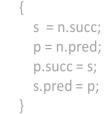
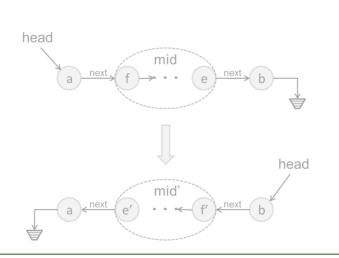
$\exists c \forall in \ Q(c, in)$

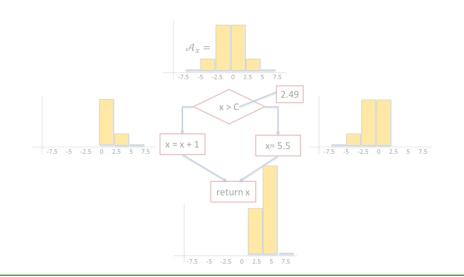
```
/* Average of x and y without using x+y (avoid overflow)*/
int avg(int x, int y) {
  int t = expr({x/2, y/2, x%2, y%2, 2 }, {PLUS, DIV});
  assert t == (x+y)/2;
  return t;
}
```

```
f_1
f_2
f_3
f_3
f_3
f_4
f_5
f_7
```



Program Synthesis







Sk[c](in)

Lecture 1 Course Overview and Introduction to Synthesis

Nadia Polikarpova

Instructor



Nadia Polikarpova

- Assistant Professor since 2017
- Before that: postdoc at MIT with Armando Solar-Lezama
- Before that: PhD at ETH Zurich
- Research areas: program synthesis and program verification
- she / her

Logistics

Lecture

- When: Tue/Thu 3:30-4:50
- Where: Zoom

Office Hours

- When: after class (4:50-5:30)
- Where: same Zoom link

Course Website

- https://github.com/nadia-polikarpova/cse291-program-synthesis
- To ask questions: use canvas

1. Understand what program synthesis can do and how

1. Understand what program synthesis can do and how

2. Use existing synthesis tools

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2. Use existing synthesis tools

3. Contribute to synthesis techniques and tools towards a publication in an academic conference

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2. Use existing synthesis tools

3. Contribute to synthesis techniques and tools towards a publication in an academic conference

lectures read and discuss research papers

project

Evaluation

Class Participation: 5%

- answer questions in class
- answer your peers' questions on canvas

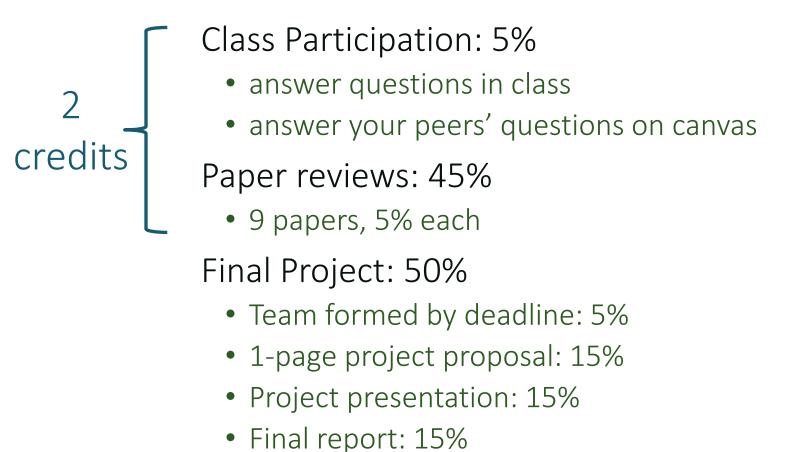
Paper reviews: 45%

• 9 papers, 5% each

Final Project: 50%

- Team formed by deadline: 5%
- 1-page project proposal: 15%
- Project presentation: 15%
- Final report: 15%

Evaluation



4 credits

Papers reviews

Due on Wed of weeks 2-10, by the end of the day

First review due next week

Posted on the Reading List at least a week before due date

Reviews submitted via a Google Form: see wiki

• Link posted on Reading List

Review content: see wiki

Project

Kinds of projects:

- re-implement a technique from a paper
- apply existing synthesis framework to a new domain
- extend/improve existing synthesis algorithm or tool
- develop a new synthesis algorithm or tool
- •

Judged in terms of

- quality of execution
- originality
- scope

Project

Team forming

Teams of 2/3

Pick a project:

 List of suggested projects on the wiki (but feel free to propose your own)

Talk to me!

