproducts and coequalizers.

Coproduct = disj. union.

Coequalizer of $A \xrightarrow{f} B$ is $B / f(a) \cap g(a)$

Prop 3.5.2. Top is cocomplete.

pf) $\cup : \text{Top} \rightarrow \text{Set} \cong \text{Top}(*, -)$

By prop 3.4.5, if Top has any limits, then its underlying set is limit in the category Set.

① Product: $\prod_{\alpha} X_{\alpha}$: Cartesian product with product topology.

(by definition, it is the coarsest topology s.t. $\prod_{\alpha} X_{\alpha} \rightarrow X_{\alpha}$ iscts.)

② Equalizer: The equalizer $E \rightarrow X \rightrightarrows Y$ with subspace topology inherited from X , i.e. the coarsest topology s.t. $E \hookrightarrow X$ iscts.

\Rightarrow Thm 3.4.12 shows that Top is complete.

Pually, ① Coproduct: disjoint union as Top space, the finest topology s.t. $X_0 \rightarrow \coprod X_{\alpha}$ iscts.

② Coequalizer of $X \rightrightarrows Y \Rightarrow Y /_{\substack{f(a) \\ g(a)}} \text{ with quotient topology}$

Hence Th 3.4.12 shows that Top B is complete \square

Ex 3.5.3. Any limit/colimit can be constructed via diagram in prop 3.5.2. by constructing (the coarsest (limit)) topology (the finest (colimit)) on the limit/colimit in set.

ex) $J: \mathcal{W}^{\text{op}} \rightarrow \text{Top}$ $p: S^1 \rightarrow S^1$ is p th power map.
 $n \mapsto S^1$
 $\downarrow \mapsto \downarrow^p$
 $n+1 \mapsto S^1$ Some have $S^1 \xrightarrow{e^{2\pi i x}} (e^{2\pi i x})^p$.

$\Rightarrow \lim J = p\text{-adic solenoid.}$

Ex 3.5.4. Space of ordered configurations of n points

: $P\text{Conf}_n(X) := \{(x_1, \dots, x_n) \in X^n : x_i \neq x_j \text{ if } i \neq j\}$
 with subspace topology $\subseteq X^n$.

Let Σ_n : symmetric group.

$\text{Conf}_n(X): B\Sigma_n \rightarrow \text{Top}$.
 $\bullet \mapsto P\text{Conf}_n(X)$
 $g \downarrow \mapsto \downarrow \text{permutation by } g$.
 $\circ \mapsto \circ$

$\Rightarrow \text{Conf}_n(X)$
 $:= \text{colim}_{P\text{Conf}_n(X)} (B\Sigma_n \xrightarrow{P\text{Conf}_n(X)} \text{Top})$

Prop. 3.5.5. \mathcal{C} is complete / cocomplete
 $\Rightarrow \mathcal{S}/\mathcal{C}$ and \mathcal{C}/\mathcal{C} are complete / cocomplete
 $\forall \mathcal{C} \in \mathcal{C}$.

Pf) Prop 3.3.8 shows that \mathcal{C} complete / cocomplete
 $\Rightarrow \mathcal{S}/\mathcal{C}$ has all limits and all connected limit.
 created by $\Pi : \mathcal{S}/\mathcal{C} \rightarrow \mathcal{C}$.

Now $1_{\mathcal{C}}$ is initial in \mathcal{S}/\mathcal{C} . Dual statement
 of Lem 3.4.16 creates coproduct in \mathcal{S}/\mathcal{C} .
 (Since pushout is connected colimit.)

And dual of Lem 3.4.17 generates coequalization.
 So, by thm 3.4.12, \mathcal{S}/\mathcal{C} is cocomplete (ii)

In this proof, coproduct of $(c \rightarrow x_a)$ is a
 map $c \rightarrow (x_0 \sqcup x_1 \sqcup \dots)$
 defined by $c \rightarrow x_a$'s. (iii)

Rmk) $\text{Ob}, \text{mor} : \text{Cat} \rightleftarrows \text{Set}$ are representable by
 Ex 2.1.5 (ix) and (x).

By Thm 3.4.6(i), Both preserves limits when
 limit of $K : J \rightarrow \text{Cat}$ exists.

Furthermore, Jonah, codomain identities are
 natural transformations between Ob and mor .

\Rightarrow They are preserved via $\text{Ob} \cong \text{mor}$.

Also composition as a $\text{Cat}(J, -) \rightarrow \text{CAT}(J, -)$

preserved. These give a candidate of product and
thus, \oplus equalizer in Cat or CAT.

Prop 3.5.6. Cat and CAT are complete.

Pf) ① Product : product of category

② Equalizer: $E \rightarrow C \xrightarrow{\begin{smallmatrix} F \\ G \end{smallmatrix}} D$

where E is a subcategory of C

$$\text{Ob } E = \{c \in \text{Ob } C : Fc = Gc\}$$

$$\text{Mor } E = \{f \in \text{Mor } C : Ff = Gf\}$$

Apply Thm 3.4.12. \square

Ex 3.5.7. Let $F: C \rightarrow \text{Set}$. Pullback in CAT
of $C \xrightarrow{F} \text{Set} \leftarrow \text{Set}_+$.

is a category $S+$.

Object: $(c, (x, \succ))$ $\in S+$. $Fc = X$.

Mor: $(c, (x, \succ)) \xrightarrow{(f, g)} (c', (x', \succ'))$

s.t. $f: c \rightarrow c' \in C$.

$g: (x, \succ) \rightarrow (x', \succ') \in \text{Set}_+$

$$Ff = Ug$$

Note that $(c, (x, \succ))$ can be denoted as (c, \succ)
with $x \in Fc$.

And morphism can be denoted as $(c, \succ) \xrightarrow{f} (c', \succ')$
s.t. $Ff(\alpha) = \alpha'$.

\Rightarrow This pullback is $\int F$.

(Claim) Cat and (AT) are cocomplete.

But proof is a little bit delicate;

① Coproduct: Just disjoint union of category

(Notes that any category is cproduct of T/Is connected components)

To use

② Lemma 3, 4, 17, we need pushout, But it is difficult.

Ex 3.5.8 $1 \sqcup 1 \rightarrow 2$ ~~We~~ we want to have pushout P.

$$\begin{array}{ccc} & & \\ \downarrow & + & \downarrow \\ 1 & \rightarrow & P \end{array}$$

Let $P \rightarrow C$ a functor.

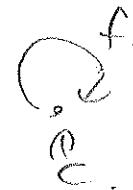
Then $1 \sqcup 1 \rightarrow 2$ gives a cocone over

$$\begin{array}{ccc} & & \\ \downarrow & & \downarrow \\ 1 & \rightarrow & P \rightarrow C \end{array}$$

diagram $1 \leftarrow 1 \sqcup 1 \rightarrow 2$ with node C.

thus, $1 \rightarrow P \rightarrow C = 1 \sqcup 1 \rightarrow 1 \rightarrow P \rightarrow C$

$$= 2 \rightarrow P \rightarrow C.$$



implies that all these map denote

thus, P should have an endomorphism without any relation on the composite f-f.

$\Rightarrow P =$ free category with one object and one non-identity morphism.

By identifying $\text{Mor}(\cdot, \cdot) = \{\mathbb{I}, f, f^2, \dots\}$

as a monoid, we know that $P = \text{BIN}$

Where IN is \checkmark natural number under addition

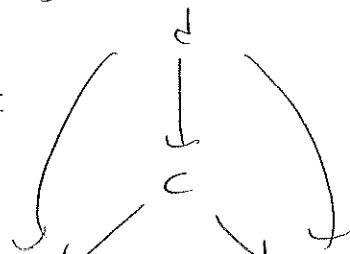
Similarly, $1 \sqcup 1 \rightarrow \mathbb{I}$ since \mathbb{D} is chosen
to working iso \mathbb{I}
 $\mathbb{I} \rightarrow \text{BIN}$ thus f should be
iso.

Completeness proved at chapter 4.

Notes that equivalence of categories preserve
complete / cocomplete

Prop. 3.5.9. P : poset B , complete and cocomplete
as a category $\Leftrightarrow P$ is a complete lattice.
i.e. every subset $A \subset P$ has both infimum and
a supremum.

pt) In any poset or preorder, if limit
(or colimit) exists, it should be infimum of
its object (or supremum of its object)

Since  implies that any $d \in \bigcap_{i \in I} U_i$
must $d \leq c$, as a term of
order.

$F_i \rightarrow F_j$ (Similar to colimit)

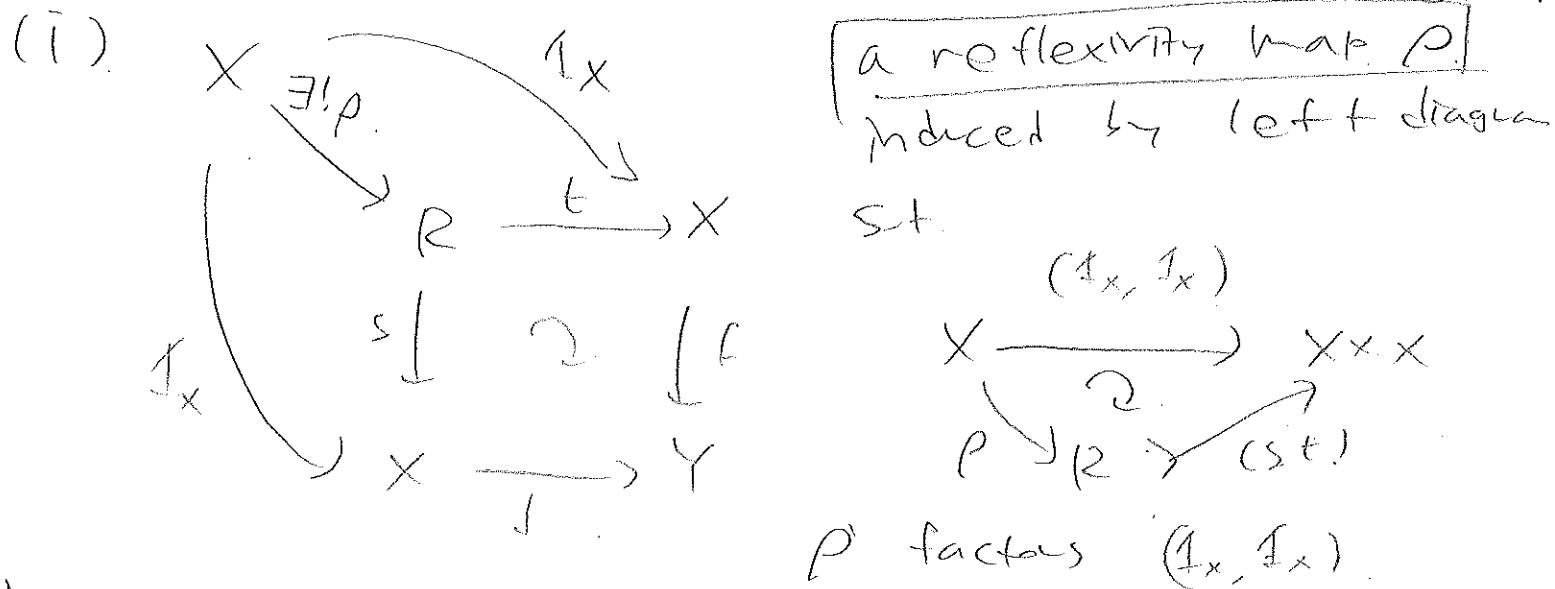
It doesn't matter when I is discrete or not.

Ex 3.5.10. Let C : category with finite limits
 Kernel pair of $f: X \rightarrow Y$ = pullback
 over f as, $R \xrightarrow{t} X$. Thus pullback
 gives $(s, t): R \rightarrow X \times X$.

