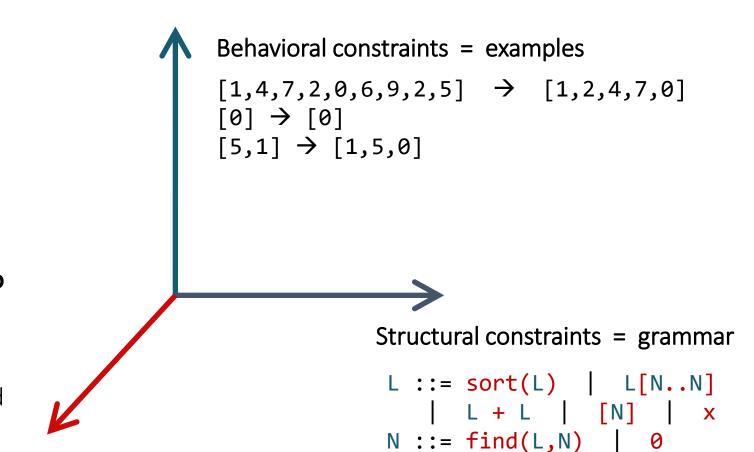
# Lecture 6 Stochastic Search

Nadia Polikarpova

# The problem statement



#### Search strategy?

Enumerative

Stochastic

Representation-based

Constraint-based

# Stochastic search in synthesis

Weimer, Nguyen, Le Goues, Forrest. *Automatically Finding Patches Using Genetic Programming*. ICSE'09

Gissurarson, Applis, Panichella, van Deursen, Sands. *PropR: Property-Based Automatic Program Repair.* ICSE'22

Schkufza, Sharma, Aiken: *Stochastic superoptimization*. ASPLOS 2013

Shi, Steinhardt, Liang: FrAngel: Component-Based Synthesis with Control Structures. POPL'19

# Stochastic search in synthesis

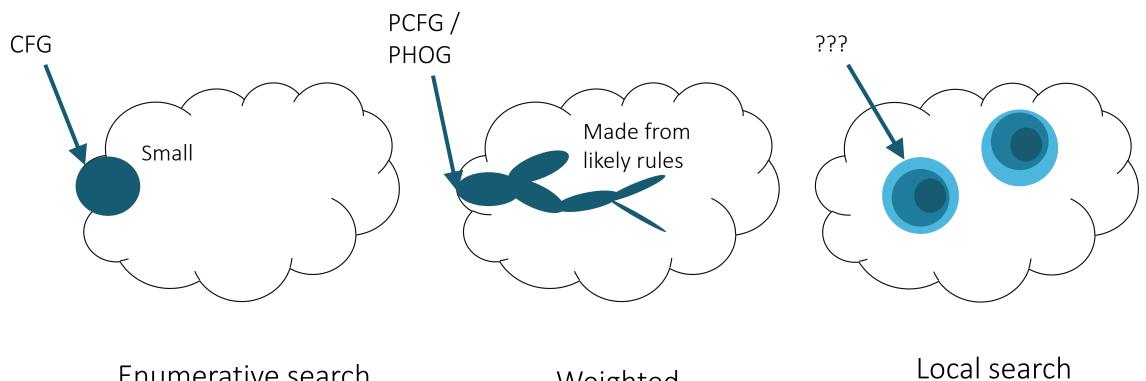
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# Search space



Enumerative search

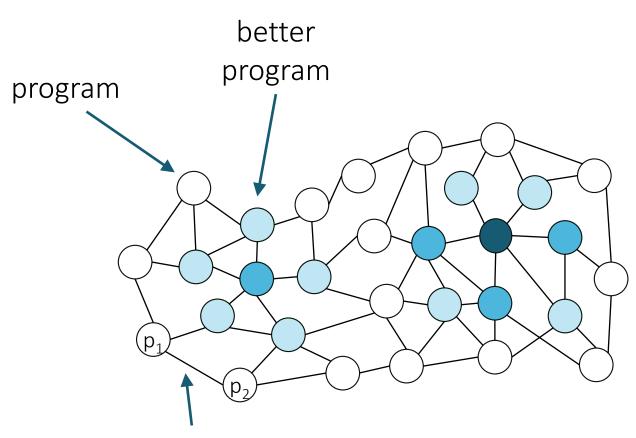
Weighted enumerative search

#### Naïve local search

To find the best program:

```
p := random()
while (true) {
   p' := mutate(p);
   if (cost(p') < cost(p))
      p := p';
}</pre>
```

