

CET 241: Day 5

Dr. Noori KIM

Agenda

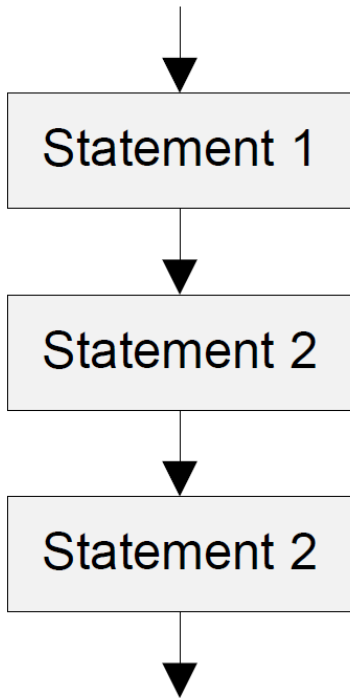
- Top down design
- Reuse of registers
- Passing Parameters to Subroutines via Register
 - “ARM Procedure Call Standard”
 - An example
- Preserve Environment via Stack

"Nothing is particularly hard if you divide it into small jobs."

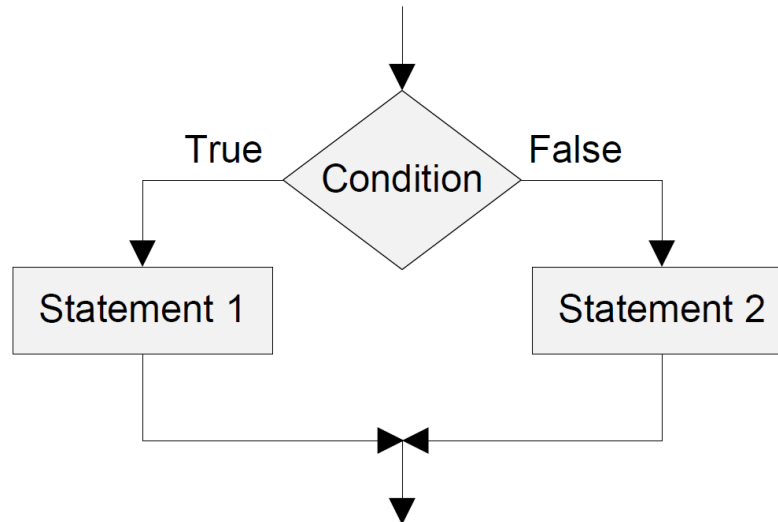
Henry Ford, Founder of Ford Motor

➔ The beauty of Divide and Conquer

Three basic control structures

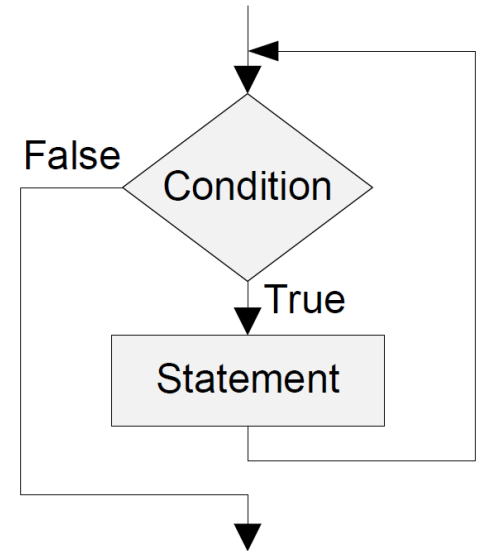


Sequence Structure



Selection Structure

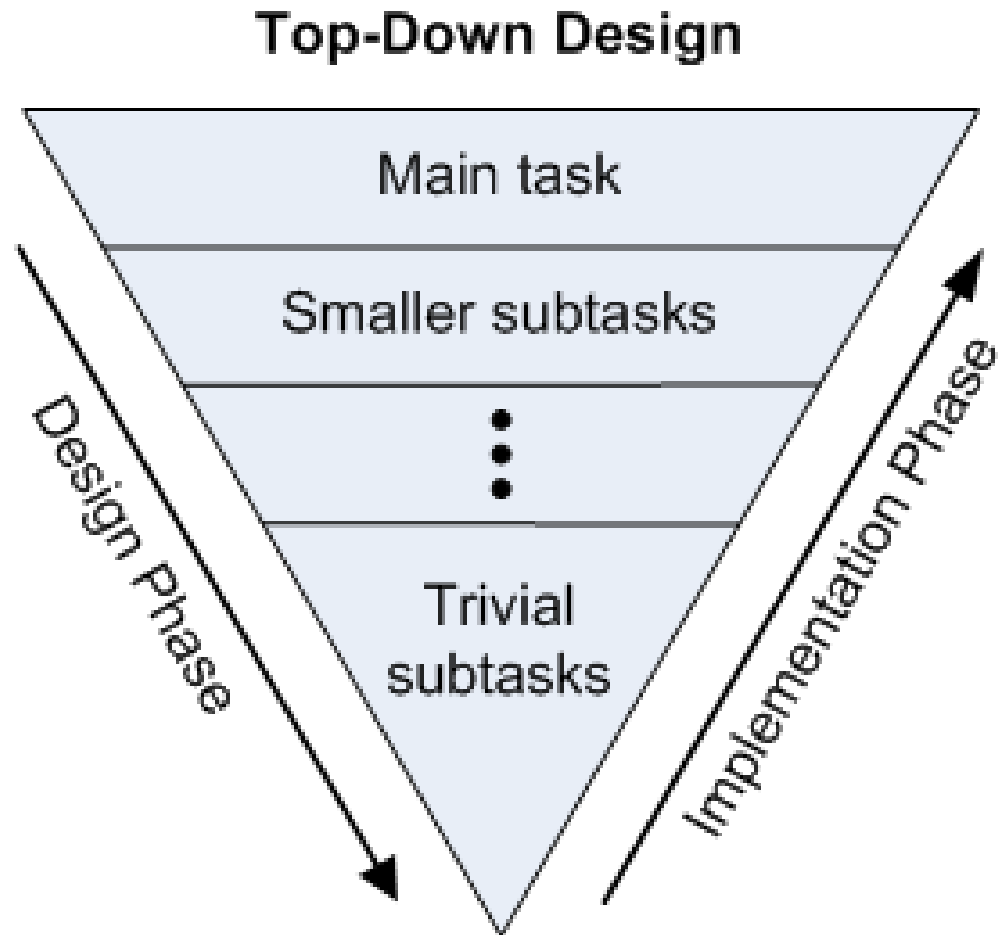
- If, else-if, else
- switch



Loop Structure

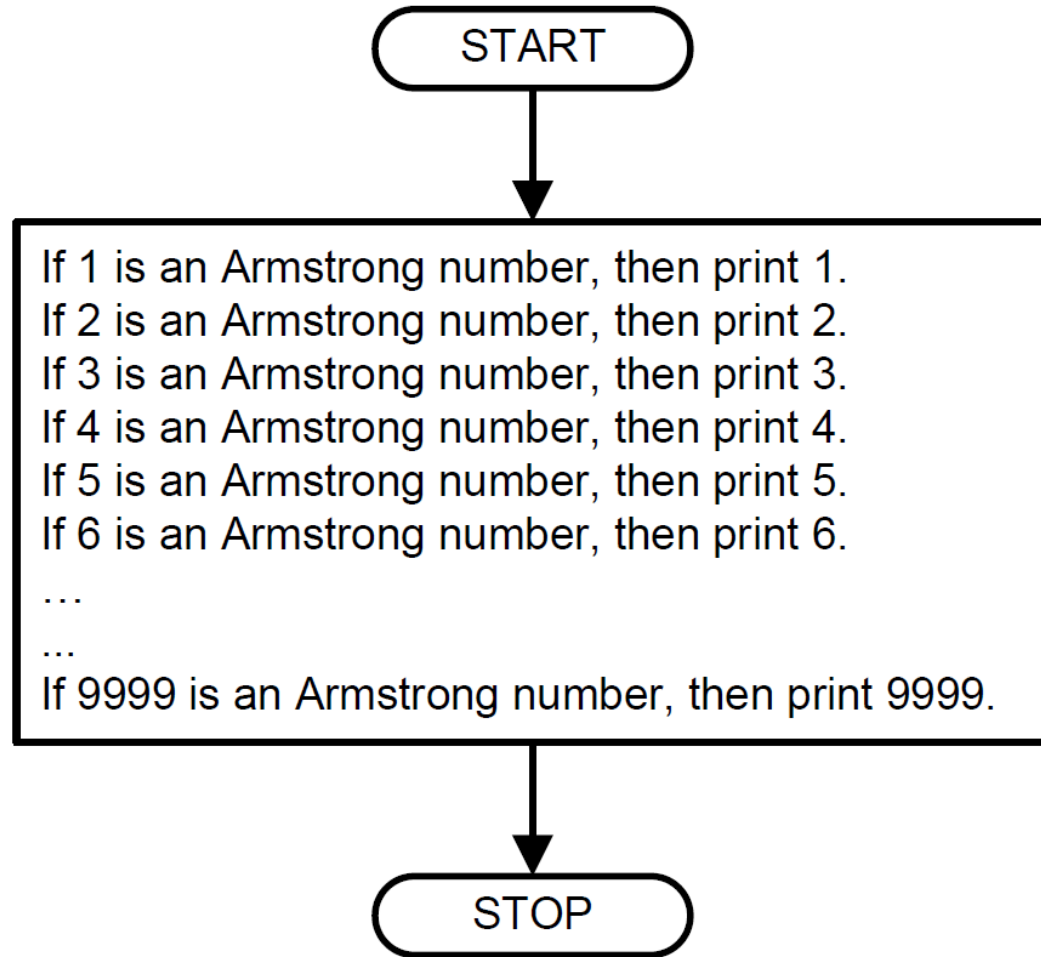
- for
- while

Top-Down Design

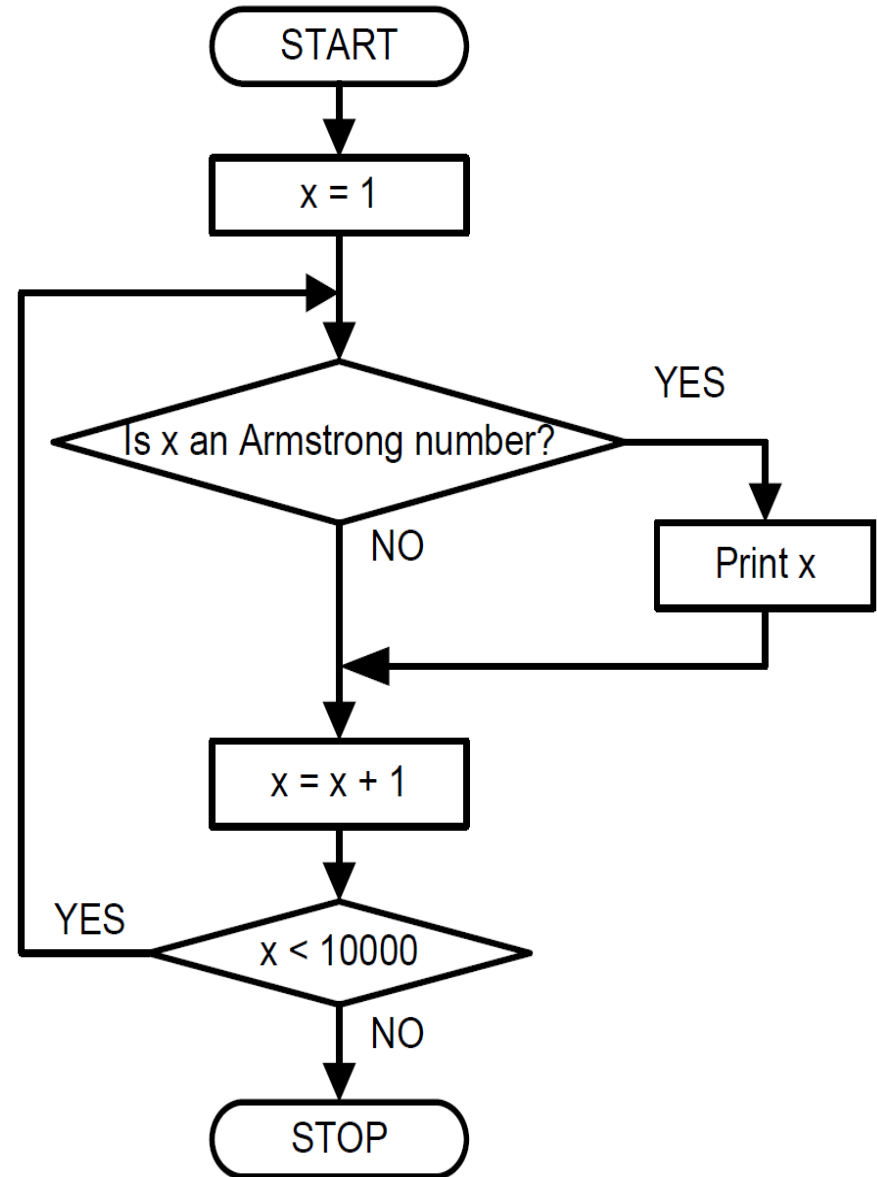
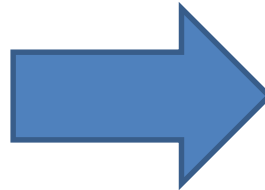
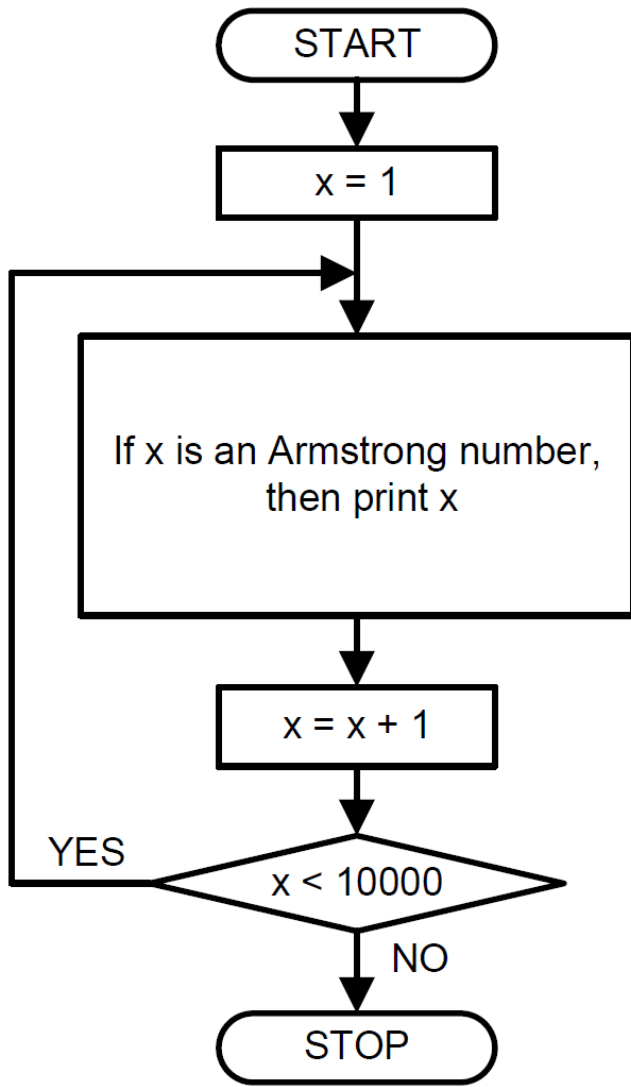


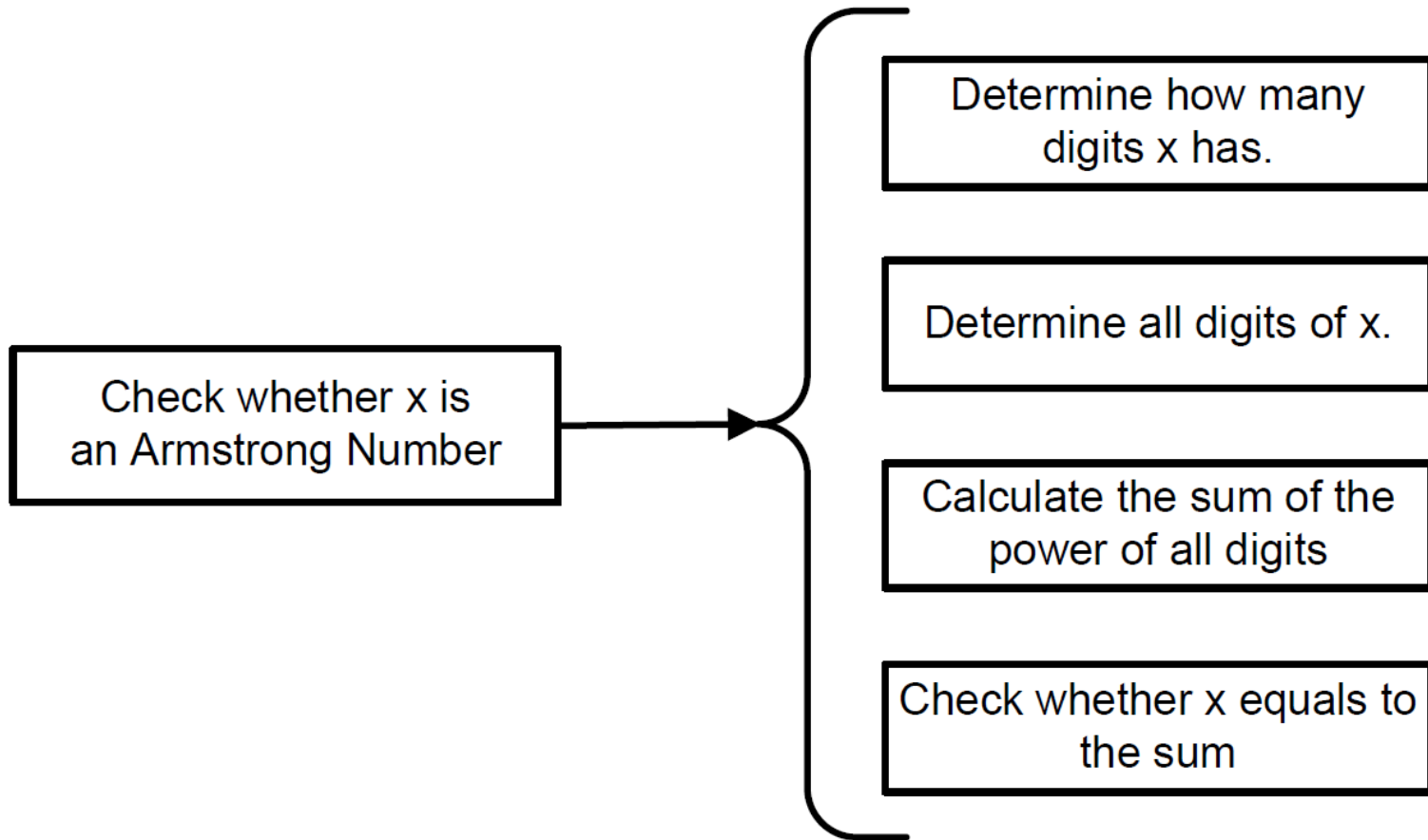
Top-Down Design Example

- Find all Armstrong numbers less than 10,000
- Given a positive integer that has n digits, it is an Armstrong number if the sum of the n th powers of its digits equals the number itself.
- For example, 371 is an Armstrong number since we have $371 = 3^3 + 7^3 + 1^3$.