$\begin{array}{ccc} s \downarrow & \cap & f \downarrow \\ X & \xrightarrow{f} & Y \end{array}$

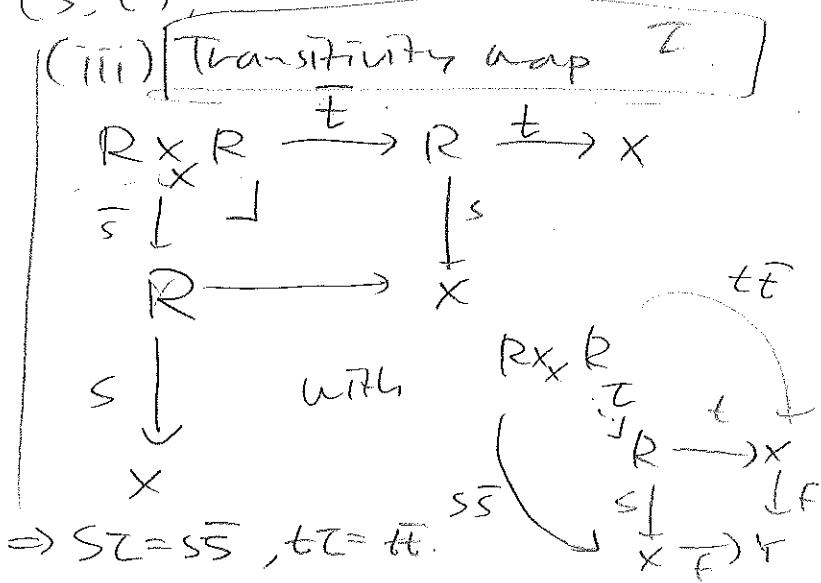
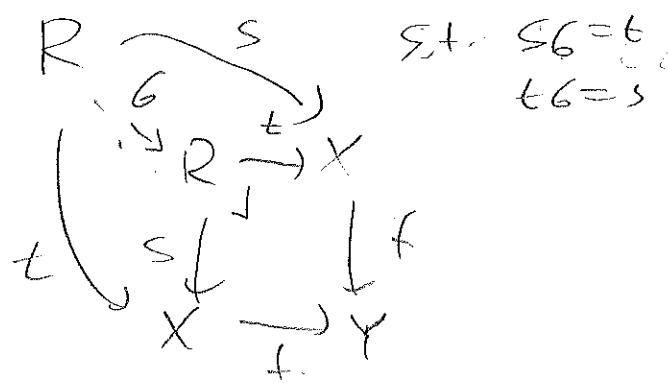
So R is always a
subset of $X \times X$
 (= monomorphism with codomain
 $X \times X$)

In Set, $R \subset X \times X$ defines relation on X .
 This is always "equivalence relation" in the below
 sense.



i.e. p is a section of (s, t)

(ii) Symmetry map σ



An equivalence relation in a category C with finite (sets)

$$= (S, t): R \rightrightarrows X \times X.$$

with ρ, δ, τ in a manner of (i), (ii) (iii).

In this case, Coequalizer of $S, t: R \rightrightarrows X \xrightarrow{e} X_{/R}$
gives $e: X \rightarrow X_{/R}$ a quotient object.

In Set, $X_{/R}$ is the set of R -equivalence classes
of elements of X . Moreover, $\exists!$ factorization

$$X \xrightarrow{f} Y \quad \text{In a good situation (Grothendieck topology)} \\ e \searrow , m \quad m \text{ is monomorphism. In this case,}$$

The factorization is called
image factorization and m is

identified as the image $X_{/R}$ as
a subobject of Y .

Ex 3.5.1) BG : 1-object category

Let $X: BG \rightarrow \text{Set}$. Colimit of X can be
decomposed as

$$\text{Colim } X \leftarrow \begin{array}{c} X \xrightarrow{\text{!-!-! ! ! ! }} X \\ \downarrow \text{!-!-! ! ! ! } \quad \downarrow \text{!-!-! ! ! ! } \\ X \xleftarrow{\text{!-!-! ! ! ! }} X \end{array} \quad \text{Hence, } \text{Colim } X = \frac{X}{\text{range}(x)} \\ = \{ G \cdot x : x \in X \}$$

a set of orbits
of G action.

Ex 3.5.ii) Use the diagram in (i) again

$$\begin{array}{ccccc}
 & & \overset{\iota_{\text{out}}}{\leftarrow} & & F(\text{out}) \\
 & \overset{\iota_{\text{out}}}{\leftarrow} & \underset{d}{\longleftarrow} & \underset{f \in \text{MorJ}}{\longleftarrow} & \downarrow \iota_f \\
 \text{Colinf} \leftarrow & \underset{j \in \text{ObJ}}{\sqcup} F(j) & & \underset{f \in \text{MorJ}}{\sqcup} F(\text{out}) & \\
 & \iota_{\text{out}} \uparrow & & & \uparrow \iota_f \\
 & F(\text{out}) & \xrightarrow{\quad Ff \quad} & &
 \end{array}$$

Then $\text{Colinf} = \{[\lambda] : \lambda \in \bigcup_{j \in \text{ObJ}} F(j)\}$

where $[\lambda] := \{x' \in \bigcup_{j \in \text{ObJ}} F_j : \exists z \in \bigcup_{f \in \text{MorJ}} F(\text{out})$
 $\text{s.t. } x' = f(z), \lambda = f(z)$
 $\text{or vice versa}\}$

Since d is just inclusion, c is Ff for each component $F(\text{out})$, this means that

$$[\lambda] := \{x' \in \bigcup_{j \in \text{ObJ}} F_j : \exists z \in \bigcup_{f \in \text{MorJ}} F(\text{out}) \text{ s.t. } x' = z, x' = Ff(z) \text{ or vice versa}\}$$

Now we claim that (j, λ) and (i, λ) is connected in \mathcal{F}
 $\Leftrightarrow [\lambda] = [z]$.

Suppose LHS. Then \exists a finite length zigzag morphism $(j, \lambda) \xrightarrow{f_1} (i_1, \lambda_1) \xleftarrow{f_2} (j_2, \lambda_2) \xrightarrow{f_3} \dots \xleftarrow{f_n} (i_n, \lambda_n)$.
 Say $j_n = i$, $\lambda_n = \lambda$. Then $[\lambda] = [\lambda_n]$ since $Ff_1(\lambda) = \lambda_1$, $[\lambda_1] = [\lambda_2]$ since $Ff_2(\lambda_2) = \lambda_1 \dots$
 $Ff_n(\lambda) = \lambda$. $[\lambda_n] = [\lambda]$ since either $Ff_n(\lambda_{n-1}) = \lambda_n$ or $Ff_n(\lambda_n) = \lambda_{n-1}$.

Is well-defined and true.

Conversely, let $[x] = [z]$. Then $\exists (x' \in x)$
 S.t. $x = x'$ and $z = Ff(x')$. for some $f: x \rightarrow z$
 or $x = x'$ and $z = Ff(x')$

This shows $i \rightarrow j \sim j \rightarrow i$, so (i, j)
 (j, i) is connected in SF .

Thus, if $\lambda: F \Rightarrow \text{Colim } F$ is a colimit cone,
 $F\bar{i} \xrightarrow{Ff} Fj$ Thus, if $x \in F\bar{i}$, then
 $\lambda_{\bar{i}}: ? \leftarrow \lambda_j$ $\lambda_{\bar{i}}(x)$ denotes the orbit of x .
 $\text{Colim } F$ under $Ff \circ Hf \circ \lambda_j$.

Thus colimit cone is just set of functions
 sending element in $F\bar{i}$ to its orbit in Set .

Ex. 3.5(iii) Let $A = \{E \xrightarrow[s]{t} V\}$ be a category
 with two parallel nontrivial morphism.

Then,

① Set^A is the category of directed multigraph.
 pf). Let $F: A \rightarrow \text{Set}$. as $E_F \xrightarrow[s]{t} V_F$
 $G: A \rightarrow \text{Set}$ $E_G \xrightarrow[s]{t} V_G$.

By thinking, $e \in E_F$ is an arrow from $s(e)$ to $t(e)$,
 gives a mult. digraph. Moreover, if $\exists X: F \Rightarrow G$

Then, $E_X: E_F \rightarrow E_G$

is

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$$E_F \xrightarrow[s]{\quad t \quad} U_F \quad \text{Thus, if } V = s(e) \text{ with}$$

$$\alpha_E \downarrow \quad ? \downarrow \alpha_G$$

$$E_G \xrightarrow[s]{\quad t \quad} U_G$$

then α_G maps $s(e)$ as $s(\alpha_E(e))$

and w as $s(t(\alpha_E(e)))$

so that incidence relation is preserved.

② Set^A is complete / cocomplete.

This is direct application of Thm 3.3.9 and the fact that Set is co/complete.

③ DirGraph is a full subcategory of Set^A .

p.f.) any morphism between digraph is a morphism as multi-digraph.

Thus $\text{DirGraph} \hookrightarrow \text{Set}^A$ is a fully faithful functor.

④ There is a functor $G: \text{Set}^A \rightarrow \text{DirGraph}$
 s.t. $\text{DirGraph} \hookleftarrow \text{Set}^A \xrightarrow{G} \text{DirGraph} = 1_{\text{DirGraph}}$

p.f.) Define $G: \text{Set}^A \rightarrow \text{DirGraph}$.

$$E \xrightarrow{\quad s \quad} V \longleftarrow E/h \xrightarrow{\quad t \quad} U$$

where $e_1, e_2 \in E \Leftrightarrow s(e_1) = s(e_2), t(e_1) = t(e_2)$

Then, \sim is equivalence relation: And for each $[e] \in E/h$, $s([e]) := s(e)$ is well-defined. Moreover, $[e] \neq [e']$, $t([e]) := t(e)$, then $s(e) \neq s(e')$ or $t(e) \neq t(e')$.

Thus, $E/\sim \xrightarrow{\cong} V$ is a directed graph.

If $E_1 \xrightarrow{\alpha_E} V_1$ is a morphism.

$\downarrow [(\alpha_E, \alpha_G)]$ notes that $\alpha_{E/\sim}([e]) := [\alpha_E(e)]$

$E_2 \xrightarrow{\beta} V_2$

gives.

$E_1/\sim \xrightarrow{\cong} V_1$

$\alpha_{E/\sim} \downarrow \quad \quad \quad \downarrow \alpha_G$

$E_2/\sim \xrightarrow{\cong} V_2$

Since if $e, e' \in [e]$, then

$$s(\alpha_E(e)) = \alpha_G(s(e)) = \alpha_G(s(e')) = s(\alpha_E(e'))$$

$$t(\text{ } \text{ } \text{ } \text{ }) = \text{ } \text{ } \text{ } (t(e)) = \text{ } \text{ } \text{ } (t(e')) = t(\text{ } \text{ } \text{ } \text{ })$$

\Rightarrow $\alpha_E(e)$ and $\alpha_E(e')$ are equivalent.

$\therefore [e] \xleftarrow{\alpha_E} t([e]) = t(e) \quad \text{So } \alpha_{E/\sim} \text{ is well-defined}$

$\downarrow \quad \quad \quad \downarrow$
 $\alpha_G(s(e)) \quad \alpha_G(t(e)) \quad \text{commutes}$

$[e] \xrightarrow{\alpha_E} s(\alpha_E(e)) \quad // \quad \dots$
 $\xrightarrow{\alpha_E} t(\alpha_E(e))$

Hence $\text{Set}^A \xrightarrow{G} \text{DirGraph}$ is a functor.

Notes that $\text{DirGraph} \hookrightarrow \text{Set}^A \xrightarrow{G} \text{DirGraph}$.

$\cong 1_{\text{DirGraph}}$

⑤) Dn Graph has colimit if Set^A has a corresponding colimit.

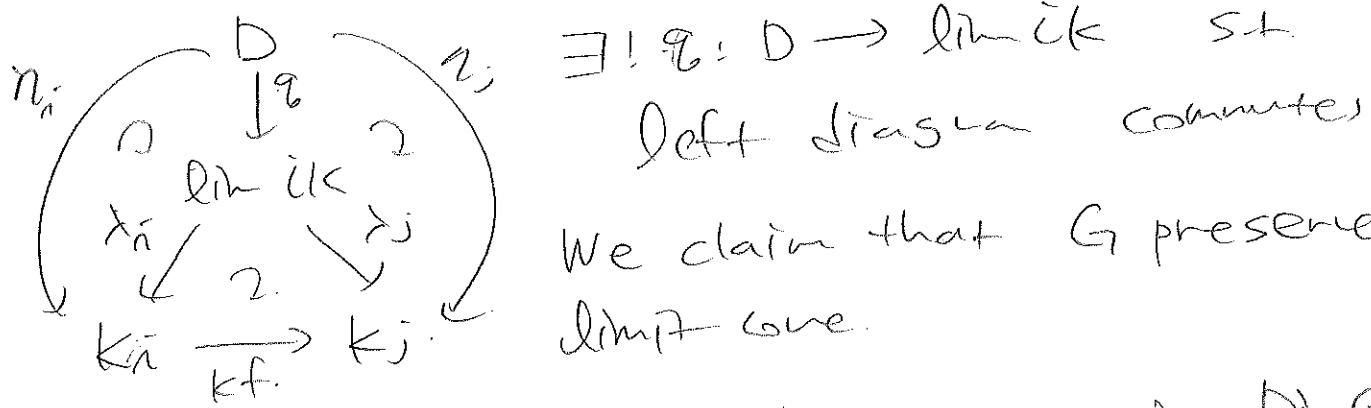
pf). Let $k: J \rightarrow$ Dn Graph.

then, $\bar{k}: J \rightarrow$ Dn Graph \longrightarrow Set^A.
is also a small diagram.

