TIC TaC Toe: Understanding the Minimax Algorithm

Published on December 13, 2013 by <u>Jason Fox</u> (//plus.google.com/109774408645971700391?rel=author) in <u>Software Engineering</u> (/category/software-engineering).

I recently built an unbeatable game of tic tac toe. It was a fun and very humbling project that taught me a ton. If you want to get totally schooled, give the <u>tic tac toe game a shot here (http://perfecttictactoe.herokuapp.com/)</u>.

In order to make the game unbeatable, it was necessary to create an algorithm that could calculate all the possible moves available for the computer player and use some metric to determine the best possible move. After extensive research it became clear that the <u>Minimax</u> (http://en.wikipedia.org/wiki/Minimax) algorithm was right for the job.

It took a little while to really fundamentally understand the algorithm and implement it in my game. I found many code examples and explanations, but none that really walked a simpleton like me through the ins and outs of the process. I hope this post will help some of you to appreciate the elegance of this algorithm.

Describing a Perfect Game of Tic Tac Toe

To begin, let's start by defining what it means to play a perfect game of tic tac toe:

If I play perfectly, every time I play I will either win the game, or I will draw the game. Furthermore if I play against another perfect player, I will always draw the game.

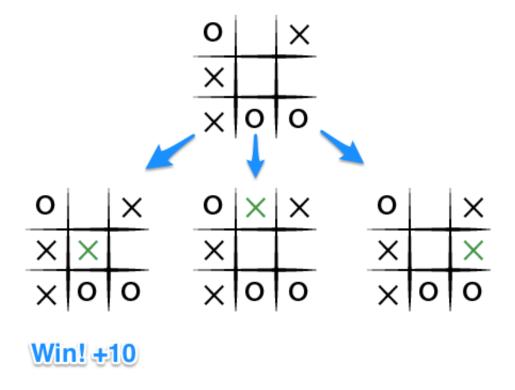
How might we describe these situations quantitatively? Let's assign a score to the "end game conditions:"

- I win, hurray! I get 10 points!
- I lose, shit. I lose 10 points (because the other player gets 10 points)
- I draw, whatever. I get zero points, nobody gets any points.

So now we have a situation where we can determine a possible score for any game end state.

Looking at a Brief Example

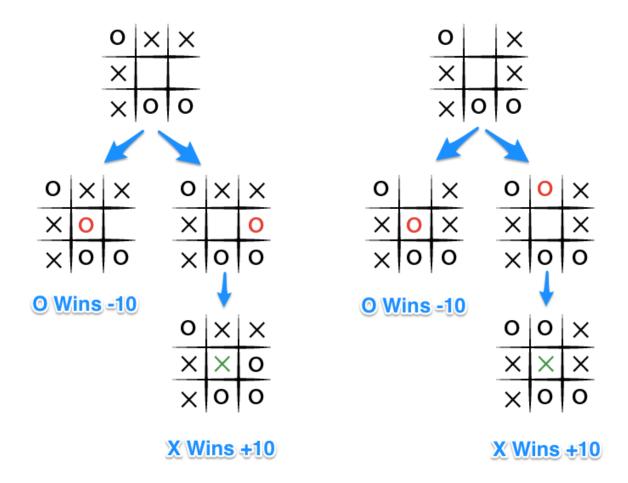
To apply this, let's take an example from near the end of a game, where it is my turn. I am X. My goal here, obviously, is to *maximize* my end game score.



If the top of this image represents the state of the game I see when it is my turn, then I have some choices to make, there are three places I can play, one of which clearly results in me wining and earning the 10 points. If I don't make that move, O could very easily win. And I don't want O to win, so my goal here, as the first player, should be to pick the maximum scoring move.

But What About O?

What do we know about O? Well we should assume that O is also playing to win this game, but relative to us, the first player, O wants obviously wants to chose the move that results in the worst score for us, it wants to pick a move that would *minimize* our ultimate score. Let's look at things from O's perspective, starting with the two other game states from above in which we don't immediately win:



The choice is clear, O would pick any of the moves that result in a score of -10.

Describing Minimax

The key to the Minimax algorithm is a back and forth between the two players, where the player whose "turn it is" desires to pick the move with the maximum score. In turn, the scores for each of the available moves are determined by the opposing player deciding which of its available moves has the minimum score. And the scores for the opposing players moves are again determined by the turn-taking player trying to maximize its score and so on all the way down the move tree to an end state.

