# Semaphore in RTEMS

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#### Outline

- Introduction of Semaphore in RTEMS
- Priority Inversion
- Priority Inheritance Protocols
- Exercises

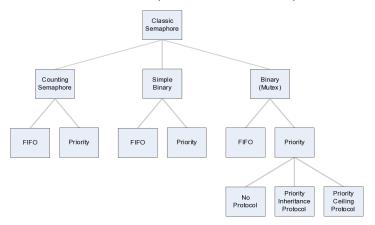






#### Semaphore in RTEMS

Semaphore Attribute Set (Possible combinations)



Source from https://docs.rtems.org/doc-current/share/rtems/html/c\_user/Semaphore-Manager-Building-a-Semaphore-Attribute-Set.html







#### Features of RTEMS

- rtems\_semaphore\_create(name, count, attribute\_set, priority\_ceiling, rtems\_id \*id)
- Some attributes:
  - RTEMS FIFO tasks wait by FIFO (default)
  - RTEMS PRIORITY tasks wait by priority
  - RTEMS COUNTING SEMAPHORE no restriction on values (default)
  - RTEMS BINARY SEMAPHORE restrict values to 0 and 1
  - RTEMS\_NO\_INHERIT\_PRIORITY do not use priority inheritance (default)
  - RTEMS\_NO\_PRIORITY\_CEILING do not use priority ceiling (default)
  - RTEMS\_LOCAL local semaphore (default)
  - ...
- For example: RTEMS\_BINARY\_SEMAPHORE | RTEMS\_FIFO |
  RTEMS\_NO\_INHERIT\_PRIORITY | RTEMS\_NO\_PRIORITY\_CEILING |
  RTEMS\_LOCAL
- Count should be larger than 1 for the normal usage of binary/counting semaphore.







# Dig into the source code cpukit/

- In fact Binary Semaphore in RTEMS is implemented by the structure of Mutex. (Binary Semaphore != Mutex)
- RTEMS interface:
  - rtems\_semaphore\_obtain() is implemented into rtems/src/semobtain.c
  - rtems\_semaphore\_release() is implemented into rtems/src/semrelease.c
- Core functions:
  - \_CORE\_mutex\_Seize\_interrupt\_blocking() in score/src/coremutexseize.c
  - \_Thread\_Raise\_priority() in score/src/threadchangepriority.c
  - inline functions in score/include/rtems/score/.



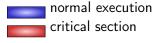


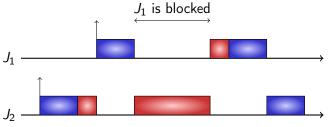


# Priority Inversion

A higher priority job is blocked by a lower-priority job.

Unavoidable when there are critical sections







### Priority Inheritance Protocol (PIP)

When a lower-priority job  $J_j$  blocks a higher-priority job, the priority of job  $J_j$  is promoted to the priority level of highest-priority job that job  $J_j$  blocks.

For example, if the priority order is  $J_1 > J_2 > J_3 > J_4 > J_5$ ,

- When job  $J_4$  blocks jobs  $J_2$  and  $J_3$ , the priority of  $J_4$  is promoted to the priority level of  $J_2$ .
- When job  $J_5$  blocks jobs  $J_1$  and  $J_3$ , the priority of  $J_5$  is promoted to the priority level of  $J_1$ .





### Exercises (10 points)

- Please build the source code and execute SEMAPHORE\_TEST example. Then, draw the diagram to check the system behaviours. (3 points)
- Remove the marked thread\_raise\_priority() in coremutexseize.c to recover PIP behaviours and draw the timing diagram. (2 points)
- Revise PIP to promote the priority of resource holder to the highest priority. What is the drawback? Please explain in the diagram. (5 points)

Hint: coremutexseize.c, taskcreate.c, grep -r "XXX" \*