One page: explain what you plan to do and give some evidence that you've started to work on it

During the exam week

• ~5-10 min per project

3-8 pages, structured like a research paper

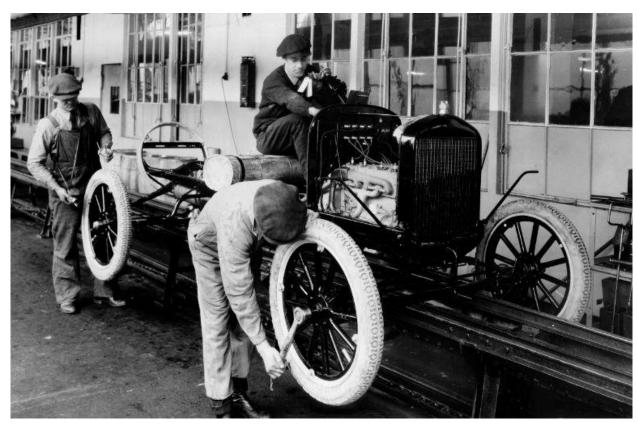
Proposal

Presentation

Report

And now the good stuff

The goal: automate programming





What is program synthesis?



What is program synthesis?



The FORTRAN Automatic Coding System

J. W. BACKUS†, R. J. BEEBER†, S. BEST‡, R. GOLDBERG†, L. M. HAIBT†, H. L. HERRICK†, R. A. NELSON†, D. SAYRE†, P. B. SHERIDAN†, H. STERN†, I. ZILLER†, R. A. HUGHES§, AND R. NUTT||

Introduction

HE FORTRAN project was begun in the summer of 1954. Its purpose was to reduce by a large factor the task of preparing scientific problems for IBM's next large computer, the 704. If it were possible for the 704 to code problems for itself and produce as

system is now complete. It has two components: the FORTRAN language, in which programs are written, and the translator or executive routine for the 704 which effects the translation of FORTRAN language programs into 704 programs. Descriptions of the FORTRAN language and the translator form the principal

```
append:
    push ebp
    mov ebp, esp
    push eax
    push ebx
    push len
    call malloc
    mov ebx, [ebp + 12]
    mov [eax + info], ebx
    mov dword [eax + next], 0
    mov ebx, [ebp + 8]
    cmp dword [ebx], 0
    je null_pointer
    mov ebx, [ebx]
next_element:
    cmp dword [ebx + next], 0
    je found_last
    mov ebx, [ebx + next]
    jmp next_element
found_last:
    push eax
    push addMes
    call puts
    add esp, 4
    pop eax
    mov [ebx + next], eax
go_out:
    pop ebx
    pop eax
    mov esp, ebp
    pop ebp
    ret 8
null_pointer:
    push eax
    push nullMes
    call puts
    add esp, 4
    pop eax
    mov [ebx], eax
    jmp go_out
```

Assembly

```
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    je null_pointer
                                             void insert(node *xs, int x) {
    mov ebx, [ebx]
                                               node *new;
                                               node *temp;
next_element:
                                               node *prev;
    cmp dword [ebx + next], 0
    je found last
                                               new = (node *)malloc(sizeof(node));
    mov ebx, [ebx + next]
                                               if(new == NULL) {
    jmp next element
                                                 printf("Insufficient memory.");
                                                 return;
found_last:
    push eax
                                               new->val = x;
    push addMes
                                               new->next = NULL;
    call puts
                                               if (xs == NULL) {
    add esp, 4
                                                 xs = new;
    pop eax
                                               } else if(x < xs->val) {
    mov [ebx + next], eax
                                                 new->next = xs;
                                                 xs = new;
go_out:
                                               } else {
    pop ebx
                                                 prev = xs;
    pop eax
                                                  temp = xs->next;
    mov esp, ebp
                                                 while(temp != NULL && x > temp->val) {
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    push nullMes
                                                 } else {
    call puts
                                                   new->next = temp;
    add esp, 4
                                                   prev->next = new;
    pop eax
    mov [ebx], eax
    jmp go_out
```

Assembly

 C

```
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  if(new == NULL) {
   printf("Insufficient memory.");
    return;
  new->val = x;
  new->next = NULL;
  if (xs == NULL) {
   xs = new;
  } else if(x < xs->val) {
    new->next = xs;
    xs = new;
  } else {
    prev = xs;
    temp = xs->next;
    while(temp != NULL && x > temp->val) {
      prev = temp;
      temp = temp->next;
    if(temp == NULL) {
      prev->next = new;
    } else {
      new->next = temp;
      prev->next = new;
```

```
insert x xs =
  match xs with
  Nil → Cons x Nil
  Cons h t →
   if x ≤ h
      then Cons x xs
      else Cons h (insert x t)
```

Assembly

(

Haskell

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    mov [ebx], eax
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```

"Any sufficiently advanced compiler is indistinguishable"

from a synthesizer"

