



Data Structures

Introduction

Design and Analysis
of Algorithms I

Data Structures

Point : organize data so that it can be accessed quickly and usefully.

Examples : lists, stacks, queues, heaps, search trees, hashtables, bloom filters, union-find, etc.

Why so Many ? : different data structures support different sets of operations => suitable for different types of tasks.

Rule of Thumb : choose the “minimal” data structure that supports all the operations that you need.



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Data Structures

Heaps and Their Applications

Heap: Supported Operations

- A container for objects that have keys
- Employer records, network edges, events, etc.

Insert: add a new object to a heap.

Running time : $O(\log(n))$

Equally well,
EXTRACT MAX



Extract-Min: remove an object in heap with a minimum key value. [ties broken arbitrarily]

Running time : $O(\log n)$ [n = # of objects in heap]

Also : **HEAPIFY** ($\begin{matrix} n \text{ batched Inserts} \\ \text{in } O(n) \text{ time} \end{matrix}$), **DELETE** ($O(\log(n))$ time)

Application: Sorting

Canonical use of heap : fast way to do repeated minimum computations.

Example : **SelectionSort** $\sim \theta(n)$ linear scans, $\theta(n^2)$ runtime on array of length n

Heap Sort : 1.) insert all n array elements into a heap
2.) Extract-Min to pluck out elements in sorted order

Running Time = $2n$ heap operations = $O(n \log(n))$ time.

=> optimal for a “comparison-based” sorting algorithm!

Application: Event Manager

“Priority Queue” – synonym for a heap.

Example : simulation (e.g., for a video game)

- Objects = event records $\left[\begin{array}{l} \text{Action/update to occur at} \\ \text{given time in the future} \end{array} \right]$
- Key = time event scheduled to occur
- Extract-Min => yields the next scheduled event

Application: Median Maintenance

I give you : a sequence x_1, \dots, x_n of numbers, one-by-one.

You tell me : at each time step i , the median of $\{x_1, \dots, x_i\}$.

Constraint : use $O(\log(i))$ time at each step i .

Solution : maintain heaps H_{Low} : supports Extract Max
 H_{High} : supports Extract Min

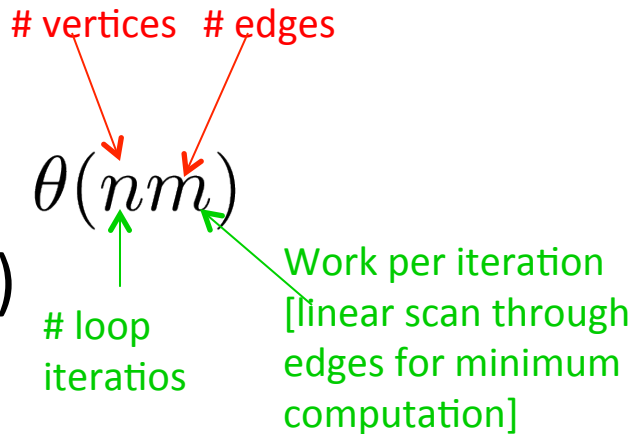
Key Idea : maintain invariant that $\sim i/2$ smallest (largest) elements in
 H_{Low} (H_{High})

You Check : 1.) can maintain invariant with $O(\log(i))$ work
2.) given invariant, can compute median in $O(\log(i))$ work

Application: Speeding Up Dijkstra

Dijkstra's Shortest-Path Algorithm

- Naïve implementation \Rightarrow runtime =
- with heaps \Rightarrow runtime = $O(m \log(n))$



Taking It To The Next Level

Level 0

- “what’s a data structure ?”

Level 1

- cocktail party-level literacy

Level 2

- “this problem calls out for a heap”

Level 3

- “I only use data structures that I create myself”



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Heaps: Some
Implementation Details

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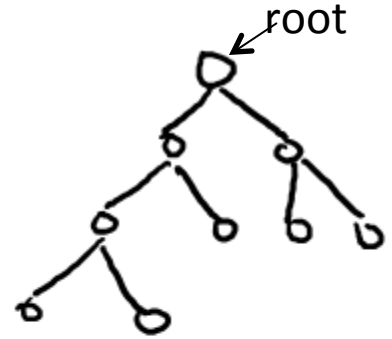
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The Heap Property

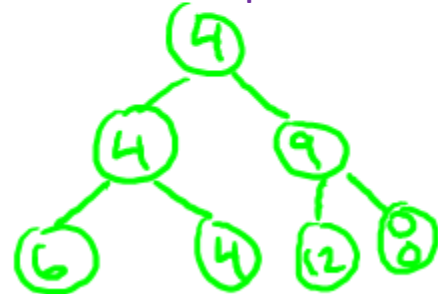
Conceptually : think of a heap as a tree.
-rooted, binary, as complete as possible

