# A Non-homogeneous Time Mixed Integer LP Formulation for Traffic Signal Control

- 3 Iain Guilliard
- 4 National ICT Australia
- 5 7 London Circuit
- 6 Canberra, ACT, Australia
- 7 iguilliard@nicta.com.au
- 8 Scott Sanner
- 9 Oregon State University
- 10 1148 Kelley Engineering Center
- 11 Corvallis, OR 97331
- 12 scott.sanner@oregonstate.edu
- 13 Felipe W. Trevizan
- 14 National ICT Australia
- 15 7 London Circuit
- 16 Canberra, ACT, Australia
- 17 felipe.trevizan@nicta.com.au
- 18 Brian C. Williams
- 19 Massachusetts Institute of Technology
- 20 77 Massachusetts Avenue
- 21 Cambridge, MA 02139
- 22 williams@csail.mit.edu
- 23 4172 words + 7 figures + 0 table + 23 citations (Weighted total words: 5922 out of 7000 + 35
- 24 references)
- 25 August 1, 2015

#### ABSTRACT

We build on the body of work in mixed integer linear programming (MILP) approaches that attempt to jointly optimize traffic signal control over an entire traffic network (rather than focus on arterial routes) and specifically on improving the scalability of these methods for large urban traffic networks. Our primary insight in this work stems from the fact that MILP-based approaches to traffic control used in a receding horizon control manner (that replan at fixed time intervals) need to compute high fidelity control policies only for the early stages of the signal plan; therefore, coarser 7 time steps can be employed to "see" over a long horizon to preemptively adapt to distant platoons and other predicted long-term changes in traffic flows. To this end, we contribute the queue transmission model (QTM) which blends elements of cell-based and link-based modeling approaches 10 11 to enable a non-homogeneous MILP formulation of traffic signal control. We then experiment with this novel QTM-based MILP control in a range of networks demonstrating the improved scalabil-12 ity possible with non-homogeneous time steps in comparison to the best homogeneous time step. 13 Our experiments also provide near-optimal traffic control policies for larger horizons and larger 14 networks than shown in previous implementations of MILP-based traffic signal control. 15

Using 204 words up to here. Maximum is 250 words.

17 <sup>1</sup>

<sup>&</sup>lt;sup>1</sup>Make sure to follow instructions and author guide: http://onlinepubs.trb.org/onlinepubs/AM/InfoForAuthors.pdf http://onlinepubs.trb.org/onlinepubs/am/2015/WritingForTheTRRecord.pdf

#### INTRODUCTION

11

17

18

19

20

21

22 23

24 25

26

27

28

29

31

33

36 37

39

40 41

42 43

44

As cities rapidly grow in population while urban traffic infrastructure often adapts at a slower pace, it is critical to maximize capacity and throughput of existing road infrastructure through optimized traffic signal control. Unfortunately, many large cities still use some degree of fixed-time control (e.g., Toronto (1)) even if they also use actuated or adaptive control methods such as SCATS (2) or SCOOT (3). However, there is further opportunity to improve traffic signal control even beyond adaptive methods through the use of *optimized* controllers as evidenced in a variety of approaches 7 ranging from mixed integer (linear) programming (4, 5, 6, 7, 8, 9) to heuristic search (10, 11) to scheduling (12) to reinforcement learning (1). While such optimized controllers hold the promise 10 of maximizing existing infrastructure capacity by finding more complex (and potentially closer to optimal) jointly coordinated intersection policies than arterially-focused master-slave approaches such as SCATS and SCOOT, such optimized methods are computationally demanding and either 12 (a) do not guarantee jointly optimal solutions over a large intersection network (often because they 13 only consider coordination of neighboring intersections or arterial routes) or (b) fail to scale to 14 large intersection networks simply for computational reasons (which is the case for many mixed 16 integer programming approaches).

In this work, we build on the body of work in mixed integer linear programming (MILP) approaches that attempt to jointly optimize traffic signal control over an entire traffic network (rather than focus on arterial routes) and specifically on improving the scalability of these methods for large urban traffic networks. In our investigation of existing approaches in this vein, namely exemplar methods in the spirit of (6, 8, 9) that use a (modified) cell transmission model (CTM) (13, 14) for their underlying prediction of traffic flows, we remark that a major drawback is the CTMimposed requirement to choose a predetermined homogeneous (and often necessarily small) time step for reasonable modeling fidelity. This need to model large number of CTM cells with a small time step leads to MILPs that are exceedingly large and intractable to solve.

Our primary insight in this work stems from the fact that MILP-based approaches to traffic control used in a receding horizon control manner (that replan at fixed time intervals) need to compute high fidelity control policies only for the early stages of the signal plan; therefore, coarser time steps can be employed to "see" over a long horizon to preemptively adapt to distant platoons and other predicted long-term changes in traffic flows. This need for non-homogeneous control in turn spawns the need for an additional innovation: we require a traffic flow model that permits non-homogeneous time steps and properly models the travel time delay between lights. To this end, we might consider CTM extensions such as the variable cell length CTM (15), stochastic CTM extensions (16, 17), extensions for better modeling freeway-urban interactions (18) including CTM hybrids with link-based models (19), assymmetric CTMs for better handling flow imbalances in merging roads (20), the situational CTM for better modeling of boundary conditions (21), and the lagged CTM for improved modeling of the flow density relation (22). However, despite the widespread varieties of the CTM and the usage of the CTM (23) for a range of applications, there seems to be no extension that permits non-homogeneous time steps as required in our novel MILPbased control approach.