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```

```
"Any sufficiently advanced compiler is indistinguishable"
```

```
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 node *temp;
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  new = (node *)malloc(sizeof(node));
  if(new == NULL) {
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   return;
  new->val = x;
  new->next = NULL;
 if (xs == NULL) {
    xs = new;
 } else if(x < xs->val) {
    new->next = xs;
    xs = new;
 } else {
    prev = xs;
    temp = xs->next;
    while(temp != NULL && x > temp->val) {
     prev = temp;
      temp = temp->next;
    if(temp == NULL) {
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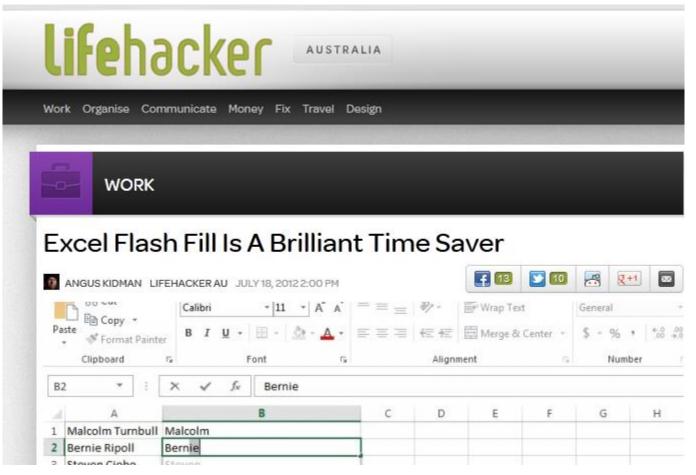


Modern program synthesis: FlashFill

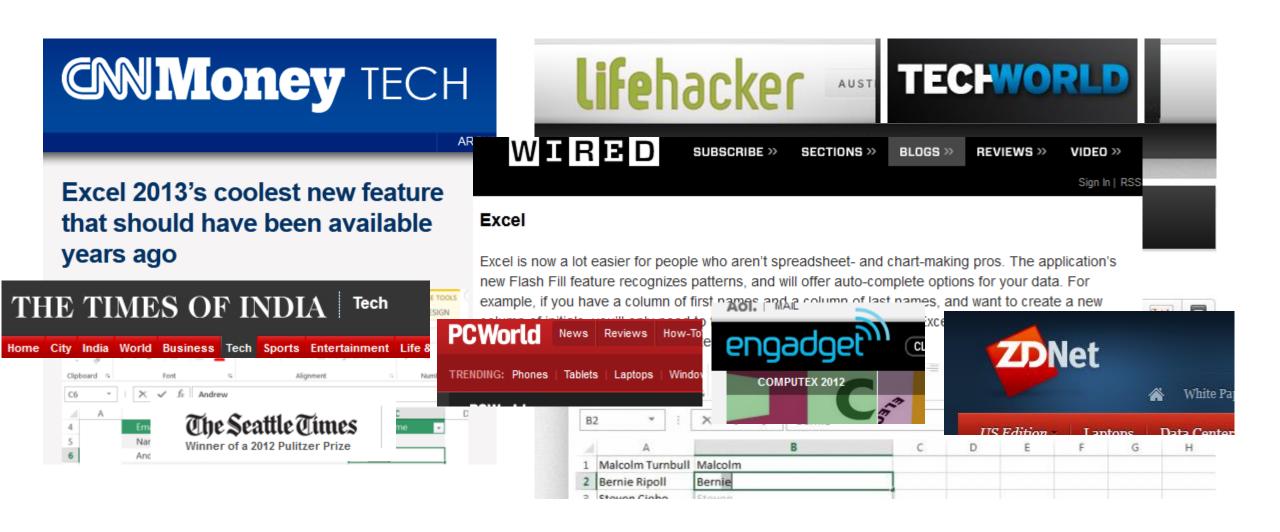


Modern program synthesis: FlashFill



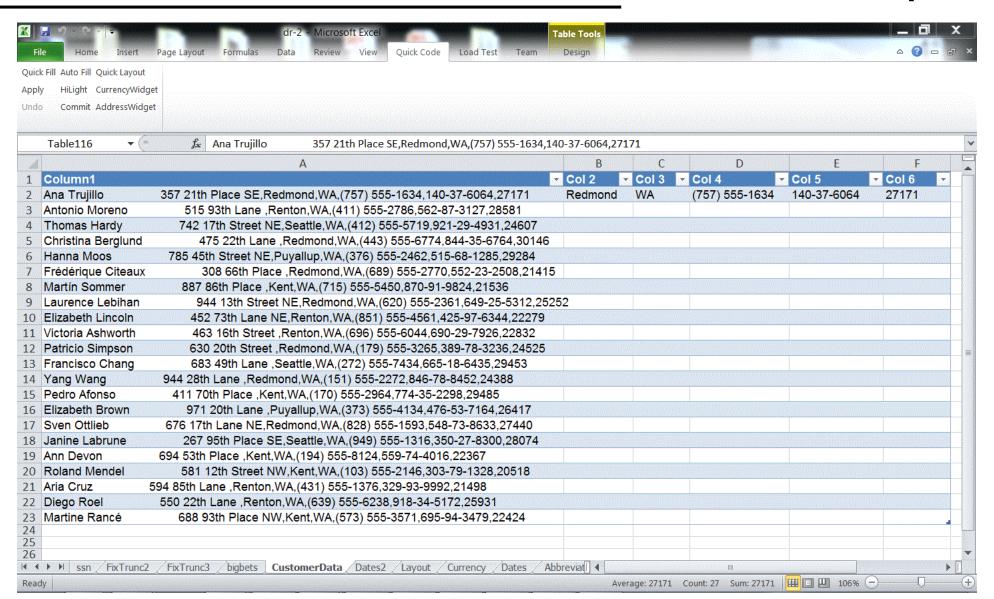


[Gulwani 2011]

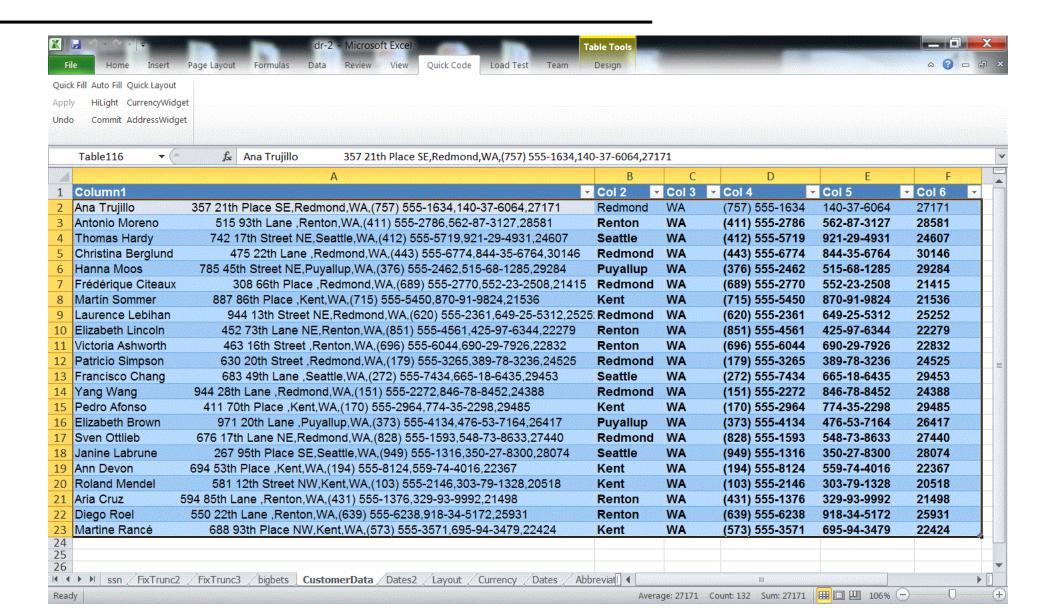


FlashFill: a feature of Excel 2013

[Gulwani 2011]



FlashFill: a feature of Excel 2013



[Solar-Lezama 2013]

Problem: isolate the least significant zero bit in a word

• example: 0010 0101 → 0000 0010

[Solar-Lezama 2013]

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Easy to implement with a loop

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Can this be done more efficiently with bit manipulation?

Problem: isolate the least significant zero bit in a word

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Easy to implement with a loop

Can this be done more efficiently with bit manipulation?

- Trick: adding 1 to a string of ones turns the next zero to a 1
- i.e. 000111 + 1 = 001000

Sketch: space of possible implementations