Heap Property: at every node x ,
 $\text{Key}[x] \leq \text{all keys of } x\text{'s children}$

Consequence : object at root must
have minimum key value



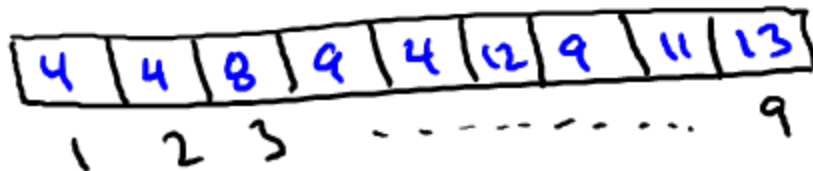
A heap



alternatively



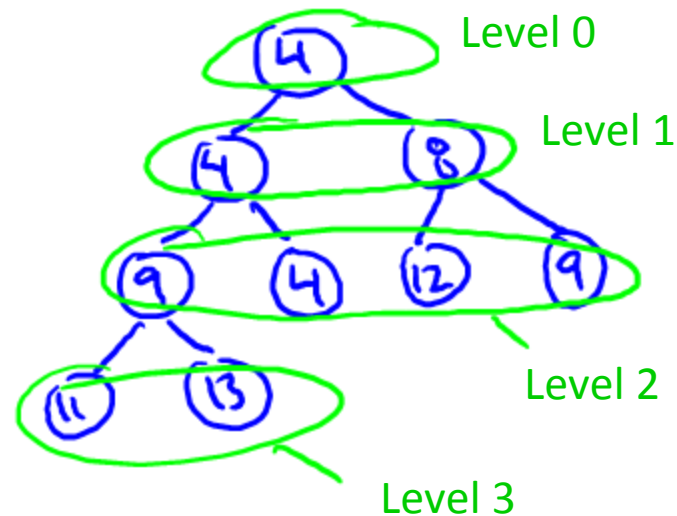
Array Implementation



Note : $\text{parent}(i) = i/2$ if i even
 $= \lfloor i/2 \rfloor$ if i odd

i.e., round down

and children of i are $2i, 2i+1$

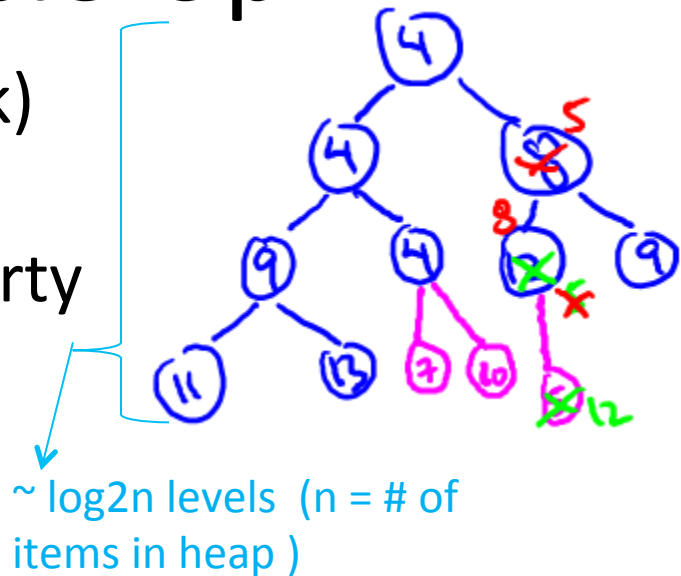


Insert and Bubble-Up

Implementation of Insert (given key k)

Step 1: stick k at end of last level.

Step 2 : Bubble-Up k until heap property is restored (i.e., key of k 's parent is $\leq k$)



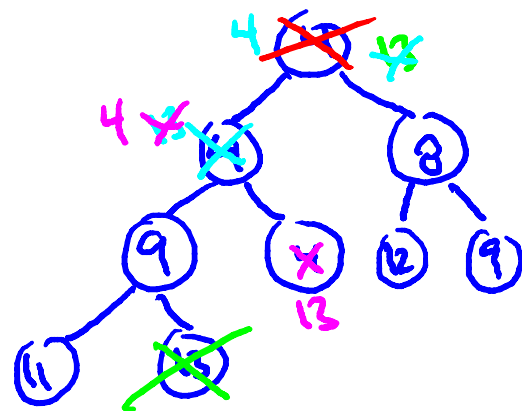
Check : 1.) bubbling up process must stop, with heap property restored
2.) runtime = $O(\log(n))$

Extract-Min and Bubble-Down

Implementation of Extract-Min

1. Delete root
2. Move last leaf to be new root.
3. Iteratively Bubble-Down until heap property has been restored

[always swap with smaller child!]



Check : 1.) only Bubble-Down once per level, halt with a heap
2.) run time = $O(\log(n))$