Will never get to  $\bigcirc$  from  $p_1!$ 



can generate p<sub>2</sub> from p<sub>1</sub> (and vice versa) via mutation

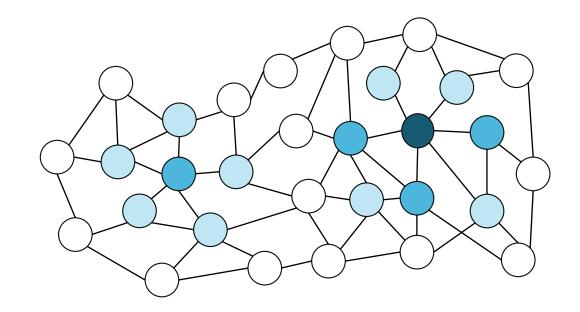
## MCMC sampling

Avoid getting stuck in local minima:

```
p := random()
while (true) {
   p' := mutate(p);
   if (random(A(p -> p'))
      p := p';
}
```

#### where

- if p is better than p:  $A(p \rightarrow p') = 1$
- otherswise:  $A(p \rightarrow p')$  decreases with difference in cost between p' and p



# MCMC sampling

Metropolis algorithm:

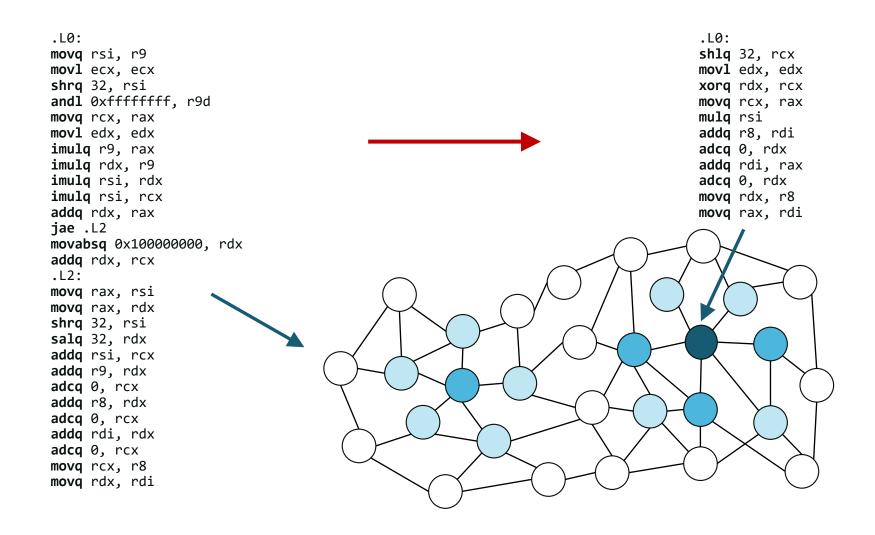
$$A(p \to p') = \min(1, e^{-\beta(C(p') - C(p))})$$

The theory of Markov chains tells us that in the limit we will be sampling with the probability proportional to

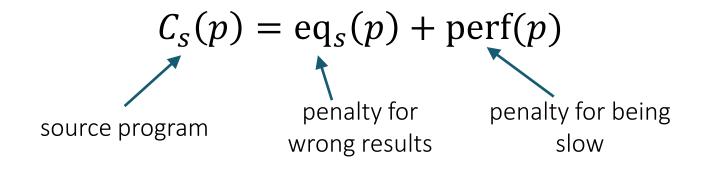
$$e^{-\beta * C(p)}$$

# MCMC for superoptimization

[Schkufza, Sharma, Aiken '13]



## Cost function



when  $eq_s(p) = 0$ , use a symbolic validator

### Cost function

$$C_S(p) = \operatorname{eq}_S(p) + \operatorname{perf}(p)$$
source program

penalty for penalty for being wrong results slow

$$perf(p) = \sum_{i \in instr(p)} latency(i)$$

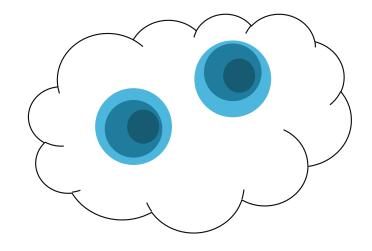
#### Local search: discussion

#### Strengths:

• can explore program spaces with no a-priori bias

#### Limitations?

- only applicable when there is a cost function that faithfully approximates correctness
- Counterexample: round to next power of two
  - 0011 -> 0100
  - 0101 -> 1000
  - 0111 -> 1000



# Stochastic search in synthesis

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• Similar but for program repair, uses genetic programming

Schkufza, Sharma, Aiken: *Stochastic superoptimization*. ASPLOS 2013

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- Samples from a grammar with bias towards partial solutions
- I assume they use stochastic just for ease of sampling

#### Next



#### Search strategy?

Enumerative Stochastic Representation-based Constraint-based

Structural constraints = grammar

Behavioral constraints = examples

 $[0] \rightarrow [0]$ 

 $[5,1] \rightarrow [1,5,0]$ 

 $[1,4,7,2,0,6,9,2,5] \rightarrow [1,2,4,7,0]$