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Agenda (3/27)
                       Reminder: Exam 2 is next Tuesday (4/2)
- more linked lists
                            · recursion
                            · sorting
- stacks
                            · strings
- queves.
                            · 2D Arrays.
    Struct Node {
        int val;
         struct Node * next
     Struct Nade * head = NULL; // list has no elements
     void insert (struct Mode * * locptr, int key) &
           4 struct Node & new Node = maller (sizeof(struct Node))
               new Mode -> val = 18 key;
               new Node -> next = * locptr ;
                * looph zaneso 4 = newAtcle;
               return;
      3
```

to add 7 to the list: insert (& nead, 7); There is no one way to write Linked Lit functions key feature: a node structure that has data and a pointer to the next Node.

to get direct access to the Noderest pointer in it wints to change. Not a copy

some times you might see linked lists that use a "dwmmy" head node — a node that acts as the head of the list without holding data.

for example:

struct Node * head = malloc (** size of (struct Node)

struct Node * head = malloc (she size of (struct Node *));

// allocate a Mode that doesn't hold data.

head duringth

// write insert in a different way:

Void insert (street Node * nead, int key) {
street Node * new Node = malloc(--)

new Node -> val = key;

new Nade -> next = heard -> next

Nead -> next = new Nade;

3 insert (head, 7)

CI STATE

new Mode

the key thing we needed was direct access to this next pointer "head" is still a pointer to the thing we want to change.

Mote: because insert works on the address of the pointer you want to change, bet's see what happens if we pass in a different pointer's address.

nead

insert (& nead, 2);

head

insert (& (nead > next -> next), 4);

looptr

new Node

1414

rewired the linked list from the <u>middle</u>.

Insert puts a new node right after the next pointer
you pass in.

II) delete a node from the linked list try one: the node head points to.

head = head -> next;

leaks memory, so let's first remember what we need to free struct Node & to Delete = head; head = head -> next; free (to Delete); VI) list that into a function.

The address of the pointer

that points to the node we

void delete (struct Nade ** locptr) { want to delete

struct Node * toDelete = * head locptr;

locptr = (locptr) > next;

Free (**Delete)

delete (& (read -> next)); // deletes the first nock from the list.

VII) delete all the nodes in a list?

while (head != NULL) I

detete (8 head) a don't ferget, we've working on

pointers to pointers

VIII) find a value in the list.

i) if the node is in the list, return the <u>address</u> of the pointer that points to it.

(if it's the first element, return & head.

if it's the second element, return & (head -> next))

ii) if the node 15nt in the list, return the address of the last next pointer (which points to NULL)

Struct Node * * find Eq (struct Node ** looptr, int key) {
while ((* looptr)! = NULL) {

If ((* locptr) -> val == key) return locptr;

3 locptr = &((*locptr) -> next);

3 return locptr;

locate it walks through next pointers

IX) delete a particular key from the list if it exists.

1) figure out if the state key is in the list struct Nucle * to Delete = find Eq. (& head, key)

2) if the key is in the list, remove it:

if (*(toDelete)]= MULL) { // key is in the list clelete (toDelete);

3

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Most common uses of lists are to build two restricted
versions:
                last -in, first -out
   1: Stack (First-in First-oct)
 posh i) add an element to the stack
  pop ii) remove an element from the front of the stock.
  peek iii) look at the element at the front.
        push is just inserting a node at head:
        pop is just removing the node head points to
         peek is just looking at the first nade.
      struct stack 3
           struct Node * head;
           int 512e;
      3
       void push ( & struct Stack . * s, int key) &
            In stodaged
             insert (& (s -> head ) &; , key);
             512644
        3
   int road pop ( strict stack & s) 1
           if (size <1) & handle error }
            int retval = Keadings 5 -> head -> val ;
            delete (&(s -> head));
             51Ze-1
             return retual;
```

3

2: Queve (first-in, first-out) add to one end, delete from the other 1) enqueve (get in line) Localetter beard add to the end 2) dequeve (leave line) just have a pointer L) removes from the Front whose job is to point to the end of the list: tail struct Queve 2 struct Nade & head; struct Node + tail; was void enqueue (struct Queue * q, int key) 2 if (size == 0) 3 q-> head = malloc (-); } special case for q-> tail = head; He First node 512e # Minsert a new Node after fail struct & new Nade = malloc (-); new Node - > val = key; q-> tail -> next = newhale; // put new node after tail q-> tail = q-> fail -> next; 512eH return 3