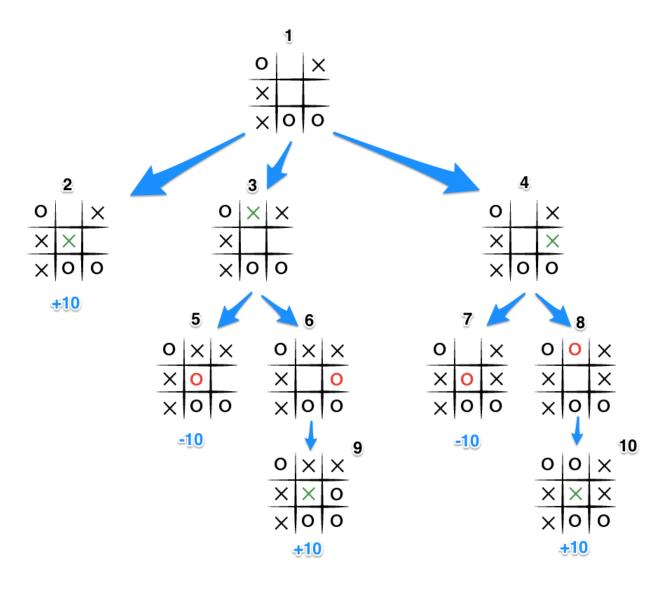
A description for the algorithm, assuming X is the "turn taking player," would look something like:

- If the game is over, return the score from X's perspective.
- Otherwise get a list of new game states for every possible move
- Create a scores list

- For each of these states add the minimax result of that state to the scores list
- If it's X's turn, return the maximum score from the scores list
- If it's O's turn, return the minimum score from the scores list

You'll notice that this algorithm is recursive, it flips back and forth between the players until a final score is found.

Let's walk through the algorithm's execution with the full move tree, and show why, algorithmically, the instant winning move will be picked:



- It's X's turn in state 1. X generates the states 2, 3, and 4 and calls minimax on those states.
- State 2 pushes the score of +10 to state 1's score list, because the game is in an end state.
- State 3 and 4 are not in end states, so 3 generates states 5 and 6 and calls minimax on them, while state 4 generates states 7 and 8 and calls minimax on them.

- State 5 pushes a score of -10 onto state 3's score list, while the same happens for state 7 which pushes a score of -10 onto state 4's score list.
- State 6 and 8 generate the only available moves, which are end states, and so both of them add the score of +10 to the move lists of states 3 and 4.
- Because it is O's turn in both state 3 and 4, O will seek to find the minimum score, and given the choice between -10 and +10, both states 3 and 4 will yield -10.
- Finally the score list for states 2, 3, and 4 are populated with +10, -10 and -10 respectively, and state 1 seeking to maximize the score will chose the winning move with score +10, state 2.

That is certainly a lot to take in. And that is why we have a computer execute this algorithm.

A Coded Version of Minimax

Hopefully by now you have a rough sense of how the minimax algorithm determines the best move to play. Let's examine my implementation of the algorithm to solidify the understanding:

Here is the function for scoring the game:

```
# @player is the turn taking player
def score(game)
    if game.win?(@player)
        return 10
    elsif game.win?(@opponent)
        return -10
    else
        return 0
    end
end
```

Simple enough, return +10 if the current player wins the game, -10 if the other player wins and 0 for a draw. You will note that *who* the player is doesn't matter. X or O is irrelevant, only who's turn it happens to be.

And now the actual minimax algorithm; note that in this implementation a **choice** or **move** is simply a row / column address on the board, for example [0,2] is the top right square on a 3x3 board.

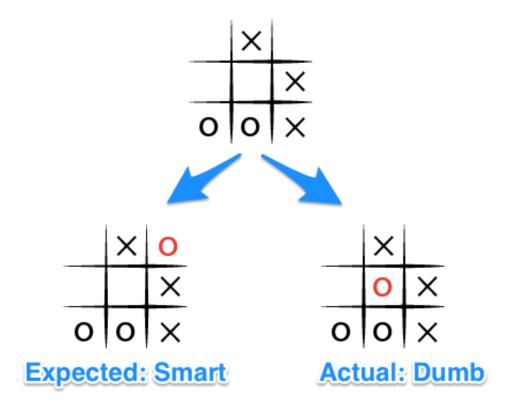
```
def minimax(game)
    return score(game) if game.over?
    scores = [] # an array of scores
    moves = [] # an array of moves
    # Populate the scores array, recursing as needed
    game.get available moves.each do |move|
        possible_game = game.get_new_state(move)
        scores.push minimax(possible game)
        moves.push move
    end
    # Do the min or the max calculation
    if game.active turn == @player
        # This is the max calculation
        max score index = scores.each with index.max[1]
        @choice = moves[max_score_index]
        return scores[max score index]
    else
        # This is the min calculation
        min score index = scores.each with index.min[1]
        @choice = moves[min_score_index]
        return scores[min score index]
    end
end
```

When this algorithm is run inside a PerfectPlayer class, the ultimate choice of best move is stored in the @choice variable, which is then used to return the new game state in which the current player has moved.