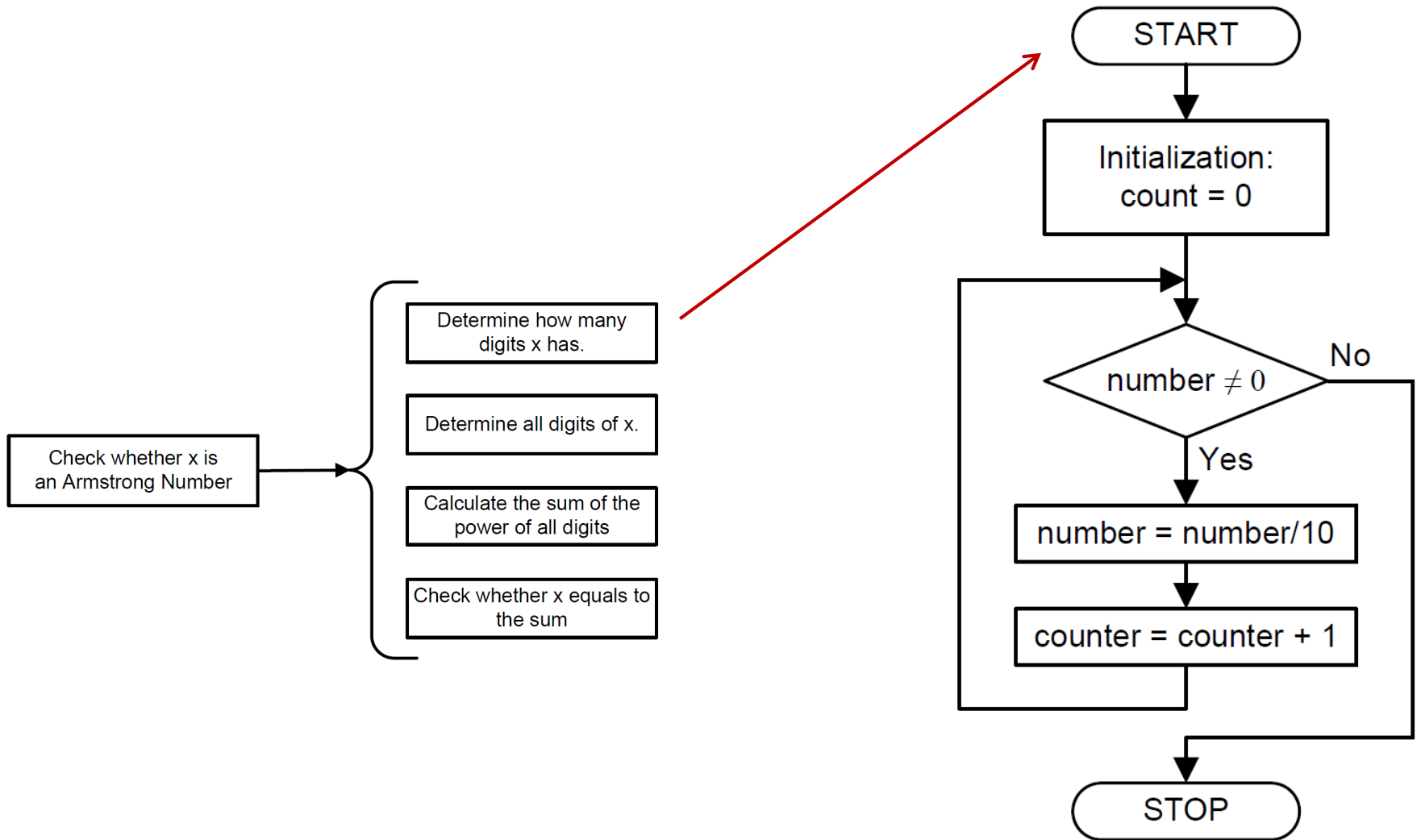


Restart!





One of the above will be tested in your quiz



Reuse of Registers

```
int A = 0; //  
int B = -1; //  
int C = -2; //  
int D = 2; //
```

?

```
void main(void){  
    A = B + C - D;  
    return;  
}
```

```
AREA myData, DATA  
A DCD 0  
B DCD -1  
C DCD -2  
D DCD 2
```

Data memory

Address	Data	
0x2000,000F	0x00	{ D = 0x0000,0002 = 2
0x2000,000E	0x00	
0x2000,000D	0x00	
0x2000,000C	0x02	{ C = 0xFFFF,FFFE = -2
0x2000,000B	0xFF	
0x2000,000A	0xFF	
0x2000,0009	0xFF	{ B = 0xFFFF,FFFF = -1
0x2000,0008	0xFE	
0x2000,0007	0xFF	
0x2000,0006	0xFF	{ A = 0x0000,0000 = 0
0x2000,0005	0xFF	
0x2000,0004	0xFF	
0x2000,0003	0x00	{
0x2000,0002	0x00	
0x2000,0001	0x00	
0x2000,0000	0x00	}

http://www.keil.com/support/man/docs/armasm/armasm_dom1361290005934.htm

```

int A = 0; // 0x00000000
int B = -1; // 0xFFFFFFFF
int C = -2; // 0xFFFFFFF0
int D = 2; // 0x00000002

void main(void){
    A = B + C - D;
    return;
}

```

8 Registers are used:
r0, r1, r2, r3, r4, r5, r6, r7

```

AREA myCode, CODE
EXPORT __main
ENTRY
__main PROC

    LDR r2, =B      ; r2 = 0x2000,0004
    LDR r3, [r2]    ; r3 = B = -1
    LDR r4, =C      ; r4 = 0x2000,0008
    LDR r5, [r4]    ; r5 = C = -2
    LDR r6, =D      ; r6 = 0x2000,000B
    LDR r7, [r6]    ; r7 = D = 2
    ADD r1, r3, r5   ; r1 = B + C
    SUB r1, r1, r7   ; r1 = B + C - D
    LDR r0, =A      ; r0 = 0x2000,0000
    STR r1, [r0]    ; Save A
    ENDP

AREA myData, DATA
A DCD 0
B DCD -1
C DCD -2
D DCD 2

END

```

```

AREA myCode, CODE
EXPORT __main
ENTRY
__main PROC

```

Lifetime of r3

```

    LDR r2, =B
    LDR r3, [r2]
    LDR r4, =C
    LDR r5, [r4]
    LDR r6, =D
    LDR r7, [r6]
    ADD r1, r3, r5
    SUB r1, r1, r7
    LDR r0, =A
    STR r1, [r0]
ENDP

```

```

AREA myData, DATA
A DCD 0
B DCD -1
C DCD -2
D DCD 2

END

```

8 registers used

```

AREA myCode, CODE
EXPORT __main
ENTRY
__main PROC

```

Lifetime of r2

```

    LDR r2, =B
    LDR r3, [r2]
    LDR r2, =C
    LDR r5, [r2]
    LDR r2, =D
    LDR r7, [r2]
    ADD r3, r3, r5
    SUB r3, r3, r7
    LDR r2, =A
    STR r3, [r2]
ENDP

```

Lifetime of r2

```

AREA myData, DATA
A DCD 0
B DCD -1
C DCD -2
D DCD 2

END

```

4 registers used

```

AREA myCode, CODE
EXPORT __main
ENTRY
__main PROC

```

```

    LDR r2, =B
    LDR r3, [r2]
    LDR r2, =C
    LDR r5, [r2]
    LDR r2, =D
    LDR
    ADD
    SUB
    LDR r2, =A
    STR r3, [r2]
ENDP

```

```

AREA myData, DATA
A DCD 0
B DCD -1
C DCD -2
D DCD 2

END

```

3 registers used

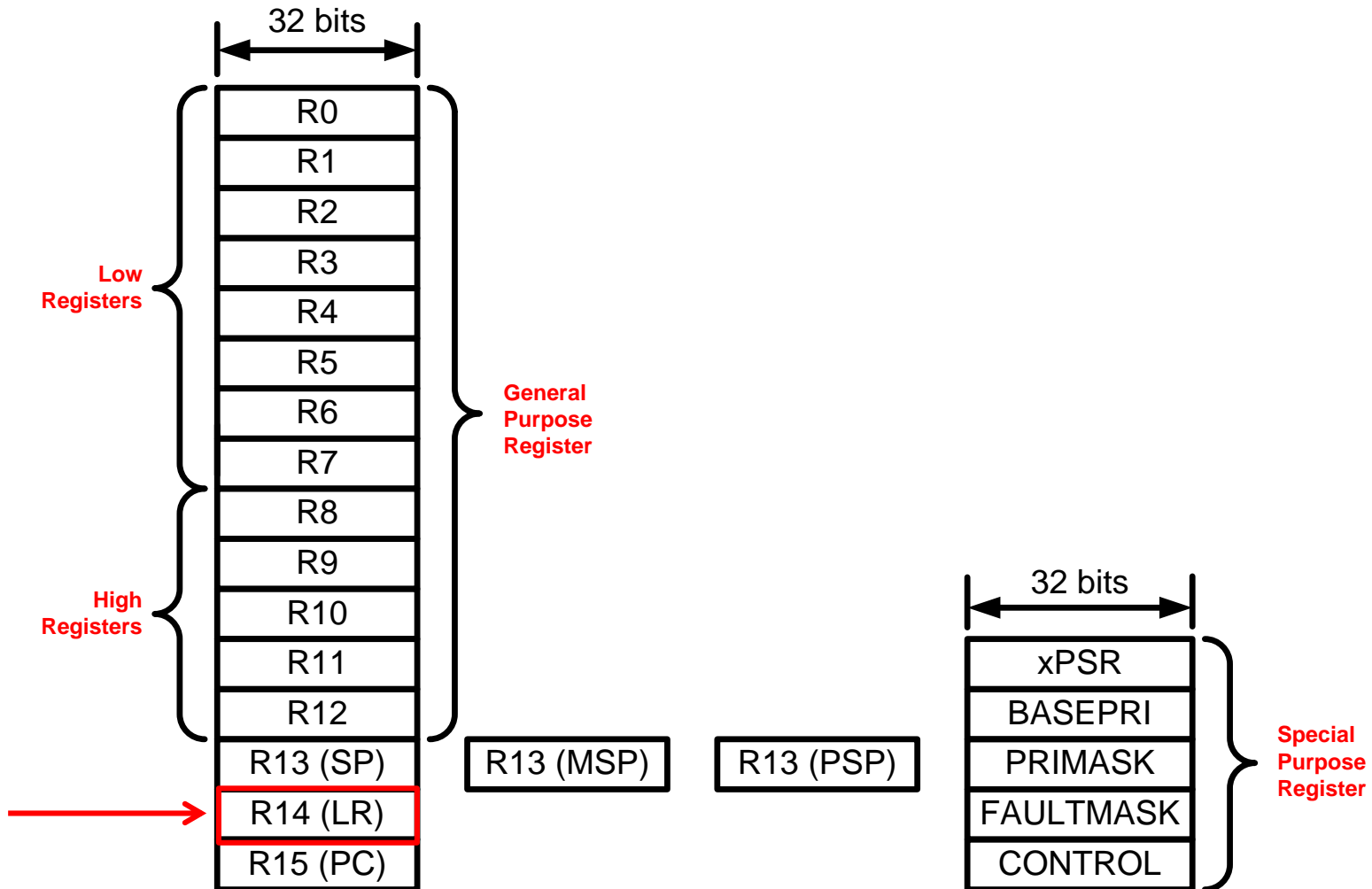
$$A = B + C - D;$$

Passing Parameters to Subroutines via
Register Preserved Environment via Stack

Subroutine

- A subroutines, also called a function or a procedure,
 - single-entry, single-exit
 - Return to caller after it exits
- When a subroutine is called, the **Link Register** (LR) holds the memory address of the next instruction to be executed after the subroutine exits.

Link Register



Calling a Subroutine

BL *label*

- Step 1: $LR = PC + 4$
- Step 2: $PC = \text{label}$
- Notes:
 - *label* is name of subroutine (address)
 - Compiler translates label to memory address
 - After call, LR holds return address (the instruction following the call)

Caller Program

```
MOV r4, #100
...
BL foo
...
```

Subroutine/Callee

```
foo PROC
...
MOV    r4, #10
...
BX     LR
ENDP
```

Exiting a Subroutine

BX LR

- PC = LR

Caller Program

```
MOV r4, #100  
...  
BL  foo  
...
```

Subroutine/Callee

```
foo PROC  
...  
MOV    r4, #10  
...  
BX    LR  
ENDP
```

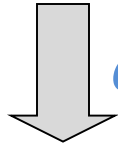
BL and BX

```
void enable(void) ;
```

• • •

```
enable() ;
```

• • •



Compiler

• • •

BL enable

• • •

export enable

enable • • •

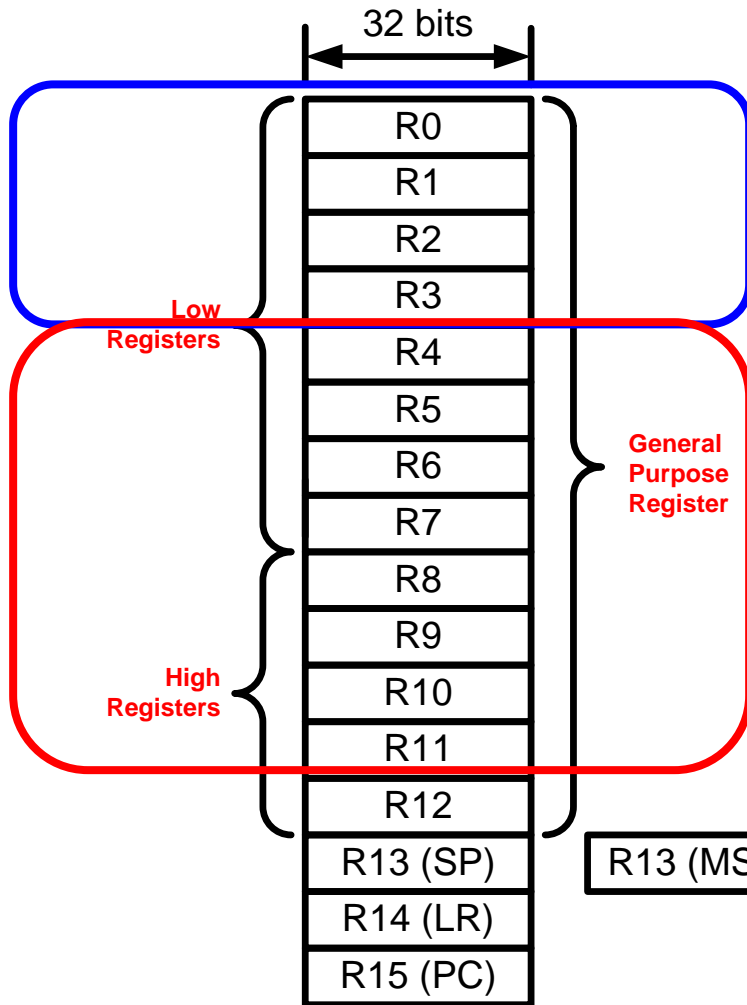
• • •

BX LR



ARM Procedure Call Standard

Register	Usage	Subroutine Preserved	Notes
r0	Argument 1 and return value	No	If return has 64 bits, then r0:r1 hold it. If argument 1 has 64 bits, r0:r1 hold it.
r1	Argument 2	No	
r2	Argument 3	No	If the return has 128 bits, r0-r3 hold it.
r3	Argument 4	No	If more than 4 arguments, use the stack
r4	General-purpose V1	Yes	Variable register 1 holds a local variable.
r5	General-purpose V2	Yes	Variable register 2 holds a local variable.
r6	General-purpose V3	Yes	Variable register 3 holds a local variable.
r7	General-purpose V4	Yes	Variable register 4 holds a local variable.
r8	General-purpose V5	YES	Variable register 5 holds a local variable.
r9	Platform specific/V6	No	Usage is platform-dependent.
r10	General-purpose V7	Yes	Variable register 7 holds a local variable.
r11	General-purpose V8	Yes	Variable register 8 holds a local variable.
r12 (IP)	Intra-procedure-call register	No	It holds intermediate values between a procedure and the sub-procedure it calls.
r13 (SP)	Stack pointer	Yes	SP has to be the same after a subroutine has completed.
r14 (LR)	Link register	No	LR does not have to contain the same value after a subroutine has completed.
r15 (PC)	Program counter	N/A	Do not directly change PC

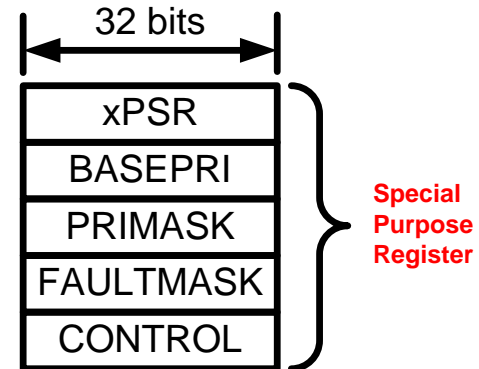


Not saved. Hold arguments, results, or temporary values. Caller doesn't expect them to be retained.