```
/**
 * Generate the set of all bit-vector expressions
 * involving +, &, xor and bitwise negation (~).
*/
generator bit[W] gen(bit[W] x){
    if(??) return x;
    if(??) return ??;
    if(??) return ~gen(x);
    if(??){
        return {| gen(x) (+ | & | ^) gen(x) |};
```

Sketch: synthesis goal

```
generator bit[W] gen(bit[W] x, int depth){
    assert depth > 0;
    if(??) return x;
    if(??) return ??;
    if(??) return ~gen(x, depth-1);
    if(??){
        return {| gen(x, depth-1) (+ | & | ^{\circ}) gen(x, depth-1) |};
bit[W] isolate0fast (bit[W] x) implements isolate0 {
     return gen(x, 3);
```

Sketch: output

```
bit[W] isolate0fast (bit[W] x) {
  return (~x) & (x + 1);
}
```

Modern program synthesis: Synquid

[Polikarpova et al. 2016]

Problem: intersection of strictly sorted lists

• example: intersect [4, 8, 15, 16, 23, 42] [8, 16, 32, 64] \rightarrow [8, 16]

Also: we want a guarantee that it's correct on all inputs!

Step 1: define synthesis goal as a *type*

```
intersect :: xs:List a → ys:List a →
    List a
```

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```
intersect :: xs:SList a → ys:SList a →
{v:SList a | elems v = elems xs n elems ys}
```

