Task1A

Here, we get the number of nodes and vertices from the input file and build an empty matrix based on them. Then we iterate through the rest of the input file to get first node, second node and the cost between them. Then this cost is set in the matrix for the two nodes in this way mat [node1] [node2] = cost

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like the previous task, we build an empty dictionary consisting of empty lists according to the number of nodes and vertices from the input file for energy configuration in the input file, the array of the node 1 in the dictionary is updated with value of node 2 and cost to go to node 2 from node 1.

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Task 2

We are given undirected unweighted graph to find path traversal. Inside the BFS function, we have a visited array and a avenue array, path array. The visited array is updated with a node when we corne across a mode that had not been visited before. The averue has the newly visited nodes while the averue is not empty, we pop a node from it and push it to the path array, we then push the neighbours of the node into the visited array or and averue if it is not visited. This continuous until the whole graph has been traversed.

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Task 3

The DFS function is a recursive function that for every unvisited neighbour, ealls DFS again, marks the node or Visited and pushes it into the path array. This way, the graph is traversed branchwise.

Task 4

The iscycle function calls the DFS function for every node in the graph. The DFS graph function maintains two arrays - visited (the overall visited nodes) and the stack (nodes found in current traversal). For every neighbour of the node, the DFS function is recursively called if It is not visited. This is to check if a eyell exists and if it possible to go back to the original node from another node. After each node is checked, the node is not included in the stack. If a cycle is detected, the DFS function returns True, which causes the iscycle function to also return True.

Task 5

Here, we find every path possible to take from the starting point and check if the end node is reached. There are two arrays-pathlist (holds all possible paths) pre visited (holds previously visited nodes to prement backtracking). For every pathlist, we take the last node, get all its children and check if the end node is among there. If not, and if the node is omisited, a new path is added to pathlist containing the previous path and this node added. It is also set an visited. In this way, we check paths till we find the shortest one. Time is just number of nodes - 1.

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We use recursive DFs to explore the jurgle. It starts from a cell and traverses the neighbour cells in all directions. The encountered diamonds are added to a diamond count variable. The visited cells are marked to prenent backtracking. # are obstacles and are avoided. The count biamonds function iterates through all cells, calling DFs FloodFILI for all empty (1) cells to find the maximum diamonds. The maximum of this result is stored.