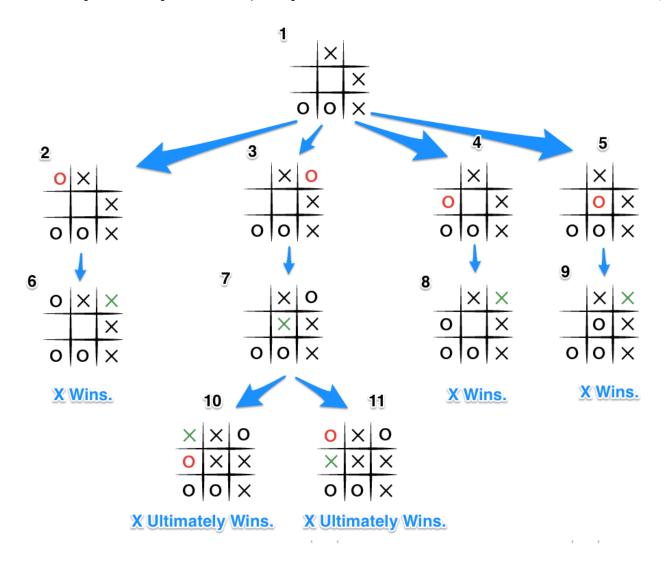
A Perfect but Fatalist Player

Implementing the above algorithm will get you to a point where your tic tac toe game can't be beat. But and interesting nuance that I discovered while testing is that a perfect player must always be perfect. In other words, in a situation where the perfect player eventually will lose or draw, the decisions on the next move are rather fatalistic. The algorithm essentially says: "hey I'm gonna lose anyway, so it really doesn't matter if I lose in the next more or 6 moves from now."

I discovered this by passing an obviously rigged board, or one with a "mistake" in it to the algorithm and asked for the next best move. I would have expected the perfect player to at least put up a fight and block my immediate win. It however, did not:



Let's see what is happening here by looking through the possible move tree (Note, I've removed some of the possible states for clarity):



- Given the board state 1 where both players are playing perfectly, and O is the computer player. O choses the move in state 5 and then immediately loses when X wins in state 9.
- But if O blocks X's win as in state 3, X will obviously block O's potential win as shown in state 7.
- This puts two certain wins for X as shown in state 10 and 11, so no matter which move O picks in state 7, X will ultimately win.

As a result of these scenarios, and the fact that we are iterating through each blank space, from left to right, top to bottom, all moves being equal, that is, resulting in a lose for O, the last move will be chosen as shown in state 5, as it is the last of the available moves in state 1. The array of moves being: [top-left, top-right, middle-left, middle-center].

What is a gosh-darn, tic tac toe master to do?

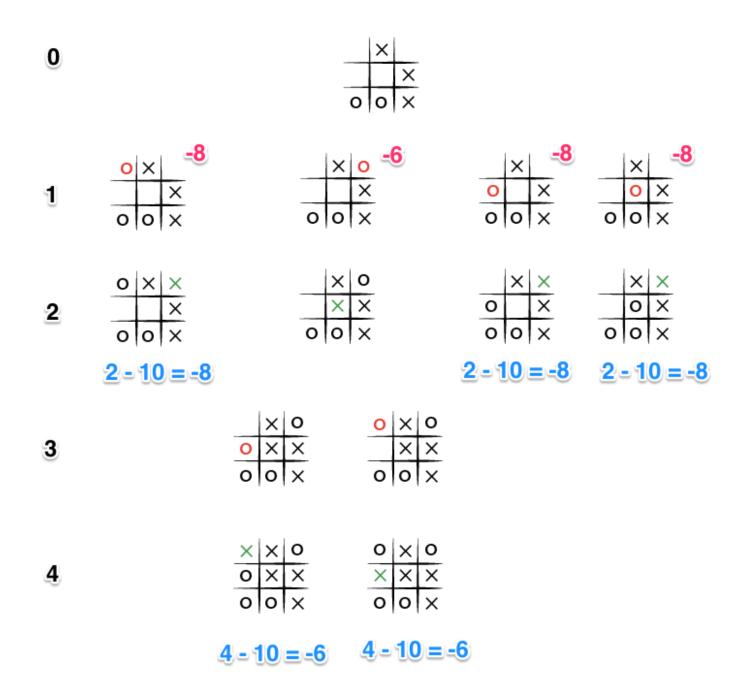
Fighting the Good Fight: Depth

The key improvement to this algorithm, such that, no matter the board arrangement, the perfect player will play perfectly unto its demise, is to take the "depth" or number of turns till the end of the game into account. Basically the perfect player should play perfectly, but prolong the game as much as possible.

