# Machine Problem 3: UNIX Processes (Due: 02/24/17)

### <u>Introduction</u>:

The /proc directory is a pseudo filesystem that allows access to kernel data structures while in user space. It allows you to view some of the information the kernel keeps on running processes. To view information about a specific process you just need to view files inside of the directory: /proc/[pid]. For more information simply view the manpage with man proc.

# The Assignment:

Using the files stored at /proc write a program/script to find information about a specific process using a user provided pid. In the following, you will find a list of the task\_struct members for which you are required to find their value. In the task\_struct a lot of the data you are finding is not represented as member values but instead pointers to other linux data structures that contain these members. All of the information you will be retrieving can be found in a process proc directory (/proc/[pid]). Your program must be able to retrieve the following data about any given process if the given process Id is existing under directory /proc/:

Table #1: Process Attributes

Category	Required Variables/Items	Description
Identifiers	PID, PPID	Process ID of the current process and its parent
	EUID, EGID	Effective user and group ID
	RUID, RGID	Real user and group ID
	FSUID, FSGID	File system user and group ID
State	R, S, D, T, Z, X	Running, Sleeping, Disk sleeping, Stopped,
		Zombie, and Dead
Thread	Thread_Info	Thread IDs of a process
Information		
Priority	Priority Number	Integer value from 1 to 99 for real time processes
	Niceness Value	Integer value from -20 to 19
Time	stime & utime	Time that a process has been scheduled
Information		in kernel/user mode
	cstime & cutime	Time that a process has waited on children
		being run in kernel/user mode
Address	Startcode & Endcode	The start and end of a process
Space	ESP & EIP	in memory
Resources	File Handles &	Number of fds used, and number of
	Context Switches	voluntary/involuntary context switches
Processors	Allowed processors and	Which cores the process is allowed to run on,
	Last used one	and which one was last used
Memory	Address range, permissions	Output a file containing the process's
Map	offset, dev, inode,	currently mappined memory regions
	and path name	

#### Refer to Machine Problem 2 for exec function family and fork

### Deliverables:

#### • Code:

- You can do this project in three different languages: C++, Bash and Python. If you do it in C++, then you are to turn in the following files: proctest.h, proctest.cpp, main.cpp and Makefile. If you do it in Bash, submit one file: proctest.sh. If you do it in Python, submit one file: proctest.py. Make sure you also submit the Makefile if you do it in C++.
- Be sure to start with the skeleton code provided with this assignment. You can implement those functions in the files in your way. However, do not change the declarations of the functions or the grading script that will test your assignment will not be able to give you the points you deserve.
- Website for learning bash: [https://www.ibm.com/developerworks/library/l-bash/]

#### • Report:

- Answer the following additional questions
  - For a process run by a user other than yourself, find the following items from Table #1: [Identifiers, State, Thread Information, Priority, Time Information, Resources, and Memory Map]
  - 2. For a process that you have created, retrieve all items enumerated in Table.
  - 3. What are the differences between the real user IDs and effective user IDs, and what is a situation where these will be different?
  - 4. Why are most of the files in /proc read only?
  - 5. Why is the task\_struct so important to the kernel and what is it used for?

# Advanced Concepts:

#### Task\_Struct Members

Linux and UNIX distributions organize all internal process data into structures named task\_struct. The kernel maintains one task\_struct member for each process running on the system. Each task\_struct member contains two pointers specifically designated to point to other task\_struct members. As such, the kernel organizes all task\_struct members into linked lists. On startup, the kernel also initializes a pid hash table whose elements are linked lists. This is done to save some time searching for process data structures. Instead of having to search one large list (for perspective, pid\_max on the author's Debian 8 system is  $2^{15}$ ), the kernel quickly hashes the process pid into one of the hash table elements and

then searches a much smaller list to find the correct task\_struct member (Bovet, 81 & Glass, 555).

A visual description of the organization of the task\_struct structure is provided in Figure 1. Also provided, for the curious reader, is the entire listing of all task\_struct members. Inside, you should be able to see all of the fields that you are required to find for the assignment (or at least pointers to other kernel data structures that actually contain those fields). For even further reading see *Bovet* and *Glass* in the bibliography section

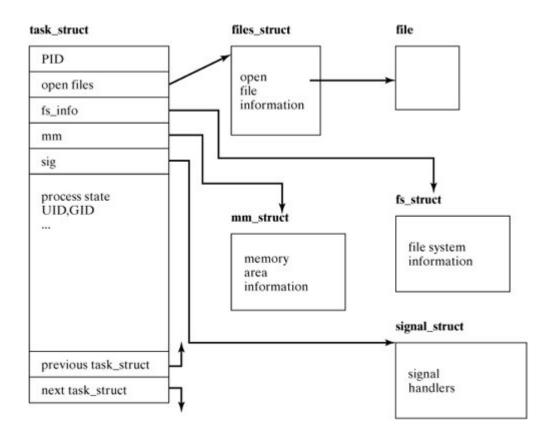
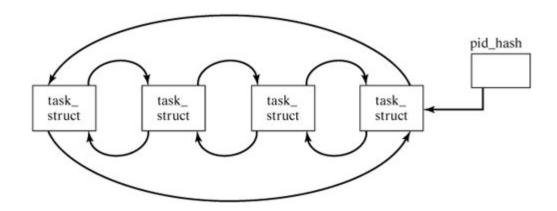


Figure #1: Task Struct Members & Fields (Glass, 554)

Figure #2: Pid Hash Table (Glass, 555)