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Step 2: define a set of components

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Step 2: define a set of components

• Which primitive operations is our function likely to use?

```
Step 1: define synthesis goal as a type
intersect :: xs:SList a → ys:SList a →
{v:SList a | elems v = elems xs n elems ys}
```

Step 2: define a set of components

- Which primitive operations is our function likely to use?
- Here: {Nil, Cons, <}

```
intersection = \xs . \ys .
  match xs with
   Nil -> xs
   Cons x xt ->
     match ys with
        Nil -> ys
        Cons y yt ->
         if x < y
          then intersection xt ys
          else
            if y < x
            then intersection xs yt
            else Cons x (intersection xt yt)
```

```
intersection = \xs . \ys .
                                                XS
                                                                    ys
                                                                                 result
  match xs with
                                     [4, 8, 15, 16, 23, 42] [8, 16, 32, 64]
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                                                                  ys
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                                                  [23, 42]
                                                                  [32, 64]
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                                                      [42]
                                                                 [32, 64]
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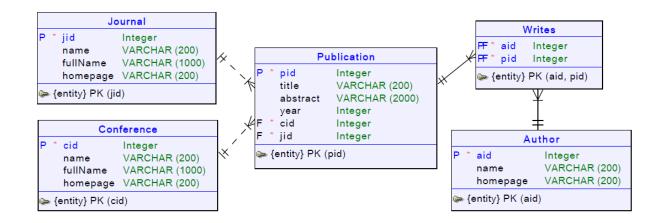
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                                                                                 [8]
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                                                                               [8, 16]
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                                                                   [32, 64]
                                                   [23, 42]
         then intersection xt ys
                                                                  [32, 64]
                                                       [42]
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                                                                                 [8]
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                                                                               [8, 16]
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                                                   [23, 42]
