

# 3\_recursion

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## 1 Recursion

### 1.1 Outline

- Call Stack
- Recursion
- Time and Space complexity of Recursion
- Mergesort
- Multiple recursive

## 2 Call Stack

### 2.1 How a Call Stack Works

- Your computer internally uses a call stack (stack ADT) to execute functions
- When you run your Python file, the `main` functions is called. `main` is pushed onto the stack
  - Sounds familiar? `if __name__ == "__main__":`
- As the main function executes, it may call other functions, each functions is pushed to the top of the stack
  - The currently executing function is at the top of the stack
- When each function is executed, it is popped from the stack
- The function may return a value, which is passes to the calling function (the function below in the stack).
- The calling function can use the return value and continue execution until the stack is empty

### 2.2 Basic Example

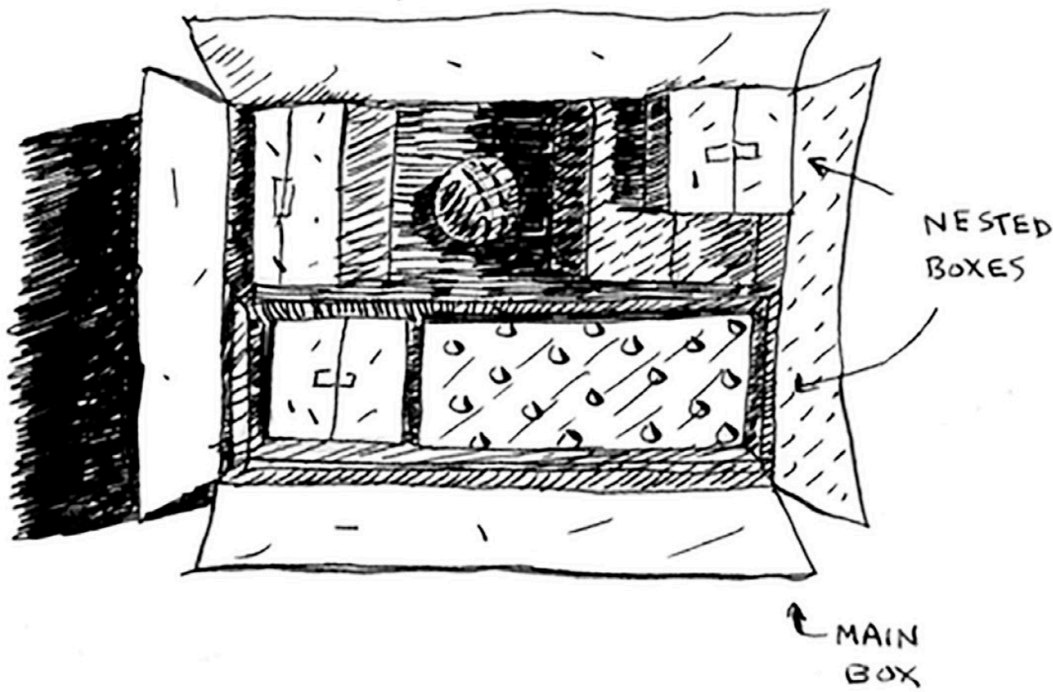
- If I run `round(float("20.24"))`, I expect 20
  - The `round` function is first to be called, it is pushed on the call stack
  - Then, `float("20.24")` is called and pushed on the call stack
- Now, we pop each function off the call stack.

- `float("20.24")` returns 20.24
- `round` uses the return value of the previous function, 20.24. It executes `round(20.24)`, which returns 20
- The stack is empty, so the program finishes

## 3 Recursion

### 3.1 Motivating Example

- Suppose you are looking for a key in a box, but the box contains more boxes!



- 2 minutes: write down the steps of the algorithm you would take to search for the key

### 3.2 Algorithm 1: Loop

1. Make a pile of all the boxes
2. Grab a box and open it
3. If it contains a box, append it to your pile of boxes
4. If it contains the key, you're done!
5. Repeat

### 3.3 Algorithm 2: Recursion

1. Grab a box and open it

2. If it contains a box, repeat step 1
3. If it contains the key, you're done!

Which algorithm do you like more?