Callee must save them. Caller expects these values are retained.

R13 (MSP)

R13 (PSP)



Example: $R2 = R0 * R0 + R1 * R1$

```
MOV R0 , #3
MOV R1 , #4
BL  SSQ
MOV R2 , R0
B  ENDL
...
SSQ  PROC
    MUL R2 , R0 , R0
    MUL R3 , R1 , R1
    ADD R2 , R2 , R3
    MOV R0 , R2
    BX  LR
    ENDP
...
```

R1: second argument

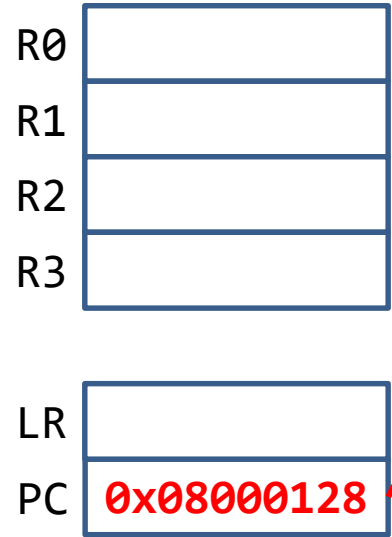
R0: first argument

```
int SSQ(int x, int y){
    int z;
    z = x*x + y * y;
    return z;
}
```

R0: Return Value

SSQ

MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL
...
PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
...
ENDL



SSQ

Memory Address	
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

SSQ

MOV R0,#3

MOV R1,#4

BL SSQ

MOV R2,R0

B ENDL

...

PROC

MUL R2,R0,R0

MUL R3,R1,R1

ADD R2,R2,R3

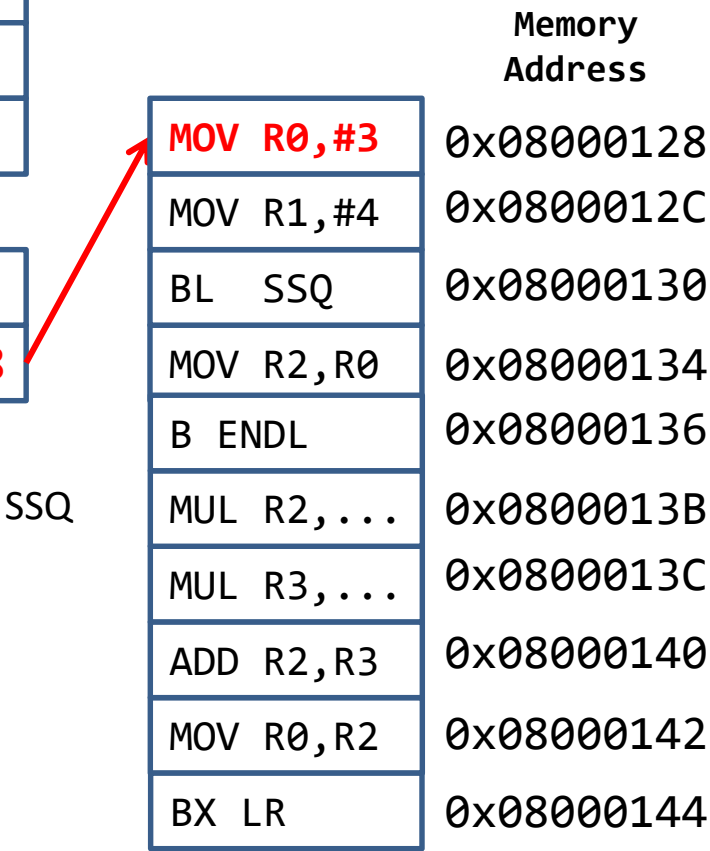
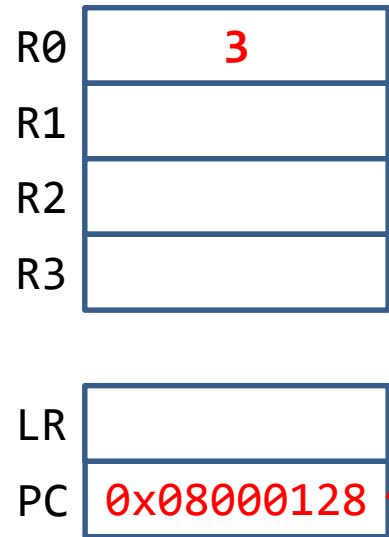
MOV R0,R2

BX LR

ENDP

...

ENDL



```
MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL

...
SSQ PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
ENDL ...
```

R0	3
R1	4
R2	
R3	
LR	
PC	0x0800012C

SSQ

	Memory Address
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

SSQ

```
MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL
...
PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
...
ENDL
```

R0	3
R1	4
R2	
R3	
LR	
PC	0x08000130

SSQ

	Memory Address
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

SSQ

MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL
...
PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
...
ENDL

R0	3
R1	4
R2	
R3	
LR	0x08000134
PC	0x0800013B

Memory Address	
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

Address of the next instruction after the branch is saved into LR.

SSQ

MOV R0,#3

MOV R1,#4

BL SSQ

MOV R2,R0

B ENDL

...

PROC

MUL R2,R0,R0

MUL R3,R1,R1

ADD R2,R2,R3

MOV R0,R2

BX LR

ENDP

...

ENDL

R0	3
R1	4
R2	9
R3	
LR	0x08000134
PC	0x0800013B

Memory Address	
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

SSQ

```
MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL
...
PROC
  MUL R2,R0,R0
  MUL R3,R1,R1
  ADD R2,R2,R3
  MOV R0,R2
  BX LR
ENDP
...
ENDL
```

R0	3
R1	4
R2	9
R3	16

LR	0x08000134
PC	0x0800013C

	Memory Address
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

```
MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL

...
SSQ PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
...
ENDL
```

R0	3
R1	4
R2	25
R3	16
LR	0x08000134
PC	0x08000140

	Memory Address
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

SSQ

MOV R0,#3

MOV R1,#4

BL SSQ

MOV R2,R0

B ENDL

...

PROC

MUL R2,R0,R0

MUL R3,R1,R1

ADD R2,R2,R3

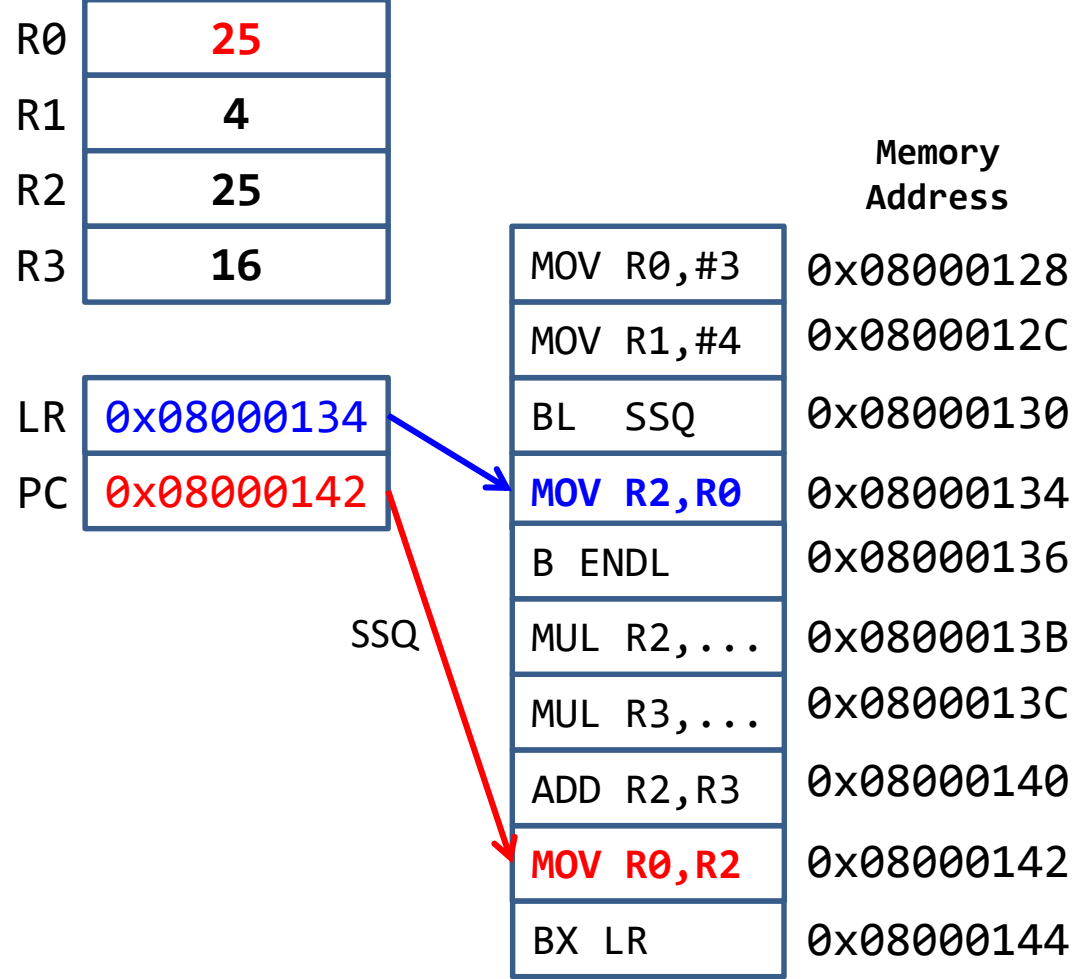
MOV R0,R2

BX LR

ENDP

...

ENDL



SSQ

MOV R0,#3

MOV R1,#4

BL SSQ

MOV R2,R0

B ENDL

...

PROC

MUL R2,R0,R0

MUL R3,R1,R1

ADD R2,R2,R3

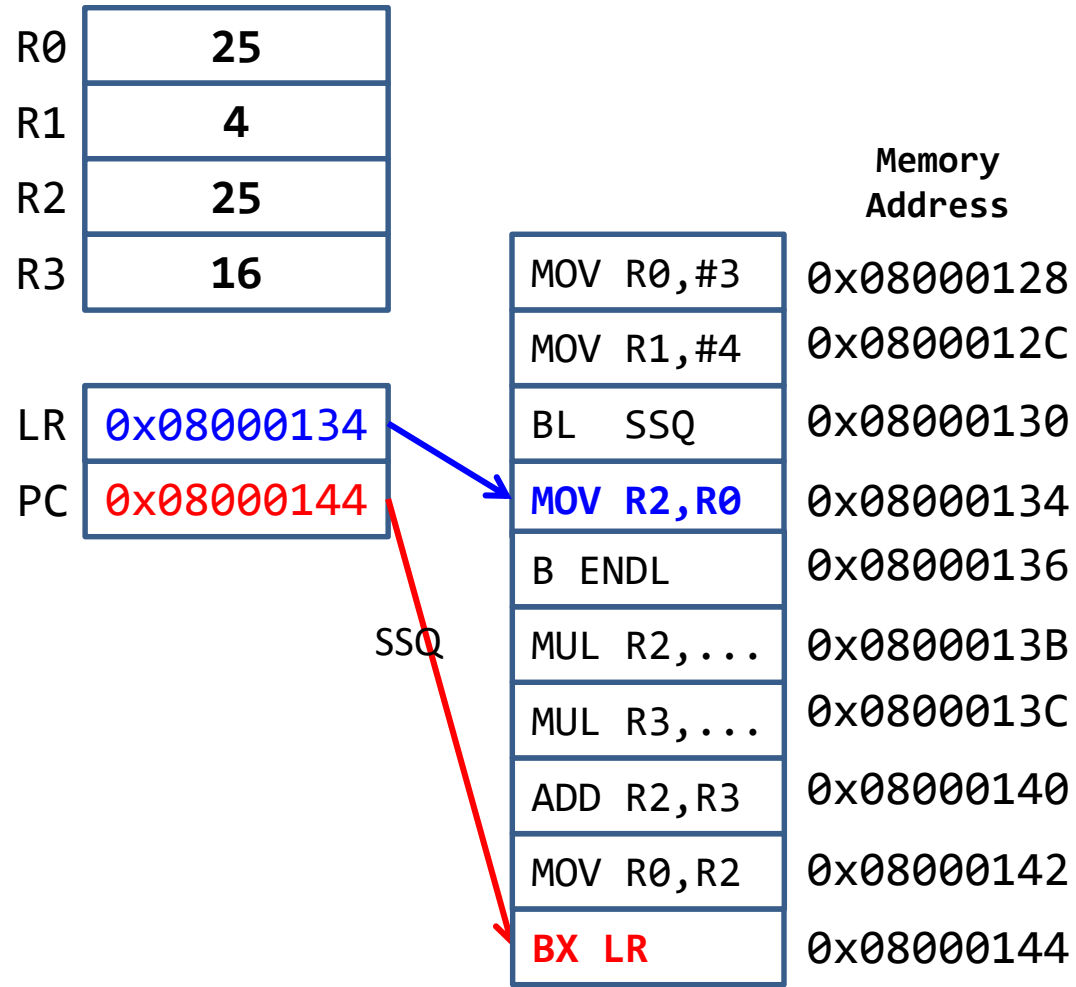
MOV R0,R2

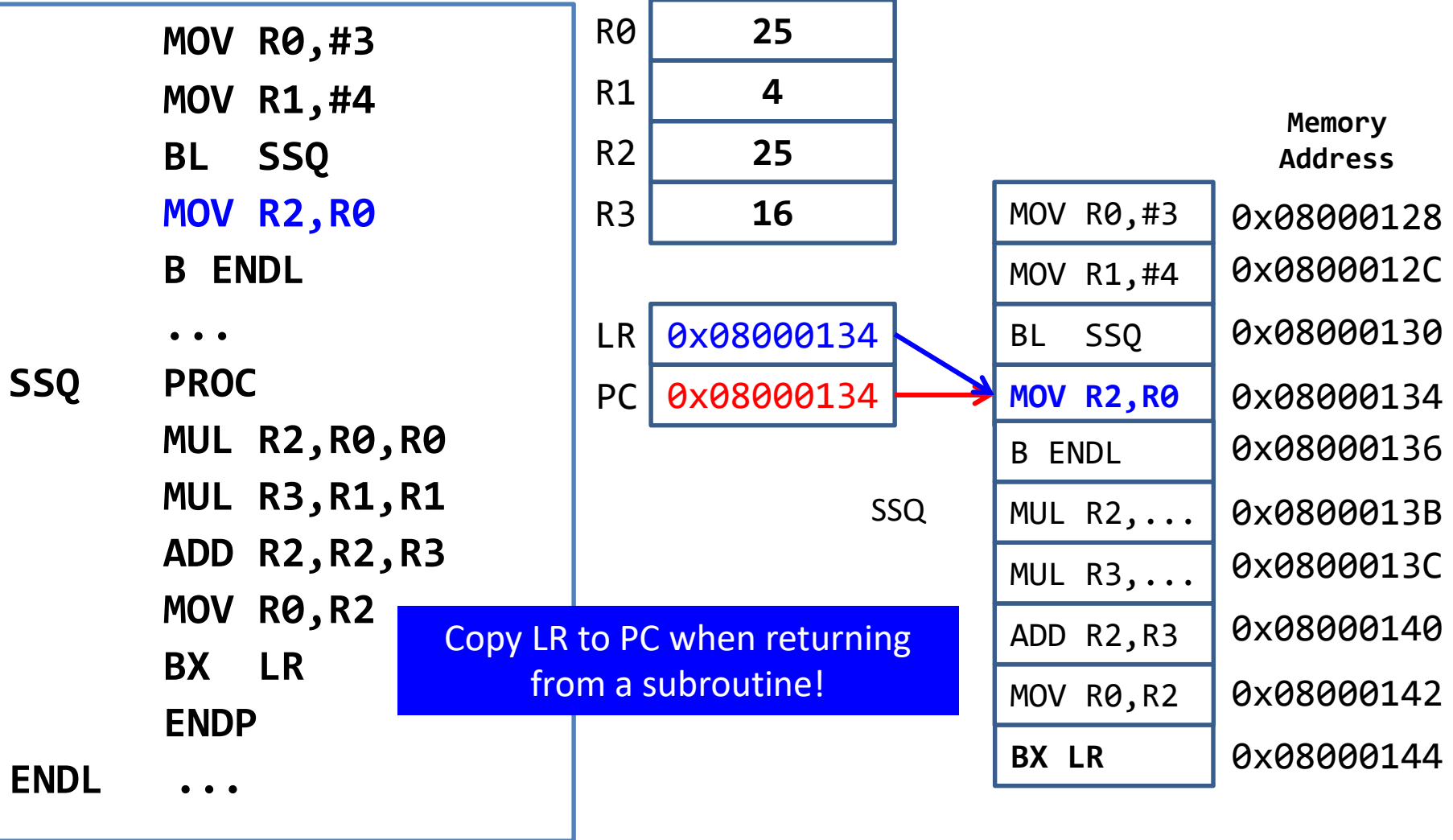
BX LR

ENDP

...