For this reason, as a major contribution of this work to enable our non-homogeneous time MILP-based model of joint intersection control, we contribute the queue transmission model (QTM) which blends elements of cell-based and link-based modeling approaches with the following key benefits:

2

3

4

5 6

7

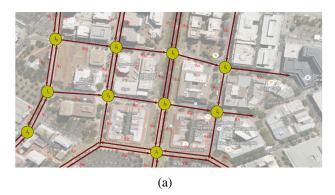
8

9

11

17

18



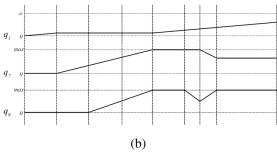


FIGURE 1 (a) Example of how a real network is modeled using QTM. (b) Volume of traffic in different queues as a function of non-homogeneous discretized time.

- unlike previous joint intersection control work (6, 8, 9), it is inherently intended for nonhomogeneous time steps that can be used for control over large horizons,
- any length of roadway with no merges or diverges can be modeled as a single queue leading to compact models of large traffic networks thus maintaining relatively compact MILPs for large traffic networks (i.e., larges numbers of cells are not required between intersections), and
- it accurately models fixed travel time delays critical to green wave coordination as in (4, 5, 7) through the use of a non-first order Markovian update model and combines this with the more global intersection signal optimization approach of (6, 8, 9).

In the remainder of this paper, we first formalize our novel QTM model of traffic flow 10 with non-homogeneous time steps and show how to encode it as a linear program for simulating traffic. We proceed to allow the traffic signals to become discrete variables subject to a delay 12 13 minimizing optimization objective and standard cycle and phase time constraints leading to our final MILP formulation of traffic signal control. We then experiment with this novel QTM-based MILP control in a range of networks demonstrating the improved scalability possible with nonhomogeneous time steps in comparison to the best homogeneous time step. These experiments also provide near-optimal traffic control policies for larger horizons and larger networks than shown in previous implementations of MILP-based traffic signal control. <sup>2 3</sup>

## THE QUEUE TRANSMISSION MODEL

- 20 A Queue Transmission Model (QTM) is the tuple  $(\mathcal{Q}, \mathcal{L}, \vec{\Delta t}, \mathbf{I})$ , where  $\mathcal{Q}$  and  $\mathcal{L}$  are, respectively,
- the set of queues and lights;  $\vec{\Delta t}$  is a vector of size N representing the discretization of the simulation
- horizon [0, T] and the duration in seconds of the n-th time interval is denoted as  $\Delta t_n$ ; and I is a 22

<sup>&</sup>lt;sup>2</sup>We could really use some pictures in the Intro to refer to here and subsequently – both a traffic network divided into queues, and the concept of the piecewise linear evolution of traffic flow with non-homogeneous (dilated) time steps, something like I had provided in my early writeup. I think these help visually explain much of the context for the paper and its approach and are critical for reviewer understanding on a time budget for reading this They may only read the first 2-3 pages and then skim!

<sup>&</sup>lt;sup>3</sup>A picture is worth a 1000 words but we only pay 250, hence a 4X ROI on pictures!

- matrix  $|Q| \times T$  in which  $I_{i,n}$  represents the flow of cars requesting to enter queue i from the outside of the network at time n.
- A traffic light  $\ell \in \mathcal{L}$  is defined as the tuple  $(\Psi_{\ell}^{\min}, \Psi_{\ell}^{\max}, \mathcal{P}_{\ell}, \vec{\Phi}_{\ell}^{\min}, \vec{\Phi}_{\ell}^{\max})$ , where:
- $\mathcal{P}_{\ell}$  is the set of phases of  $\ell$ ;
- $\Psi_\ell^{min}$  ( $\Psi_\ell^{max}$ ) is the minimum (maximum) allowed cycle time for  $\ell$ ; and
- $\vec{\Phi}_{\ell}^{\min}$  ( $\vec{\Phi}_{\ell}^{\max}$ ) is a vector of size  $|\mathcal{P}_{\ell}|$  and  $\Phi_{\ell,k}^{\min}$  ( $\Phi_{\ell,k}^{\max}$ ) is the minimum (maximum) allowed time for phase  $k \in \mathcal{P}_{\ell}$ .
- A **queue**  $i \in \mathcal{Q}$  represents a segment of road that vehicles traverse at free flow speed; once traversed, the vehicles are vertically stacked in a stop line queue. Formally, a queue i is defined by the tuple  $(Q_i, T_i^{prop}, F_i^{out}, \vec{F_i}, \vec{Pr_i}, \mathcal{Q}_i^{\mathcal{P}})$  where:
- $Q_i$  is the maximum capacity of i;
- $T_i^{\text{prop}}$  is the time required to traverse i and reach the stop line;
- $F_i^{\text{out}}$  represents the maximum traffic flow from i to the outside of the modeled network;
- $\vec{F}_i$  and  $\vec{Pr}_i$  are vectors of size  $|\mathcal{Q}|$  and their j-th entry (i.e.,  $F_{i,j}$  and  $Pr_{i,j}$ ) represent the maximum flow from queue i to j and the turn probability from i to j ( $\sum_{j \in \mathcal{Q}} Pr_{i,j} = 1$ ), respectively; and
- $\mathcal{Q}_i^{\mathcal{P}}$  denotes the set of traffic light phases controlling the outflow of queue i.
- Differently than CTM (8, 13), QTM does not assume that  $\Delta t_n = T_i^{\text{prop}}$  for all  $n \in \{1, \dots, N\}$ ,
- 19 that is, the QTM can represent non-homogeneous time intervals. The only requirement over  $\Delta t_n$  is
- 20 that no traffic light maximum phase time is smaller than any  $\Delta t_n$  since phase changes occur only
- 21 between time intervals; formally,  $\Delta t_n \leq \min_{\ell \in \mathcal{L}, k \in \mathcal{P}_{\ell}} \Phi_{\ell, k}^{\max}$  for all  $n \in \{1, \dots, N\}$ .