         then intersection xt ys
                                                                   [32, 64]
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                                                                       [64]
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```

Modern program synthesis: SQLizer

[Yaghmazadeh et al. 2017]

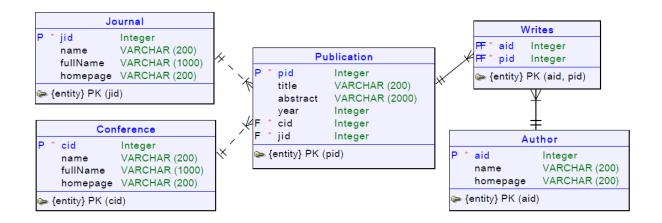
Problem: "Find the number of papers in OOPSLA 2010"



Modern program synthesis: SQLizer

[Yaghmazadeh et al. 2017]

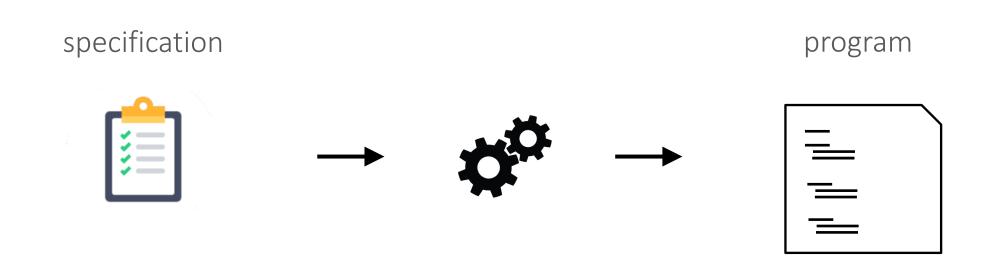
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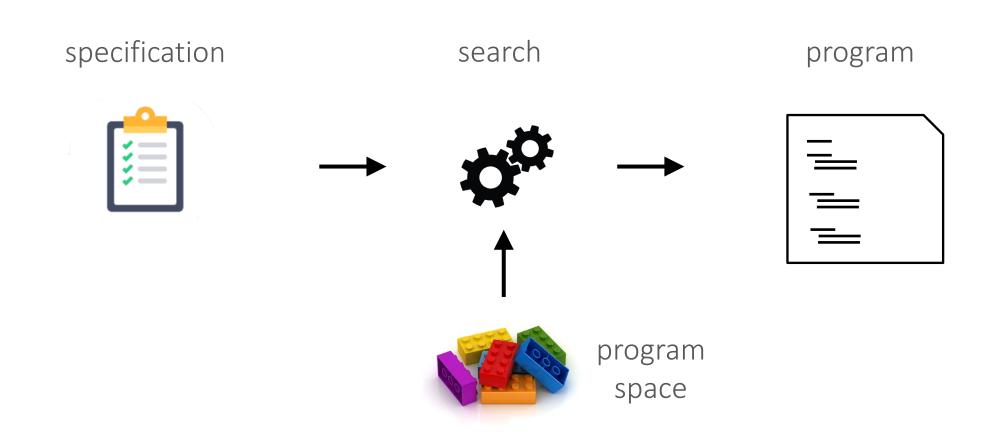
•Output:

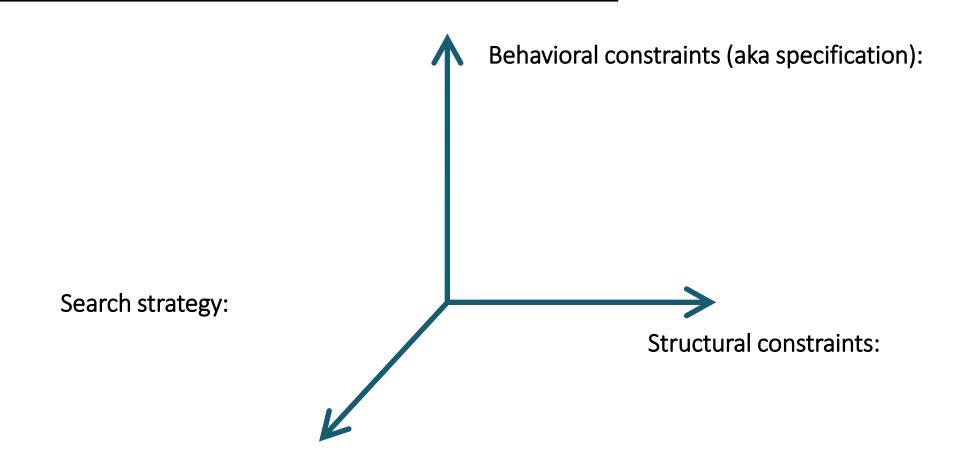
```
SELECT count(Publication.pid)
FROM Publication JOIN Conference ON Publication.cid = Conference.cid
WHERE Conference.name = "OOPSLA" AND Publication.year = 2010
```

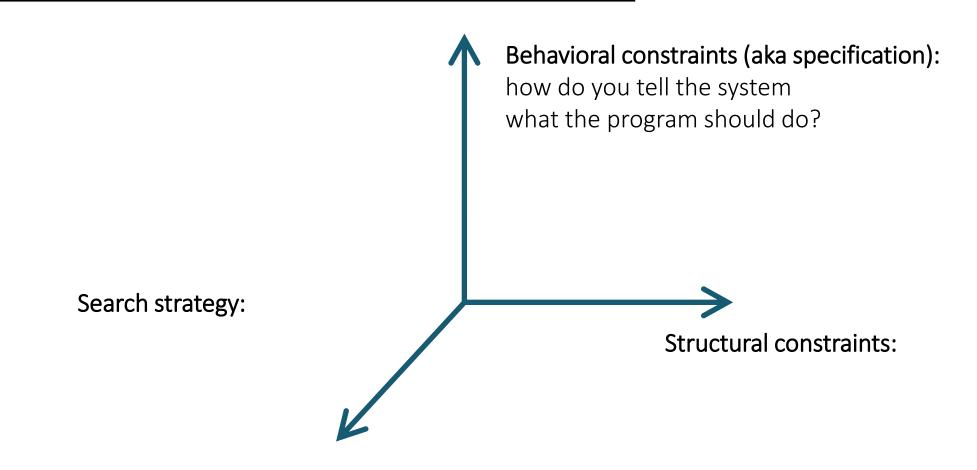
What is program synthesis?

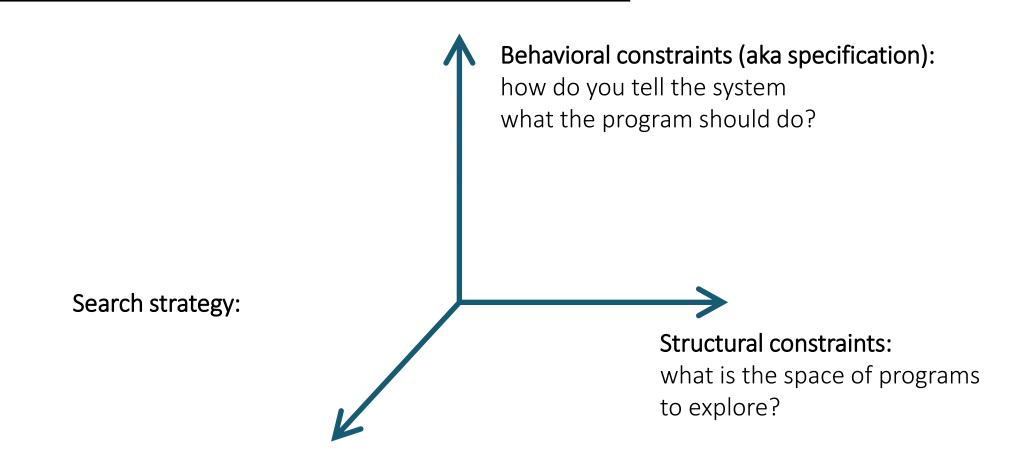


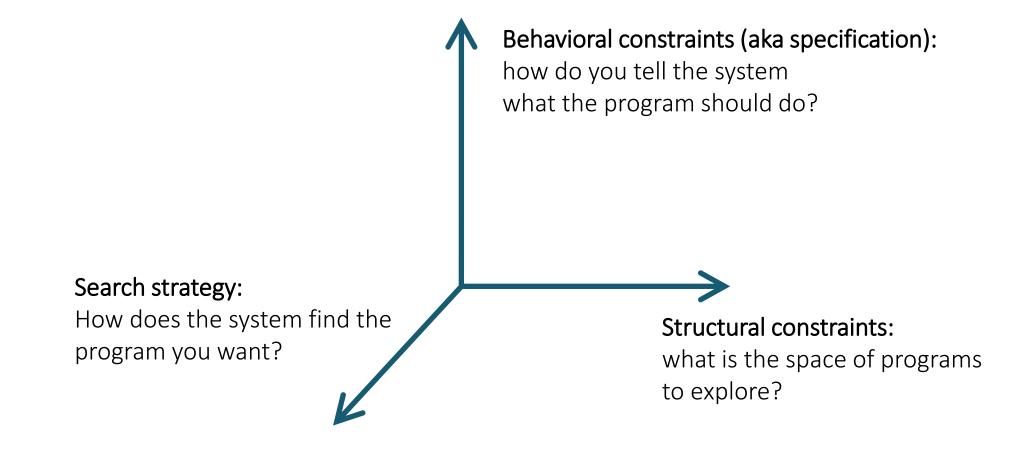
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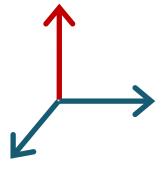








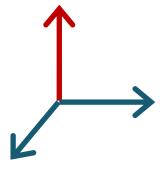
Behavioral constraints



How do you tell the system what the program should do?

- What is the input language / format?
- What is the interaction model?
- What happens when the intent is ambiguous?

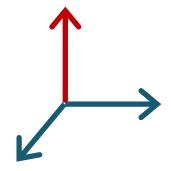
Behavioral constraints



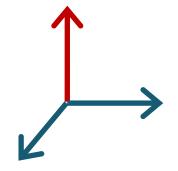
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Q: What did behavioral constraints look like in FlashFill / Sketch / Synquid / SQLizer?

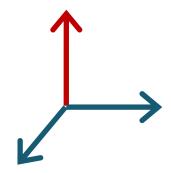


Input/output examples



Input/output examples

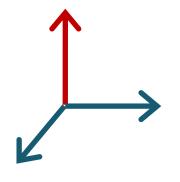
Equivalent program



Input/output examples

Equivalent program

Formal specifications (pre/post conditions, types, ...)



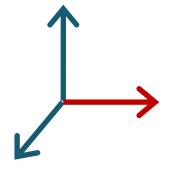
Input/output examples

Equivalent program

Formal specifications (pre/post conditions, types, ...)

Natural language

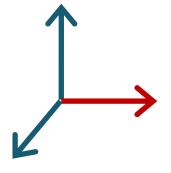
Structural constraints



What is the space of programs to explore?

- Large enough to contain interesting programs, yet small enough to exclude garbage and enable efficient search
- Built-in or user defined?
- Can we extract domain knowledge from existing code?

Structural constraints

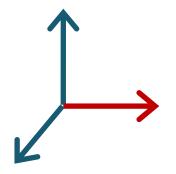


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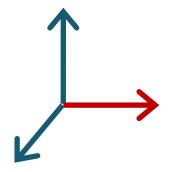
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Structural constraints: examples



Built-in DSL

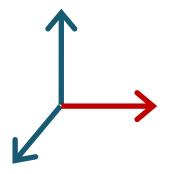
Structural constraints: examples