In order to achieve this we will subtract the depth, that is the number of turns, or recursions, from the end game score, the more turns the lower the score, the fewer turns the higher the score. Updating our code from above we have something that looks like this:

```
def score(game, depth)
    if game.win?(@player)
        return 10 - depth
    elsif game.win?(@opponent)
        return depth - 10
    else
        return 0
    end
end
def minimax(game, depth)
    return score(game) if game.over?
    depth += 1
    scores = [] # an array of scores
    moves = [] # an array of moves
    # Populate the scores array, recursing as needed
    game.get available moves.each do |move|
        possible game = game.get new state(move)
        scores.push minimax(possible game, depth)
        moves.push move
    end
    # Do the min or the max calculation
    if game.active turn == @player
        # This is the max calculation
        max score index = scores.each with index.max[1]
        @choice = moves[max score_index]
        return scores[max score index]
    else
        # This is the min calculation
        min score index = scores.each with index.min[1]
        @choice = moves[min score index]
        return scores[min score index]
    end
end
```

So each time we invoke minimax, depth is incremented by 1 and when the end game state is ultimately calculated, the score is adjusted by depth. Let's see how this looks in our move tree:



This time the depth (Shown in black on the left) causes the score to differ for each end state, and because the level o part of minimax will try to maximize the available scores (because O is the turn taking player), the -6 score will be chosen as it is greater than the other states with a score of -8. And so even faced with certain death, our trusty, perfect player now will chose the blocking move, rather than commit honor death.

In Conclusion

I hope all of this discussion has helped you further understand the minimax algorithm, and perhaps how to dominate at a game of tic tac toe. If you have further questions or anything is confusing, leave some comments and I'll try to improve the article. All of the source code for my tic tac toe game (http://perfecttictactoe.herokuapp.com/) can be found on github (https://github.com/jasonrobertfox/tictactoe).

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Billy • 6 months ago

It's not a minimax, actually. Minimax presupposes evaluation function for any game situation. And here you can estimate only the result of the very end of the game (win, draw, lose). It's just no more than something like decision tree. Since that it's not right to call it minimax.



Jason Fox Mod → Billy · 6 months ago

Here the goal is to make the game unbeatable, however "This can be extended if we can supply a heuristic evaluation function which gives values to non-final game states without considering all possible following complete sequences" (Wikipedia) This implementation seems to follow all the other minimax implementations I researched in finding all end states.



Zarhock → Jason Fox · a month ago

I think what you are talking about would be the answer to a problem I'm working on right now.

Let's say I am asked to develop a program for a player using minimax algorithm. The depth limit is 2, which means when the program-player plays (max), he can develop the level of his opponent (min), and his next level (max). At that level, it has to estimate a value for every configuration.

What heuristic function could I use for this value? Thanks.



David brear · / months ago

This is by far the best explanation I've seen on minimax for tic tac toe. Everywhere else gave very obscure, often enterprisey solutions but this article explains in a language that anyone can understand. Great article Jason, thank you!



Jason Fox Mod → David Brear ⋅ 6 months ago Thanks!



Randall · 2 months ago

In the last code section, should "return score(game) if game.over?" be "return score(game, depth) if game.over?"?



Jason Fox Mod → Randall • 2 months ago

Yeah possibly, check out the linked repository on github though, the code has likely been updated, a little cleaner, and more complete.



Abouzar Kamaee · 3 months ago

Wonderful Post!



Lorenzo Rossi · 7 months ago

what is the programming language??? java? c? c++? c#????



Jason Fox Mod → Lorenzo Rossi · 7 months ago

Ruby is shown above.



Ankit Singh ⋅ a month ago

thanks jason but why it is not working when i return 1 for win and -1 for lose and 0 for draw ... def score(game, depth)

if game.win?(@player)

return 1 - depth

elsif game.win?(@opponent)

return depth - 1

else

return 0

end end



Jason Fox Mod → Ankit Singh · a month ago

what exactly is not working? i think you have to have a starting score greater than the possible depth, or else you well get negative values.