#### 22 Traffic Flow Simulation with QTM

- 23 In this section, we present how to simulate traffic flow in a network using QTM and non-homogeneous
- 24 time intervals  $\Delta t$ . We assume for the remainder of this section that a *valid* control plan for
- 25 all traffic lights is fixed and given as parameter; formally, for all  $\ell \in \mathcal{L}$ ,  $k \in \mathcal{P}_{\ell}$ , and interval
- 26  $n \in \{1, ..., N\}$ , the binary variable  $p_{\ell,k,n}$  is known a priori and indicates if phase k of light  $\ell$  is
- 27 active (i.e.,  $p_{\ell,k,n} = 1$ ) or not on interval n.

32

- We represent the problem of finding the flow between queues as a Linear Program (LP) over the following variables defined for all interval  $n \in \{1, ..., N\}$  and queues i and j:
- of  $q_{i,n} \in [0, Q_i]$ : traffic volume waiting in the stop line of queue i at the beginning of interval n;
  - $f_{i,n}^{\text{in}} \in [0, I_{i,n}]$ : inflow to the network via queue i during interval n;

<sup>&</sup>lt;sup>4</sup>**To Iain**: Maybe bring forward a small network and any other figure that would help illustrate the model and comment about it.

17

18

19

21

22

23

26

27

28

29

- $f_{i,n}^{\text{out}} \in [0, \mathbf{F}_i^{\text{out}}]$ : outflow from the network via queue i during interval n; and
- $f_{i,j,n} \in [0, F_{i,j}]$ : flow from queue i into queue j during interval n.

The maximum traffic flow from queue i to queue j is enforced by constraints (C1) and (C2). (C1) ensures that only the fraction  $\Pr_{i,j}$  of the total internal outflow of i goes to j, and (C2) forces the flow from i to j to be zero if all phases controlling i are inactive (i.e.,  $p_{\ell,k,n} = 0$  for all  $k \in \mathcal{Q}_i^{\mathcal{P}}$ ). If more than one phase  $p_{\ell,k,n}$  is active, then (C2) is subsumed by the domain upper bound of  $f_{i,j,n}$ .

$$f_{i,j,n} \le \Pr_{i,j} \sum_{k=1}^{|\mathcal{Q}|} f_{i,k,n} \tag{C1}$$

$$f_{i,j,n} \le F_{i,j} \sum_{p_{\ell,k,n} \in \mathcal{Q}_i^{\mathcal{P}}} p_{\ell,k,n} \tag{C2}$$

To simplify the presentation of remainder of the LP, we define the helper variables  $q_{i,n}^{\rm in}$  (C3),  $q_{i,n}^{\rm out}$  (C4), and  $t_n$  (C5) to represent the volume of traffic to enter and leave queue i during interval n, and the time elapsed since the beginning of the simulation until the end of interval  $\Delta t_n$ .

13 
$$q_{i,n}^{\text{in}} = \Delta t_n (f_{i,n}^{\text{in}} + \sum_{j=1}^{|\mathcal{Q}|} f_{j,i,n})$$
 (C3)

14 
$$q_{i,n}^{\text{out}} = \Delta t_n (f_{i,n}^{\text{out}} + \sum_{j=1}^{|\mathcal{Q}|} f_{i,j,n})$$
 (C4)

$$t_n = \sum_{x=1}^n \Delta t_x \tag{C5}$$

In order to account for the misalignment of the different  $\Delta t$  and  $T_i^{\text{prop}}$ , we need to find the volume of traffic that entered queue i between two arbitrary points in time x and y ( $x \in [0,T], y \in [0,T]$ , and x < y), i.e., x and y might not coincide with any  $t_n$  for  $n \in \{1,\ldots,N\}$ . This volume of traffic, denoted as  $V_i(x,y)$ , is obtained by integrating  $q_{i,n}^{\text{in}}$  over [x,y] and is defined in (1) where m and w are the index of the time intervals s.t.  $t_m \le x < t_{m+1}$ , and  $t_w \le y < t_{w+1}$ . Because the QTM dynamics is piecewise linear,  $q_{i,n}^{\text{in}}$  is a step function w.r.t. time and this integral reduces to the sum of  $q_{i,n}^{\text{in}}$  over the intervals contained in [x,y] and the appropriate fraction of  $q_{i,m}^{\text{in}}$  and  $q_{i,w}^{\text{in}}$  representing the misaligned beginning and end of [x,y].

$$V_i(x,y) = (t_{m+1} - x) \frac{q_{i,m}^{\text{in}}}{\Delta t_m} + \left(\sum_{k=m+1}^{w-1} q_{i,k}^{\text{in}}\right) + (y - t_w) \frac{q_{i,w}^{\text{in}}}{\Delta t_w}$$
(1)

Using these helper variables, (C6) represents the flow conservation principle for queue i where  $V_i(t_{n-1} - T_i^{\text{prop}}, t_n - T_i^{\text{prop}})$  is the volume of cars that reached stop line during  $\Delta t_n$ . Since  $\Delta t$  and  $T_i^{\text{prop}}$  for all queues are known a priori, the indexes m and w used by  $V_i$  can be precomputed in order to encode (1); moreover, (C6) represents a non-first order Markovian update because the update considers the previous w-m time steps. To insure that the total volume of

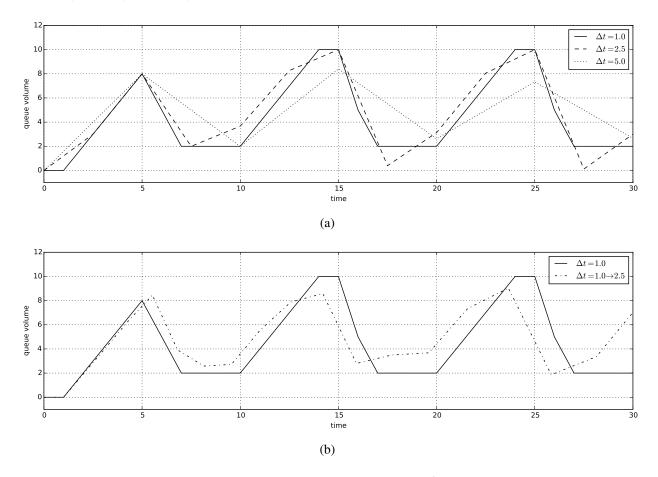


FIGURE 2 (a) Convergence with increasing refinement of  $\Delta t$  from 5.0 down to 1.0. (c) Dilation of  $\Delta t$  from 1.0 to 2.5 compared to a fixed  $\Delta t$  of 1.0.