Since set^A is complete / complete, so
 $\lim \bar{k}$ and $\text{colim } \bar{k}$ exists.

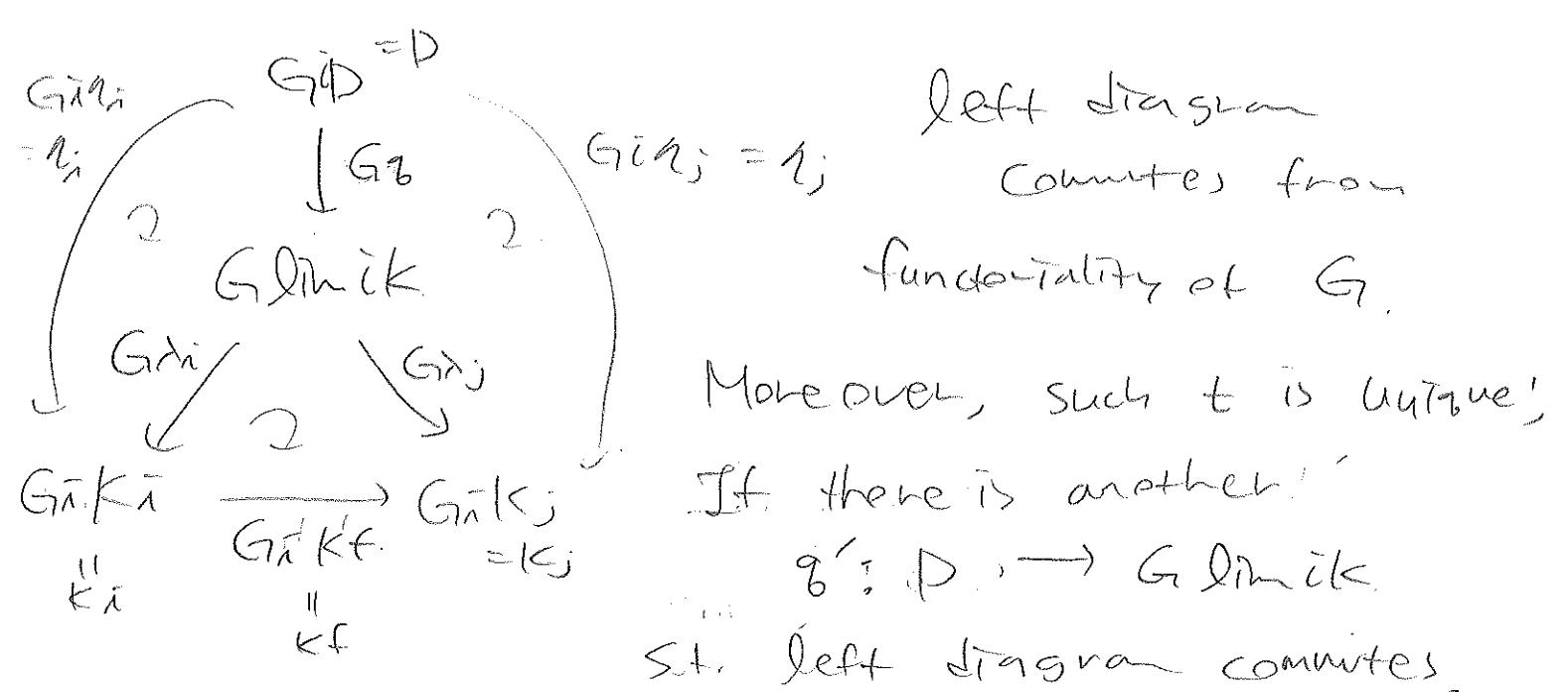
Now we claim that $G(\lim \bar{k}) \cong \lim (G\bar{k})$
 $= \lim k$.

Let $\lambda: \lim \bar{k} \Rightarrow \bar{k}$ be a limit cone. Then,
for any cone $\eta: D \Rightarrow \bar{k}$, $\forall i \in J$, we have.



We claim that G preserves
limit cone.

Let $\eta: D \Rightarrow G\bar{k} = k$ be a cone in Dn Graph.
Then, $\bar{\eta}: D \Rightarrow \bar{k}$ is a cone in Set^A
since the function $\bar{\cdot}$ does not change component
of natural transformations $\eta_i: D_i \rightarrow \bar{k}_i$, as a
sub category. Thus, $\exists! g$ s.t. above diagram holds.
By applying G ,



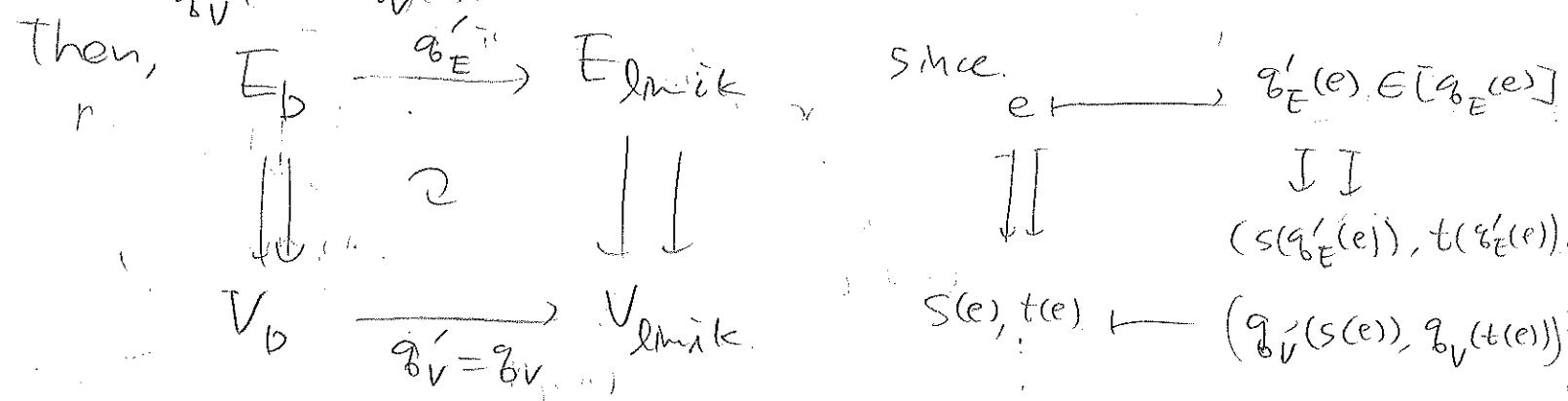
Moreover, such t is unique!

If there is another:
 $g': D \rightarrow G^F_i K_i$

St. left diagram commutes.

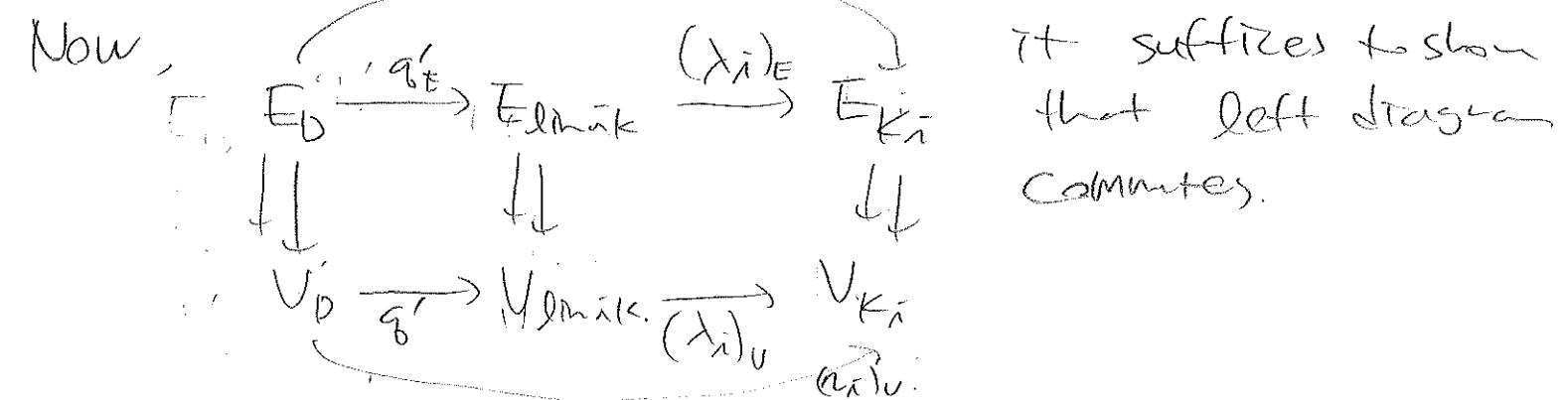
Then, define $g': D \rightarrow G^F_i K_i$ as a section:

i.e., $g'_E(x) = e'$ for some fixed representation $e' \in [e(x)]_E$.
 $g'_V(x) = g_V(x)$.



but $g'_E(e) \in [g_E(e)] \Rightarrow s(g'_E(e)) = s(g_E(e)) = g_V(s(e))$
 $t(g'_E(e)) = t(g_E(e)) = g_V(t(e))$

gives commutivity property. Hence g' is well-defined natural transformation. $(n_i)_E$



For $e \in E_p$,

$$e \mapsto g'_E(e) \in [g_E(e)] \implies (\lambda_i)_E(g'_E(e)).$$

Then, $s(g'_E(e)) \neq s(g_E(e))$, $t(g'_E(e)) = t(g_E(e))$

implies that $(\lambda_i)_E(g_E(e)) = (\lambda_i)_E(g'_E(e))$

Since K_1 is not a multigraph, so for any two ordered vertices, \exists at most 1 arrow.

By universal property, $(\lambda_i)_E(g_E(e)) = (\lambda_i)_E(e)$

$$\Rightarrow g'_E \circ (\lambda_i)_E = (\lambda_i)_E.$$

And by definition, $g'_U \circ (\lambda_i)_U = g_U \circ (\lambda_i)_U = (\lambda_i)_U$.

$$\Rightarrow g' \circ \lambda = \lambda.$$

The universal property shows that $g' = g$.

Thus, $g = Gg' = Gg$. $\therefore Gg$ is unique.

This shows that $G\lambda : G\text{Dm } \bar{A} \rightarrow G\bar{A}$
is a limit cone in Dm Graph .

$\Rightarrow G$ preserves limit cone.

Thus, $G(\text{lim } \bar{A}) \cong \text{lim } (G\bar{A}) = \text{lim } \bar{A}$.

$\Rightarrow \text{Dm Graph}$ admits limits if it is admissible set A
 $\Rightarrow \text{Dm Graph}$ is complete since set A is complete.

By the duality, doing the all argument dually,
we can get Dm Graph admits colimit if it is set A .

$\Rightarrow \text{Dm Graph}$ is cocomplete.