Built-in DSL

User-defined DSL (grammar)

Structural constraints: examples

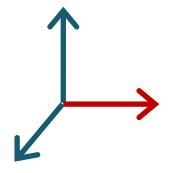


Built-in DSL

User-defined DSL (grammar)

User-provided components

Structural constraints: examples



Built-in DSL

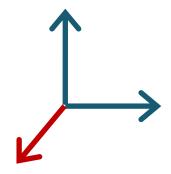
User-defined DSL (grammar)

User-provided components

Languages with synthesis constructs

• e.g. generators in Sketch

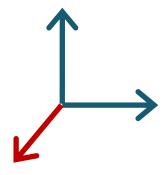
Search strategies



Synthesis is search:

• Find a program in the space defined by *structural constraints* that satisfies *behavioral constraints*

Search strategies



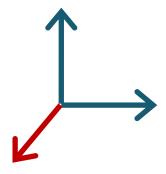
Synthesis is search:

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Challenge: the space is astronomically large

• The search algorithm is the heart of a synthesis technique

Search strategies



Synthesis is search:

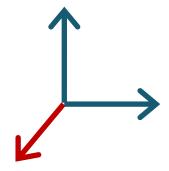
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Challenge: the space is astronomically large

• The search algorithm is the heart of a synthesis technique

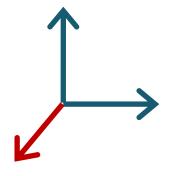
How does the system find the program you want?

- How does it know it's the program you want?
- How can it leverage structural constraints to guide the search?
- How can it leverage behavioral constraints to guide the search?



Enumerative (explicit) search

• exhaustively enumerate all programs in the language in the order of increasing size

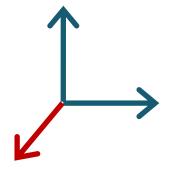


Enumerative (explicit) search

 exhaustively enumerate all programs in the language in the order of increasing size

Stochastic search

• random exploration of the search space guided by a fitness function



Enumerative (explicit) search

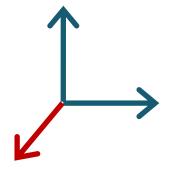
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Representation-based search

• use a data structure to represent a large set of programs



Enumerative (explicit) search

 exhaustively enumerate all programs in the language in the order of increasing size

Stochastic search

random exploration of the search space guided by a fitness function

Representation-based search

• use a data structure to represent a large set of programs

Constraint-based search

translate to constraints and use a solver

Structure of the Course

Module 1: Synthesis of Simple Programs

- Easy to decide when a program is correct
- Challenge: search in a large space

Module 2: Synthesis of Complex Programs

- Deciding when a program is correct can be hard
- Search in a large space is still a problem

Module 3: Advanced Topics

Quantitative synthesis, human aspects, applications

Example: FlashFill

specification