traffic traversing i (i.e.,  $V_i(t_n - T_i^{\text{prop}}, t_n)$ ) and waiting at the stop line does not exceed the capacity of the queue, we apply (C7).

$$q_{i,n} = q_{i,n-1} - q_{i,n-1}^{\text{out}} + V_i(t_{n-1} - T_i^{\text{prop}}, t_n - T_i^{\text{prop}})$$
 (C6)

$$V_i(t_n - \mathcal{T}_i^{\text{prop}}, t_n) + q_{i,n} \le \mathcal{Q}_i$$
(C7)

As with MILP formulations of CTM (e.g. Lin and Wang (8)), QTM is also susceptible to withholding traffic, i.e., the optimizer might prevent cars from moving from i to j even though the associated traffic phase is active and j is not full. We address this issue through our objective function (O1) by maximizing the total outflow  $q_{i,n}^{\text{out}}$  (i.e., both internal and external outflow) of i plus the inflow  $f_{i,n}^{\text{in}}$  from the outside of the network to i. This quantity is weighted by the remaining time until the end of the simulation horizon T to force the optimizer to allow as much traffic volume as possible into the network and move traffic to the outside the network as soon as possible. (O1) is analogous to minimizing delay in CTM models, e.g., (O1) is equivalent to the objective function (O3) in Lin and Wang (8) for their parameters  $\alpha = \beta = 1$ . Figure 5(d) shows the delay experienced by each vehicle travelling along an avenue, where delay is the horzontal difference between the cumulative departure and arrival curves at each point, less the free flow travel time

6

10

11 12

<sup>&</sup>lt;sup>5</sup>To Iain: Add a paragraph linking the plots with the objective function.

5

6

21

23

24

26

27

28 29

along the avenue. The objective function tries to maximise the arrival curve by pushing it up closer to the departure curve, which also has the effect of minimising the horizontal distance, or delay. <sup>6</sup>

$$\max \sum_{n=1}^{N} \sum_{i=1}^{|\mathcal{Q}|} (T - t_n + 1) (q_{i,n}^{\text{out}} + f_{i,n}^{\text{in}})$$
 (O1)

The objective function (O1) and constraints (C1–C7) form the LP representing the dynamic, piecewise linear model of flow in a QTM network over time when a control plan  $p_{\ell,k,n}$  is given as an input parameter.

Figures 2(a) and 2(b) show the results of applying the LP formulation to a simple model with a fixed signal plan, using both homogeneous  $\vec{\Delta t}$  and non-homogeneous  $\vec{\Delta t}$ .

### 9 TRAFFIC CONTROL WITH QTM AS AN MILP

In this section, we remove the assumption that a valid control plan for all traffic lights is given and extend the LP (O1, C1–C7) to an Mixed-Integer LP (MILP) that also computes the optimal control plan. Formally, for all  $\ell \in \mathcal{L}$ ,  $k \in \mathcal{P}_{\ell}$ , and interval  $n \in \{1, \ldots, N\}$ , the phase activation parameter  $p_{\ell,k,n} \in \{0,1\}$  becomes a free variable to be optimized. In order to obtain a valid control plan, we enforce that one phase of traffic light  $\ell$  is always active at any interval n (C8) and that phase changes happen sequentially (C9), i.e., if phase k was active during interval n-1 and has become inactive in interval n, then phase k+1 must be active in interval n. (C9) assumes that k+1 equals 1 if  $k=|\mathcal{P}_{\ell}|$ .

18 
$$\sum_{k=1}^{|\mathcal{P}_{\ell}|} p_{\ell,k,n} = 1$$
 (C8)

$$p_{\ell,k,n-1} \le p_{\ell,k,n} + p_{\ell,k+1,n} \tag{C9}$$

Next, we enforce the minimum and maximum phase durations (i.e.,  $\Phi_{\ell,k}^{\min}$  and  $\Phi_{\ell,k}^{\max}$ ) for each phase  $k \in \mathcal{P}_{\ell}$  of traffic light  $\ell$ . To encode these constraints, we use the helper variable  $d_{\ell,k,n} \in [0,\Phi_{\ell,k}^{\max}]$  defined by constraints (C10–C14) that: (i) holds the elapsed time since the start of phase k when  $p_{\ell,k,n}$  is active (C10,C11); (ii) is constant and holds the duration of the last phase until the next activation when  $p_{\ell,k,n}$  is inactive (C12,C13); and (iii) is restarted when phase k changes from inactive to active (C14). Notice that (C10–C14) employs the big-M method to turn the cases that should not be active into subsumed constraints based on the value of  $p_{\ell,k,n}$ . We use  $\Phi_{\ell,k}^{\max}$  as our large constant since  $d_{\ell,k,n} \leq \Phi_{\ell,k}^{\max}$  and  $\Delta t_n \leq \Phi_{\ell,k}^{\max}$  by assumption (Section 2.1). Similarly, constraint (C15) ensures the minimum phase time of k and is not enforced while k is still active. Figures 3(a) to 3(c) present an example of how (C10–C15) work together as a function of the time n for  $d_{\ell,k,n}$ ; the domain constraint  $0 \leq d_{\ell,k,n} \leq \Phi_{\ell,k}^{\max}$  for all  $n \in \{1,\dots,N\}$  is omitted

<sup>&</sup>lt;sup>6</sup>Show diagrams with traffic predictions converging as time increment gets smaller. Validates that large time-steps are rough approximations while model behavior converges for small time steps.