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Ex. 3.5, N) Since we proved Cat is complete/cocomplete w.r.t its set structure so we can make J^Δ concretely. Define $i_0: J \rightarrow J \times 2$, $j \mapsto (j, 0)$.

$$\text{Obj} \xrightarrow{!} \text{obj}^1.$$

i_0 is a pushout in Set,

$$\begin{array}{ccc} i_0 \downarrow & \quad \downarrow & \\ \text{Obj} \times 2 & \longrightarrow & \text{Obj}^0 \end{array} \Rightarrow \text{Obj}^{\Delta} = \text{Obj} \times 2 \amalg \frac{1}{\sim},$$

where $(j, 0) \sim 1 \forall j \in J$.

$$\Rightarrow \text{Obj}^{\Delta} \cong \{(j, 1) : j \in \text{Obj}\} \cup \{1\}.$$

Now by more forgetful functor,

$\text{Mor } J \xrightarrow{!} \text{Mor } 1$ induces that.

$$\begin{array}{ccc} i_0 \downarrow & \quad \downarrow & \text{Mor } J^{\Delta} = \text{Mor}(J \times 2) \not\sim \\ (\text{Mor}(J \times 2)) & \longrightarrow & \text{Mor } J^{\Delta} \end{array}$$

where $(f, 1_0) = 1_1$.

$$\begin{aligned} \text{Thus } \text{Mor } J^{\Delta} &\cong \{1 \xrightarrow{=} (j, 1) : \forall j \in \text{Obj}\} \\ &\cup \{f \circ \rightarrow 1 : f \in \text{Mor } J\} \\ &\cup \{f, 1_1 : f \in \text{Mor } J\}. \end{aligned}$$

Now, we claim that J^{Δ} is a cone over J with sumit 1 , by identifying $\{(j, 1)\}$ as base of J .

To see that, let $f: i \rightarrow j \in \text{Hom}(J)$. Then its image in J^2 is $(i, 1) \xrightarrow{(f, 1_1)} (j, 1)$. It suffices to show that

$$a \begin{array}{c} \swarrow \\ (i, 1) \end{array} \begin{array}{c} \searrow \\ (j, 1) \end{array} b \quad \text{commutes. where}$$

$$(i, 1) \xrightarrow{(f, 1_1)} (j, 1) \quad b \quad \parallel \quad (j, 0) \rightarrow (j, 1).$$

Since $(i, 0) \xrightarrow{(1_i, 0 \rightarrow 1)} (i, 1)$ in $J \times 2$,

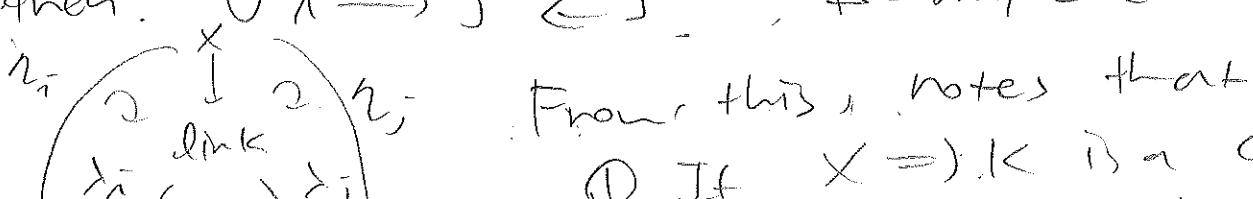
$$(f, 1_1) \downarrow \quad \quad \quad \downarrow (f, 1_2)$$

$$(j, 0) \xrightarrow{(1_j, 0 \rightarrow 1)} (j, 1) \quad \text{we have.}$$

$$(f, 1_1) \circ a = (f, 1_1) \circ (1_{i, 0 \rightarrow 1}) = (1_{j, 0 \rightarrow 1}) \circ (f, 1_0)$$

$$= b \circ f_2 = b. \quad \square$$

Ex. 3.5.v) Let $k: J \rightarrow ((N, \leq))$ be a small diagram. If $\lambda: \text{link} \Rightarrow k$ a limit cone exists, then $\forall \pi \rightarrow j \in J$, for any cone $\lambda: X \Rightarrow k$



From this, note that

① If $X \Rightarrow k$ is a cone,

then $X(F\pi) \xrightarrow{\cong} F(j)$.

② If link exists, then $X(\text{link})$.

Thus, link is $\text{g.c.d}(F\pi; \pi \in \text{ob } J)$.

Conversely, suppose $\chi = \text{g.c.d}(F\pi; \pi \in \text{ob } J)$.

Then, $\chi(F\pi) \xrightarrow{\cong} F(j)$ exists.

Since (N, \leq) is a poset as a category

these $\{\lambda_i : i \in \text{Obj}\}$ form a core., since every division commutes:

If $X \Rightarrow F$ is a core, then $X(F, \forall i =) X / Y$ since Y is gcd thus the above Init clause holds $\Rightarrow Y$ is l.h.s.

In case of colimit, every arrows are reversed,
so, by duality argument, colim k is l.c.m $(F_i, i \in J)$

Ex. 3.5. vi)

$F_{\text{Inj}_{\text{non}}}^{\text{op}} : \text{Obj}(\text{finite set}) \xrightarrow{\text{Morphisms: injections.}}$ Fix $X \in \text{Top}$.

$F_{\text{Inj}_{\text{non}}}^{\text{op}} \rightarrow \text{Top}$.

$[n] \longmapsto P(\text{Conf}_n(X)) : \{(x_1, \dots, x_n) \in X^n : x_i \neq x_j \forall i \neq j\}$
 $(x_{f(1)}, \dots, x_{f(n)})$ partial function

$f \downarrow \longmapsto \uparrow : \uparrow f^{-1} = [m] \xrightarrow{\text{for } i \longmapsto f^{-1}(i)}$

$[m] \longmapsto P(\text{Conf}_m(X)) : (x_1, \dots, x_m)$

this is well-defined since f^{-1} is just a projection map. (with permutation of coordinate.)

However, $[n] \longmapsto \text{Conf}_n(X)$ cannot be a functor!

since $[(x_1, \dots, x_n)] \longmapsto [(x_{f(1)}, \dots, x_{f(n)})]$ is not well-defined. For example: $f : \{1, 2\} \hookrightarrow \{1, 2, 3\}$ with
canonical injection cannot determine $[(4, 5, 6)] \longmapsto [(4, 5)]$
since $[(4, 5, 6)] = [(5, 6, 4)] = [(4, 6)] \dots \longmapsto [(5, 6)]$
 $\longmapsto [(4, 6)]$.

Ex 3.5. vii) a morphism.

Fiber Space : $p: E \rightarrow B \in \text{Top}$

Map between fiber spaces: a commutative square, i.e., $(p: E \rightarrow B) \longrightarrow (p': E' \rightarrow B')$ is $(g: E \rightarrow E', f: B \rightarrow B')$ s.t.

$$\begin{array}{ccc} E & \xrightarrow{g} & E' \\ p \downarrow & \cong & \downarrow p' \\ B & \xrightarrow{f} & B' \end{array}$$

thus, Category of fiberspace is \subseteq diagram category
 Top^2 .

And one of its non-full subcategory is

Top/B , a slice category over B .

i.e. objects are $p: E \rightarrow B$ fiber space over B
and morphisms are $E \xrightarrow{\cong} E'$ identified as $(g, 1_B)$ in Top^2 .

1) In a morphism,

$$\begin{array}{ccc} E' & \xrightarrow{g} & E \\ p' \downarrow & & \downarrow p \\ B' & \xrightarrow{f} & B \end{array}$$

let $p^{-1}(b)$ a fiber of $b \in B$.
Then, $g^{-1}(p'(b)) = (p')^{-1}(f^{-1}(b))$
by commutative diagram
 $\Rightarrow g^{-1}: p^{-1}(b) \longrightarrow f^{-1}(b)$
is a canonical map.

2) WTS

fiber of product of fiber spaces = product of fibers.

First of all we need to figure out product of fiber spaces in Top^2

Let $p: E \rightarrow B$, $p': E' \rightarrow B'$.

then product is $g: E'' \rightarrow B''$ with projection,

$\pi_p: g \rightarrow p$, $\pi_{p'}: g \rightarrow p'$ s.t.

$$\begin{array}{ccc} & s & \\ p & \swarrow \quad \downarrow \quad \searrow & \\ & g & \\ p' & \leftarrow \quad \rightarrow & p' \end{array}$$

$\forall s \in \text{Top}^2$

with $s \rightarrow p$, $s \rightarrow p'$.

Likewise. Ex 3.5 iv, we can generate g via set.

$$\begin{array}{ccccc} & E_s & & & \\ & \downarrow s & & & \\ E & \xleftarrow{(\pi_p)_E} & E'' & \xrightarrow{(\pi_{p'})_E} & E' \\ p \downarrow & & \downarrow z & & \downarrow f' \\ B & \xleftarrow{(\pi_p)_B} & B'' & \xrightarrow{(\pi_{p'})_B} & B' \end{array}$$

From left picture
which holds a set.
we know $E'' = E \times E'$
 $B'' = B \times B'$.
Then $g = (p, p')$.

Thus, the unique map $s \mapsto g$ is

$(E_s \rightarrow E'', B_s \rightarrow B'')$ a tuple of unique map

induced by E'' and B'' as a product.
Then it is just trivial to see that above diagram commutes.

Now, take a fiber in g over $(b, b') \in B''$.
Then it is $p^{-1}(b) \times (p')^{-1}(b')$ which is product of fibers.

(iii) Let $\pi_1: B \times F \rightarrow B$, then fiber of π_1 over $b \in B$ is $\pi_1^{-1}(b) = \{b\} \times F$.

As a subspace $\{b\} \times F$ is homeomorphic to F .
 Since it's finite product, box topology = product topology.
 So any basic open set in $B \times F$ is $O_B \times O_F$ form,
 thus $(\{b\} \times F) \cap (O_B \times O_F) = \{b\} \times O_F$ is basic open set
 in $\{b\} \times F$, thus $f: \{b\} \times F \rightarrow \mathbb{R}$ induces $f^{-1}(O_F) = \{b\} \times O_F$.
 If all open set $O_F \in \mathcal{F}$.
 $\Rightarrow \{b\} \times F \cong F$ in Top.

(iv). Let $\pi_E: B \times F \rightarrow B$, $\pi_G: B \times G \rightarrow B$ be two trivial fiber space. Then, if it is iso.

$$\begin{array}{ccc} \exists f, g \text{ s.t.} & \begin{array}{c} \pi_1 \nearrow^B \\ \downarrow f \\ B \times F \end{array} & \begin{array}{c} \pi_G \searrow^G \\ \downarrow g \\ B \times G \end{array} \end{array}$$

Let $\pi_G: B \times G \rightarrow G$
 $\pi_E: B \times E \rightarrow E$
 canonical morphism.

Then, $f(b, f) = (b, y)$ for some $y \in G$
 $g(b, y) = (b, y')$

$\Rightarrow f' = \pi_G \circ f$, $g' = \pi_E \circ g$ are OTS.

Conversely, we claim that given $f': B \times F \rightarrow G$
 and $g': B \times G \rightarrow E$, s.t. $g'(b, f'(b, x)) = x$

and $f'(b, g'(b, y)) = y$, this induces a map
 f and g satisfies the above diagram.

$$\text{Notes that } f := B \times F \xrightarrow{f'} D \times G$$

$$(b, x) \longmapsto (b, f'(b, x))$$

$$g := D \times G \longrightarrow B \times F$$

$$(b, y) \longmapsto (b, g'(b, y))$$

are cts; notes that for open set $O_B \times O_G \ni B \times G$,

$$f^{-1}(O_B \times O_G) = (f')^{-1}(O_G) \cap (O_B \times F)$$

Since if $f(b, x) \in O_B \times O_G$, then

$$f'(b, x) \in O_G \text{ and } b \in O_B$$

$$\Rightarrow (b, x) \in f^{-1}(O_G) \cap (O_B \times G)$$

Conversely, if $(b, x) \in (f')^{-1}(O_G) \cap (O_B \times F)$

$$\text{then } f(b, x) = (b, f'(b, x)) \in O_B \times O_G$$

So the equality holds.

$$\text{Moreover, } g(f(b, x)) = g(b, f'(b, x)) = (b, g'(b, f'(b, x)))$$

$$= (b, x)$$

$$\text{and } f(g(b, y)) = f(b, g'(b, y)) = (b, f'(b, g'(b, y)))$$

$$= (b, y)$$

$\Rightarrow f, g$ satisfy the diagram

(v) Let $p: E \rightarrow B$ a fiber space.

