

Test4-Canvas-Graphs Results for Caden Roberts

❗ Correct answers are hidden.

Score for this quiz: **10** out of 10

Submitted Aug 30 at 9:02pm

This attempt took 28 minutes.

Question 1

4 / 4 pts

A common recommendation feature used in social networking websites is the number of common friends. In an undirected graph G , find a pair (u, v) with the largest number of common neighbors.

Your Answer:

Main Idea: A common neighbor of two nodes u and v in a graph G is a node that has edges from it to both u and v . We have to iterate through all possible node pairs and count their common neighbors to determine the pair with the largest count. For each pair of nodes (u, v) , count the # of common neighbors by checking the intersection of their adjacency lists. The pair with the maximum # of common neighbors is the result.

Pseudocode:

1. Initialize `maximumCommon = 0` and `resultPair = (Null, Null)`
 2. For each pair of vertices u and v in the graph:
 - 2a. If u is not v , find the intersection of the adjacency lists of u and v .
 - 2b. Let `commonCount = size of this intersection`.
 - 2c. If `commonCount > maximumCommon`:
 - set `maxCommon = commonCount`.
 - set `pair = (u, v)`.
- return `pair`.

Running time analysis:

Initializing variables takes $O(1)$. The loop iterates all vertices, giving $O(V^2)$ iterations, with V as the number of vertices.

To compute the intersection of their adjacency lists takes $O(\text{minimum of the degrees of } u \text{ or } v)$. So, the worst case time would be $O(V)$ time per pair. Returning the pair is $O(1)$. Overall running time is $O(V^3)$.

Question 2

6 / 6 pts

Biologists often construct a food network of species in an ecosystem. The vertices represent species, and a directed edge (u, v) means species u eats species v . An apex species is one that is not eaten by another species. Suppose you are given a list of all possible "eating" relationships (so a list of " u eats v "). Determine all apex species.

Your Answer:

Main Idea:

We have to identify apex species in a directed graph where vertices represent species and directed edges represent 'eats' relationships (u eats v for (u, v)). An apex species is a species that is not eaten by any other species. We have to find all vertices with no incoming edges. We will construct a directed graph from the list of ' u eats v ' relationships. For each vertex in the graph we will count the number of incoming edges. We collect all vertices with an incoming edge count of 0 and they are the apex species.

Pseudocode:

1. Initialize an empty dictionary graph to represent the adjacency list of the graph.
2. Initialize a dictionary degree to store the # of in-degrees of each vertex
3. Initialize a set apexSpecies to store the apex species.
4. For each pair (u, v) in the list of relationships:
 - 4a. If v is not in degree, set $\text{degree}[v]$ to 0.
 - 4b. Increment $\text{degree}[v]$ by 1.

4c. If u is not in degree, set $\text{degree}[u]$ to 0.

5. For each vertex u in the graph:

5a. If $\text{degree}[u]$ equals 0:

 Add u to apexSpecies.

return apexSpecies.

Running time analysis:

Constructing the graph using the adjacency list graph takes $O(E)$ because each edge is processed once. Initializing/updating degree dictionary takes $O(E)$ as well. Finding the apex species takes $O(V)$ to check each vertex. Overall running time then is $O(V+E)$.

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