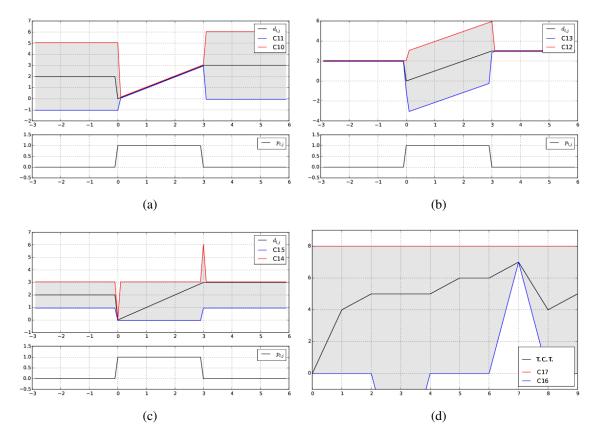


FIGURE 3 Visualization of constraints (C10–C17) for a traffic light  $\ell$  as a function of time. (a–c) present, pairwise, the constraints (C10–C15) for phase k ( $d_{\ell,k,n}$  as the black line) and the activation variable  $p_{\ell,k,n}$  in the attached plot. (d) presents the constraint for the cycle time of  $\ell$  (C16 and C17) where T.C.T. is the total cycle time and is defined as the left hand side of both constraints. For this example,  $\Phi_{\ell,k}^{\min}=1$ ,  $\Phi_{\ell,k}^{\max}=3$ ,  $\Psi_{\ell}^{\min}=7$ , and  $\Psi_{\ell}^{\max}=8$ .

1 for clarity.

9

10

11

12

13

$$d_{\ell,k,n} \le d_{\ell,k,n-1} + \Delta t_{n-1} p_{\ell,k,n-1} + \Phi_{\ell,k}^{\max} (1 - p_{\ell,k,n-1})$$
(C10)

$$d_{\ell,k,n} \ge d_{\ell,k,n-1} + \Delta t_{n-1} p_{\ell,k,n-1} - \Phi_{\ell,k}^{\max} (1 - p_{\ell,k,n-1})$$
(C11)

$$d_{\ell,k,n} \le d_{\ell,k,n-1} + \Phi_{\ell,k}^{\max} p_{\ell,k,n-1}$$
 (C12)

5 
$$d_{\ell,k,n} \ge d_{\ell,k,n-1} - \Phi_{\ell,k}^{\max} p_{\ell,k,n}$$
 (C13)

6 
$$d_{\ell,k,n} \le \Phi_{\ell,k}^{\max} (1 - p_{\ell,k,n} + p_{\ell,k,n-1})$$
 (C14)

$$d_{\ell,k,n} \ge \Phi_{\ell,k}^{\min}(1 - p_{\ell,k,n}) \tag{C15}$$

Lastly, we constrain the sum of all the phase durations for light  $\ell$  to be within the cycle time limits  $\Psi_{\ell}^{\min}$  (C16) and  $\Psi_{\ell}^{\max}$  (C17). In both (C16) and (C17), we use the duration of phase 1 of  $\ell$  from the previous interval n-1 instead of the current interval n because (C14) forces  $d_{\ell,1,n}$  to be 0 at the beginning of each cycle; however, from the previous end of phase 1 until n-1,  $d_{\ell,1,n-1}$  holds the correct elapse time of phase 1. Additionally, (C16) is enforced right after the end of the each cycle, i.e., when its first phase is changed from inactive to active. The value (C16) and (C17)

1 over time for a traffic light  $\ell$  is illustrated in Figure 3(d).

2 
$$d_{\ell,1,n-1} + \sum_{k=2}^{|\mathcal{P}_{\ell}|} d_{\ell,k,n} \ge \Psi_{\ell}^{\min}(p_{k,1,n} - p_{k,1,n-1})$$
 (C16)

$$d_{\ell,1,n-1} + \sum_{k=2}^{|\mathcal{P}_{\ell}|} d_{\ell,k,n} \le \Psi_{\ell}^{\max}$$
(C17)

5 The MILP that encodes the problem of finding the optimal traffic control plan in a QTM network 6 is defined by (O1, C1–C17).

#### 7 EMPIRICAL EVALUATION

In this section we compare the solutions for traffic networks modeled as a QTM using homogeneous and non-homogeneous time intervals in two aspects: the quality of the solution and convergence to the optimal solution. We compare the quality of solutions based on the total travel time and we also consider the third quartile and maximum of the observed delay distribution. Our hypotheses are: (i) the quality of the non-homogeneous solutions is at least as good as the homogeneous ones when the number of time intervals N is fixed; and (ii) the non-homogeneous approach requires less time intervals (i.e., smaller N) than the homogeneous approach to converge to the optimal solution. In the remainder of this section, we present the traffic networks considered in the experiments, our methodology, and the results.