Then a section $s: B \rightarrow E$ is acts map

$$\text{S.t. } p \circ s = 1_B.$$

Now define $\text{Sect}: \text{Top}/B \longrightarrow \text{Set}$

$$p: E \rightarrow B \longmapsto \{\text{sections of } p\}$$

$$f: E \rightarrow E' \downarrow \longmapsto \begin{cases} f_* \text{ post comp.} \\ \end{cases}$$

$$p': E' \rightarrow B' \longmapsto \{\text{sections of } p'\}$$

then ① $\forall s: B \rightarrow E$ a fiber of $p: E \rightarrow B$,

$f \circ s: B \rightarrow E \rightarrow E'$ is a fiber of $p': E' \rightarrow B'$

Since $p' \circ f \circ s(x) = p \circ s(x) = x, \forall x \in B$.

thus, as a set map, f_* is well-defined.

② Functionality trivially holds by argument in ①

⇒ Sect is well-defined function.

(vi). Notes that

$$\text{Top}_{pB}^2: \text{obj}: p: E \rightarrow B \quad \text{any fiber space} \\ \text{Mor}: (p: E \rightarrow B) \xrightarrow{(f, g)} (p': E' \rightarrow B')$$

If the commutes square is actually

a pull back; i.e.

$$\begin{array}{ccc} E & \xrightarrow{f} & E' \\ p \downarrow & \square & \downarrow p' \\ B & \xrightarrow{g} & B' \end{array}$$

Now, $\text{Sect} : (\text{Top}_{\text{pb}}^2)^{\text{op}} \rightarrow \text{Set}$.

$p : E \rightarrow B \longmapsto \text{Section of } p$

$(f, g) \downarrow \quad \uparrow \phi$

$p' : E' \rightarrow B' \longmapsto \text{Section of } p'$

from the fact that for any $s' : B' \rightarrow E'$ of section of p' ,

$$\begin{array}{c} B \xrightarrow{s} E \xrightarrow{f} E' \\ \downarrow p \qquad \downarrow f' \\ B \xrightarrow{s} B' \xrightarrow{p'} E' \end{array}$$

$\exists! s : E \rightarrow B$ induced by $s' : B' \rightarrow E'$.
in Difft diagram.
and s is a section of p .

$\text{pos} = 1_B$.

thus, ϕ is well-defined as the unique map induced by 1_B and $(-\log)$ on the pullback of f and p .

Also, if $B \xrightarrow{(f, g)} P' \xrightarrow{(f', g')} P''$, then,

$B \xrightarrow{s_0} E \xrightarrow{f} E' \xrightarrow{f'} E''$ induces $s_0 : B \rightarrow E$.

$$\begin{array}{ccc} B & \xrightarrow{s_0} & E \\ \downarrow p & \downarrow f & \downarrow f' \\ B & \xrightarrow{s_1} & E' \\ \downarrow p'' & \downarrow f'' & \downarrow f'' \\ B & \xrightarrow{g} & B' \end{array}$$

s_0, s_1 are maps.

$\text{and } B \xrightarrow{s_2} E' \xrightarrow{f'} E'' \text{ and } B \xrightarrow{s_2 \circ g} B' \xrightarrow{g} B'$

From notes that

$$s'' \circ s' = f' \circ s_0, \quad s_2 \circ g = f \circ s_1.$$

$$\Rightarrow s'' \circ s' \circ g = f' \circ s_0 \circ g = (f' \circ f) \circ s_1 \text{ and } \text{pos}_0 = 1_B = ps_1.$$

thus universal property of pull back $E \rightarrow E''$
 shows that $S_0 = S_2$ $\downarrow \downarrow$
 $B \rightarrow B''$

This implies functoriality of Sect.

($I_p \hookrightarrow 1$ Sect(p) is trivial)

(VII) $E' \xrightarrow{l} E$ From $F' \rightarrow F$
 $P' \downarrow \quad \downarrow g$ $\downarrow \quad \downarrow$
 $F' \xrightarrow{r} F$ $B' \rightarrow B$
 $\downarrow g' \quad \downarrow p$ with $g \circ l : E' \rightarrow E$
 $B' \xrightarrow{f} B$ and $p' : E' \rightarrow B'$.

We set unique map $g' : E' \rightarrow F'$.

St. $g' \circ g = p'$ and $r \circ g' = lg$, from the
 universal property of pull back.

So, g and g' can be identified as a
 morphism in Top/B and Top/B' .

Also, (r, f) and (l, f) are morphisms
 in $\text{Top}^2_{B'}$. Now it suffices to show that.

$$\text{Sect}(p') \xrightarrow{\text{Sect}_{p'}(l, f)} \text{Sect}(F')$$

$$g_* \downarrow \quad \cap \quad \left\{ g'_*$$

$$\text{Sect}(g) \xrightarrow{\text{Sect}_{p'}(r, f)} \text{Sect}(g')$$

Let $s \in \text{Sect}(A)$. Then, $\text{Sect}_{pl}(l, f)(s) = s_g$.

$\xrightarrow{s \circ f}$ in the left diagram.

$$\begin{array}{ccccc}
 & & & & \\
 & B' & \xrightarrow{\exists! s_g} & E' & \xrightarrow{l} E \\
 & \downarrow & & \downarrow & \\
 & F' & \xrightarrow{\exists! s_g} & F & \xrightarrow{r} F' \\
 & \downarrow & & \downarrow & \\
 I_{B'} & B' & \xrightarrow{f} & B &
 \end{array}
 \quad \text{So, } g' \circ s_g : B' \xrightarrow{?} F'$$

$I_{B'}$ is in $\text{Sect}(g')$

On the other hand, $g \circ s \in \text{Sect}(g)$, thus,

$$\begin{array}{ccccc}
 & & & & \\
 & B' & \xrightarrow{\exists! s'_g} & F' & \xrightarrow{r} F \\
 & \downarrow & & \downarrow & \\
 & F' & \xrightarrow{\exists! s'_g} & F & \xrightarrow{r} F' \\
 & \downarrow & & \downarrow & \\
 I_{B'} & B' & \xrightarrow{f} & B &
 \end{array}
 \quad s'_g \in \text{Sect}(g')$$

by left diagram

$\Rightarrow s \mapsto s_g$ Now note that

$$\begin{array}{ccc}
 & & \xrightarrow{g \circ s_g} \\
 & \downarrow & \downarrow \\
 & s'_g & \xrightarrow{?} \\
 & \downarrow & \downarrow \\
 & B' & \xrightarrow{\exists! s'_g} & F' & \xrightarrow{r} F \\
 & \downarrow & & \downarrow & \\
 & F' & \xrightarrow{\exists! s'_g} & F & \xrightarrow{r} F' \\
 & \downarrow & & \downarrow & \\
 I_{B'} & B' & \xrightarrow{f} & B &
 \end{array}$$

Since $f' \circ g' \circ s_g = I_{B'}$ because $s'_g \circ s_g$ is a section.

And $r \circ g' \circ s_g = g \circ l \circ s_g = g \circ s \circ f$. Thus

$$\begin{array}{c}
 \text{from} \\
 \begin{array}{ccc}
 & E' & \xrightarrow{l} E \\
 & \downarrow & \\
 & F' & \xrightarrow{r} F
 \end{array}
 \end{array}$$

$$\begin{array}{c}
 \text{from} \\
 \begin{array}{ccc}
 & B' & \xrightarrow{\exists! s_g} & E' & \xrightarrow{l} E \\
 & \downarrow & & \downarrow & \\
 & F' & \xrightarrow{\exists! s_g} & F & \xrightarrow{r} F' \\
 & \downarrow & & \downarrow & \\
 I_{B'} & B' & \xrightarrow{f} & B &
 \end{array}
 \end{array}$$

universal
properties
says that
 $s'_g = g' \circ s_g$.

3.6. Functionality of limit and colimits.

Functionality \neq canonicity.

Lots of non-unique objects can define limit.

Prop 3.6.1. Suppose \mathcal{C} has all J -shape limits.
Then choice of limit for each diagram

defines $\lim_J : \mathcal{C}^J \rightarrow \mathcal{C}$ as a functor.

Part 1) \lim_J is not canonically defined.

pf) By choice of limit and limit cone, we define \lim_J .
on $\text{Ob } \mathcal{C}$. To define \lim_J on morphism,

Suppose $F, G \in \mathcal{C}^J$ and $\lambda : \lim_J F \Rightarrow F$ are
chosen limit cone,
 $\alpha : \lim_J G \Rightarrow G$.

then if $\alpha : F \Rightarrow G \in \text{Hom}(F, G)$, then
 $\alpha_i : F_i \Rightarrow G_i$

$$\begin{array}{ccc} \lambda_i : \lim_J F_i & \xrightarrow{\quad} & \alpha_i \circ \lambda_i \text{ is a cone over } G_i \\ F_i \xrightarrow{\quad} & \downarrow & \text{thus } \exists ! k : \lim_J F_i \rightarrow \lim_J G_i \\ \alpha_{ij} : F_i \xrightarrow{\quad} & \downarrow & \text{s.t. } \alpha_i \circ \lambda_i = \lambda_j \circ k \\ G_i \xrightarrow{\quad} & \downarrow & \end{array}$$

thus, define $\lim_J(\alpha) := k$
The uniqueness of k implies functionality.

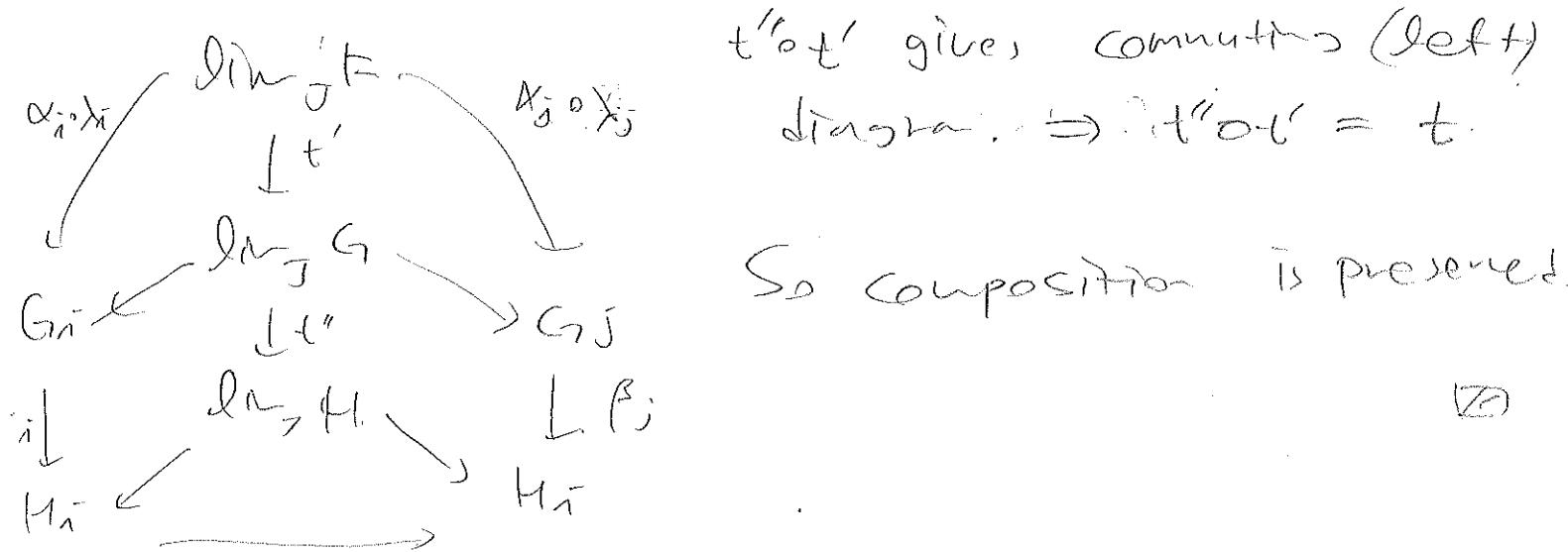
Notes that $1_F : F \Rightarrow F$ is sent to $1_{\lim_J F}$.