```
1: "Polikarpova, Nadia" → "Nadia"
```

2: "Van Damme, Jean Claude" → "Jean"

Example: FlashFill

```
program space
 (( ))
```

constant string: or substring of input: between("...", "...")

specification

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2: "Van Damme, Jean Claude" → "Jean"
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between("...", "...")



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between(" ", " ")

between(", ", " ")
```



```
Example: FlashFill
                                                        specification
                                              1: "Polikarpova, Nadia" → "Nadia"
                                              2: "Van Damme, Jean Claude" → "Jean"
  program space
 constant string:
   (( ))
                                       (())
                                                           "hello"
                                                                                "Nadia"
 or substring of input:
   between("...", "...")
                                       between(" ", " ")
                                                                      between(", ", " ")
                                                              too many
```

```
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too many

Example: Synquid

specification

```
Example: Synquid
                                                                            program
                                                        intersection = \xs . \ys .
                                                          match xs with
          specification
                                                            Nil -> xs
                                                            Cons x xt ->
                                                             match ys with
intersect :: xs:SList a →
                                                               Nil -> ys
  ys:SList a →
  \{v: SList a \mid elems v = elems xs n\}
                                                               Cons y yt ->
                                elems ys}
                                                                 if x < y
                                                                 then intersection xt ys
                                                                 else
                                                                   if y < x
                                                                   then intersection xs yt
                                                                   else Cons x (intersection xt yt)
```

Example: Synquid

intersection = \xs . \ys .

match xs with

Nil -> xs

Cons x xt ->

intersect :: xs:SList a ->

ys:SList a ->

Mil -> ys

elems ys}

How do we know this program always produces a sorted list that is the intersection?

 $\{v: SList a \mid elems v = elems xs n\}$

```
Cons y yt ->
 if x < y
 then intersection xt ys
  else
    if y < x
    then intersection xs yt
    else Cons x (intersection xt yt)
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Module 3: Advanced Topics

Synthesis with Quantitative Objectives

• Find the shortest / fastest / most likely program

Module 3: Advanced Topics

Synthesis with Quantitative Objectives

Find the shortest / fastest / most likely program

Synthesis as a Programming Tool

- How can synthesis help programmers?
- What is the right user interaction model?

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Synthesis with Quantitative Objectives

Find the shortest / fastest / most likely program

Synthesis as a Programming Tool

- How can synthesis help programmers?
- What is the right user interaction model?

Domain-Specific Synthesis

- Optimization
- CAD models
- Cryptographic schemes
- SQL / Regex

Weeks 1-2

Topic: Enumerative synthesis from examples

Paper: Alur, Radhakrishna, Udupa. <u>Scaling Enumerative Program</u> <u>Synthesis via Divide and Conquer</u>

- Review due Wednesday
- Link to PDF on the course wiki
- Submit through Google Form (link on course wiki)

Project:

- Teams due next Friday
- Submit through a Google Sheet (check email for invite and instructions)