ENDL





SSQ

MOV R0,#3

MOV R1,#4

BL SSQ

MOV R2,R0

B ENDL

...

PROC

MUL R2,R0,R0

MUL R3,R1,R1

ADD R2,R2,R3

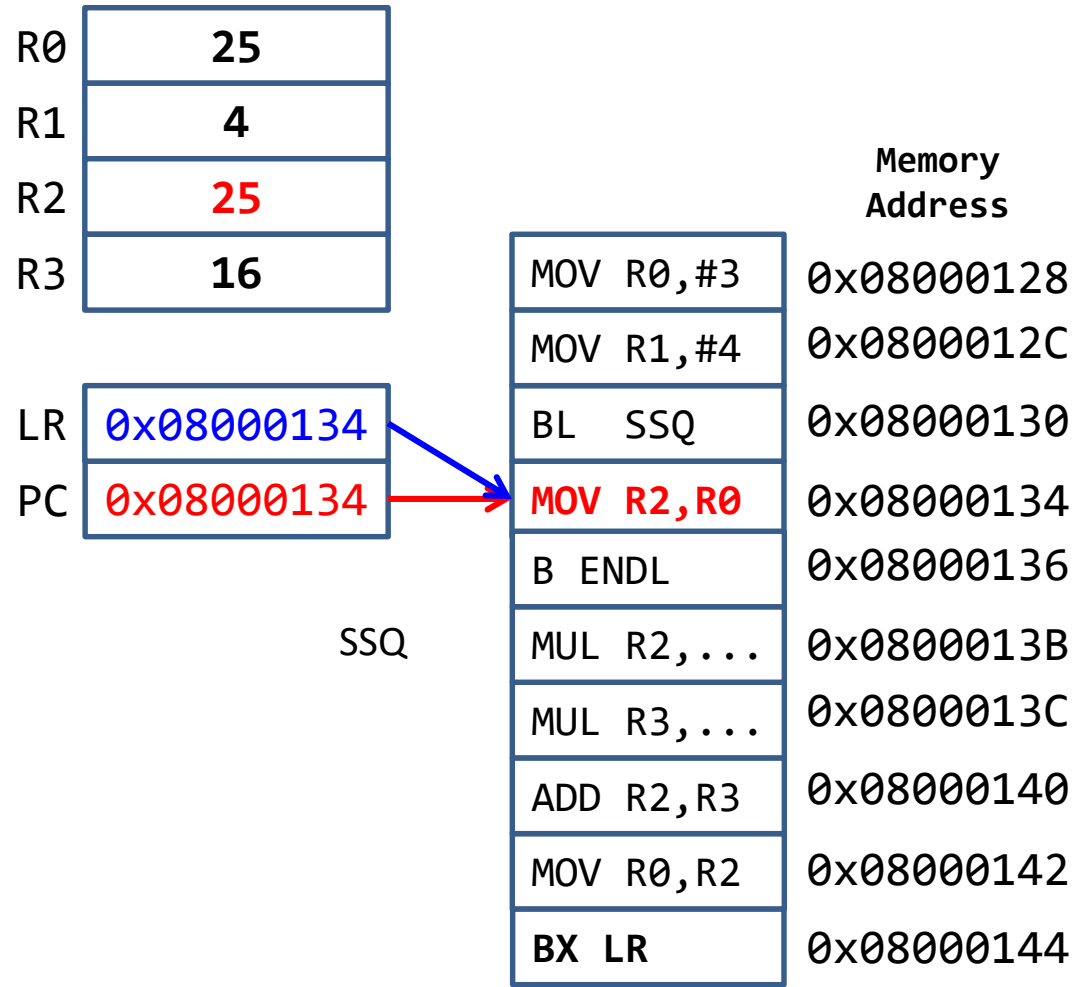
MOV R0,R2

BX LR

ENDP

...

ENDL



SSQ

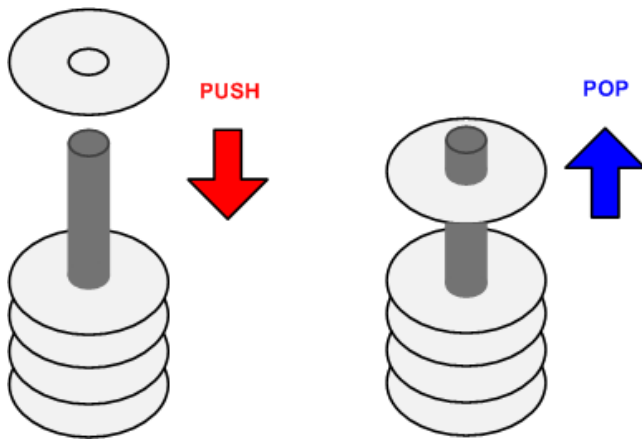
MOV R0,#3
MOV R1,#4
BL SSQ
MOV R2,R0
B ENDL
...
PROC
MUL R2,R0,R0
MUL R3,R1,R1
ADD R2,R2,R3
MOV R0,R2
BX LR
ENDP
ENDL ...

R0	25
R1	4
R2	25
R3	16
LR	0x08000134
PC	0x08000136

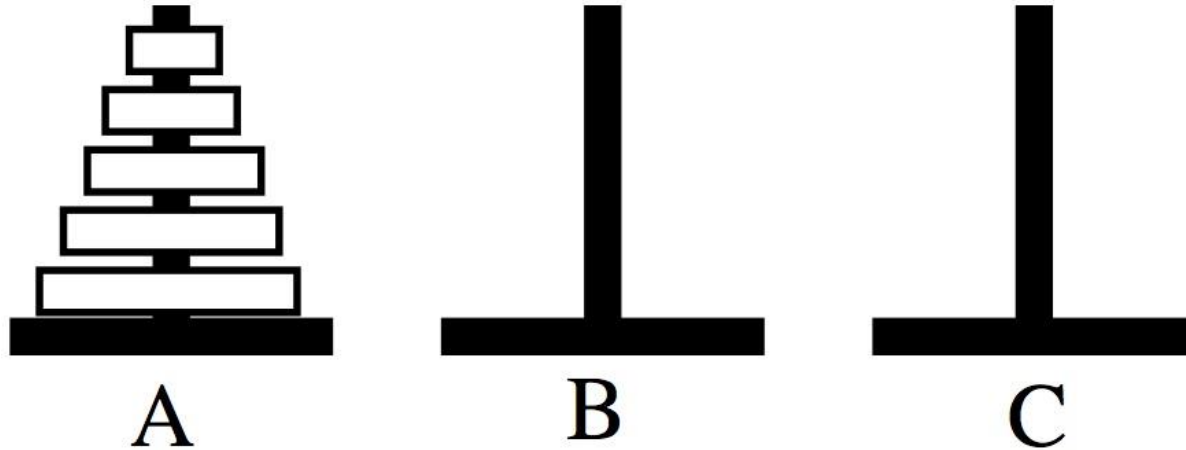
Memory Address	
MOV R0,#3	0x08000128
MOV R1,#4	0x0800012C
BL SSQ	0x08000130
MOV R2,R0	0x08000134
B ENDL	0x08000136
MUL R2,...	0x0800013B
MUL R3,...	0x0800013C
ADD R2,R3	0x08000140
MOV R0,R2	0x08000142
BX LR	0x08000144

Stack

- A **Last-In-First-Out** data structure
- Only allow to access the most recently added item
 - Also called the top of the stack
- Key operations:
 - push (add item to stack)
 - pop (remove top item from stack)



Stack ex: Tower of Hanoi



- Only one disk may be moved at a time.
- Each move consists of taking the upper disk from one of the rods and sliding it onto another rod, on top of the other disks that may already be present on that rod.
- **No disk may be placed on top of a smaller disk.**

STACK: Last In First Out



Stack 1

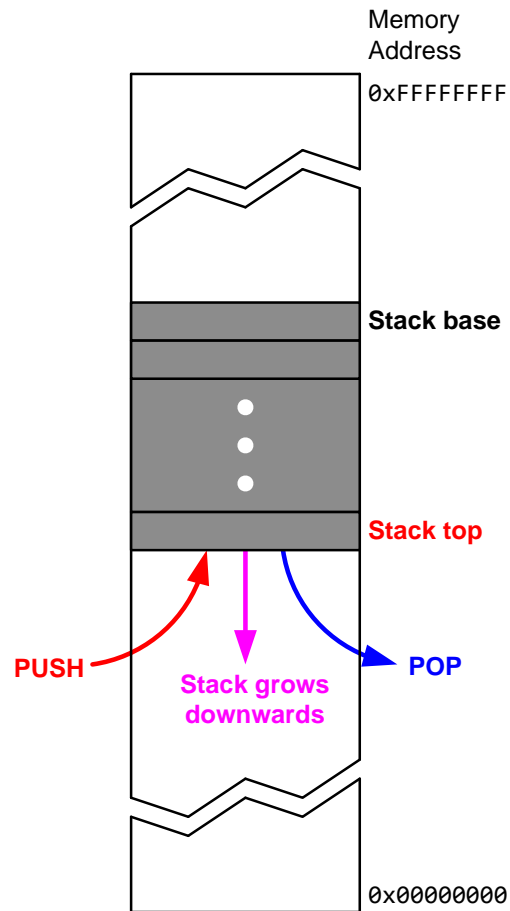
Stack 2

Stack 3

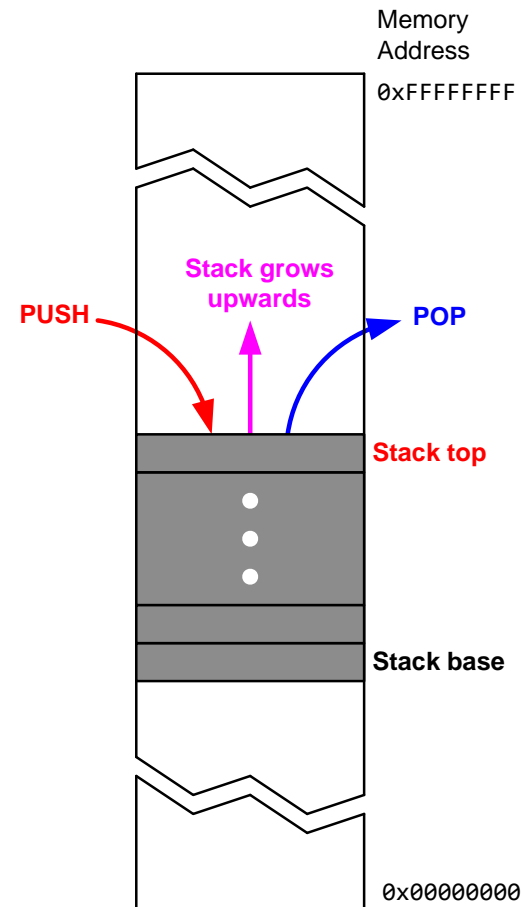
http://en.wikipedia.org/wiki/File:Tower_of_Hanoi_4.gif

Stack Growth Convention:

Ascending vs Descending



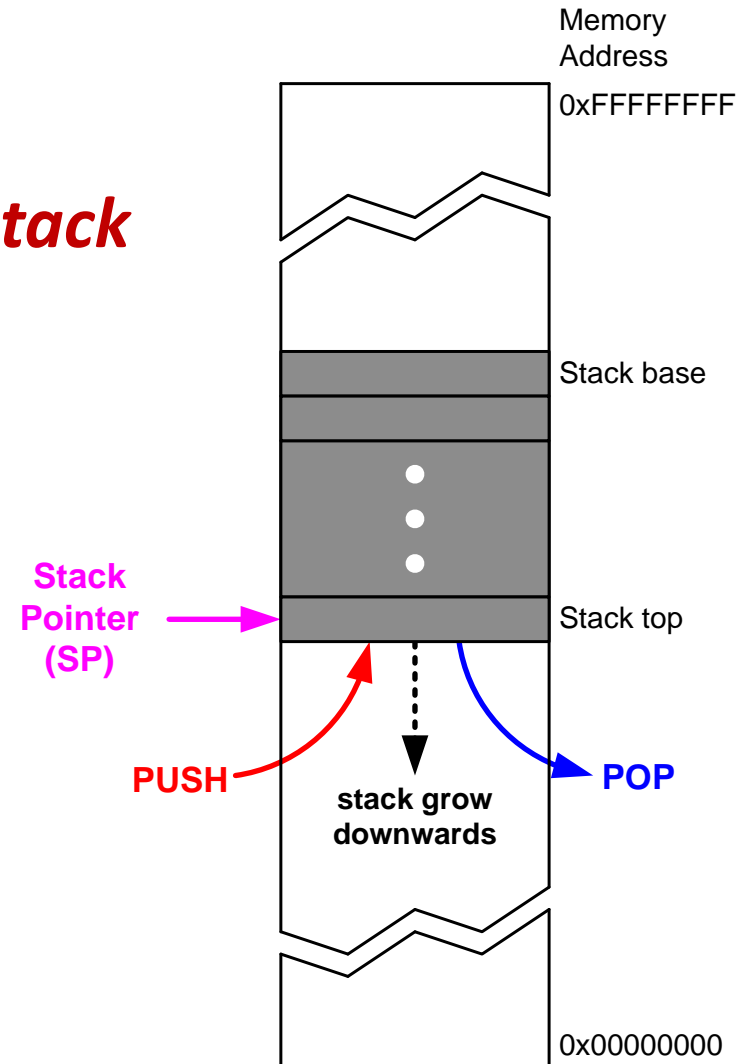
Descending stack: Stack grows towards low memory address



Ascending stack: Stack grows towards high memory address

Cortex-M Stack

- stack pointer (SP) = R13
- Cortex-M uses *full descending stack*
- stack pointer
 - decremented on **PUSH**
 - incremented on **POP**





PUSH {Rd}

– $SP = SP - 4 \rightarrow$ descending stack

– $(*SP) = Rd \rightarrow$ full stack

(SP points to the last item pushed onto the stack)

Push multiple registers *They are equivalent.
(push large first)*

PUSH {r6, r7, r8}  **PUSH {r8, r7, r6}**  **PUSH {r8}**
PUSH {r7}
PUSH {r6}

PUSH {r6-r8}

- The order in which registers listed in the register list does not matter.
- When pushing multiple registers, these registers are automatically **sorted by name** and **the lowest-numbered register** is stored to the lowest memory address, *i.e.* **is stored last**.

POP {Rd}

- $Rd = (*SP) \rightarrow$ full stack
- $SP = SP + 4 \rightarrow$ Stack shrinks

Pop multiple registers

They are equivalent.

Pop small first

POP {r6, r7, r8}



POP {r8, r7, r6}



POP {r6}
POP {r7}
POP {r8}

- The order in which registers listed in the register list does not matter.
- When popping multiple registers, these registers are automatically **sorted by name** and **the lowest-numbered register** is loaded from the lowest memory address, *i.e.* **is loaded first**.