And for $F, G \xrightarrow{\alpha} H$, $\lim_J F$ we get to

$$\begin{array}{ccc} \beta_{F,G,H} / \text{It is } \beta_{F,G,H} & \xrightarrow{\quad} & \text{by universal property} \\ \lim_J F & \xrightarrow{\quad} & \\ H_i \xrightarrow{\quad} & \xrightarrow{\quad} & H_j \end{array}$$

And if $t': \text{Im}_J F \rightarrow \text{Im}_J G$, $t'': \text{Im}_J G \rightarrow \text{HJ}$.

derived from universal property of limit,



Ex 3.6.2 Path : Top \rightarrow Set $\cong \text{Top}(I, -)$
 $X \mapsto \text{Top}(I, X)$. with $I = [0, 1]$

Let $0, 1, * \rightrightarrows I$ endpoints. Define

$P: \text{Top} \xrightarrow{\quad} \text{Set}$

$$X \mapsto \left(\text{Top}(I, X) \xrightarrow[\text{ev}_1]{\text{ev}_0} \text{Top}(t, X) \right)$$

(1) (2)

$$\text{Path}(X) \xrightarrow{\quad} P(X)$$

This is a functor. And its coequalizer is.

$$\text{Path}(X) \xrightarrow[\text{ev}_1]{\text{ev}_0} P(X) \longrightarrow \pi_0(X)$$

the set of path components of X , since $\pi_0(X)$ is quotient of $P(X)$ by equivalence relation s.t.
 $x \sim y$ if x and y are path connected.

$\Rightarrow \pi_0: \text{Top} \xrightarrow{P} \text{Set} \xrightarrow{\quad} \text{Set}$ is well-defined
 by Prop 3.6.1. (We just choose specific coequalizer $\pi_0(X)$)

(Notes that any set with the same cardinality with $T_0(Y)$ can be chosen. But we choose \mathbb{N} , and it defines \mathcal{X}_0 as a functor.)

Other choice gives another functor.)

Rmk) Prop 3.6.1 shows any natural transformation in \mathcal{C}^J are sent to morphism between limits (or colimits) regardless of complete/completeness of domain category.

Smallary 3.6.3) A natural iso between diagram induces a naturally defined iso between their limits or colimits whenever they exists.

b) Lem. 3.6.1 can be deduced from the full subcategory of \mathcal{C}^J spanned by functors having limits in \mathcal{C} . Now apply Lem 1.3.8. (all functors preserve iso.)

Rmk) Cor 3.6.3 is actually motivation of Eilenberg and MacLane.

Ex 3.6.4) G : sp. $X, Y : BG \Rightarrow \text{Set}$

X and Y are objectswise isomorphic $\Rightarrow X \cong Y$ as a set.
" naturally " \Rightarrow There exists

an equivariant bijection $t : X \xrightarrow{\cong} Y$ st.

$$\begin{array}{ccc} X & \xrightarrow{t} & Y \\ \exists f & \exists g & \\ x & \xrightarrow{f} & y \end{array}$$

Let $\{a, b\}$ has the nontrivial \mathbb{Z}_2 action by
sending $a \xrightarrow{\text{act}} b$.

and $\{c, d\}$ has $=$ 64.

sends $c \xrightarrow{\text{act}} d$

Then $\{a, b\} : B(\mathbb{Z}_2) \rightarrow \text{Set}$ are obcture β .
 $\{c, d\} : " \rightarrow \text{Set}$ but not natural iso.

since there is no function $f : \{a, b\} \rightarrow \{c, d\}$
satisfying natural iso.

(p1) If $f(b) = c$, then $\begin{array}{c} b \xrightarrow{\text{act}} c \\ \downarrow \cancel{x} \quad \downarrow \\ b \xrightarrow{\text{act}} d \end{array}$

If $f(b) = d$, then $\begin{array}{c} b \xrightarrow{\text{act}} d \\ \downarrow \cancel{x} \quad \downarrow \\ b \xrightarrow{\text{act}} c \end{array}$

Or in viewpoint of limit, $\lim B(G \rightarrow X)$ is X^G .

$\lim \{a, b\} = \{a\}$ $\lim B(G \rightarrow X)$ is G -orbit sets.

$\lim \{c, d\} = \emptyset$, $\lim \{c, d\} = \{*\}$.

Ex. 3.6.5. X : finite set.

$\text{Sym}(X)$: all permutations of $X \Rightarrow |X|!$
elements

$\text{Ord}(X)$: all total ordering on X .

But $\text{Sym}(F_m)$ \rightarrow F_n acts by conjugation.

$\text{Ord} : \{ \cdot \} \rightarrow \{ \cdot \}$ acts by translation.

ex) $(132) \in \text{Sym}(\{1, 2, 3\})$ $| \geq 3 \geq 2$ $\text{Ord}(\{1, 2, 3\})$

$\begin{array}{c} \boxed{1} \\ (bc)a \end{array} \in \text{Sym}(\{b, a, c\})$ $| \dots$ $\text{Ord}(\{b, a, c\})$

But this is not natural. (even objective iso?)
⇒ limit (comit) need not be iso.

Species := a functor $F: \text{Fin}_{\text{iso}} \rightarrow \text{Fin}$.

labeled F-structures on n := $F(n)$.

unlabeled " := \lim_{isom} of

$$\begin{array}{c} \text{BI}_n \hookrightarrow \text{Fin}_{\text{iso}} \xrightarrow{F} \text{Fin} \\ \circ \longmapsto [n] \longrightarrow F(n) \end{array}$$

objective iso \Rightarrow labeled Sym-structure

= labeled Ord =

But unlabeled Sym-structure: set of conjugacy classes of permutations of n elements

(Since any bijection is a permutation)

unlabeled Ord-structure: trivial.

Since all linear orders on n are isomorphic.

Ques) ① Even limit is unique, there might be distinct limit cons.

② choice of limit cone is usually not compatible.

ex) Freyd shows that C^{WDF} pullback has a "canonical pullback" with "horizontal composite". But it seems not possible to find such canonical one for "vertical composite".

③ Even dozen limit objects are equal, natural iso might not be identity.

Lemma 3.6.6 For any triple $Y, Y, Z \in C$
where C has binary product,

$\exists ! \alpha : X \times (Y \times Z) \rightarrow (X \times Y) \times Z$ natural
is. Commutes with projections to X, Y, Z .

$$\text{pf). } X \times (Y \times Z) \xrightarrow{\pi_{1,X}} X \quad \text{The universal property of } \\ \text{int } X \times Y \text{ induces the unique} \\ \pi_{1,Y \times Z} \downarrow \qquad \qquad \qquad t : X \times (Y \times Z) \rightarrow X \times Y \\ Y \times Z \xrightarrow{\pi_{1,Y}} Y \qquad \qquad \qquad \pi_{2,Y} \downarrow \qquad \qquad \qquad Y \\ \pi_{1,Z} \downarrow \qquad \qquad \text{s.t.} \qquad \qquad \qquad X \\ Z \qquad \qquad \qquad \qquad \qquad \pi_{1,X} = \pi_{2,X} \circ t \\ \pi_{1,Y} \circ \pi_{1,Y \times Z} = \pi_{2,Y} \circ t.$$

Now, t and $\pi_{1,Z} \circ \pi_{1,Y \times Z}$ induce,

$$\alpha : X \times (Y \times Z) \rightarrow (X \times Y) \times Z \quad \begin{matrix} \pi_{3,Y} \downarrow & \pi_{3,Z} \downarrow \\ \alpha \circ \pi_{1,Y} & Z \end{matrix} \\ \text{s.t.} \\ t = \pi_{3,X \times Y} \circ \alpha.$$

$$\pi_{1,Z} \circ \pi_{1,Y \times Z} = \pi_{3,Z} \circ \alpha.$$

Conversely, $\pi_{2,Y} \circ \pi_{3,X \times Y}$ and $\pi_{3,Z}$ induce,

$$s : (X \times Y) \times Z \rightarrow Y \times Z \quad \begin{matrix} \pi_{1,Y} \downarrow & \pi_{1,Z} \downarrow \\ Y & Z \end{matrix}$$

$$\text{s.t. } \pi_{2,Y} \circ \pi_{3,X \times Y} = \pi_{3,Y} \circ s.$$

$$\pi_{3,Z} = \pi_{4,Z} \circ s.$$

Also, s and $\pi_{2,X} \circ \pi_{3,X \times Y}$ induce,

$$\beta : (X \times Y) \times Z \rightarrow X \times (Y \times Z) \quad \text{s.t. } \pi_{1,X} \circ \beta = \pi_{2,Y} \circ \pi_{3,X \times Y} \\ \pi_{1,Y \times Z} \circ \beta = s.$$

$$\text{Then, } \pi_{3,XXY}(\alpha \circ \beta) = t \circ \beta : (XY) \times Z \rightarrow (XY)$$

$$\text{Since } \pi_{2,X}(t \circ \beta) = \pi_{1,X}(\beta) = \pi_{2,X} \circ \pi_{3,XY}$$

$$\begin{aligned} \pi_{2,Y}(t \circ \beta) &= \pi_{1,Y} \circ \pi_{1,Y \times Z} \circ \beta = \pi_{1,Y} \circ s = \pi_{3,XY} \\ &= \pi_{2,Y} \circ \pi_{3,XXY} \end{aligned}$$

$$\Rightarrow t \circ \beta = \pi_{3,XXY} \text{ by the universal property of } X \times Y$$

$$\Rightarrow \pi_{3,XXY}(\alpha \circ \beta) = \pi_{3,XXY}$$

$$\text{Also, } \pi_{3,Z}(\alpha \circ \beta) = \pi_{1,Z} \circ \pi_{1,YZ} \circ \beta \\ = \pi_{1,Z} \circ s = \pi_{3,Z}$$

$$\Rightarrow \alpha \circ \beta = 1_{(XY) \times Z} \text{ by the univ. property of } (XY) \times Z$$

Conversely,

$$\pi_{1,X}(\beta \circ \alpha) = \pi_{2,X} \circ \pi_{3,XXY} \circ \alpha = \pi_{2,X} \circ t = \pi_{1,X}$$

$$\pi_{1,Y \times Z} \circ (\beta \circ \alpha) = s \circ \alpha \text{ and}$$

$$\pi_{1,Y} \circ (s \circ \alpha) = \pi_{2,Y} \circ \pi_{3,XXY} \circ \alpha = \pi_{2,Y} \circ t = \pi_{1,Y} \circ \pi_{1,Y \times Z}$$

$$\pi_{1,Z} \circ (s \circ \alpha) = \pi_{3,Z} \circ \alpha = \pi_{1,Z} \circ \pi_{1,Y \times Z}$$

$$\Rightarrow \pi_{1,Y \times Z} = s \circ \alpha \text{ by universal property}$$

$$\text{thus } \pi_{1,Y \times Z} \circ (\beta \circ \alpha) = \pi_{1,Y \times Z}$$

$$\Rightarrow \beta \circ \alpha = 1_{X \times (Y \times Z)} \text{ by the universal property}$$

$\Rightarrow \alpha, \beta$ are isomorphisms in C .

This is natural iso if we see this. Let $A = \Sigma_2 - 3$.

$$\text{Let } (X, \dashrightarrow) : \mathcal{C} \rightarrow \mathcal{C}^A.$$

Category
 without nontrivial
 morphism.

$$Y \dashrightarrow (X, Y) : A \rightarrow \mathcal{C}$$

$$\downarrow \dashrightarrow \downarrow$$

$$Z \dashrightarrow (X, Z) : A \rightarrow \mathcal{C}.$$

Then, $Y \dashrightarrow Z$ induces natural transformation between (X, Y) and (X, Z) since A has no nontrivial morphism.

Then, $(X, \lim_A(Y, -)) \xrightarrow{\alpha} (\lim_A(X, Y), -)$
 by our construction of α, β .

∴ Natural! □

Ex. 3.6, 7. $\text{sk}(\text{Set})$ a skeletal category of Set .
 Then actually for each cardinal k , there is only one object in $\text{sk}(\text{Set})$.

$\text{sk}(\text{Set})$ is complete; (Just use isomorphisms for finite sets)

\Rightarrow It has a product.

Let C : ctble infinite $\Rightarrow \text{Ob}(C = C)$.

as $\pi_1, \pi_2 : C \rightarrow C$ are epic.

