

## Assignment 3 – Sample Solution

Q1. [6 marks] Parts (a), (b), and (c) are 0.5 mark each meaning 1.5 marks for each of (i), (ii), (iii), and (iv).

R (ABCDEF)

i.  $AD \rightarrow B$ ,  $C \rightarrow D$ ,  $BC \rightarrow A$ ,  $B \rightarrow D$

a) Candidate keys are: ACEF and BCEF.

b) All FDs are BCNF violations.

c)  $C \rightarrow D$  violates  $\Rightarrow$  decompose R into R1(CD) and R2(ABCEF)

R1 is BCNF since it has 2 attributes

For R2, FD is  $BC \rightarrow A$  only, candidate key is BCEF

$BC \rightarrow A$  violates  $\Rightarrow$  decompose R2 into R3(BCA) and R4(BCEF)

R3 is in BCNF since FD is  $BC \rightarrow A$  only and candidate key is BC and  $BC \rightarrow A$  does not violate

R4 is in BCNF since no FDs

Hence the final decomposition is: CD, ABC, and BCEF

If you start with a different FD, you will get the following decompositions:

AB, BD, CD, and ACEF

or

BD, ABC, and BCEF

ii.  $BC \rightarrow E$ ,  $C \rightarrow AB$ ,  $AF \rightarrow CD$

a) Candidate keys are: AF and CF.

b)  $BC \rightarrow E$  and  $C \rightarrow AB$  are the BCNF violations.

c)  $C \rightarrow AB$  violates  $\Rightarrow$  decompose R to R1(CAB) and R2(CDEF)

No more violation  $\Rightarrow$  final decomposition: ABC, CDEF

If you start with  $BC \rightarrow E$ , your final decomposition will be BCE, CAB, and CDF.

iii.  $ABF \rightarrow D$ ,  $CD \rightarrow E$ ,  $BD \rightarrow A$

a) Candidate keys are: ABCF and BCDF.

b) All FDs are BCNF violations.

c)  $CD \rightarrow E$  violates  $\Rightarrow$  decompose R to R1(CDE) and R2(ABCF)

R1 is BCNF since it has two attributes

For R2,  $ABF \rightarrow D$  and  $BD \rightarrow A$  are the FDs and candidate keys are ABCF BCDF

$BD \rightarrow A$  violates  $\Rightarrow$  decompose R2 to R3(ABD) and R4(BCDF)

R3 is in BCNF since FD is  $BD \rightarrow A$  only and candidate key is BD and  $BD \rightarrow A$  does not violate

R4 is in BCNF since no FDs

Hence, final decomposition is CDE, ABD, and BCDF

If you start with a different FD, you will get following decompositions:

CDE, ABD, BDF, and ABCF

or

ABD, BDF, and ABCEF

iv.  $AB \rightarrow D$ ,  $BCD \rightarrow EF$ ,  $B \rightarrow C$

a) Candidate key is: AB.

b)  $BCD \rightarrow EF$  and  $B \rightarrow C$  are the BCNF violations.

c)  $BCD \rightarrow EF$  violates  $\Rightarrow$  decompose R to R1(BCDEF) and R2(ABCD)

$B \rightarrow C$  violates  $\Rightarrow$  decompose R1 and R2 to: BC, BDEF, and ABD

No more violation  $\Rightarrow$  final decomposition: BC, ABD, and BDEF

If you choose  $B \rightarrow C$  to start with, the final decomposition will be BC, ABD, and **ABEF**.

Q2. [4 marks] One (1) mark for each of the four RA queries.

Note that there are many correct answers to this question. The following are just sample answers.

i.

Proj[name](Sel[category.sector = 'Technology'](category Join company))

ii.