C:\Users\Noori\Desktop\CS100\CS100C_Fall2018\2_LectureLabNote\Week11\PushPop\test.uvprojx - µVision

File Edit View Project Flash Debug Peripherals Tools SVCS Window Help

main

Registers

Register	Value
Core	
R0	0x00000012
R1	0x00000034
R2	0x00000056
R3	0x00000000
R4	0x00000000
R5	0x00000000
R6	0x00000000
R7	0x00000000
R8	0x00000000
R9	0x00000000
R10	0x00000000
R11	0x00000000
R12	0x00000000
R13 (SP)	0x10000200
R14 (LR)	0xFFFFFFFF
R15 (PC)	0x00000114
xPSR	0x01000000
Banked	
System	
Internal	
Mode	Thread
Privilege	Privileged
Stack	MSP
States	14
...	0x00000117

Disassembly

0x00000114 E7FE B 0x00000114

0x00000116 0000 MOVS r0,r0

0x00000118 0000 MOVS r0,r0

0x0000011A 0000 MOVS r0,r0

test.s startup_LPC17xx.s system_LPC17xx.c

```
4 AREA Main, CODE, READONLY ; Area of code named "Main" that is read only
5 ALIGN 2 ; Align the data boundary to a multiple of 2
6 ENTRY ; Entry into the code segment
7
8 ; Label main which is branched to on reset
9 _main
10 MOV R0, #0x12
11 MOV R1, #0x34
12 MOV R2, #0x56
13
14 ; Learning Focus
15 PUSH{R0,R2,R1} ; 1. Disassembly instructions on push and pop
16 POP{R0,R1,R2} ; 2. R0-R2 don't change after pop
17 ; 3. Change of SP, it decrements
18 ; initial SP = 0x10000200
19 ; After push SP = 0x100001F4
20 ; After pop SP = 0x10000200
21 ; Where Stack is located (usually) beginning of RAM (0x10000000)
22 ; Size of stack is defined at Startup.s, we have defined 0x200
23 ; therefore initial SP = 0x10000200 makes sense
24
```

Memory 1

0x100001F4

0x100001F4: 00000012

0x100001F8: 00000034

0x100001FC: 00000056

0x10000200: 00000000

0x10000204: 00000000

0x10000208: 00000000

0x1000020C: 00000000

0x10000210: 00000000

0x10000214: 00000000

0x10000218: 00000000

0x1000021C: 00000000

0x10000220: 00000000

0x10000224: 00000000

0x10000228: 00000000

0x1000022C: 00000000

0x10000230: 00000000

0x10000234: 00000000

0x10000238: 00000000

0x1000023C: 00000000

0x10000240: 00000000

0x10000244: 00000000

0x10000248: 00000000

0x1000024C: 00000000

0x10000250: 00000000

0x10000254: 00000000

0x10000258: 00000000

Call Stack + L... Memory 1

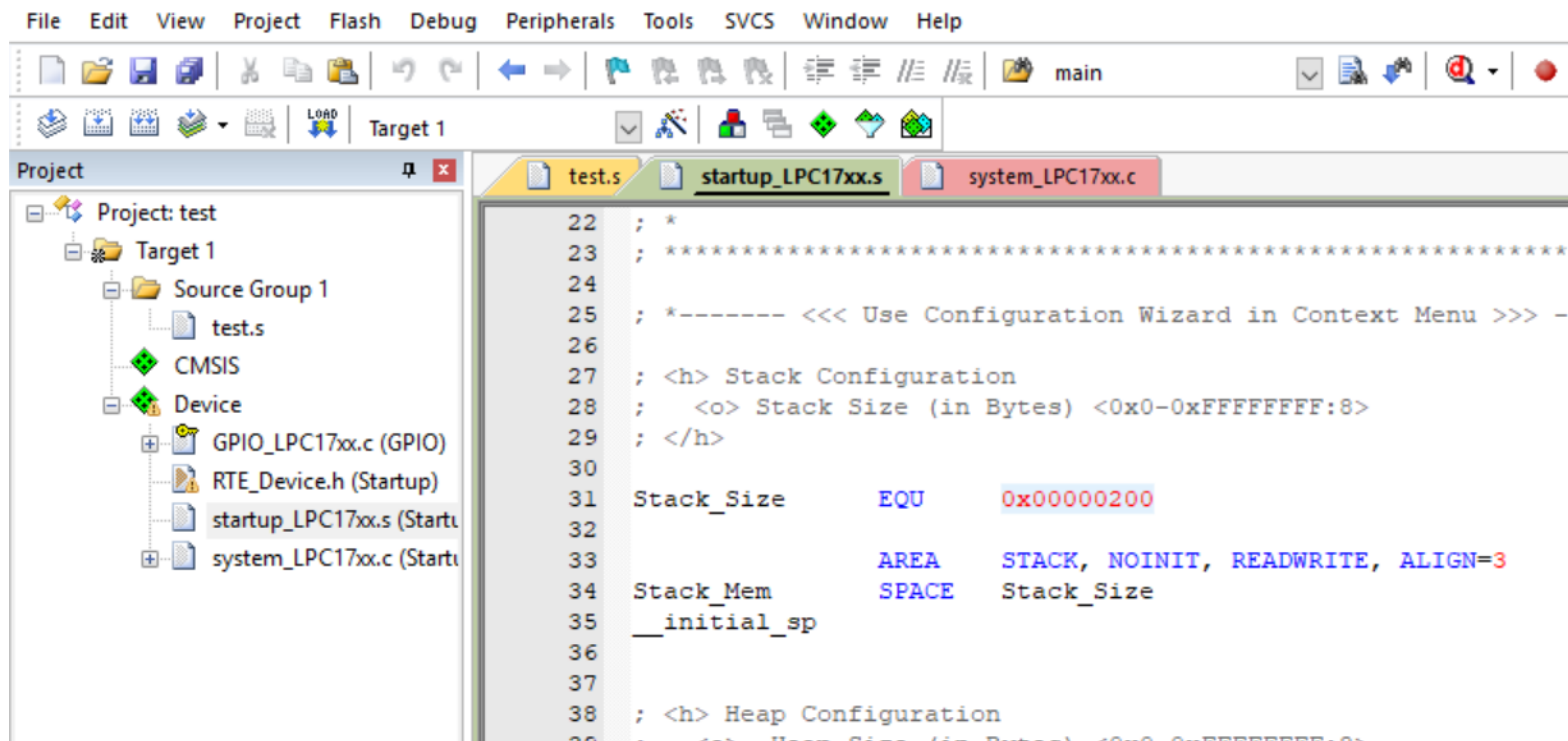
Command

BS 0x00000106, 1

Simulation t1: 0.00000117 sec L:22 C:71 CAP NUM SCRL OVR RA

Initializing the stack pointer (SP)

- Before using the stack, software has to define stack space and initialize the stack pointer (SP).
- The assembly file **startup.s** defines stack space and initialize SP.

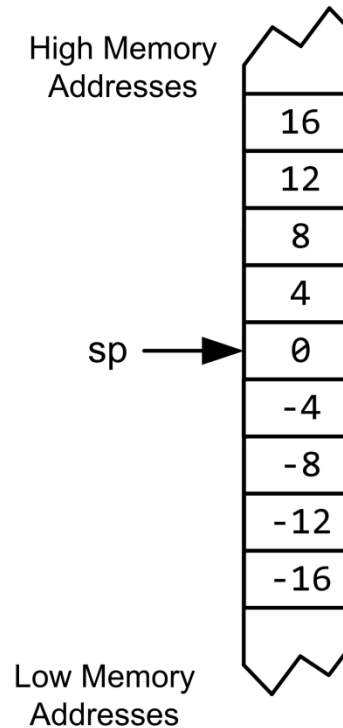
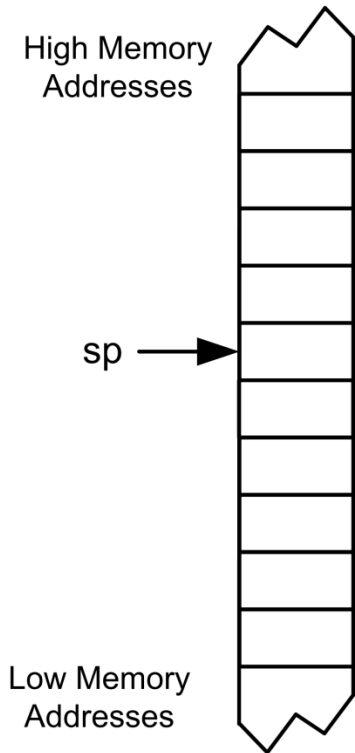


R11	0x00000000	11
R12	0x00000000	12
R13 (SP)	0x10000200	13
R14 (LR)	0xFFFFFFFF	14
R15 (PC)	0x00000114	15
+ xPSR	0x01000000	16
+ Banked		