#### 17 Networks

We consider three networks of increasing complexity (Figure 5): an avenue crossed by three side streets; a 2-by-3 grid; and a 3-by-3 grid with a diagonal avenue. The queues receiving cars from outside of the network are marked in Figure 5 and we refer to them as input queues. The maximum capacity ( $Q_i$ ) is 60 cars for non-input queues and infinity for input queues to prevent interruption of the input demand due to spill back from the stop line. The traversal time of each queue i ( $T_i^{\text{prop}}$ ) is set at 9s (a distance of about 100m with a free flow speed of 50km/h). Flows are defined from the head of each queue i into the tail of the next queue j; there is no turning traffic ( $\Pr_{i,j} = 1$ ), and the maximum flow rate between queues,  $F_{i,j}$ , is set at 5 cars/s. All traffic lights have two phases, north-south and east-west, and lights 2, 4 and 6 of network 3 have the additional northeast-southwest phase to control the diagonal avenue. For networks 1 and 2,  $\Phi_{\ell,k}^{\min}$  is 1s,  $\Phi_{\ell,k}^{\max}$  is 3s,  $\Psi_{\ell}^{\min}$  is 2s, and  $\Psi_{\ell}^{\max}$  is 6s, for all traffic light  $\ell$  and phase  $\ell$ . For network 3,  $\Phi_{\ell,k}^{\min}$  is 1s and  $\Phi_{\ell,k}^{\max}$  is 6s for all  $\ell$  and  $\ell$ ; and  $\ell$  and  $\ell$  is 2s and  $\ell$  and  $\ell$  is 12s for all lights  $\ell$  except for lights 2, 4 and 6 in which  $\Psi_{\ell}^{\min}$  is 3s and  $\Psi_{\ell}^{\max}$  is 18s.

#### 1 Experimental Methodology

For each network, a constant background level traffic is injected in the network in the first 55s to allow the solver to settle on a stable policy. Then a spike in demand is introduced in the queues marked as (Figure 5) from time 55s to 70s to trigger a policy change. From time 70s to 85s, the demand is returned to the background level, and then reduced to zero for all input queues. We extend the problem horizon T until all cars have left the network. By clearing the network, we can

<sup>7</sup>FWT: I don't think that optimal is the best word here since we arbitrarily fixed a value of  $\Delta t$ . Also, there is the technical problem that Gurobi might not have found the true optimal.

7

10

11

12

15

17

22

23

24

26

2728

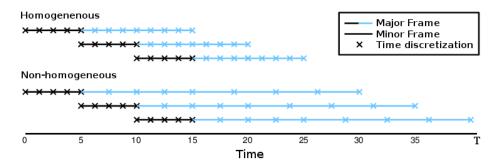


FIGURE 4 Receding horizon control. For this figure, the problem horizon T is 40s. The major frames are discretized in 12 time intervals (N=12) and they span 15s and 30s for homogeneous and non-homogeneous discretizations, respectively.

easily measure the total travel time for all the traffic as the area between the cumulative arrival and departure curves measured at the boundaries of the network. The background level for the input queues are 1, 4 and 2 cars/s for queues marked as  $\diamondsuit$ ,  $\clubsuit$  and  $\spadesuit$  (Figure 5), respectively; and during the high demand period, the queues  $\spadesuit$  receive 4 cars/s.

For both homogeneous and non-homogeneous intervals, we use the MILP QTM formulation (Section 3) in a receding horizon manner: a control plan is computed for a pre-defined horizon (smaller than T) and only a prefix of this plan is executed before generating a new control plan. Figure 4 depicts our receding horizon approach and we refer to the planning horizon as a major frame and its executable prefix as a minor frame. Notice that, while the plan for a minor frame is being executed, we can start computing the solution for the next major frame based on a forecast model.<sup>9</sup>

To perform a fair comparison between the homogeneous and non-homogeneous discretizations, we fix the size of all minor frames to 10s and force it to be discretized in homogeneous intervals of 0.25s. For the homogeneous experiments,  $\Delta t$  is kept at 0.25s throughout the major frame; therefore, given N, the major frame size equals N/4 seconds for the homogeneous approach. For the non-homogeneous experiments,  $\Delta t$  linearly increases from 0.25s at the end of the minor frame to 1.0s at the end of the major frame; therefore, the major frame size used by the non-homogeneous approach is 10.375 + 0.625(N-40) seconds for a given N>40. Once we have generated a series of minor frames, we concatenate them into a single plan and simulate the flow through the network using the QTM LP formulation with a fixed (homogeneous)  $\Delta t$  of 0.25s. We also compare both receding horizon approaches against the optimal solution obtained by computing a single control plan for the entire control horizon (i.e., [0,T]) using a fixed  $\Delta t$  of 0.25s.

For all our experiments, we used Gurobi<sup>TM</sup> as MILP solver with 12 threads on a 3.1GHz AMD Opteron<sup>TM</sup> 4334 processor with 12 cores. We limit MIP gap accuracy to 0.1% and the time cutoff for solving a major frame to 3000s for the receding horizon approaches and unbounded for the optimal plan. All our results are averaged over five runs to account for Gubori's stochastic strategies.

<sup>&</sup>lt;sup>8</sup>FWT: is this explanation of how to compute the total travel time still necessary?

<sup>&</sup>lt;sup>9</sup>FWT: not sure if we should mention this.

<sup>&</sup>lt;sup>10</sup>Do we need to justify why we use the QTM as the simulator over say a micro simulator?

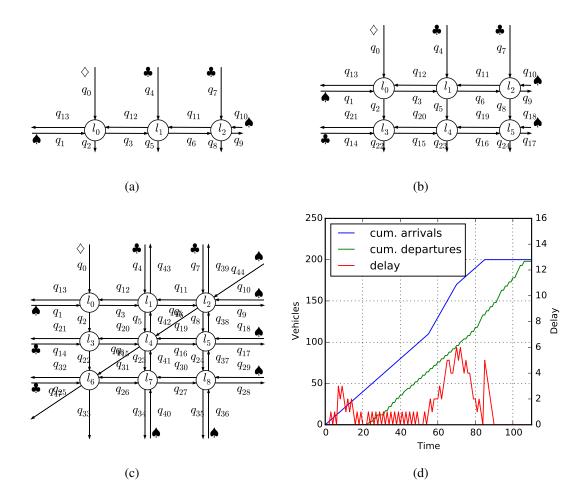


FIGURE 5 Networks used to evaluate the QTM performance. Queues markerd as  $\diamondsuit$ ,  $\clubsuit$ , and  $\spadesuit$  recieve traffic from the outside of the network. Now that we have plenty of space for figures, is Fig (d) here still necessary?

