

## Exchange Economy with Externalities and Complex Utility Functions

Consider a pure exchange economy with two agents and two goods. Let  $x_{ij}$  be the consumption of good  $i$  by agent  $j$

The agents' utility functions are:

$$u_1(x_{11}, x_{21}) = \min\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\}$$

$$u_2(x_{12}, x_{22}) = 3x_{12} + x_{21}$$

The initial endowments are:

$$\omega_1 = (0, \beta) \quad \text{and} \quad \omega_2 = (1, 0)$$

- (a) Find the set of Pareto-optimal allocations in terms of  $\beta$
- (b) Solve the utility maximization problem for both agents to determine their demands as functions of prices and  $\beta$

## Solution

### (a) Pareto Optimal Allocations

First, observe that the total endowments are

$$x_{11} + x_{12} = 1 \quad \text{and} \quad x_{21} + x_{22} = \beta$$

so any feasible allocation must satisfy those equations and non-negativity of consumptions

**Agent 1's utility:**

$$u_1(x_{11}, x_{21}) = \min\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\}$$

**Agent 2's utility:**

$$u_2(x_{12}, x_{22}) = 3x_{12} + x_{22} \quad (\text{note that } x_{21} \text{ is Agent 1's consumption of good 2 according to the problem statement})$$

### Checking Pareto improvements from the initial endowment

The initial endowment is

$$\omega_1 = (0, \beta), \quad \omega_2 = (1, 0)$$

which implies the allocation

$$(x_{11}, x_{21}) = (0, \beta), \quad (x_{12}, x_{22}) = (1, 0)$$

Under this allocation

$$u_1(0, \beta) = \min\{0 + 3\beta, 4 \cdot 0 + \beta\} = \min\{3\beta, \beta\} = \beta \quad (\text{assuming } \beta > 0)$$

$$u_2(1, 0) = 3 \cdot 1 + 0 = 3$$

Hence, at the initial endowment,

$$u_1 = \beta \quad \text{and} \quad u_2 = 3$$

A direct examination shows that any attempt to reallocate goods in a way that increases one agent's utility strictly ends up decreasing the other's utility. For example:

- Giving good 1 from Agent 2 to Agent 1 reduces  $x_{12}$  and thus reduces  $3x_{12}$  in  $u_2$
- Giving good 2 from Agent 1 to Agent 2 reduces  $x_{21}$ , which directly appears as  $+x_{21}$  in  $u_2$ , so that also hurts Agent 2
- Interior reallocations end up lowering either  $u_1$  or  $u_2$  below their initial values

Therefore, the initial endowment itself is already Pareto optimal

### Generalizing to all possible reallocations

To see the full range of Pareto-optimal allocations, consider a general allocation  $(x_{11}, x_{21}; x_{12}, x_{22})$  subject to

$$x_{11} + x_{12} = 1 \quad x_{21} + x_{22} = \beta \quad x_{ij} \geq 0$$

We have

$$u_1(x_{11}, x_{21}) = \min\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\} \quad \text{and} \quad u_2(x_{12}, x_{22}) = 3x_{12} + x_{22}$$

**(i) All of good 2 goes to Agent 1.** In  $u_2 = 3x_{12} + x_{22}$ , the partial derivative with respect to  $x_{21}$  is  $+1$ , so Agent 2 strictly prefers that any amount of good 2 be in the hands of Agent 1 (since it appears positively in  $u_2$ ). Meanwhile, Agent 1 is not harmed if  $x_{21}$  is larger, because  $x_{21}$  also enters positively in both expressions inside the  $\min\{\}$  of  $u_1$ . Consequently, any amount of good 2 given to Agent 2 can be reallocated to Agent 1 to strictly increase  $u_2$  without lowering  $u_1$ . Thus Pareto efficiency requires

$$x_{21} = \beta \quad \text{and} \quad x_{22} = 0$$

**(ii) Splitting good 1.** Since  $x_{21} = \beta$  is fixed for Pareto efficiency, the only remaining choice is how to split the total of 1 unit of good 1 between the two agents. Let  $t \in [0, 1]$  be the fraction of good 1 that goes to Agent 1:

$$x_{11} = t \quad x_{12} = 1 - t$$

so the allocation is

$$(x_{11}, x_{21}; x_{12}, x_{22}) = (t, \beta; 1 - t, 0)$$

Their utilities become

$$u_1(t, \beta) = \min\{t + 3\beta, 4t + \beta\}$$

$$u_2(1 - t, \beta) = 3(1 - t) + \beta = 3 + \beta - 3t$$

No further reallocation of good 2 is beneficial because of the positive effect on both agents of having  $x_{21} = \beta$ . Thus the entire Pareto frontier is spanned by  $t \in [0, 1]$ .

Hence, for any  $\beta > 0$ , the full set of Pareto-optimal allocations is  $\{(t, \beta; 1 - t, 0) : 0 \leq t \leq 1\}$ , and the initial endowment corresponds to  $t = 0$ .