If $\text{Ob}(C \times C) \xrightarrow{i} (\text{Ob}C) \times C$ is identity, then.

$$\text{Ob}(C \times C) = ((\text{Ob}C) \times C) \quad \text{for my } f, g, h : C \rightarrow C.$$

$$\begin{aligned} f \times (g \times h) &\vdash \vdash \vdash (f \times g) \times h \\ \text{Ob}(C \times C) &= ((\text{Ob}C) \times C) \end{aligned}$$

Now, we set below commutative diagram

$$\begin{array}{ccc}
 C & \xrightarrow{\quad g \times h \quad} & C \\
 \pi_2 \downarrow & & \uparrow \pi_2 \\
 (X_C) & \xrightarrow{f \times g} & (Y_C) \\
 \pi_1 \downarrow & f & \downarrow \pi_1 \\
 C & \xrightarrow{\quad f \times g \quad} & C
 \end{array}
 \Rightarrow \text{Since } \pi_1, \pi_2 \text{ are epi, } f = f \times g, \quad g \times h = h.$$

$f, g : C \rightrightarrows C$. equal.
 Since $\begin{array}{c} f \times f \\ \parallel \\ C \xrightarrow{f \times g} C \end{array}$
 $\therefore f = g.$

But this is absurd! since $\text{sk}(\text{Set}) (C, C) = \text{Set}(C, C)$

Ex 3.6.1) Special case of prop 3.6.1

((iff) shown in the above.)

3.7. Size Matters

Lem 3.7.1. C : category. $I_C : C \rightarrow C$ admits limit $\Leftrightarrow C$ has initial object.

Pf) $(\lambda_c : I \rightarrow C)_{c \in C}$ a limit cone.

We can say that I is "weakly initial object"

In the sense that $\forall c \in C, \exists \lambda_c : I \rightarrow c$.

To see λ_c is only morphism, let $f : I \rightarrow c$.

Then, $I \xrightarrow{f} c \Rightarrow f \lambda_c = \lambda_c$. So WTS λ_c is identity.

$$\begin{array}{ccc}
 \lambda_l & \swarrow & \searrow \lambda_c \\
 I & \xrightarrow{f} & c
 \end{array}$$

To see (iii), we have $\lambda \circ \lambda_c : (\lambda_c \circ \lambda_c) \rightarrow \lambda_c$

thus, $\lambda \circ \lambda_c : (\lambda_c \circ \lambda_c) \rightarrow \lambda_c$.

It is another cone. By the universal property of λ_c , we have

$$\begin{array}{ccc} & \lambda & \\ \lambda_c \circ \lambda & \swarrow \quad \downarrow & \searrow \lambda_c \circ \lambda \\ & \lambda_c = \lambda_c & \end{array}$$

since both λ_c and λ
make the left diagram
commutes,

$$\Rightarrow f = \lambda_c.$$

Conversely, if C has initial element, l , then, $\lambda_c : l \rightarrow c$ form a cone over I_c , and it is limit cone. Let $\eta : t \rightarrow I_c$.

Then, $\forall f : d \rightarrow c \in C$, $\eta_d \circ f = \eta_c$. This implies,
 $\eta_d \circ f = \eta_c$.
 $\Rightarrow \eta_d \circ f = \eta_c$. $\eta_c \Rightarrow \eta_d$. $\eta_d : t \rightarrow \lambda_c$ exists.

So, λ is a limit cone. □

Def 3.2.7. Cardinality of a small category.

Ref 3.2.7. The set of its morphism

a K -small category = a category whose cardinality is less than K .

k -small diagram: a diagram whose indexing category is k -small.

Prop 3.17.3 (Freyd) Any k -small category admits all k -small limits or k -small colimits is a preorder.

Pf). Let C be a k -small category with $(C = \lambda)$.
 $f, g: D \rightarrow A \in \text{Mor}(C)$ with $f \neq g$.

By Ex 3.4.4, $B \rightarrow A^\lambda$ is generated by λ indexed morphisms $B \rightarrow A$.

\Rightarrow The collection of all $B \rightarrow A^\lambda$ whose component is either f or g has cardinality 2^λ .

$$|C(B, A^\lambda)| \geq 2^\lambda > \lambda = |\text{Mor}(C)|.$$

↑ from Cantor's diagonalization argument.

But this contradicts the fact that $C(B, A^\lambda) \subseteq \text{Mor}(C)$.

\Rightarrow Only hom set in C is singleton.

$\Rightarrow C$ is preorder.

Do the same thing using cotensor, for k -complete case.

2nd) Prop 3.5.9 shows that there is a preorder to have all limit/colimit of unindexed size.

Since P poset has this property

Ex 3.17.1) done in the above. $\hookrightarrow P = \text{Complete lattice}$.

8. Interaction between limits and colimits.

Let $F: I \times J \rightarrow C$. By Ex 1.7. VII), F can be regarded as $F: I \rightarrow C^J$ or $F: J \rightarrow C^I$.

If $\forall i \in I, \exists \lim_{j \in J} F(i, j)$, then Prop 3.3.9

gives $\lim_{j \in J} F(-, j): I \rightarrow C$. Moreover,

Thm 3.8.1. If $\lim_{i \in I} \lim_{j \in J} F(i, j)$, $\lim_{i \in I} \lim_{j \in J} F(i, j)$ exist.

then there are iso and define the $\lim_{I \times J} F$.

Pf) From Yoneda Lemma, it suff to prove.

$$((X, \lim_{i \in I} \lim_{j \in J} F(i, j))) \cong ((X, \lim_{(i, j) \in I \times J} F(i, j))) \cong ((Y, \lim_{(i, j) \in I \times J} F(i, j)))$$

By Thm 3.4.6(i), covariant representable functor preserves limit.

$$\Rightarrow ((X, \lim_{i \in I} \lim_{j \in J} F(i, j))) \cong \lim_I ((X, \lim_{j \in J} F(i, j))) \cong \lim_I \lim_{j \in J} ((X, F(i, j)))$$

$$((X, \lim_{(i, j) \in I \times J} F(i, j))) \cong \lim_{(i, j) \in I \times J} ((X, F(i, j)))$$

$$((X, \lim_{j \in J} \lim_{i \in I} F(i, j))) \cong \lim_J \lim_I ((X, F(i, j)))$$

thus, it suffices to show that $\lim_I \lim_{j \in J} H: I \times J \rightarrow \text{Set}$ a set valued functor, the limits.

$$\lim_I \lim_{j \in J} H: \cong \lim_{I \times J} H \cong \lim_I \lim_{j \in J} H \text{ is o.}$$

Is this iso of category $I \times J \cong J \times I$,
only left iso is needed to show.

By 3.2.3, let $\lambda: \mathbb{A} \Rightarrow \lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \in \lim_{\substack{\text{I} \\ \rightarrow}} \lim_{\substack{\text{J} \\ \rightarrow}} H$.

Then, $\begin{array}{ccc} 1 & \xrightarrow{i} & i' \\ \downarrow & \nearrow & \searrow \\ 2 & & \end{array} \quad \forall i \rightarrow i' \in I$
 $\lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \rightarrow \lim_{\substack{\text{J} \\ \rightarrow}} H(i', j)$

By prop 3.6.1, since we already chosen limit one,

at $\lim_{\substack{\text{J} \\ \rightarrow}} H(i, j)$, $\forall i \in I$, $i \rightarrow i'$ also a natural
transformation of $H(i, -) \xrightarrow{\sim} H(i', -)$.

determines $\lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \rightarrow \lim_{\substack{\text{J} \\ \rightarrow}} H(i', j)$.

Denote $(\eta_i): \lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \Rightarrow H(i, -)$.

Then, $\forall i \in I, \phi_{i,j}: 1 \xrightarrow{i} \lim_{\substack{\text{I} \\ \rightarrow}} H(i, -) \xrightarrow{\eta_{i,j}} H(i, j)$

is a cone over $F(I \times J) \rightarrow \mathcal{C}$; To see this
notes that $\begin{array}{ccc} 1 & \xrightarrow{i} & i' \\ \downarrow & \nearrow & \searrow \\ 2 & & \end{array} \quad \forall f: i \rightarrow i'$,

$\lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \rightarrow \lim_{\substack{\text{J} \\ \rightarrow}} H(i', j)$ commutes by prop 3.6.1.

$\begin{array}{ccc} 1 & \xrightarrow{i} & i' \\ \downarrow & \nearrow & \searrow \\ 2 & & \end{array} \quad \int \eta_{i,j}: H(i, j) \xrightarrow{\sim} H(i', j)$

$H(f, -) \xrightarrow{\sim} H(f, j)$ (Actually, $\lim_{\substack{\text{I} \\ \rightarrow}} H(i, j) \rightarrow \lim_{\substack{\text{J} \\ \rightarrow}} H(i', j)$)

Thus, over is given by the universal property of

$\begin{array}{ccc} 1 & \xrightarrow{i} & i' \\ \downarrow & \nearrow & \searrow \\ 2 & & \end{array} \quad \text{cone } H(i, -) \Rightarrow H(f, -) \circ \eta_{i,j} = k \circ \eta_{i',j} \quad)$

$\Rightarrow H(i, j) \rightarrow H(i', j) \rightarrow H(i', j')$ (Last triangle is from η .)

This shows that $\phi: I \Rightarrow F \in \text{LinF}_{\overline{J} \times \overline{J}}$.

Conversely, if $\phi \in \text{LinF}_{\overline{J} \times \overline{J}}$, then, for fixed i ,

$\phi_i: I \Rightarrow H(i, -)$, thus, $\phi_i \in \lim_{\overline{J}} H(i, j)$

So we get $\lambda_i: I \rightarrow \lim_{\overline{J}} (H(i, j))$

$$I \mapsto \phi_i$$

Hence, this makes $\lambda: I \Rightarrow \lim_{\overline{J}} H(i, -)$. (elements, colimit)

These show that $\lim_{\overline{I}} \lim_{\overline{J}} H \cong \lim_{\overline{I} \times \overline{J}} H$.

By duality, column case also holds. \square

Rmk). This shows that \exists natural iso.

$$(C^{\overline{I} \times \overline{J}} \xrightarrow{\text{Im}_{\overline{J}}}, C^{\overline{I}}) \cong (C^{\overline{I} \times \overline{J}} \xrightarrow{\text{colim}_{\overline{J}}}, C^{\overline{I}})$$

$$\begin{array}{ccc} \text{Im}_{\overline{I}} & \cong & \text{colim}_{\overline{J}} \\ \downarrow & \cong & \downarrow \\ C^{\overline{J}} & \xrightarrow{\text{Im}_{\overline{J}}} & C^{\overline{I}} \end{array} \quad \begin{array}{ccc} \text{colim}_{\overline{I}} & \cong & \text{colim}_{\overline{J}} \\ \downarrow & \cong & \downarrow \\ C^{\overline{I}} & \xrightarrow{\text{colim}_{\overline{J}}} & C^{\overline{J}} \end{array}$$

for any choice of $\text{Im}_{\overline{I}}$ and $\text{colim}_{\overline{J}}$, only.

Lemma 3.8.3. $F: \overline{I} \times \overline{J} \rightarrow C$ bi-fuctor s.t.

$\text{Im } F, \text{colim } F$ exist.

$\Rightarrow \exists k: \text{colim } \text{Im } F(i, j) \rightarrow \text{Im } \text{colim } F(i, j)$

$$\begin{matrix} i \in \overline{I} & j \in \overline{J} \end{matrix}$$

$$\begin{matrix} j \in \overline{J} & i \in \overline{I} \end{matrix}$$

a canonical map exists.

pf) By the univ. prop. of colimit, if k is canonical, then \exists

$$k_i : \left(\underset{j \in J}{\text{lim}} F(i, j) \longrightarrow \underset{\bar{j} \in J}{\text{colim}} F(i, \bar{j}) \right)_{i \in I}$$

that defines a cone.

By the univ prop of limit, k_i is defined by

$$(k_{i,j} : \underset{j \in J}{\text{lim}} F(i, j) \rightarrow \underset{i' \in I}{\text{colim}} F(i, j))_{j \in J}$$

that define a cone over $\underset{i \in I}{\text{colim}} F(i, -)$.

So, define $k_{\bar{i}, j} : \underset{j \in J}{\text{lim}} F(\bar{i}, j) \xrightarrow{\pi_{\bar{i}, j}} F(\bar{i}, j) \rightarrow \underset{\bar{i}' \in I}{\text{colim}} F(\bar{i}', j)$

where $\pi_{\bar{i}}, i_j$ are limit and colimit cone over $F(\bar{i}, -)$ and $F(-, j)$ respectively.