#### Results

Figures 6(a), 6(c) and 6(e) show, for each network, the increase in the total travel time w.r.t. the optimal solution as a function of N. As we hypothesized, the non-homogeneous discretization requires less time intervals to converge. This is important because the size of the MILP, including the number of binary variables, scales linearly with N; therefore, the non-homogeneous approach can scale up better than the homogeneous one (e.g., Figure 6(e)).

The distribution of the total delay observed by each car while traversing the network is shown in Figures 6(b), 6(d) and 6(f). Each group of box plots represents a different value of N: when the non-homogeneous  $\Delta t$  first converges to the optimum solution; when the homogeneous  $\Delta t$  first converges on the optimum solution; and the optimum solution itself. In all networks, the quality of the solution obtained using non-homogeneous  $\Delta t$  is better or equal than using homogeneous  $\Delta t$  for fixed N in both the total travel time and *fairness*, i.e., smaller third quartile and maximum delay.

FWT: In the paragraphs above, we need to address network 2 because it is the exception in both cases: in the end of Figure 6(c), homogeneous is better, and the homogeneous delay in

#### Figure 6(d) is also better.

Finally, Figure 7 shows the cumulative arrival and departure curves and the how delay evolves over time for  $q_1$  of network 2. Figure 7(a) shows the comparison at the point where the non-homogeneous  $\Delta t$  first converges and shows that with the longer major frame time of the non-homogeneous  $\Delta t$ , the solver is able to find a coordinated signal policy along the avenue to dissipate the extra traffic that arrives at the 55s point, while the homogeneous  $\Delta t$  with its shorter major frame. Once the homogeneous  $\Delta t$  has converged in Figure 7(b), both solutions are close to the optimum solution which is shown in Figure 7(c).

#### 9 **CONCLUSION**

In this paper, we showed how to formulate a novel queue transmission model (QTM) model of traffic flow with non-homogeneous time steps as a linear program. We then proceeded to allow the traffic signals to become discrete variables subject to a delay minimizing optimization objective and standard traffic signal constraints leading to a final MILP formulation of traffic signal control. We experimented with this novel QTM-based MILP control in a range of networks and demonstrated that by exploiting the non-homogeneous time steps supported by the QTM, we are able to scale the model up to larger networks whilst maintaining the same quality of a homogeneous solution using more binary variables. Altogether, this work represents a major step forward in the scalability of MILP-based jointly optimized traffic signal control via the use of a non-homogeneous traffic models and thus helps pave the way for fully optimized joint urban traffic signal controllers

#### 21 ACKNOWLEDGMENT

20

- 22 This work is part of the Advanced Data Analytics in Transport programme, and supported by Na-
- 23 tional ICT Australia (NICTA) and NSW Trade&Investment. NICTA is funded by the Australian
- 24 Government through the Department of Communications and the Australian Research Council
- 25 through the ICT Centre of Excellence Program. NICTA's role is to pursue potentially economi-
- 26 cally significant ICT related research for the Australian economy. NSW Trade&Investment is the
- 27 business development agency for the State of New South Wales.

as an improved successor technology to existing signal control methods.

## 28 **REFERENCES**

- 29 [1] El-Tantawy, S., B. Abdulhai, and H. Abdelgawad, Multiagent reinforcement learning for 30 integrated network of adaptive traffic signal controllers (MARLIN-ATSC): methodology 31 and large-scale application on downtown toronto. *Intelligent Transportation Systems, IEEE* 32 *Transactions on*, Vol. 14, No. 3, 2013, pp. 1140–1150.
- 33 [2] Sims, A. G. and K. W. Dobinson, SCAT–The Sydney co-ordinated adaptive traffic system: 34 Philosophy and benefits. *IEEE Transactions on Vehicular Technology*, Vol. 29, 1980.
- [3] Hunt, P. B., D. I. Robertson, R. D. Bretherton, and R. I. Winton, SCOOT–A traffic responsive
   method of coordinating signals. Transportation Road Research Lab, Crowthorne, U.K., 1981.
- [4] Gartner, N., J. D. Little, and H. Gabbay, *Optimization of traffic signal settings in networks by mixed-integer linear programming*. DTIC Document, 1974.

<sup>&</sup>lt;sup>11</sup>FWT: although Figure 7 is a nice illustration of how homogeneous and non-homogeneous differ, it is currently not backing up any of our claims.

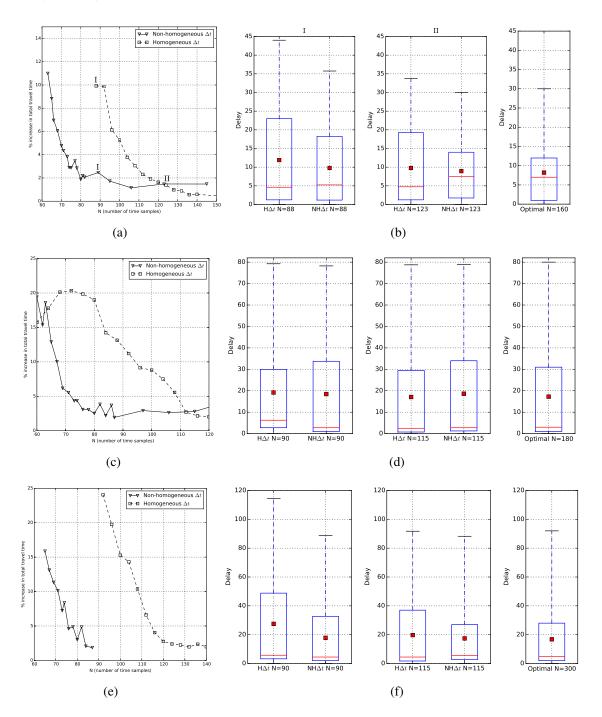


FIGURE 6 Increase in the total travel time w.r.t. the optimal solution as a function of N (Figures a, c, and e) and distribution of the total delay of each car for different values of N (Figures b, d, and f). The mean of the total delay is presented as a red square in box plots. Plots in the i-th row correspond to the results for the i-th network in Figure 5.