Example #1: task\_struct Declaration

```
struct task_struct {
/* these are hardcoded - don't touch */
 volatile long
                                      /*-1 unrunnable, 0 runnable, >0 stopped */
                       state;
 long
                      counter;
 long
                       priority;
 unsigned
                      long signal;
                                      /* bitmap of masked signals */
 unsigned
                      long blocked;
                                      /* per process flags, defined below */
 unsigned
                      long flags;
 int errno;
                      debugreg[8];
                                      /* Hardware debugging registers */
 long
 struct exec_domain *exec_domain;
/* various fields */
 struct linux_binfmt *binfmt;
 struct task_struct
                      *next_task, *prev_task;
                       *next_run, *prev_run;
 struct task_struct
                       saved_kernel_stack;
 unsigned long
 unsigned long
                      kernel_stack_page;
 int
                      exit_code, exit_signal;
 /*Other features of a process. personality is for setting process execution domain;
      dumpable:1 means process can be dumpable and so on.*/
 unsigned long
                       personality;
 int
                      dumpable:1;
                      did_{exec}:1;
 int
 int
                      pid;
 int
                      pgrp;
 int
                       tty_old_pgrp;
                       session;
 /* boolean value for session group leader */
                      leader;
 int
                      groups[NGROUPS];
 int
 /*
```

```
* pointers to (original) parent process, youngest child, younger sibling,
   * older sibling, respectively. (p-) father can be replaced with
   * p->p_pptr->pid)
 struct task_struct
                       *p_opptr, *p_pptr, *p_cptr,
                       *p_ysptr, *p_osptr;
 struct wait_queue
                       *wait_chldexit;
 unsigned short
                       uid, euid, suid, fsuid;
 unsigned short
                       gid, egid, sgid, fsgid;
 unsigned long
                       timeout, policy, rt_priority;
                       it_real_value , it_prof_value , it_virt_value ;
 unsigned long
 unsigned long
                       it_real_incr , it_prof_incr , it_virt_incr ;
 struct timer_list
                       real_timer;
                       utime, stime, cutime, cstime, start_time;
 long
/* mm fault and swap info: this can arguably be seen as either
  mm-specific or thread-specific */
 unsigned long
                       min_flt, maj_flt, nswap, cmin_flt, cmaj_flt, cnswap;
 int swappable:1;
 unsigned long
                       swap_address;
 unsigned long
                                       /* old value of maj_flt */
                       old_maj_flt;
                                       /* page fault count of the last time */
 unsigned long
                       dec_flt;
 unsigned long
                       swap_cnt;
                                       /* number of pages to swap on next pass */
/* limits */
 struct rlimit
                       rlim[RLIM_NLIMITS];
                       used_math;
 unsigned short
                       comm[16];
 char
/* file system info */
 int
                       link_count;
                                       /* NULL if no tty */
 struct ttv_struct
                       *ttv:
/* ipc stuff */
 \mathbf{struct}\ \mathrm{sem\_undo}
                       *semundo:
 struct sem_queue
                      *semsleeping;
/* ldt for this task - used by Wine. If NULL, default_ldt is used */
 struct desc_struct *ldt:
/* tss for this task */
 struct thread_struct tss;
/* filesystem information */
 struct fs_struct
                       *fs:
/* open file information */
 struct files_struct
                       * files ;
/* memory management info */
 struct mm_struct
                       *mm;
/* signal handlers */
 struct signal_struct *sig;
#ifdef __SMP__
```

# UNIX man Pages:

One incredibly useful feature of UNIX operating systems that many new developers do not know about is the built in manual system. Using the command 'man', you can access information about most aspects of the operating system from general commands all the way to system call APIs.

The structure of man pages in UNIX are organized into sections by number follows (Wikimedia Foundation):

- 1. General Commands
- 2. System Calls
- 3. Library Functions (Specifically, the C standard library)
- 4. Special Files
- 5. File Formats
- 6. Games
- 7. Miscellanea
- 8. System Administration

You may (or may not, that's fine too) find the following manual pages useful when creating this assignment. Each one of these lines can be executed as a valid shell command to open a particular manual page. Note, the number indicates the manual section that that function resides.

- man 3 exec
- man 2 fork
- man 2 chdir
- man 2 pipe
- man 2 dup
- man 2 wait

# Grading Rubric:

Points will be assigned according to the following rubric. The code portion of the assignment is worth 160 points, divided as shown below. The report is worth an additional 40 points. Consequently, the point total for the assignment is 200 points.

- getpid(): 5 points
- getppid(): 5 points
- geteuid(): 5 points
- getegid(): 5 points
- getruid(): 5 points
- getrgid(): 5 points
- getfsuid(): 5 points
- getfsgid(): 5 points
- getstate(): 5 points
- getthread\_count(): 5 points
- getpriority(): 5 points
- getniceness(): 5 points
- getstime(): 5 points
- getutime(): 5 points
- getcstime(): 5 points
- getcutime(): 5 points
- getstartcode(): 5 points
- getendcode(): 5 points
- getesp(): 5 points
- geteip(): 5 points
- getfiles(): 20 points
- getvoluntary(): 5 points
- getnonvol(): 5 points
- getallowed\_cpus(): 5 points
- getlast\_cpu(): 5 points
- getmemory\_map(): 20 points

# Bibliography:

[1] Bovet, Daniel Pierce. Understanding the Linux Kernel: From I/O Ports to Process Management. U.S.A: O'Reilly, 2003. Print.

[2] Glass, Graham. Linux for Programmers and Users. Upper Saddle River, NJ: Pearson Prentice Hall, 2006. Print.