## (b) Utility Maximization and Demands

### Agent 2's problem

Agent 2 has the endowment  $(1, 0)$ , so their income is

$$I_2 = p_1 \cdot 1 + p_2 \cdot 0 = p_1$$

Agent 2 maximizes

$$u_2(x_{12}, x_{22}) = 3x_{12} + x_{22}$$

subject to

$$p_1 x_{12} + p_2 x_{22} \leq p_1, \quad x_{12}, x_{22} \geq 0$$

Notice that  $x_{21}$  is *not* chosen by Agent 2; it is Agent 1's consumption of good 2. Therefore, for a standard Walrasian demand, only the term  $3x_{12}$  is actually controlled by Agent 2. From Agent 2's perspective,  $x_{21}$  is a constant (an externality determined by Agent 1)

Focusing on the part they *do* control, Agent 2 solves

$$\max_{x_{12}, x_{22} \geq 0} 3x_{12}$$

subject to

$$p_1 x_{12} + p_2 x_{22} \leq p_1$$

This is a linear program in  $(x_{12}, x_{22})$ . Since only  $x_{12}$  yields positive marginal utility (equal to 3), Agent 2 spends the entire budget on good 1. Hence

$$x_{12} = \frac{p_1}{p_1} = 1, \quad x_{22} = 0$$

Thus, the demand for Agent 2 is

$$\text{Agent 2's Demand: } (x_{12}^*, x_{22}^*) = (1, 0)$$

### Agent 1's problem

We want to solve the following utility maximization problem:

$$\max_{x_{11}, x_{21} \geq 0} \min\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\} \quad \text{subject to} \quad p_1 x_{11} + p_2 x_{21} \leq p_2 \beta$$

### Step 1: Check for an interior solution

For an interior solution that makes both expressions inside the  $\min\{\cdot, \cdot\}$  equal, we set:

$$x_{11} + 3x_{21} = 4x_{11} + x_{21} \implies 3x_{21} - x_{21} = 4x_{11} - x_{11} \implies 2x_{21} = 3x_{11} \implies x_{21} = \frac{3}{2}x_{11}$$

Plug this relation into the budget constraint:

$$p_1 x_{11} + p_2 \left(\frac{3}{2}x_{11}\right) = x_{11} \left(p_1 + \frac{3}{2}p_2\right) \leq p_2 \beta$$

At optimum (assuming an interior solution is indeed optimal), the budget is fully spent, so

$$x_{11} = \frac{p_2 \beta}{p_1 + \frac{3}{2}p_2} \quad , \quad x_{21} = \frac{3}{2} \frac{p_2 \beta}{p_1 + \frac{3}{2}p_2}$$

### Step 2: Corner solutions

Because utility is the minimum of two linear expressions, we must also check corners where the consumer spends all income on a single good:

1. **Corner 1:**  $x_{21} = 0$ . Then

$$u_1 = \min\{x_{11}, 4x_{11}\} = x_{11}$$

and the budget constraint becomes  $p_1 x_{11} \leq p_2 \beta$ . Hence

$$x_{11} = \frac{p_2 \beta}{p_1} \quad \text{and} \quad u_1 = \frac{p_2 \beta}{p_1}$$

2. **Corner 2:**  $x_{11} = 0$ . Then

$$u_1 = \min\{3x_{21}, x_{21}\} = x_{21}$$

(because  $x_{21} \leq 3x_{21}$  for  $x_{21} \geq 0$ ), and the budget constraint is  $p_2 x_{21} \leq p_2 \beta$ . Thus

$$x_{21} = \beta \quad \text{and} \quad u_1 = \beta$$

### Step 3: Compare utilities

Let

$$U_{\text{int}} = \min\left\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\right\} \quad (\text{under the interior solution}) = \frac{11}{2}x_{11} = \frac{11p_2\beta}{2p_1 + 3p_2}$$

$$U_{\text{corner1}} = \frac{p_2\beta}{p_1} \quad , \quad U_{\text{corner2}} = \beta$$

By comparing  $U_{\text{int}}$  with  $U_{\text{corner1}}$  and  $U_{\text{corner2}}$ , we find the Marshallian demands are piecewise according to the ratio  $\frac{p_1}{p_2}$ :

$$\begin{aligned} &\text{If } \frac{p_1}{p_2} < \frac{1}{3}, \quad \text{Corner 1 is better} \\ &\text{If } \frac{1}{3} \leq \frac{p_1}{p_2} \leq 4, \quad \text{Interior solution is better} \\ &\text{If } \frac{p_1}{p_2} > 4, \quad \text{Corner 2 is better} \end{aligned}$$

#### Step 4: Final Marshallian demand

$$x_{11}^*(p_1, p_2, \beta) = \begin{cases} \frac{p_2 \beta}{p_1} & \text{if } \frac{p_1}{p_2} < \frac{1}{3} \\ \frac{p_2 \beta}{p_1 + \frac{3}{2}p_2} & \text{if } \frac{1}{3} \leq \frac{p_1}{p_2} \leq 4 \\ 0 & \text{if } \frac{p_1}{p_2} > 4 \end{cases}$$

$$x_{21}^*(p_1, p_2, \beta) = \begin{cases} 0 & \text{if } \frac{p_1}{p_2} < \frac{1}{3} \\ \frac{3}{2} \frac{p_2 \beta}{p_1 + \frac{3}{2}p_2} & \text{if } \frac{1}{3} \leq \frac{p_1}{p_2} \leq 4 \\ \beta & \text{if } \frac{p_1}{p_2} > 4 \end{cases}$$

Therefore, the Marshallian demand for the agent with utility  $\min\{x_{11} + 3x_{21}, 4x_{11} + x_{21}\}$  is given by the above piecewise functions