Ankit Singh → Jason Fox · a month ago

so the problem is the negative value right..?

it means i have to put return value greater than possible game depth..!!



Edward John Laruya ⋅ 3 months ago

do you have a downloadable application of this game?



Jason Fox Mod → Edward John Laruya · 3 months ago

Should be able to set up an Heroku instance and host it there for yourself, or run it on a server locally:

https://github.com/jasonrobert...



Nikhil • 5 months ago

i am not able to download the source code, please help, thanxxxxxxx in advance



Jason Fox Mod → Nikhil • 5 months ago

Should be able to clone or download a zip from here: https://github.com/jasonrobert... if you have a github account.



Christophe De Troyer ⋅ 5 months ago

Excellent article! :) I did notice however I have to implement it for OXO, which is a "more complex" version of tic-tac-toe. This was excellent to get my feet wet :)



jani · 5 months ago

Thanxs ... :-)

_ _ _ _

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Icefire57 · 6 months ago

the score awarded to the win/loss (10) needs to exceed the max possible depth of the game tree, correct?

∧ V · Reply · Share ›



Jason Fox Mod → Icefire57 · 6 months ago
Yep

∧ V · Reply · Share ›



Zheng Luo • 6 months ago Looks Good.



Ahmad Hassanat ⋅ 6 months ago

the time complexity seems to be very high, since each time minimax needs to go down to a goal state, if the state space is very large say 9x9 tic-tac-toe, the algorithm would be very slow.

∧ V · Reply · Share ›



Anish · 7 months ago

Hi Jason. I am working on similar game but it's little complicated than Tic Tac Toe. I was not aware of Al Logic. But your explanation gives me the idea which would be very useful to implement in my game. Thanks:-)

∧ V · Reply · Share ›



jusopi · 7 months ago

Hi Jason. This is a great article. I was recently given the challenge to write a tic-tac-toe game in under 24 hrs. I decided that betw. my day job and family life I was only going to be able to give it 8 hours tops. So after reading several articles on the minimax, and being left scratching my head, I finally found your article. It explained things so simply. Thank you for that. I still didn't quite get the algorithm down pat as I can beat it in some cases (actually it just misses obvious wins). I put it up on github for folks that stumble upon your article and want to see a JavaScript implementation. https://github.com/jusopi/t3



Rinku Saru · 7 months ago

at what moment in game i need to calculate tha min-max? from starting or when 3 move are done?

in this problem only 3 no. r generated i.e -10,0,10, is that all we need in a program or can u

explain the the heurstic i.e statescore generated algorithm

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Jason Fox Mod → Rinku Saru · 7 months ago

In this case I was building an unbeatable game, so I had to examine the game to the end states. I the choice of score was arbitrary at first. It could have been -1 0 or 1 as well. When I added the depth as an additional factor it needed to be at least as large as the maximum depth. In fact the initial score is determined dynamically in the ultimate version of the game, as it allows a larger board (4x4) to be easily implemented.

A more complicated heuristic would be involved if you did not want to explore the move tree to an end state, or make a game where the computer player was smart, but not perfect. You might, for example, look 3 or 4 moves ahead and if the game is not finished, calculate a score based on the state of the board. (One way I read about is the number of "possible lines" that exist for a player.)

Hope this helps:)

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Rinku Saru → Jason Fox · 2 months ago

thanks for help I'm able to implement this in java ,you are great help thanks again .:P



Ryan · 7 months ago

You mean seppuku, not bushido:)

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Jason Fox Mod → Ryan · 7 months ago

You're right! I've changed that.

∧ | ∨ · Reply · Share ›



Soham Bhattacharjee · 7 months ago

it is a bit difficult to understand without the images please check



Jason Fox Mod → Soham Bhattacharjee · 7 months ago

Sorry about this, yeah I accidentally deleted the images like a fool, it's on my list to

ובטומטב נוובווו.



Soham Bhattacharjee → Jason Fox · 7 months ago

i checked the site but the images are still not there; i know its a lot to ask but can you upload them a lil bit faster :-)



Jason Fox Mod → Soham Bhattacharjee · 7 months ago Images are back up.



Soham Bhattacharjee → Jason Fox · 7 months ago thank you man awesome help really appreciate it !!!



nis26 · 8 months ago

Good Job! Very well explained.



Jason Fox Mod → nis26 · 7 months ago Thanks!





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