[5] Gartner, N. H. and C. Stamatiadis, Arterial-based control of traffic flow in urban grid networks. *Mathematical and computer modelling*, Vol. 35, No. 5, 2002, pp. 657–671.

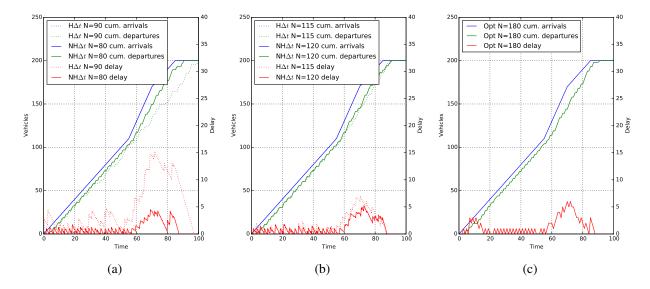


FIGURE 7 Cumulative arrival and departure curves and delay for queue 1 in the 2-by-3 network (Figure 5(b). The value of N in plots (a) and (b) corresponds, respectively, to the convergence point of the non-homogeneous and homogeneous approaches (Figure 6(c)). (c) presents the same curves for the optimal solution.

- 1 [6] Lo, H. K., A novel traffic signal control formulation. *Transportation Research Part A: Policy and Practice*, Vol. 33, No. 6, 1998, pp. 433–448.
- 3 [7] He, Q., K. L. Head, and J. Ding, PAMSCOD: Platoon-based Arterial Multi-modal Signal Control with Online Data. *Procedia-Social and Behavioral Sciences*, Vol. 17, 2011, pp. 462–489.
- 6 [8] Lin, W.-H. and C. Wang, An enhanced 0-1 mixed-integer LP formulation for traffic signal control. *Intelligent Transportation Systems, IEEE Transactions on*, Vol. 5, No. 4, 2004, pp. 238–245.
- 9 [9] Han, K., T. L. Friesz, and T. Yao, A link-based mixed integer LP approach for adaptive traffic signal control. *arXiv preprint arXiv:1211.4625*, 2012.
- 11 [10] Lo, H. K., E. Chang, and Y. C. Chan, Dynamic network traffic control. *Transportation Research Part A: Policy and Practice*, Vol. 35, No. 8, 1999, pp. 721–744.
- 13 [11] He, Q., W.-H. Lin, H. Liu, and K. L. Head, Heuristic algorithms to solve 0–1 mixed integer
  LP formulations for traffic signal control problems. In *Service Operations and Logistics and*Informatics (SOLI), 2010 IEEE International Conference on, IEEE, 2010, pp. 118–124.
- 16 [12] Smith, S., G. Barlow, X.-F. Xie, and Z. Rubinstein, SURTRAC: Scalable Urban Traffic Control. In *Transportation Research Board 92nd Annual Meeting Compendium of Papers*, Transportation Research Board, 2013.

- 1 [13] Daganzo, C. F., The cell transmission model: A dynamic representation of highway traffic
- 2 consistent with the hydrodynamic theory. Transportation Research Part B: Methodological,
- 3 Vol. 28, No. 4, 1994, pp. 269–287.
- 4 [14] Daganzo, C. F., The cell transmission model, part II: network traffic. *Transportation Research Part B: Methodological*, Vol. 29, No. 2, 1995, pp. 79–93.
- 6 [15] Xiaojian, H., W. Wei, and H. Sheng, Urban traffic flow prediction with variable cell trans-
- 7 mission model. Journal of Transportation Systems Engineering and Information Technology,
- 8 Vol. 10, No. 4, 2010, pp. 73–78.
- 9 [16] Sumalee, A., R. Zhong, T. Pan, and W. Szeto, Stochastic cell transmission model (SCTM): A
- 10 stochastic dynamic traffic model for traffic state surveillance and assignment. *Transportation*
- 11 *Research Part B: Methodological*, Vol. 45, No. 3, 2011, pp. 507–533.
- 12 [17] Jabari, S. E. and H. X. Liu, A stochastic model of traffic flow: Theoretical foundations.
- 13 Transportation Research Part B: Methodological, Vol. 46, No. 1, 2012, pp. 156–174.
- 14 [18] Huang, K. C., Traffic Simulation Model for Urban Networks: CTM-URBAN. Ph.D. thesis,
- 15 Concordia University, 2011.
- 16 [19] Muralidharan, A., G. Dervisoglu, and R. Horowitz, Freeway traffic flow simulation using the
- link node cell transmission model. In American Control Conference, 2009. ACC'09., IEEE,
- 18 2009, pp. 2916–2921.
- 19 [20] Gomes, G. and R. Horowitz, Optimal freeway ramp metering using the asymmetric cell trans-
- 20 mission model. Transportation Research Part C: Emerging Technologies, Vol. 14, No. 4,
- 21 2006, pp. 244–262.
- 22 [21] Kim, Y., Online traffic flow model applying dynamic flow-density relation. Int. At. Energy
- 23 Agency, 2002.
- 24 [22] Lu, S., S. Dai, and X. Liu, A discrete traffic kinetic model-integrating the lagged cell trans-
- 25 mission and continuous traffic kinetic models. Transportation Research Part C: Emerging
- 26 *Technologies*, Vol. 19, No. 2, 2011, pp. 196–205.
- 27 [23] Alecsandru, C., A. Quddus, K. C. Huang, B. Rouhieh, A. R. Khan, and Q. Zeng, An as-
- sessment of the cell-transmission traffic flow paradigm: Development and applications. In
- 29 Transportation Research Board 90th Annual Meeting, 2011, 11-1152.