R1 = executive

R2 = executive

R3 = executive

R4 = executive

R5 = executive

R6 = executive (since we need more than 5)

Proj[R1.code](Sel[R1.code=R2.code=R3.code=R4.code=R5.code=R6.code and R1.person!=R2.person!=R3.person!=R4.person!=R5.person!=R6.person](R1xR2xR3xR4xR5xR6))

Or

Proj[code] ( Sel[count>5] ( GroupBy[code, Count[code]](executive)) )

iii.

R1 = executive

R2 = executive

Proj[R1.person] ( Sel[R1.person = R2.person and R1.code != R2.code](R1 x R2) )

Or

Proj[person] ( Sel[count>1] ( GroupBy[person, Count[person]](executive)) )

iv.

R1 = Proj[code, industry](category)

R2 = Proj[code, industry](category)

R3 = Proj[R1.code, R1.industry] ( Sel[R1.industry = R2.industry & R1.code != R2.code](R1 x R2) )

Result = R1 - R3

Or

R1 = Proj[industry] ( Sel[count=1] ( GroupBy[industry, Count[industry]](category)) )

Result = Proj[code, industry] (R1 Join category)

Q3. [3 marks] One (1) mark for each expression (0.5 mark for min and 0.5 mark for max)

i.  $R \text{ UNION } (S \text{ INTERSECT } T)$

Min:

Minimum of INTERSECT is 0

Minimum of  $A \text{ UNION } B$  is  $\max(|A|, |B|)$

$\Rightarrow \min(\text{expression}) = \max(r, 0) = r$

Max:

Maximum of  $A \text{ INTERSECT } B$  is  $\min(|A|, |B|)$

Maximum of  $A \text{ UNION } B$  is  $|A| + |B|$

$\Rightarrow \max(\text{expression}) = r + \min(s, t)$

ii.  $\text{SEL}_{[c]} (R \times S)$ , for some condition  $c$

Min:

Minimum is selecting nothing

$\Rightarrow \min(\text{expression}) = 0$

Max:

Maximum is selecting all  $R \times S$  rows

$\Rightarrow \max(\text{expression}) = r * s$

iii.  $\text{PROJ}_{[a]} (R) - \text{PROJ}_{[a]} (R \text{ JOIN } S)$ , for some list of attributes  $a$

Min:

Minimum of  $A \text{ JOIN } B$  is 0

Maximum of  $A \text{ JOIN } B$  is  $|A| * |B|$

$\Rightarrow \min(\text{expression}) = r - (\max \text{ of JOIN})$

$\Rightarrow \min(\text{expression}) = r - r * s$

$\Rightarrow \min(\text{expression}) = r - r = 0$  (if  $r * s = r$ )

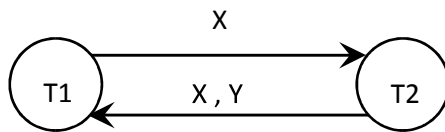
Max:

Maximum of expression is  $r - (\min \text{ of JOIN})$

$\Rightarrow \max(\text{expression}) = r - 0 = r$

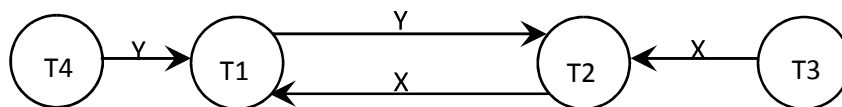
Q4. [2 marks] One (1) mark for each of (i) and (ii) (0.5 mark for the precedence graph and 0.5 mark for the explanation)

i. T1:R(X) T2:R(X) T1:W(X) T2:W(X) T2:R(Y) T1:R(Y) T1:W(Y) T2:W(X)



There is a cycle: [T1] ----X----> [T2] ----X----> [T1] (or [T1] ----X----> [T2] ----Y----> [T1]) in the precedence graph, hence the schedule is not serialisable.

ii. T3:R(X) T4:W(Y) T4:W(Z) T1:W(Y) T2:R(Y) T3:R(D) T2:W(X) T1:R(X)



There is a cycle: [T1] ----Y----> [T2] ----X----> [T1] in the precedence graph, hence the schedule is not serialisable.