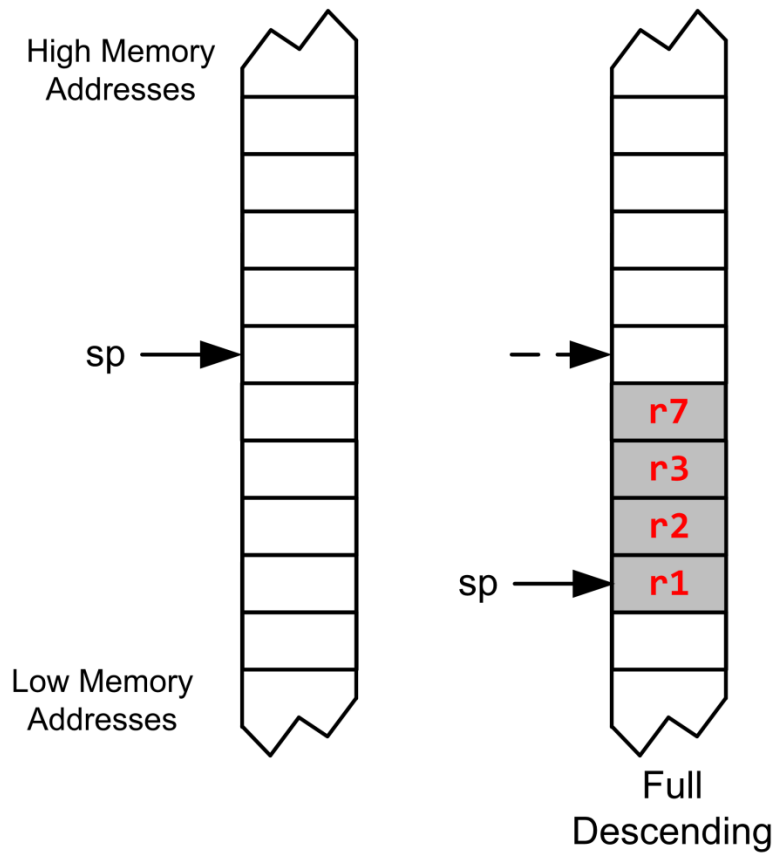
Full Descending Stack

PUSH {r3, r1, r7, r2}

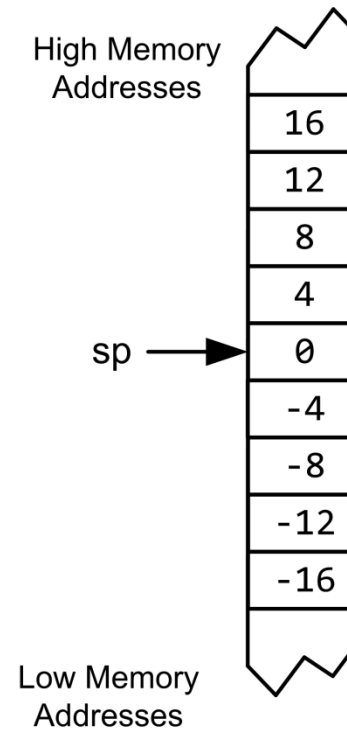
POP {r3, r1, r7, r2}



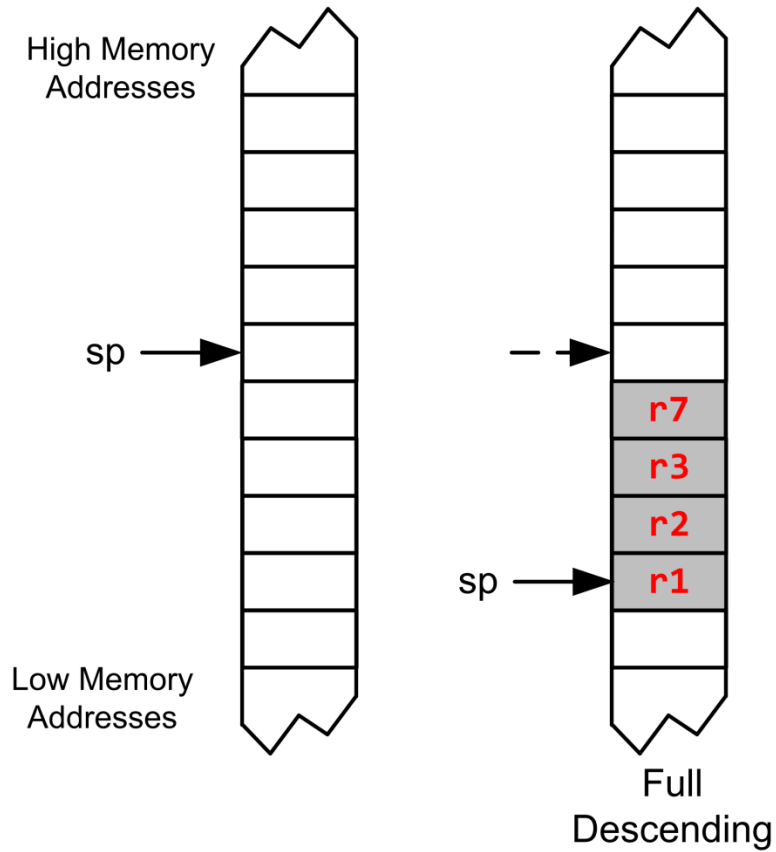
PUSH {r3, r1, r7, r2}



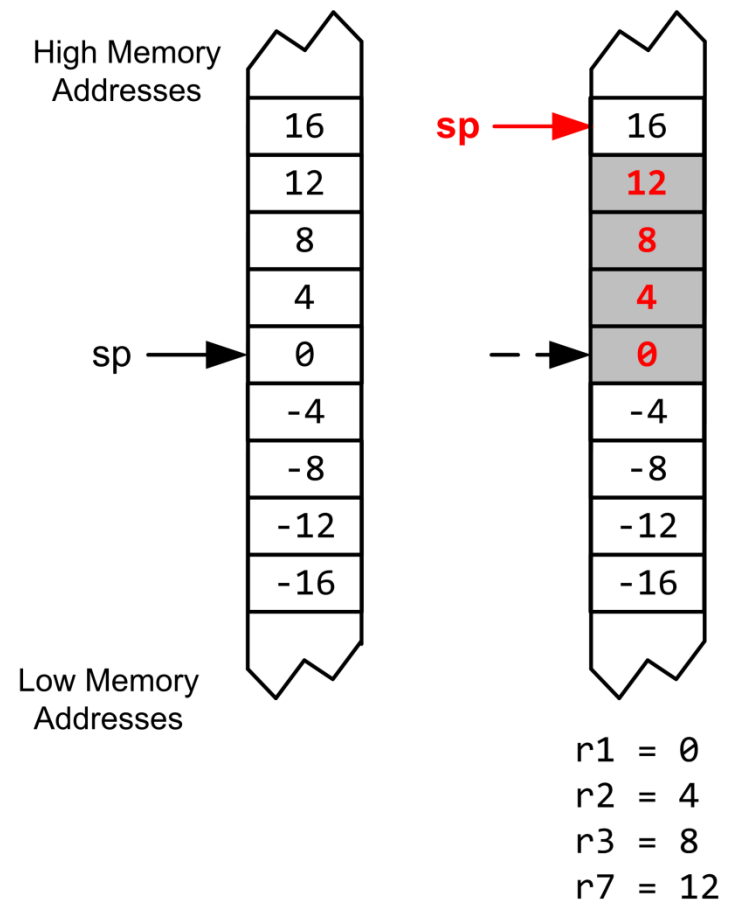
POP {r3, r1, r7, r2}



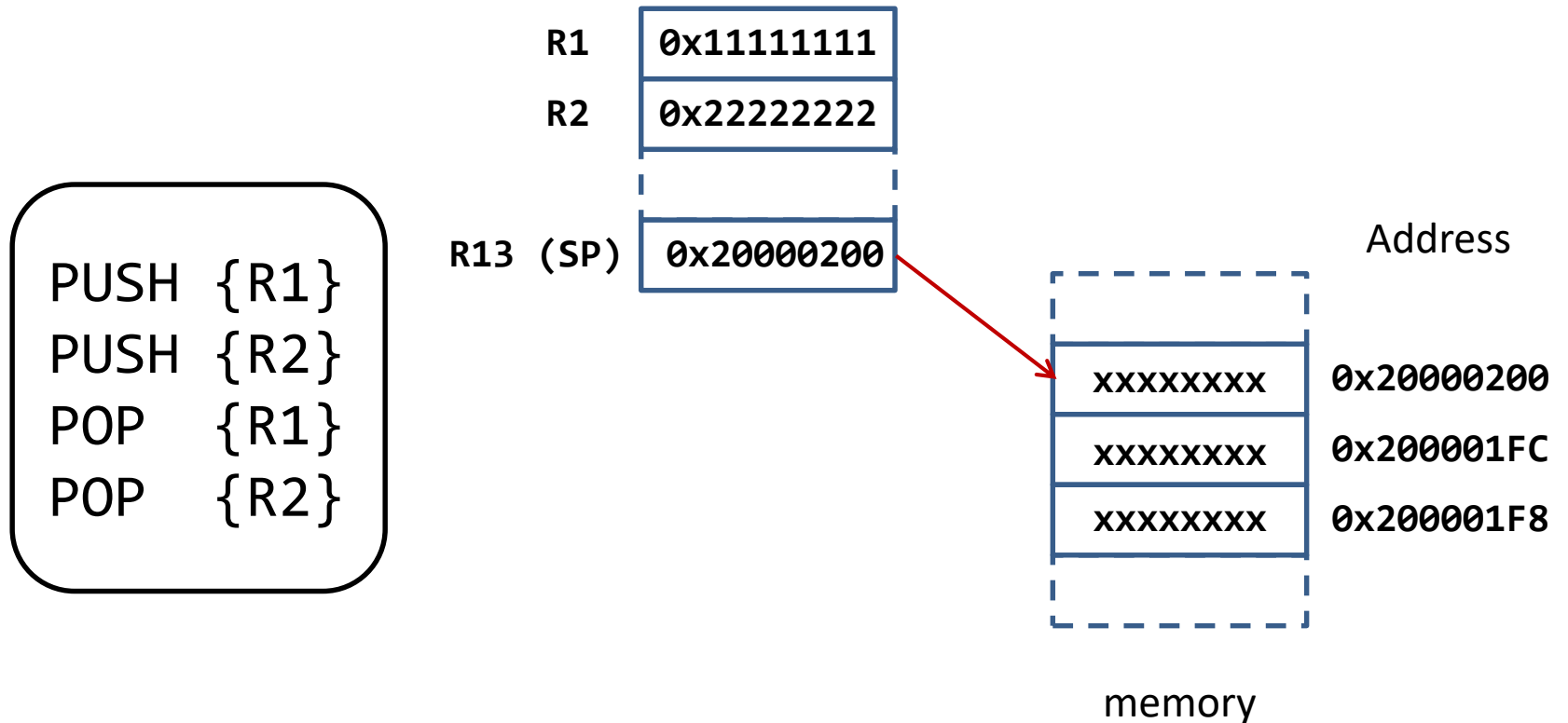
PUSH {r3, r1, r7, r2}

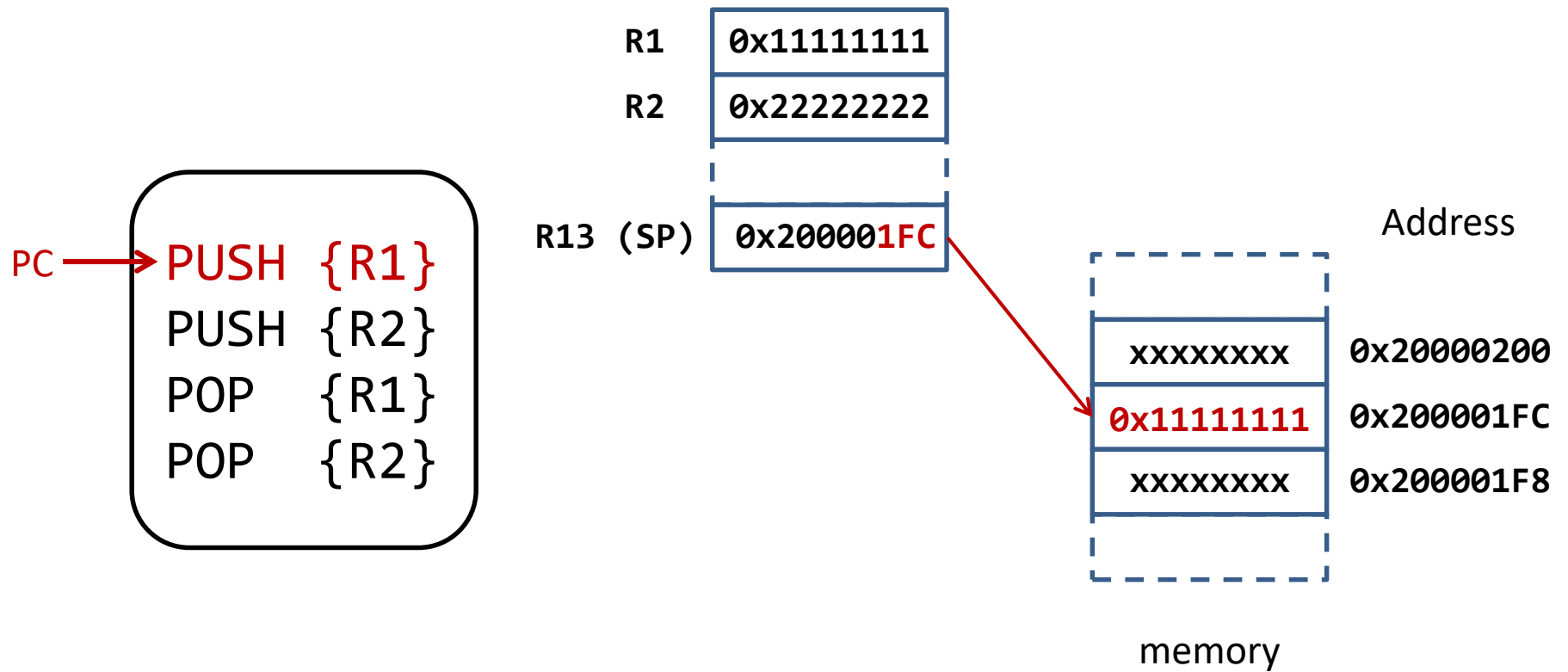


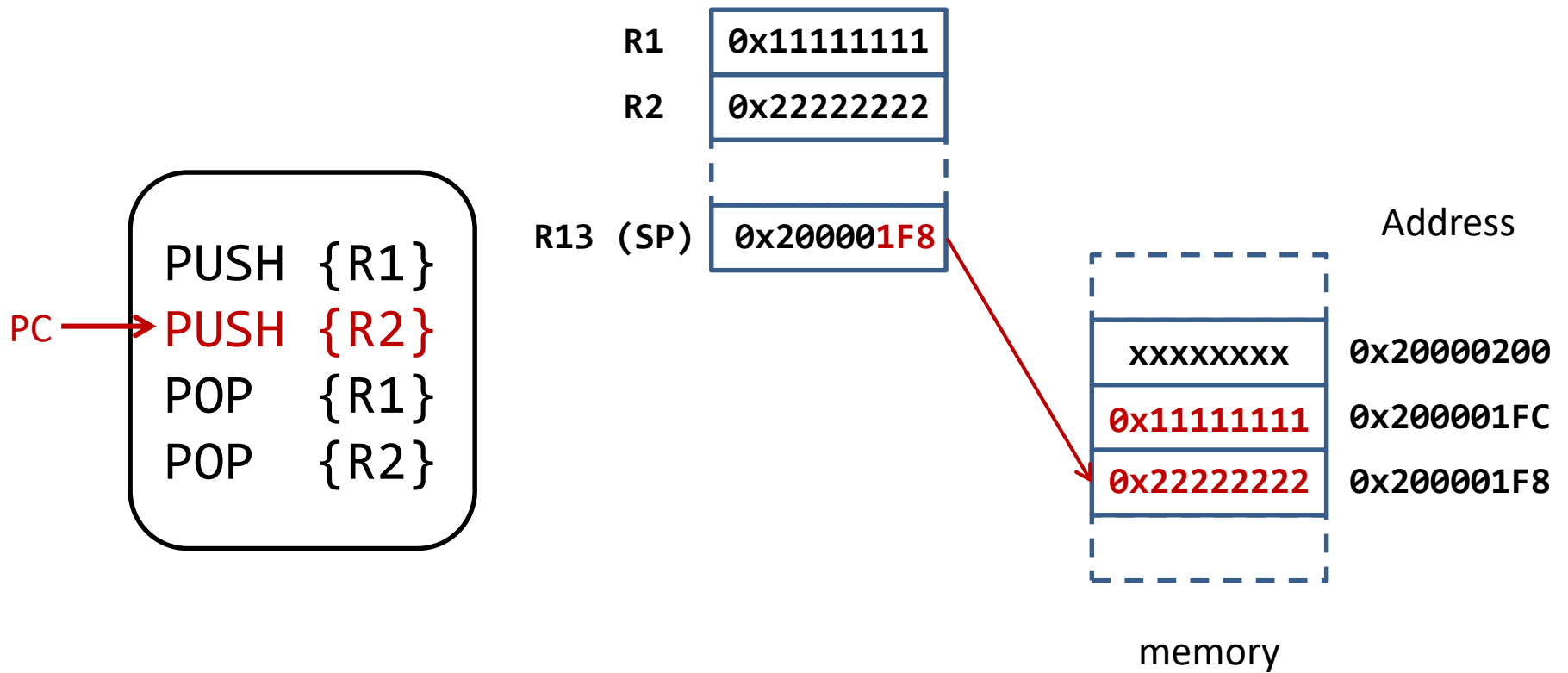
POP {r3, r1, r7, r2}

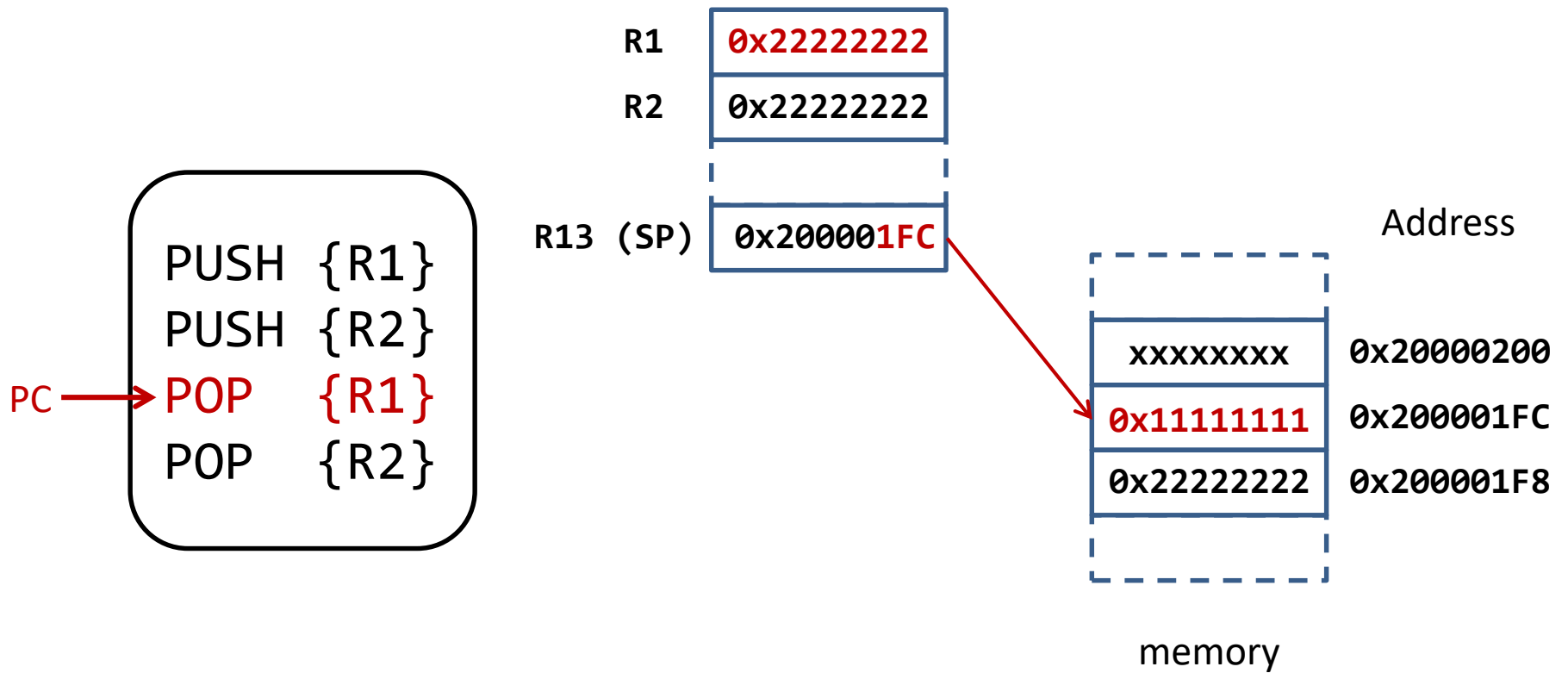


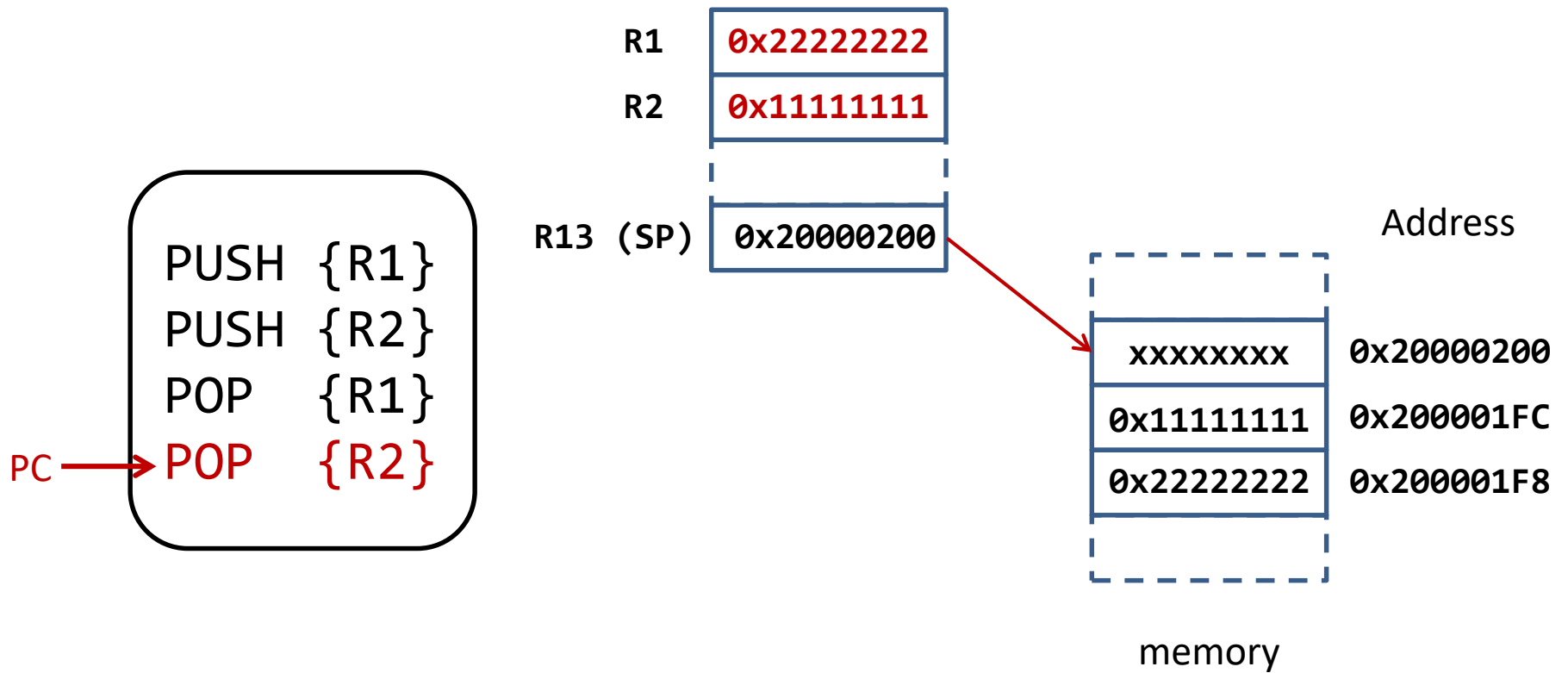
Example: swap R1 & R2











Are the values of R1 and R2 swapped?

PUSH {R1, R2}

POP {R2, R1}

Answer: No.

But you can:



Recap: Call a Subroutine I (has a problem)

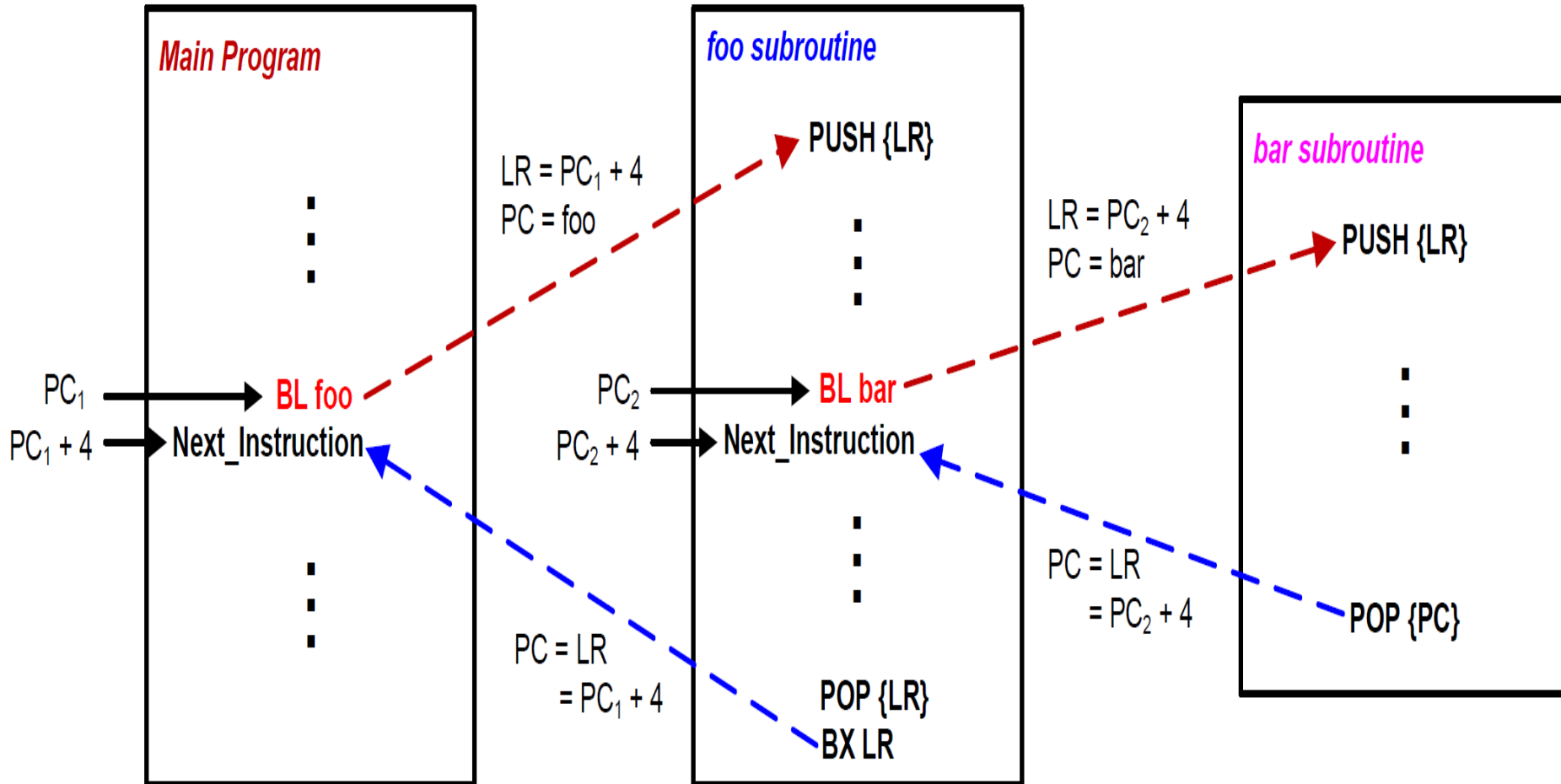
Caller Program	Subroutine/Callee
<pre>MOV r4, #100 ... BL foo ... ADD r4, r4, #1 ; r4 = 101, not 11</pre>	<pre>foo PROC ... MOV r4, #10 ; foo changes r4 ... BX LR ENDP</pre>

