This proof was inspired by [Grebinski and Kucherov (optimally...additive model)]. This holds for  $k \geq n$ .

Claim:

$$\sum_{i=0}^{\frac{x}{2}} \Pr[\operatorname{Eq}(q, v_1) = i \wedge \operatorname{Eq}(q, v_2) = i] \le \left(1 - \frac{1}{k}\right)^x$$

Case 1:  $x \le k - 3$ . We know

$$\Pr[\text{Eq}(q, v_1) = i \land \text{Eq}(q, v_2) = i] = \Pr[\text{Eq}(q, v_2) = i] \cdot \Pr[\text{Eq}(q, v_1) = i | \text{Eq}(q, v_2) = i] \\
= \binom{x}{i} \left(\frac{1}{k}\right)^i \left(\frac{k-1}{k}\right)^{x-i} \cdot \binom{x-i}{i} \left(\frac{1}{k-1}\right)^i \left(\frac{k-2}{k-1}\right)^{x-2i} \\
= \frac{x!}{(k-2)^{2i}(i!)^2 (x-2i)!} \left(\frac{k-2}{k}\right)^x \\
\leq \frac{x^{2i}}{(k-2)^{2i}(i!)} \left(\frac{k-2}{k}\right)^x.$$

Now we have

$$\sum_{i=0}^{\frac{x}{2}} \Pr[\operatorname{Eq}(q, v_1) = i \wedge \operatorname{Eq}(q, v_2) = i] \left(\frac{k}{k-1}\right)^x \leq \sum_{i=0}^{\frac{x}{2}} \frac{x^{2i}}{(k-2)^{2i}(i!)} \left(\frac{k-2}{k}\right)^x \left(\frac{k}{k-1}\right)^x$$

$$= \sum_{i=0}^{\frac{x}{2}} \frac{x^{2i}}{(k-2)^{2i}(i!)} \left(1 - \frac{1}{k-1}\right)^x$$

$$< \left(1 - \frac{1}{k-1}\right)^x \sum_{i=0}^{\infty} \left(\frac{x^2}{(k-2)^2}\right)^i \frac{1}{i!}$$

$$< e^{-\frac{x}{k-1}} \cdot e^{\frac{x^2}{(k-2)^2}}$$

$$< 1 \quad \text{for } x \leq (k-3).$$

Multiplying both sides of this inequality by  $(1-1/k)^x$  results in the claim.

Case 2:  $k-2 \le x \le k$ 

$$\sum_{i=0}^{\lfloor x/2 \rfloor} {x \choose i} {x-i \choose i} (n-2)^{-2i} = \sum_{i=0}^{\lfloor x/2 \rfloor} \frac{\frac{x(x-1)\cdots(x-2i+1)}{(n-2)^{2i}}}{i!^2}$$

For x = n - 2, n - 1, or n:

For  $i \ge 3$ ,  $x(x-1)\cdots(x-2i+1) \le n(n-1)\cdots(n-2i+1) \le (n-2)^{2i}$ . Thus,

$$\sum_{i=3}^{\lfloor x/2\rfloor} {x \choose i} {x-i \choose i} (n-2)^{-2i} \le \sum_{i=3}^{\lfloor x/2\rfloor} \frac{1}{(i!)^2}$$

$$\le \sum_{i=3}^{\infty} \frac{1}{(i!)^2}$$

$$< .0296$$

$$\sum_{i=0}^{2} {x \choose i} {x-i \choose i} (n-2)^{-2i} = 1 + \frac{x(x-1)}{(n-2)^2} + \frac{x(x-1)(x-2)(x-3)}{4(n-2)^4}$$

$$\leq 1 + \frac{n(n-1)}{(n-2)^2} + \frac{n(n-1)(n-2)(n-3)}{4(n-2)^4}$$

$$< 2.5536 \text{ for } n \geq 14$$

Thus putting the two together gives

$$\sum_{i=0}^{\lfloor x/2 \rfloor} {x \choose i} {x-i \choose i} (n-2)^{-2i} \le 2.5832$$

$$< \left(1 + \frac{1}{n-2}\right)^{n-2} \text{ for } n \ge 14$$

$$\le \left(1 + \frac{1}{n-2}\right)^x$$

So the sum of the probability that two vectors have the same response on all s questions is:

 $\sum_{x=1}^{n} (\# \text{ of reduced pairs that disagree in } x \text{ spots})(\text{Probability this pair is in the same bucket})$ 

$$= \sum_{x=1}^{n} \binom{n}{x} n^{x} (n-1)^{x} \left(1 - \frac{1}{k}\right)^{s \cdot x}$$

$$\leq \sum_{x=1}^{n} n^{3x} \left(1 - \frac{1}{k}\right)^{s \cdot x}$$

$$= \sum_{x=1}^{n} n^{3x} \left(1 - \frac{1}{k}\right)^{(4k \log n)x}$$

$$< \sum_{x=1}^{n} n^{3x} \left(\frac{1}{e}\right)^{(4 \log n)x}$$

$$= \sum_{x=1}^{n} n^{3x} \left(\frac{1}{n}\right)^{4 \cdot x}$$

$$\leq \sum_{x=1}^{n} \frac{1}{n}$$

$$\leq 1$$