Then, $k_{\bar{i}, j}$ is a cone over $\underset{\bar{i}' \in I}{\text{colim}} F(\bar{i}', j)$!

thus we have $k_{\bar{i}, j}$; thus k by above argument. \square

Corollary 3.8.4 $f : X \times Y \rightarrow \mathbb{R}$ a function

$$\Rightarrow \sup_{x \in X} \inf_{y \in Y} f(x, y) \leq \inf_{y \in Y} \sup_{x \in X} f(x, y)$$

(f) In a poset category, $\text{colim}_X = \sup_X$, $\text{inf}_X = \inf_X$. \square

Buk. 3.8.5. Direct proof of the claim, the limit analysis is actually categorical.

Pf) $\sup_x \inf_y f$ is bdd \Leftarrow $\inf_x f(x, -)$ bdd above
above by M. $\forall x \in X$
by N.

$\inf_y \sup_x f$ is bdd \Leftarrow $\sup_x f(-, y)$ bdd below by N.
below by N. $\forall y \in Y$

Then we have $\inf_x f(x, -) \leq f(x, y) \leq \sup_x f(-, y)$

(Ex) 3.8.6. Let $\widehat{\mathbb{R}} = (\mathbb{R} \cup \{-\infty, \infty\})$ with $-\infty < x < \infty$ poset.

$\Rightarrow \widehat{\mathbb{R}}$ is complete and ∞ complete. is a category for poset.
any function gives diagram, and its \lim_{\leftarrow} = \inf_{lim}
 $\text{co } \alpha = \sup_{\text{lim}}$.

Now let $x: \mathbb{N} \rightarrow \widehat{\mathbb{R}}$ be a sequence.

Then, we have: $\mathbb{N} \times \mathbb{N} \xrightarrow{+} \mathbb{N} \xrightarrow{x} \widehat{\mathbb{R}}$
 $(m, n) \longmapsto m+n \longmapsto x_{m+n}$.

So. $\liminf x_n = \sup_{n \geq 0} \inf_{m \geq n} x_m = \sup_{n \geq 0} \inf_{m \geq n} x_{n+m} = \text{colim}_{(n, k)} x_n$

By Lec. 3.8.3 $\Rightarrow \liminf x_n \leq \limsup x_n$.

In Set and some related categories there exists
an invertible k .

Def: 3.8.7. J is filtered if there is a cone.

$J' \rightarrow J$ & J' finite index \Rightarrow category

Ex) $W, \circ \rightarrow \circ$ (filtered), $\circ \leftarrow \circ \rightarrow \circ, \circ \rightrightarrows \circ$, discrete: not filtered category

Thm 3.8.9. Filtered colimit commute with finite limit in Set

Observation: $F: J \rightarrow \text{Set}$, J is small filtered category. Then, from diagram in p. 97.

$$\text{colim}_{j \in J} F_j = \coprod_{j \in J} F_j / \sim \quad \text{where } x \sim y \text{ for } x \in F_j, y \in F_k$$

$$\Leftrightarrow \exists f: j \rightarrow t \text{ s.t. } Ff(x) = Ff(y)$$

$$g: k \rightarrow t. \quad = Ff(y).$$

Thus, let A be a finite subset of $\coprod_{j \in J} F_j$ s.t. $A/\sim := \{\bar{x}_i\}_{i \in I}$, where I is finite. Then, $\exists t \in J$, s.t. $A \subseteq F_t$ and $A/\sim = A_t$.

(*) A is a finite subset of $\coprod_{j \in J} F_j$, so

$\exists \{j \in J : F_j \cap A \neq \emptyset\}$ is a finite set. Let \overline{B} be a full subcategory spanned by B . By definition, $\overline{B} \hookrightarrow J$ has a cocone with nadir, say t . Thus, by sending $F_j \cap A \rightarrow F_t$ using one map, we get such A' .

(Pf of Thm). WTS $k: \varinjlim_J \varprojlim_I F \rightarrow \varprojlim_I \varinjlim_J F$

\Rightarrow a bijection when I : finite, J : small filtered category

By def 3.2.3, as a set,

$$\varinjlim_J \varprojlim_I F = \{ \lambda : \lambda = \varprojlim_I F(-, J), \exists$$

Since I is finite, $\lambda = \{\lambda_i\}_{i \in I}$ is finite.

Also, only finitely many compatibility conditions satisfy.

Then, each $\lambda_i : I \rightarrow \operatorname{colim}_J F(i, J)$ as a element
of $\operatorname{colim}_J F(i, J)$. can be viewed

Now fix $i \in I$, $t \in J$. Choose $\lambda'_i \in F(i, t)$ s.t.
as equivalence class. Then, if $i' \in I$ and $i \sim i'$,
Since $\operatorname{colim}_J F(f, J)(\lambda_i) = \lambda'_i$,
 $F(f, J)(\lambda'_i) \in \lambda_{i'}$ as an equivalence
class.

Thus we can choose $\lambda'_i \in F(i', t)$ s.t.

$$\begin{array}{ccc} \lambda'_i & \xrightarrow{1} & \lambda'_i \\ \downarrow & \cap & \downarrow \\ F(i, t) & \xrightarrow[F(f, t)]{} & F(i', t) \end{array}$$

So if I is connected, then
we have a cone $\{X : I \rightarrow F(t, t)\}$.

Otherwise, if i'' is not connected with i in I ,
then still $\lambda'_{i''} \in F(i'', t')$ for some t' s.t.

$[\lambda'_{i''}] = \lambda''_i$. Since J is filtered, \exists a one

under t ad t' , say t'' , i.e. $t \xrightarrow{\cong} t'' \xrightarrow{\cong} t'$.

\Rightarrow we may regard λ'_i as its related λ''_i
and λ''_i as an element of $F(i, t'')$, $F(i', t'')$ and
 $F(i'', t'')$ respectively via the map π_1, π_2 .

Thus, we can say that $\exists t \in J$ s.t.

$X : I \rightarrow F(\bar{i}, t)$ and $[\lambda_i] = \lambda_{\bar{i}}$ $\forall \bar{i} \in I$.

Now this X' is inside of $\lim_{\lambda \in I} F(\bar{\lambda}, t)$.

By p. 97 diagram, dank is

$$\text{colim}_{\bar{\lambda} \in J} \lim_I F(\bar{\lambda}, i) = \frac{(\lim_{\bar{\lambda} \in J} \lim_I F(\bar{\lambda}, i))}{n}.$$

thus, $[X']$ is in dank and by

construction, k is induced by $k_{j, \bar{\lambda}} : \lim_I F(\bar{\lambda}, j) \rightarrow F(\bar{\lambda}, j)$

and by setting $j = t$, we set $k_{j, \bar{\lambda}}(X) = [X']$
in $\text{colim}_{\bar{\lambda} \in J} F(\bar{\lambda}, j)$.

$\Rightarrow k$ is surjective.

Now, to see k is injective,

let $(1 \xrightarrow{\alpha_i} F(\bar{\lambda}, i))_{\bar{\lambda} \in I}$, $(1 \xrightarrow{\beta_i} F(\bar{\lambda}, k))_{\bar{\lambda} \in I}$
a cone $\alpha : 1 \dashv \lim_I F(-, i)$, $\beta = 1 \dashv \lim_I F(-, k)$

for some $i, k \in \text{OLJ}$.

If $k([t]) = k([t'])$, then for each $\bar{\lambda} \in I$,
 $\exists t_{\bar{\lambda}} \in J$ s.t. $t_{\bar{\lambda}}$ and $\beta_{\bar{\lambda}}$ have a common nse
 in $F(\bar{\lambda}, t_{\bar{\lambda}})$ since I is finite and J is filtered,
 by the same argument we did for λ , (the size
 $|t|$ is finite) $\exists t \in J$ s.t. $t_{\bar{\lambda}}, \alpha_{\bar{\lambda}}$ and $\beta_{\bar{\lambda}}$ have
 a common nse in $F(\bar{\lambda}, t)$. Then, as an element
 of $F(\bar{\lambda}, i)$ or $F(\bar{\lambda}, k)$

~~α_i~~ image of α_i in $F(\bar{a}, t) = \text{image of } \beta_i$
in $F(\bar{a}, t)$

$\Rightarrow [\alpha] = [\beta]$ in $\text{dom } k$.

$\therefore k$ is injective.

Rmk) Proof is not categorical, since choosing \bar{a} is arbitrary. In large class of category (Ch 5) these are created by Set.

Ex. 3.8.i) Let ε be a coequalizer map

$$\begin{array}{ccc} A & \xrightarrow[e]{\quad i \quad} & A \\ & \searrow e & \downarrow r \\ & & A \end{array} \quad \begin{array}{l} \text{Then the universal property} \\ (\text{left diagram}) \text{ gives} \\ \text{unique } r \\ \text{s.t. } r\varepsilon = e. \end{array}$$

Thus we can set below diagram since

$$\begin{array}{ccc} A & \xrightarrow[e]{\quad i \quad} & A \\ & \searrow s & \downarrow r \\ & & A/e \end{array} \quad \begin{array}{l} sr\varepsilon = se = \varepsilon. \\ \text{The universal property} \\ \text{gives } sr = 1_{A/e}. \end{array}$$

Thus, A/e is a retract of A . Conversely, any retract of the form $B \xrightarrow[r]{\quad} A \xrightarrow[\varepsilon]{\quad} B$ gives an idempotent sr , which is split by B .

To see any functor preserves idempotent, if F is a functor, then $F(e \cdot e) = F(e) \cdot F(e) = Fe$.

Now suppose that $j: FA \rightarrow B$ is cod F .
s.t. $j = j \circ Fe$.

Then we have a diagram.

$$\begin{array}{ccc} FA & \xrightarrow{\text{Fe}} & FA \\ & \downarrow & \downarrow F_{\text{Fr}} \\ & 2. & \downarrow j \\ & FA & \downarrow j \\ & B & \end{array}$$

We claim that

Left triangle is commutative

Since

$$j \circ F_{\text{Fr}} = j \circ \text{Fe} = j.$$

Moreover, since F_S is still split epi,
(since functor preserves split epi)

thus epi, therefore if $\exists b: F(A/e) \rightarrow B$
s.t. $b \circ F_S = j$, then $(j \circ F_{\text{Fr}}) \circ F_S = j$
implies $b = j \circ F_{\text{Fr}}$.

Thus, $F(A/e)$ is a regularizer in $\text{cod } F$.
Now, by duality, idempotent as a limit also
preserved by any functor.

Thus, by Thm 3.8.1, they commute with
both limits and colimits of any shape.

Ex 3.8.ii) It suff to show that $\underset{H, G}{\text{colim}} (\lim X)$
 $= \lim_{G, H} (\text{colim } X)$.

And from previous lecture, we know that

LHS = The set of H -orbits on $X^{G, H}$

RHS = The set of // consists of G -fixed elements

Suppose $\exists G \times H \rightarrow X$ exists.

$$\begin{array}{ccc} (g, h) & \downarrow & X \\ \text{I}_G & \xrightarrow{\quad} & \text{I}_H \\ & \cdot & \cdot \\ & & X \end{array}$$

At least $X = G \times H$ exists in Set.

Then, by identifys $g \circ c = (g, I_H) \circ c$
 $h \circ c = (I_G, h) \circ c$,

\bullet X is G -Set and H -Set.

Moreover, if x, y are in the same H -orbit,
then $x = (d, h)y$ for some $h \in H$, thus

$$\begin{aligned} \forall g \in G, \quad g \circ x &= (g, I_H)x = (g, I_H)(d, h)y = (g, h)x \\ &= (d, h)(g, I)x. \end{aligned}$$

$\Rightarrow g^*x$ and g^*y are in the same orbit.

thus, $\exists G$ action on the set of H -orbits.

Let A be an H -orbit s.t. $gA = A$ $\forall g$.

Then ~~for any~~ $\exists x \in A, \forall g \in G \exists h \in H$ s.t.

$g \circ x = h \circ x$ for some h .

$g^n \circ x = x$. $\forall n \in \mathbb{N} \Rightarrow \exists n > 0$ s.t.

But this implies $n \mid |g|$ and $n \mid |h|$.

$\Rightarrow n=1$ since $|g|$ and $|h|$ are coprime.

\therefore all elements in A are in X^G .

Conversely, if we take an H -orbit on X^G , then these are orbits of X consisting of X^G elements. the