Call a Subroutine II:

Preserve Runtime Environment via Stack

Caller Program	Subroutine/Callee
<pre>MOV r4, #100 ... BL foo ... ADD r4, r4, #1 ; r4 = 101, not 11</pre>	<pre>foo PROC PUSH {r4} ; preserve r4 ... MOV r4, #10 ; foo changes r4 ... POP {r4} ; Recover r4 BX LR ENDP</pre>

Stacks and Subroutines



Subroutine Calling Another Subroutine

```
MAIN  
    MOV R0, #2  
    BL QUAD  
ENDL    ...
```

Function **MAIN**



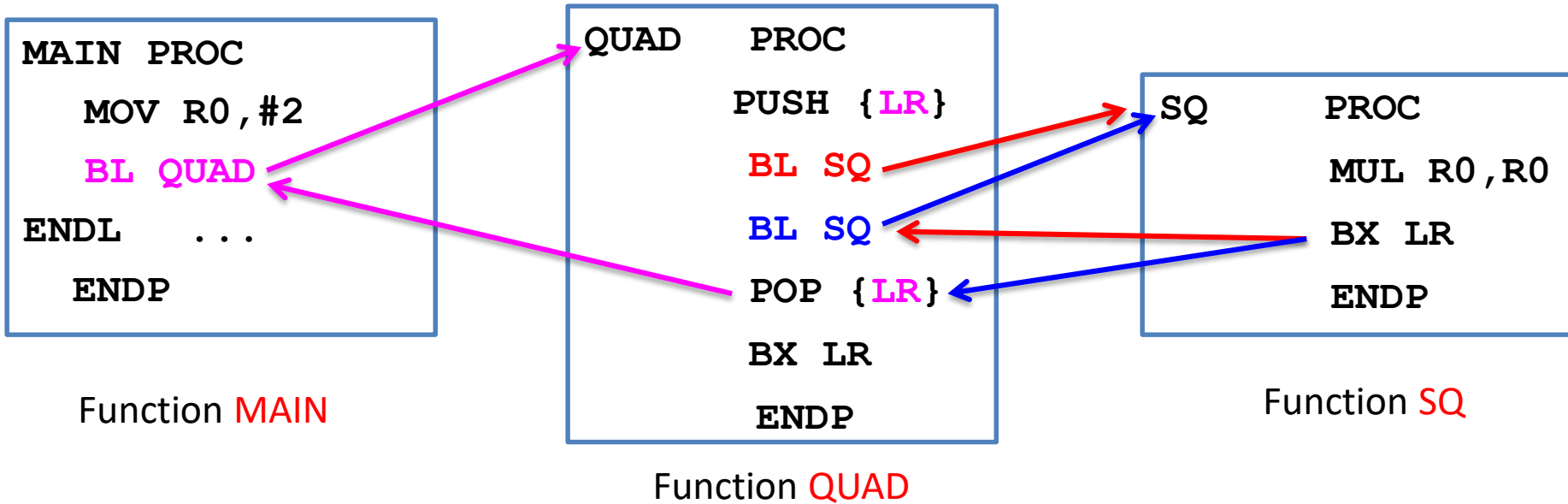
```
QUAD    PUSH {LR}  
        BL SQ  
        BL SQ  
        POP {LR}  
        BX LR
```

Function **QUAD**



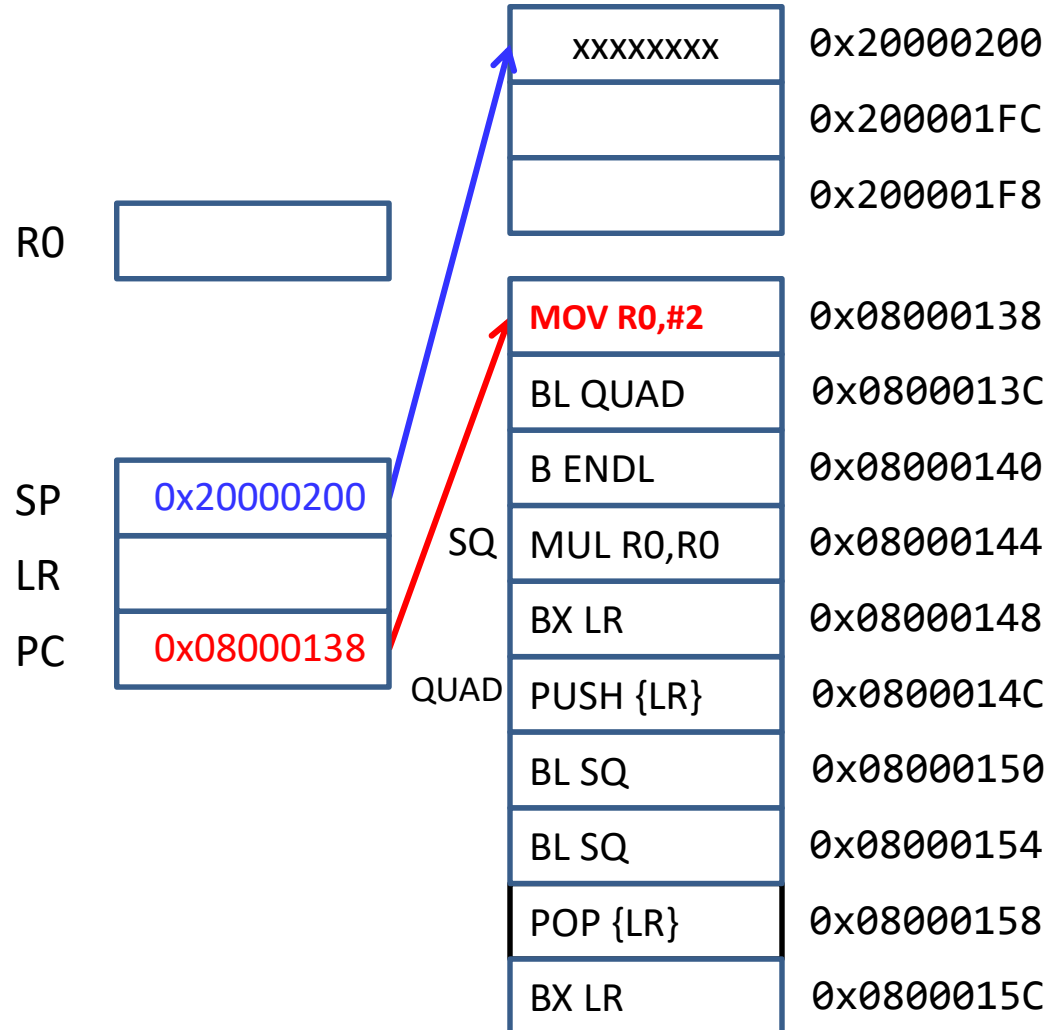
```
SQ      MUL R0, R0  
        BX LR
```

Function **SQ**

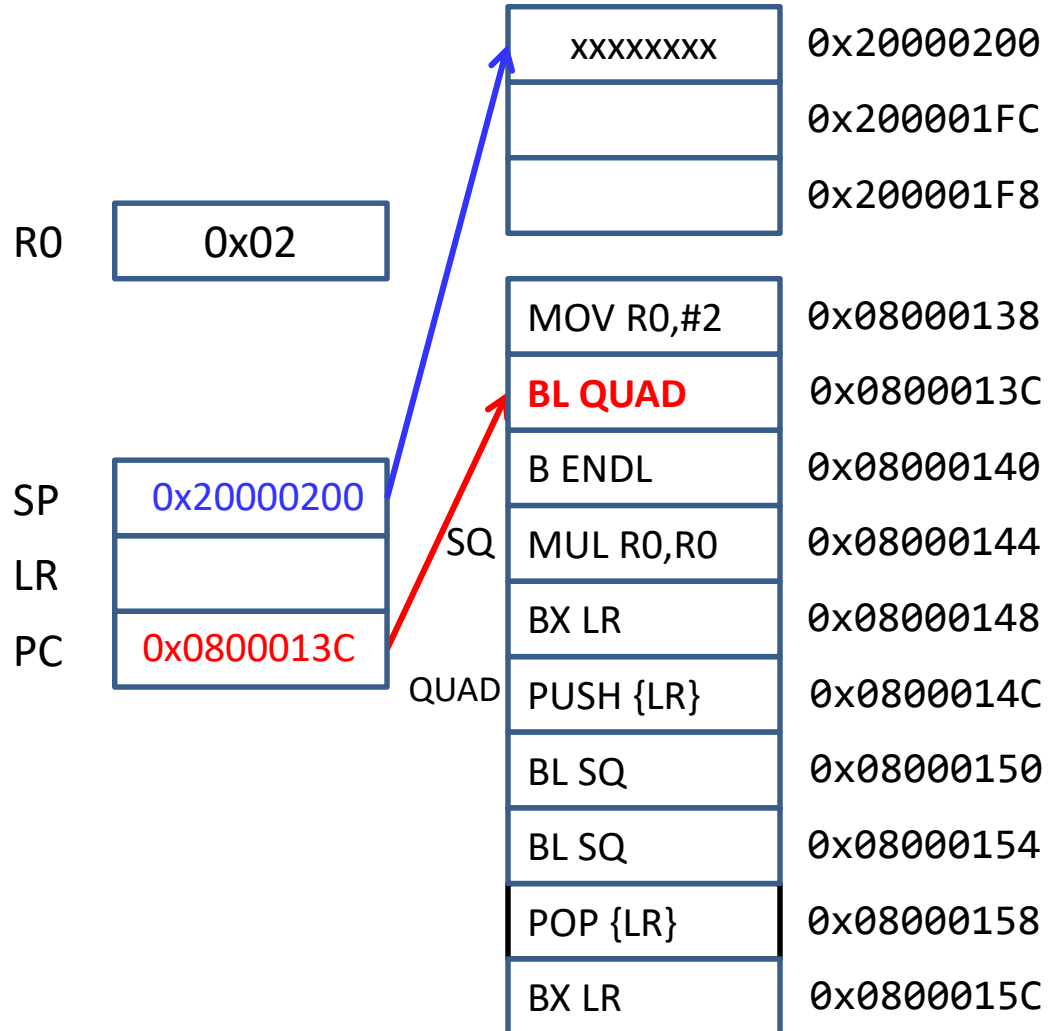


Example: R0 = R0⁴

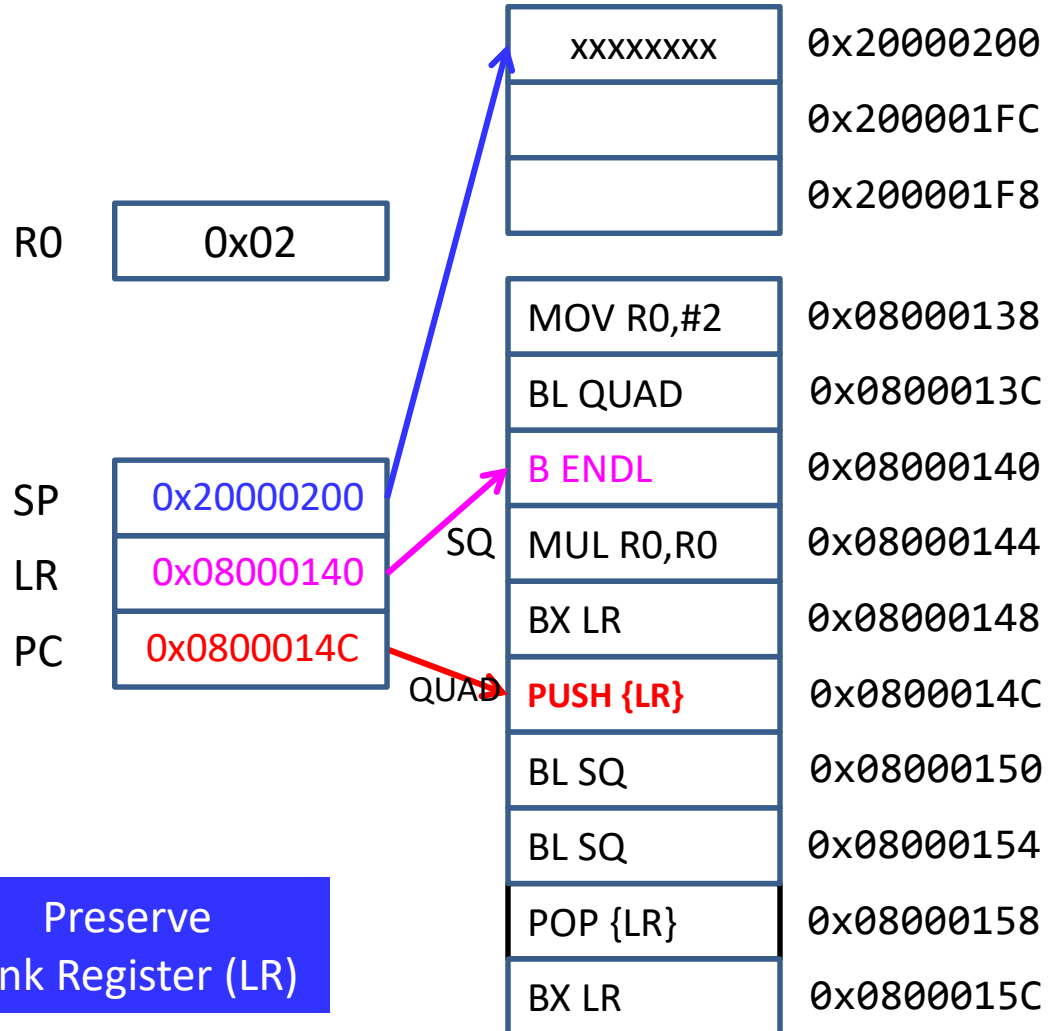
	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...



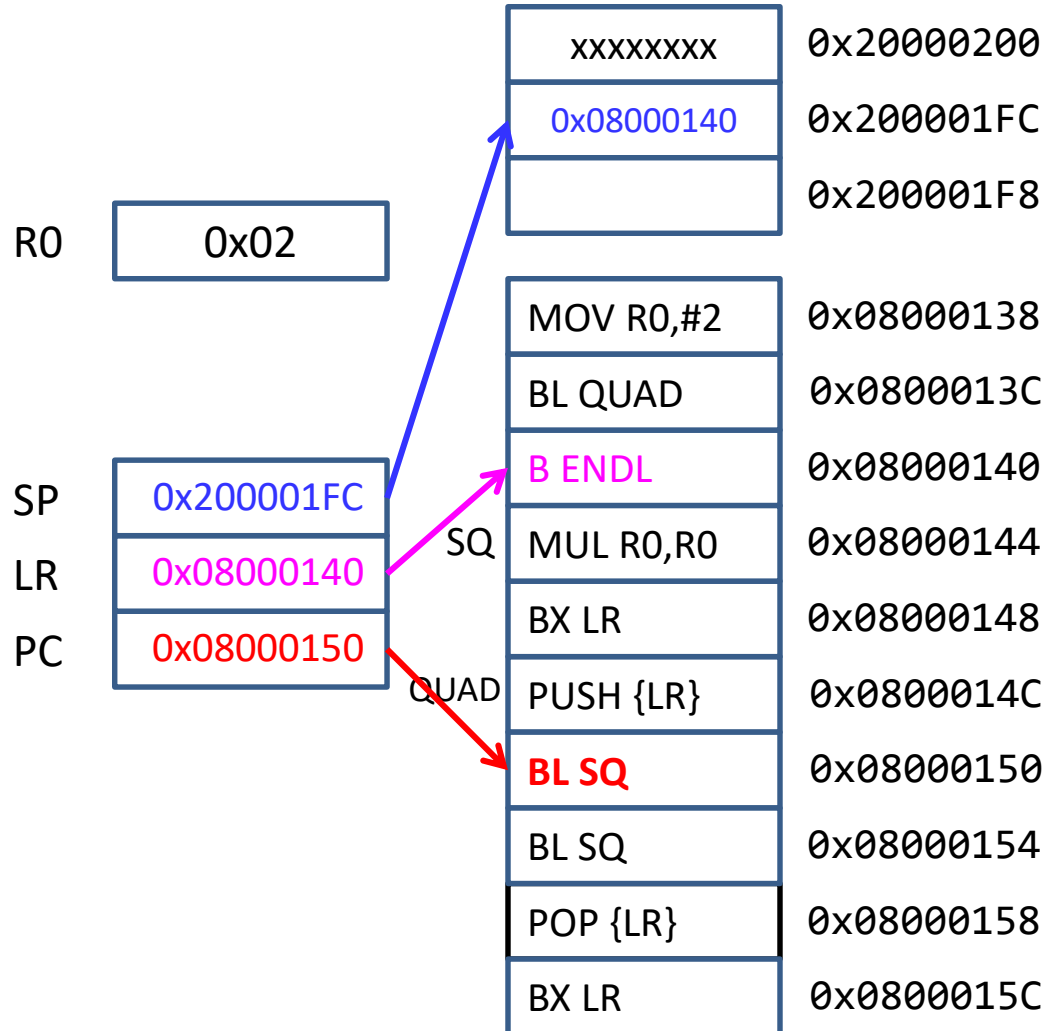
	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...