- Notice the function is recursive because it calls itself
- Both algorithms achieve the same thing, but recursion is clearer (to me)

```
[ ]: # Write your answer here
```

### 3.4 Formula to write a recursive function

- Since recursive functions call themselves, its easy to write an infinite loop
- Let's write a function that does a countdown

```
[1]: def countdown(i):  
    print(i)  
    countdown(i - 1)
```

- This runs forever, so we need a *base case* to tell the code when to stop

```
[2]: def countdown(i):  
    print(i)  
    if i <= 0:  
        return  
    else:  
        countdown(i - 1)
```

### 3.5 Factorial

- The *factorial* is the product of all positive integers less than or equal to the given integer
  - $5! = 5 \times 4 \times 3 \times 2 \times 1 = 120$
  - We define  $1! = 1$
- Let's use recursion to calculate factorials

```
[3]: def factorial(n):  
    if x == 1:  
        return 1  
    else:  
        return x * factorial(x - 1)
```

- Let's examine the call stack when we call `factorial(3)`

### 3.6 Recursion and the Call Stack

Code		Call Stack	
fact(3)		fact	
		x	3
if x == 1:		fact	
		x	3
else:		fact	
		x	3
return x*fact(x-1)		fact	
		x	2
		fact	
		x	3

### 3.7 Recursion and the Call Stack

if x == 1:		fact	
		x	2
		fact	
		x	3
else:		fact	
		x	2
		fact	
		x	3
return x*fact(x-1)		fact	
		x	1
		fact	
		x	2
		fact	
		x	3

### 3.8 Recursion and the Call Stack

if x == 1		fact	
		x	1
		fact	
		x	2
		fact	
		x	3
return 1		fact	
		x	1
		fact	
		x	2
		fact	
		x	3
return x*fact(x-1)		fact	
		x	2
		fact	
		x	3
return x*fact(x-1)		fact	
		x	3

### 3.9 Multiple Recursive Calls: Fibonacci Sequence

- In calculating the factorial, each recursion only calls itself once. This doesn't have to be the case
- The Fibonacci Sequence is a sequence of numbers where the first two numbers are 0 and 1, with each subsequent number being the sum of the previous two numbers in the sequence.
  - Notice how the problem is defined recursively

```
[4]: def fib(n):  
    if n <= 1:  
        return n  
    else:  
        return fib(n - 1) + fib(n - 2)
```

### 3.10 2 minutes: what is its time and space complexity?

```
[ ]: # Your answer here
```

## 4 Time and Space Complexity of Recursion

### 4.1 Time Complexity of Recursion

- Generally, recursion doesn't have performance benefits compared to loops (in problems like finding a key in nested boxes)
  - However, it is simpler to understand
- The time complexity of recursion depends on the number of times the function calls itself (branches)
  - Factorial: the **fact** is called  $n$  times before reaching the base case so its  $O(1^n) = O(n)$
  - If a recursive function called itself twice, then its  $(2^n)$
- When a recursive function makes multiple calls, the run time will often be  $O(\text{branches}^{\text{depth}})$

### 4.2 Tricky Example

```
[5]: def recursive(n):  
    for i in range(n):  
        # Something happens  
        i += 2  
    if n <= 0:  
        return 1  
    else:  
        return 1 + recursive(n - 3)
```

- Loop takes  $n/2$  steps, because we increase `i` by 2

- Recursion takes  $n/3$  steps **and** the loop is called recursively.
  - In other words, for each recursion, run the loop.
- The time complexity is  $n/2 \times n/3 = \frac{n^2}{6} = O(n^2)$

### 4.3 Space complexity of recursion

- Notice the call stack takes up space in memory. How much depends on the depth of the recursion
- Think about the maximum amount of space the call stack will need
  - Factorial:  $O(n)$ , when recursion reaches the base case
- Even when you have multiple branches, it's possible only 1 branch at depth  $n$  is in memory at a time

### 4.4 2 minutes: to find the key in nested boxes, what is the memory complexity of the recursive approach versus the loop approach?

