

$$\text{binding halt} \rightarrow \frac{x+y+z}{2} \gg 1$$

$$\frac{x+y+z}{2} \gg 1$$

① Binary search -

Given → sorted array → 

as the array is sorted.

lower indices → smaller numbers.

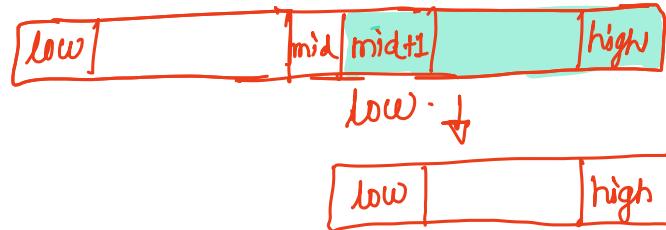
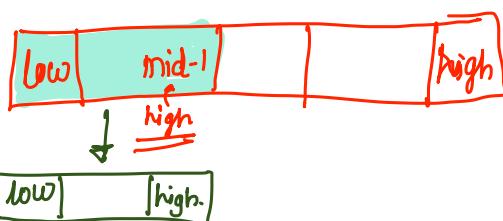
higher indices → greater numbers.



$\text{target} < \text{mid}$

$\text{target} = \text{mid}$
return

$\text{target} > \text{mid}$.



② Search in 2D matrix

Given - 2D matrix -

1	2	4	8
10	11	12	13
15	16	17	18
22	24	26	28

Given → each row is sorted.

we traverse each row in matrix

pass each row as an array in binary search.

BinarySearch (row, target).

BinarySearch ([1, 2, 4, 8], 17)

BinarySearch ([10, 11, 12, 13], 17)

BinarySearch ([15, 16, 17, 18], 17) → true

BinarySearch ((22, 24, 26, 28), 17)

③ KOKO eating banana -

Given \rightarrow array \rightarrow

1	1	1	1	1	1
---	---	---	---	---	---

hrs \rightarrow x.

we have to decide

speed at KOKO to eat Banana

tricky thing

once KOKO starts eating a pile

it has to spend full hr

Suppose -

1 3) 8) 10) 11) 14 speed = 4

even if $3 < 4 \rightarrow$ KOKO would spend full one hr here.

for $10 \rightarrow 10/4 \rightarrow 2.5 \rightarrow$ so KOKO would spend 3 hrs

that's why we take ceiling of division

As we know

1 | 4 | 3 | 2

once KOKO starts with one pile

it has to spend 1 hr on it no matter what.

so

maximum speed it could achieve is max in array

lowest speed it could get is 1

so, we have to do binary search from 1 to max.

& our target would be given hrs

Also, if we don't get matched with exacts hrs
we have to return max speed.

④ Find minimum in rotated sorted array

Given $\rightarrow [3, 4, 5, 6, 1, 2]$

3, 4, 5, 6, 1, 2
low mid high

6, 1, 2, 3, 4, 5
low mid high

mid > high. $\{$ low = mid + 1

3, 4, 5, 6, 1, 2
low high.

6, 1, 2, 3, 4, 5
low high

6, 1, 2
low high

We apply binary search
just a tweek

⑤ Search in rotated sorted array:

Binary search -

if we find element in middle \rightarrow just return
check if left part of array is sorted \rightarrow

we see if the number is between
the left part of array

if yes \rightarrow reduce the workspace

else

see if it is in right part of an array.

If it is \rightarrow reduce work space.

target = 6

3, 4, 5, 6, 1, 2.
low mid high

⑥ Time based Key-value store -

key → Integers
value → strings.

we want → get (If no match, give smaller close value)
→ set.

suppose,

we are given '7' like get(7)

but we don't have '7'

we would search for any value less than 7

like 6, 5, 4,

for that we need one map to be sorted.

so, we make a TreeMap

foo → TreeMap {
Key

1 → "apple"

4 → "banana"

7 → "cherry"

}

for every word, we have TreeMap.

foo →

treeMap	
key	value
1	apple
4	Banana
5	Kiwi

drink →

treeMap	
key	value
10	7up
12	Fanta
14	Maza

⑦ Median of two sorted arrays.

l_1 x_1
2, 4, 9, 12
 l_2 x_2
5, 6, 8, 13

In order to have correct partition

$$l_1 \leq x_2$$

$$l_2 \leq x_1$$

we find the smaller array

& try to apply binary search on it.

for applying binary search on first array

$$\text{mid}_1 = (\text{low} + \text{high}) / 2$$

2 mid_2 = would be the remaining element of that half.

2, 4, 9, 12
5, 6, 8, 13

If we choose these two as first half
→ we have to choose these two
because partition should be 4-4

① find mid of both array

② decide $l_1, l_2, \sigma_1, \sigma_2$

\downarrow \downarrow \nearrow \searrow

mid₁₋₁ mid₂₋₁ mid₁ mid₂.

③ Apply condition for perfect partition

$$l_2 \leq \sigma_2$$

$$l_2 \geq \sigma_1$$

If satisfied -
odd elements $\rightarrow \max(l_1, l_2)$
even elements $\rightarrow \max(l_1, l_2) + \max(\sigma_1, \sigma_2)$

④ If the partition is Not perfect

$$l_1 > \sigma_2$$

$$\text{high} = \text{mid} - 1.$$

else

$$\text{low} = \text{mid} + 1.$$