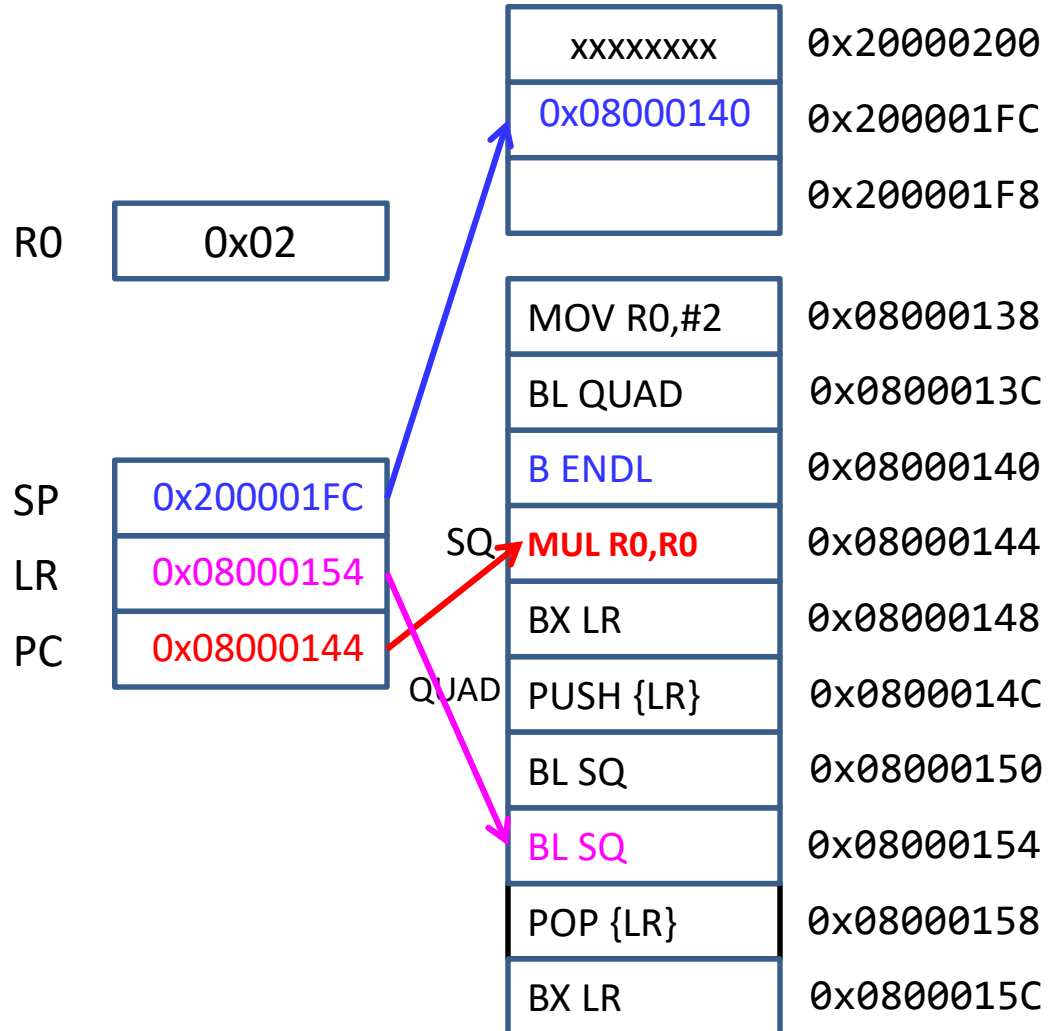
	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...

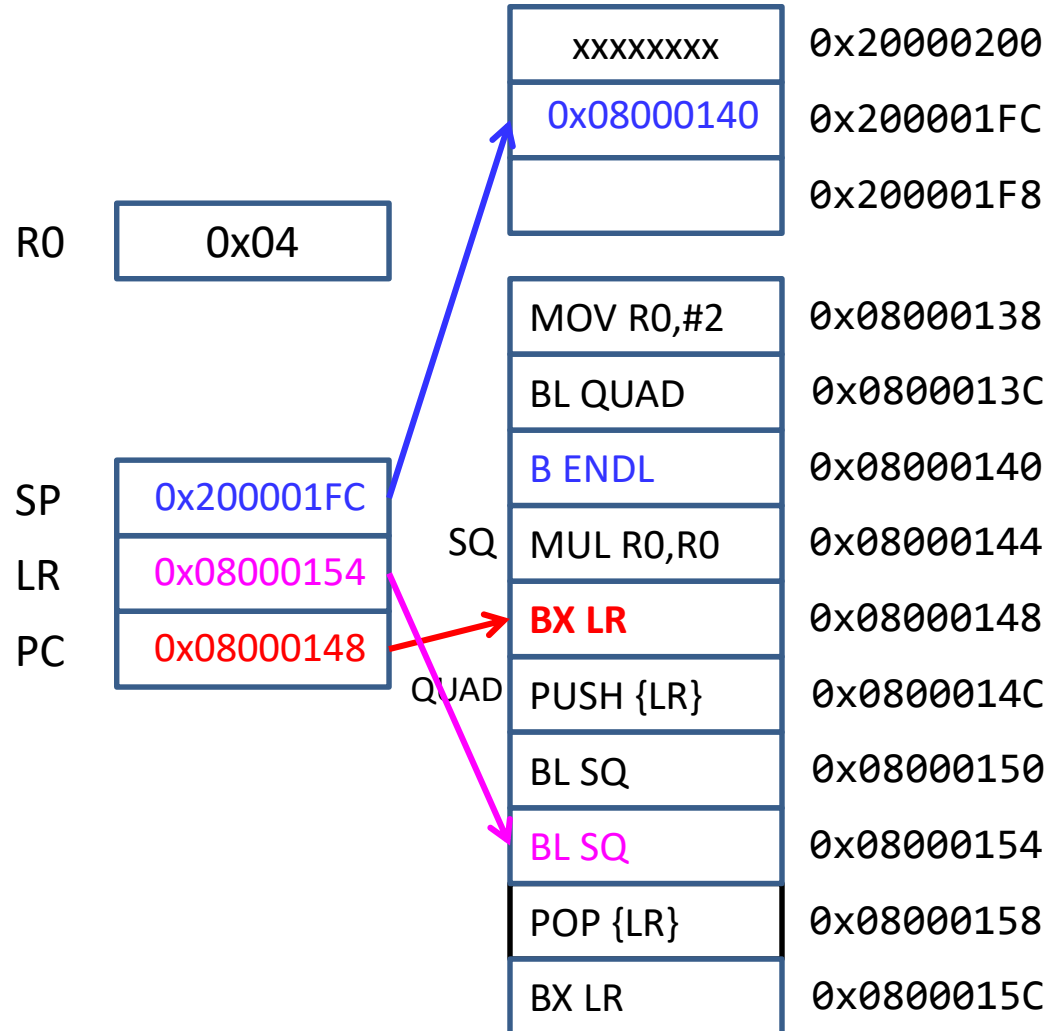
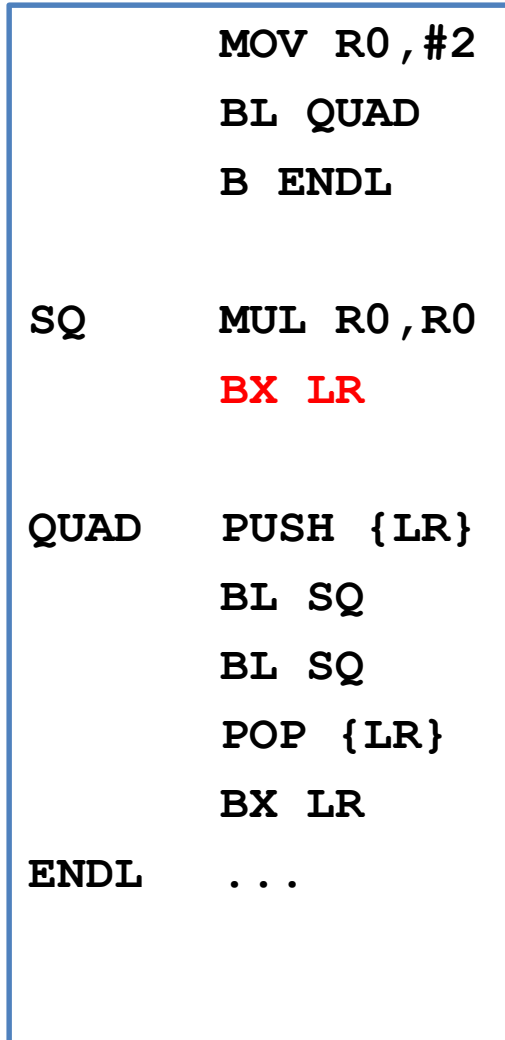


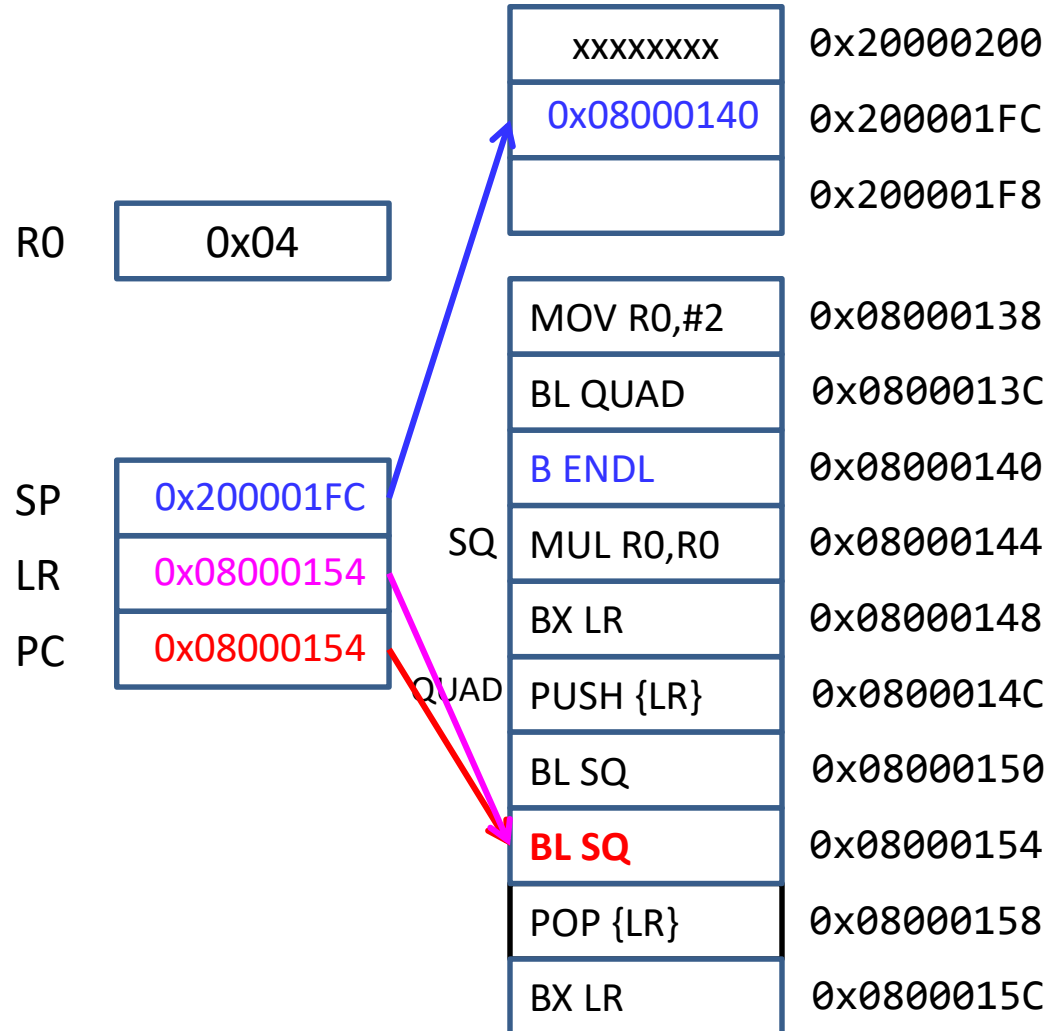
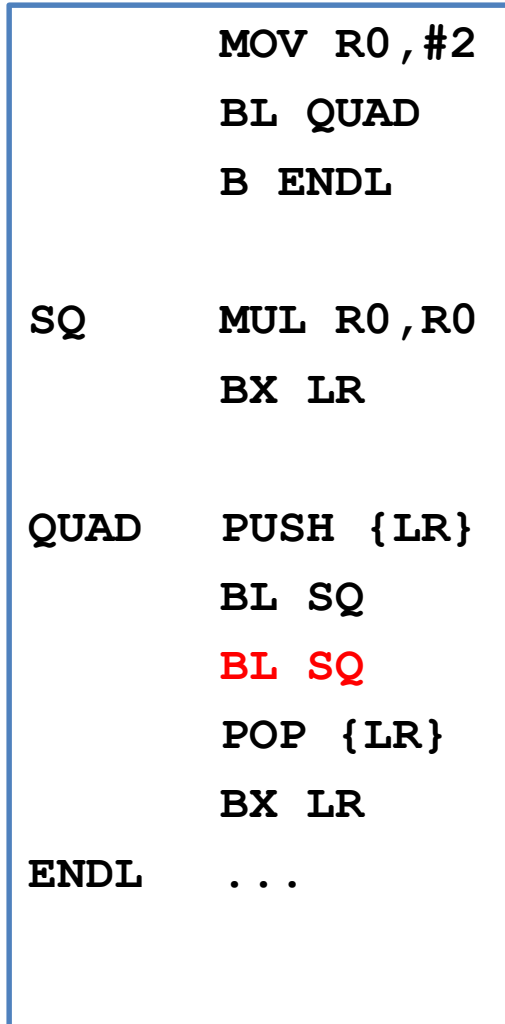
	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...



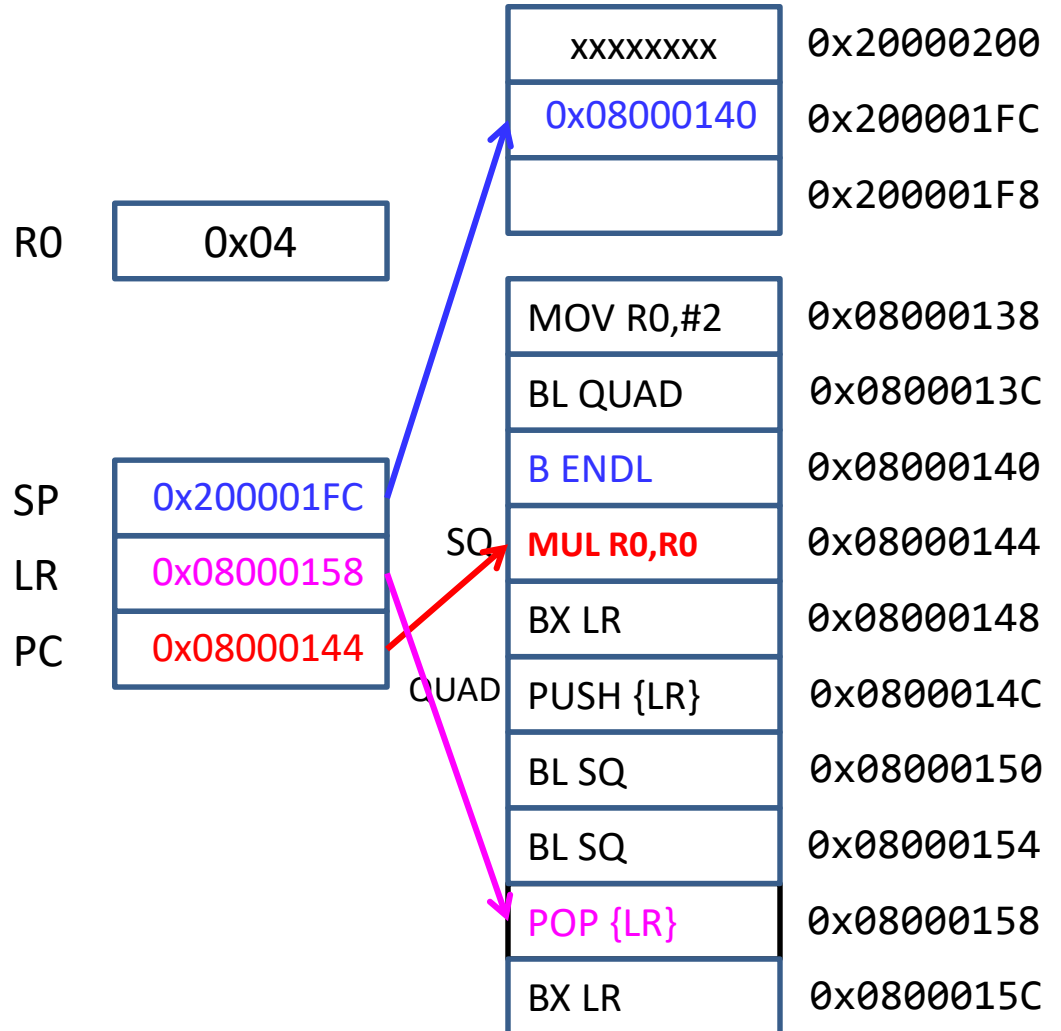
	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...







	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...



	MOV R0,#2
	BL QUAD
	B ENDL
SQ	MUL R0,R0
	BX LR
QUAD	PUSH {LR}
	BL SQ
	BL SQ
	POP {LR}
	BX LR
ENDL	...