[ ]: *# Your answer here*

### 4.5 Live Coding

Given an list of positive integers and an integer x, we want to find all unique combinations in the list where the sum is equal to x. A number in the list can be used multiple times.

#### 4.5.1 Example

```
[6]: # INPUT
lst = [1,2,5,6]
x = 6
# OUTPUT
[1, 5]
[6]
```

[ ]: *# Your code here*

## 5 Mergesort

### 5.1 Divide and Conquer Algorithms

- Divide and Conquer (D&C) is a general method to solve problems utilizing recursion.
  - Figure out the simplest case and use it as the base case
  - Figure out how to reduce your problem to the base case
- Let's start with a trivial example: how would you sum a list of integers?
  - Solution is obvious with a loop



- Let's do it recursively

## 5.2 Divide and Conquer Algorithms

Step 1

- What is the simplest array to sum?
- Arrays with no elements or 1 element
  - sum of [] is 0, sum of [8] is 8

Step 2

- How can we reduce all arrays to empty array?
- Notice  $\text{sum}[2, 4, 5] = 2 + \text{sum}[4, 5]$ , but the second version reduced the problem

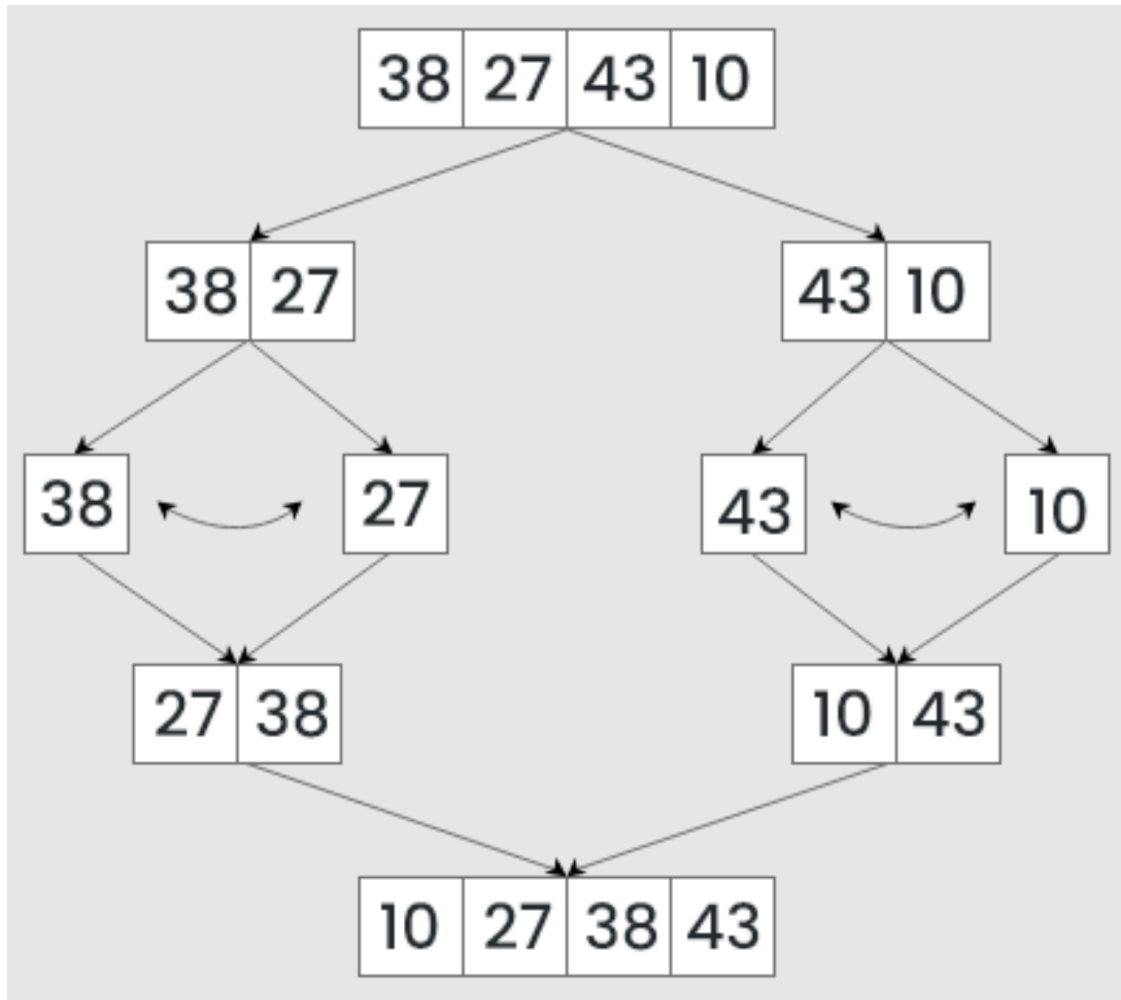
## 5.3 Divide and Conquer Algorithms

```
[7]: def rec_sum(lst):  
    if not lst:  
        return 0  
    else:  
        return lst[0] + rec_sum(lst[1:])
```

