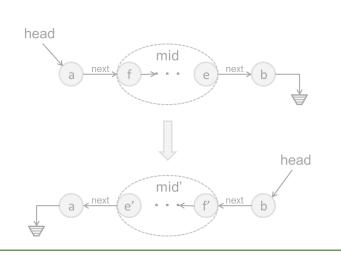
$\exists c \forall in \ Q(c, in)$

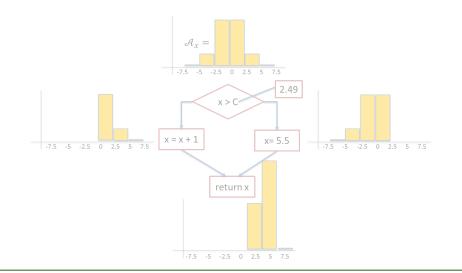
```
/* Average of x and y without using x+y (avoid overflow)*/
int avg(int x, int y) {
  int t = expr({x/2, y/2, x%2, y%2, 2 }, {PLUS, DIV});
  assert t == (x+y)/2;
  return t;
}
```

```
f_1
f_2
f_3
f_3
f_4
f_5
f_7
```

```
{
    s = n.succ;
    p = n.pred;
    p.succ = s;
    s.pred = p;
}
```

Module II: Synthesizing Complex Programs





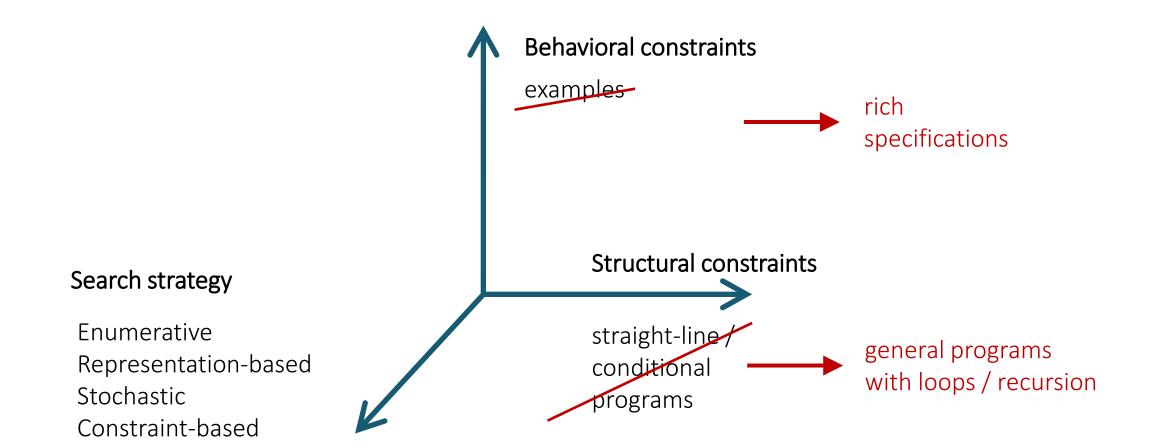


Sk[c](in)

Lecture 9 Specifications

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Module I vs Module II



Examples of rich specifications

Reference implementation

Assertions

Pre- and post-condition

Fancy types

Reference Implementation

Easy to compute the result, but hard to compute it efficiently or under structural constraints

```
bit[W] AES_round (bit[W] in, bit[W] rkey)
{
    ... // Transcribe NIST standard
}
bit[W] AES_round_sk (bit[W] in, bit[W] rkey) implements AES_round
{
    ... // Sketch for table lookup
}
```

Assertions

Hard to compute the result, but easy to check its properties

```
split_seconds (int totsec) {
  int h := ?;
  int m := ?;
  int s := ?;
  assert totsec == h*3600 + m*60 + s;
  assert 0 <= h && 0 <= m < 60 && 0 <= s < 60;
}</pre>
```

Pre-/post-conditions

Hard to compute the result; need correctness guarantees

```
sort (int[] in, int n) returns (int[] out) requires n \ge 0 ensures \forall i \ j. \ 0 \le i < j < n \Rightarrow out[i] \le out[j] \forall i. \ 0 \le i < n \Rightarrow \exists j. \ 0 \le j < n \land in[i] = out[j] {
```

Refinement types

Same as pre-/post-conditions but logic goes inside the types

```
binary search tree
                                          red nodes have
data RBT a where
                                          black children
  Empty :: RBT a
  Node :: x: a ->
    black: Bool ->
                                  !black ==> isBlack
    left: { RBT {a
                      || v < x|
    right: { RBT \{a \mid x < v\}
                                  (!black ==> isBlack
                                                         v) &&
                 (blackHeight v == blackHeight left) €
    RBT a
                                                                        same number of
                                                                        black nodes on
insert :: x: a -> t: RBT a -> {RBT a | elems _v == elems t + [x]}
                                                                        every path to leaves
insert = ?
```

Why go beyond examples?

Might need too many

- Example: Myth needs 12 for insert_sorted, 24 for list_n_th
- Examples contain too little information
- Successful tools use domain-specific ranking

Output difficult to construct

- Example: AES cypher, RBT
- Examples also contain too much information (concrete outputs)

Need strong guarantees

• Example: AES cypher

Reasoning about non-functional properties

• Example: security protocols

Why is this hard?

```
gcd (int a, int b) returns (int c)
                                                                infinitely many inputs
  requires a > 0 \land b > 0
                                                               cannot validate by testing
  ensures a \% c = 0 \land b \% c = 0
             \forall d \cdot c < d \Rightarrow a \% d \neq 0 \lor b \% d \neq 0
  int x , y := a, b;
                                                            infinitely many paths!
  while (x != y) {
                                                            hard to generate constraints
     if (x > y) x := ?;
     else y := ?;
}}
```

Why is this hard?

Synthesis from examples



validation was easy!

Synthesis from specifications

SEE IF YOU CAN FIND ANY KLINGON FRUIT!

validation is hard! (and search is still hard)

Module II

