ext4 文件系统 调研汇报

2023 春 OS 大实验交流

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上周主要工作

- 调研 ext4 文件系统的定义与设计
- 调研 ext4 文件系统的实现 (Rust 版本与其他语言版本)
- ArceOS 组件化操作系统 Hello World



ext4 Filesystem: Introduction

- ext4 is a journaling file system developed as an improvement over ext3. It is the *default file system for many popular Linux distributions*. It provides better performance and reliability than its predecessor.
- The <u>Linux kernel documentation</u> provides a detailed description of ext4's data structures and algorithms. :)
- We can also refer to this link for information about general usage and ondisk artifacts written by ext4. A wiki site is also available here.



ext4 Filesystem: Design > Blocks

• ext4 allocates storage space in units of "blocks", and the filesystem is split into a series of block groups.

For 64-bit filesystems, limits are as follows:

Item	1KiB	2KiB	4KiB	64KiB	
Blocks	2^64	2^64	2^64	2^64	
Inodes	2^32	2^32	2^32	2^32	
File System Size	16ZiB	32ZiB	64ZiB	1YiB	
Blocks Per Block Group	8, 192	16, 384	32, 768	524, 288	
Inodes Per Block Group	8, 192	16, 384	32, 768	524, 288	
Block Group Size	8MiB	32MiB	128MiB	32GiB	
Blocks Per File, Extents	2^32	2^32	2^32	2^32	
Blocks Per File, Block Maps	16, 843, 020	134, 480, 396	1, 074, 791, 436	4,398,314,962,956 (really 2^32 due to field size limitations)	
File Size, Extents	4TiB	8TiB	16TiB 256TiB		
File Size, Block Maps	16GiB	256GiB	4TiB	256TiB	



ext4 Filesystem: Design > Layout

The layout of a standard block group is approximately as follows (each of these fields is discussed in a separate section below):

Group 0 Padding	ext4 Super Block	Group Descriptors	Reserved GDT Blocks	Data Block Bitmap	inode Bitmap	inode Table	Data Blocks
1024 bytes	1 block	many blocks	many blocks	1 block	1 block	many blocks	many more blocks



ext4 Filesystem: Design > Layout

- Group 0 Padding: to allow for the installation of x86 boot sectors and other oddities.
- Super Block: The superblock records various information about the enclosing filesystem, such as block counts, inode counts, supported features...
 - We need redundancy to better protect the superblock. If the *sparse_super* feature flag is set, redundant copies of the superblock and group descriptors are kept only in the groups whose group number is either 0 or a power of 3, 5, or 7. If the flag is not set, redundant copies are kept in all groups.



ext4 Filesystem: Design > Layout

- Group Descriptors: Each block group on the filesystem has one of these descriptors associated with it.
 - The group descriptor records the location of both bitmaps and the inode table.
 - Two mechanisms: flex bg and meta bg.
 - The flex bg mechanism groups several block groups into a flex group and lay out all of the groups' bitmaps and inode tables into one long run in the first group of the flex group.
 - The meta bg mechanism enforces that only the first and last two block groups within the larger meta group contains group descriptors for the groups inside the meta group.



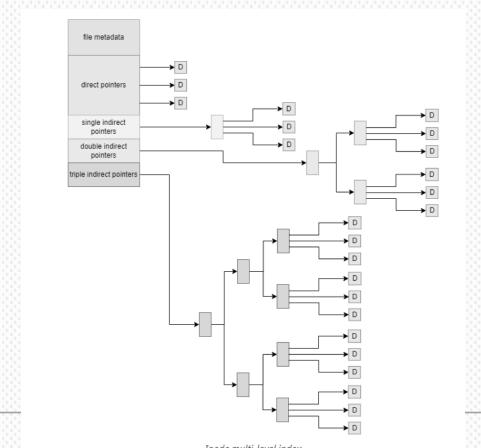
ext4 Filesystem: Design > Layout

- Block and inode Bitmaps: The *data block bitmap* tracks the usage of *data blocks* within the block group. The *inode bitmap* records which *entries* in the inode table are in use.
- Inode Table: The inode table contains the *inode* structures for all the files and directories in the block group.
- Data Blocks: The data blocks are used to store the actual file data.



ext4 Filesystem: Design > Inode Table Entries

• Either represent a file or a directory.





ext4 Filesystem: Design > Directory Entries

- In an ext4 filesystem, a directory is more or less a flat file that maps an arbitrary byte string (usually ASCII) to an inode number on the filesystem.
 - Linear (Classic) Directories
 - Hash Tree Directories: Use B-Trees to store directory entries.

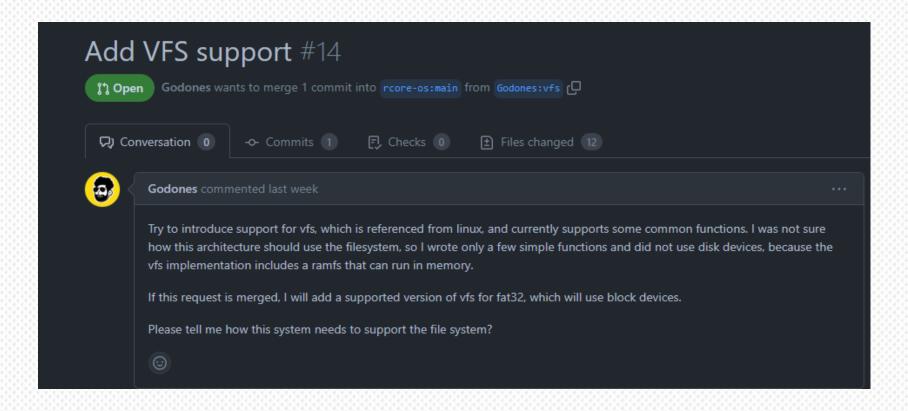


ext4 文件系统:实现

- Read files directly from ext4 filesystem images: https://github.com/FauxFaux/ext4-rs
- Standard ext4 filesystem implementation in Linux (written in C): https://github.com/torvalds/linux/tree/master/fs/ext4



AcreOS 组件化操作系统



• 和助教学长交流后等待 fs 的重构, 之后再开始实现:)



下周计划

- 仔细阅读 ext4 文件系统的实现 Code
- 如果 fs 已经并入 ArceOS, 那么就开始上手实现 ext4 文件系统的数据结构与相关算法:)