- Let's work on a real problem next!

## 5.4 Mergesort

- Some lists don't need to be sorted
  - Lists of size 1! This is our base case
- We can split lists in half until they contain 1 element, then merge all of the sub-lists
- Python's sort function uses a hybrid of merge and insertion sort, both of which you've learned!



## 5.5 Big-O of Merge Sort

First consider the non-recursive part of the code

- The “divide” step takes linear time, since slicing operations take roughly  $n/2$  steps to make a left and right copy respectively.
- The merge operation also takes  $n$  steps approximately
- All other operations are constants
- Together, the non-recursive part of this algorithm is  $O(n)$

Next consider the recursive calls

- Recall the big-O of recursion depends on the recursion depth and number of calls.  $O(\text{branches}^{\text{depth}})$
- The depth in Merge Sort is the number of times you need to divide to get to a list of length 1.
- Mathematically,  $2^{\text{depth}} = n$ , then  $\text{depth} = \log n$ . So there are approximately  $\log n$  levels

## 5.6 Big-O of Mergesort

- Since the  $O(n)$  steps must be performed each recursion, the total run time is  $O(n \log n)$ . Our analysis only depended on the size of the list, so the best and worst case of mergesort is the same
- This is much faster than insertion sort!
- 2 minutes: does it have less space complexity than insertion sort?

## 6 Recommended Problems and References

### 6.1 Recommended Problems

- Bhargava: Chapter 4 exercises
  - 4.1 to 4.8
- Write a recursive function that produces the `RecursionError: maximum recursion depth exceeded` error.
- Write a iterative function to calculate the  $n$ th Fibonacci number. What is its time and space complexity?
- Write a recursive function to determine if a string is a palindrome. What is its time and space complexity?
- Write a recursive function to check if a given positive integer is a prime number. What is its time and space complexity?

### 6.2 Recommended Problems

- Suppose you have a plot of land and want to divide the land into even square plots, while keeping the plots as big as possible. How would you do this using D&C? See Bhargava pg. 52.
- Explain why the “merge” step in mergesort is  $O(n)$
- Implement mergesort. You might find using helper functions useful.
- Write a recursive function to perform binary search on a sorted list

### 6.3 Bonus Readings

- You may be interested in learning more about quicksort in Bhargava chapter 4 or [here](#). Quicksort is another recursive sorting method

### 6.4 References

- Bhargava, A. Y. (2016). *Grokking algorithms: An illustrated guide for programmers and other curious people*. Manning. Chapter 3 and 4.
- Cormen, T. H. (Ed.). (2009). *Introduction to algorithms* (3rd ed). MIT Press. Chapter 4.