

# HOW LONG WILL MY SYNTHESIZER RUN FOR?

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Loris D'Antoni

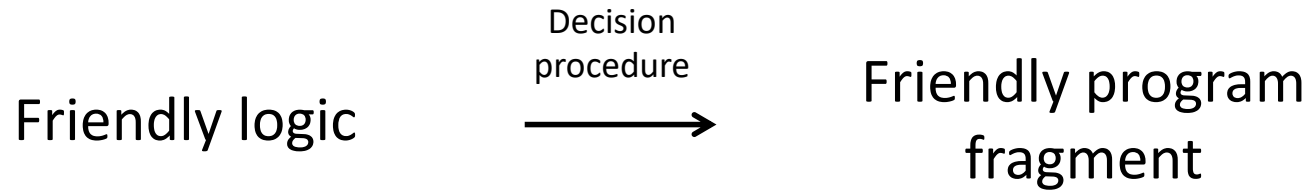
*University of Wisconsin Madison*



Joint work with Aws Albarghouthi and Samuel Drews



# The good old days of synthesis



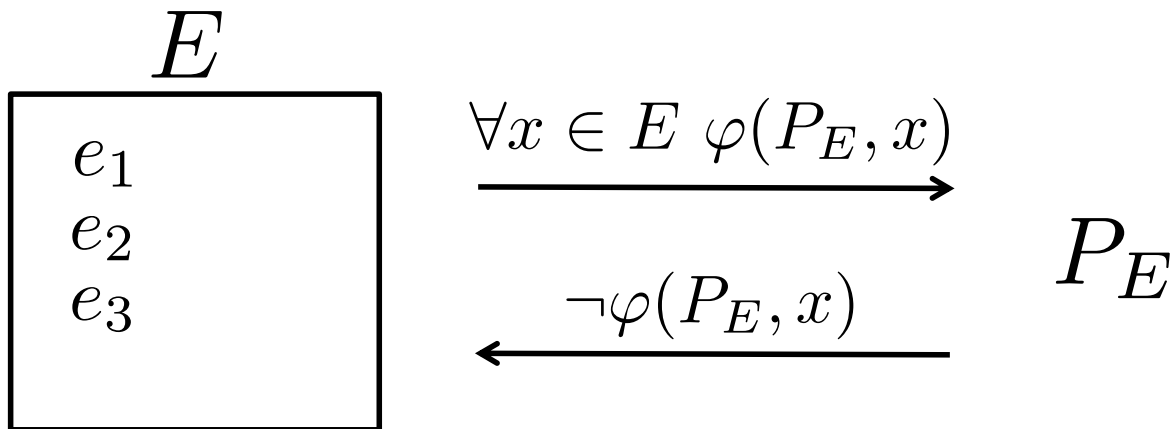
# Synthesis: Dreams $\Rightarrow$ Programs

ZOHAR MANNA AND RICHARD WALDINGER

$$\exists P \forall x \varphi(P, x)$$

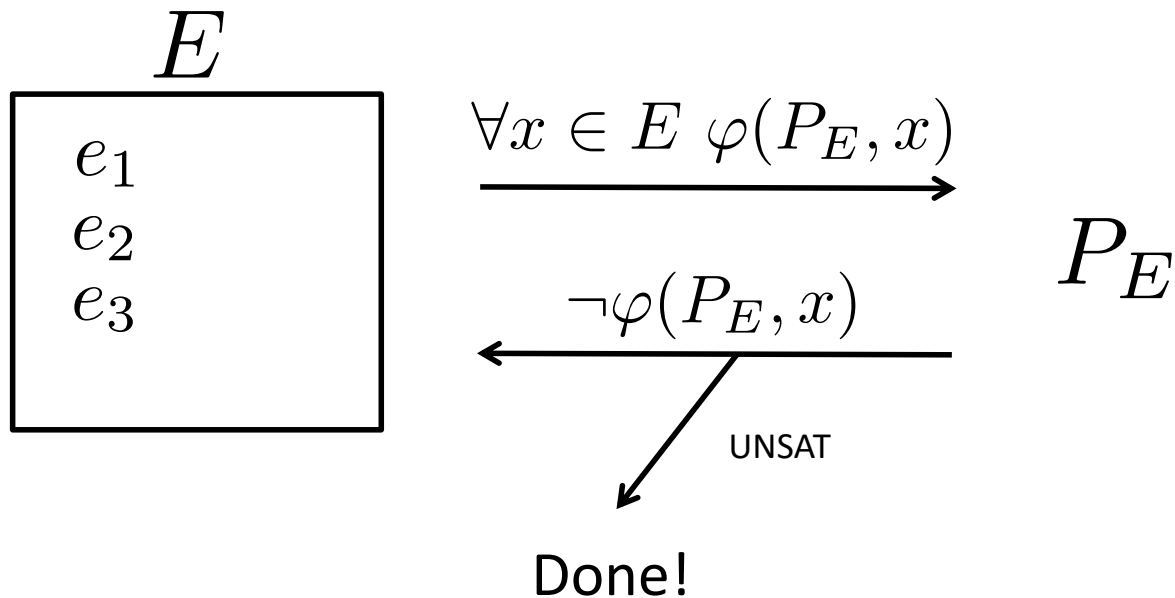
Inductive synthesis

$$\exists P \in \mathcal{S} \ \forall x \ \varphi(P, x)$$



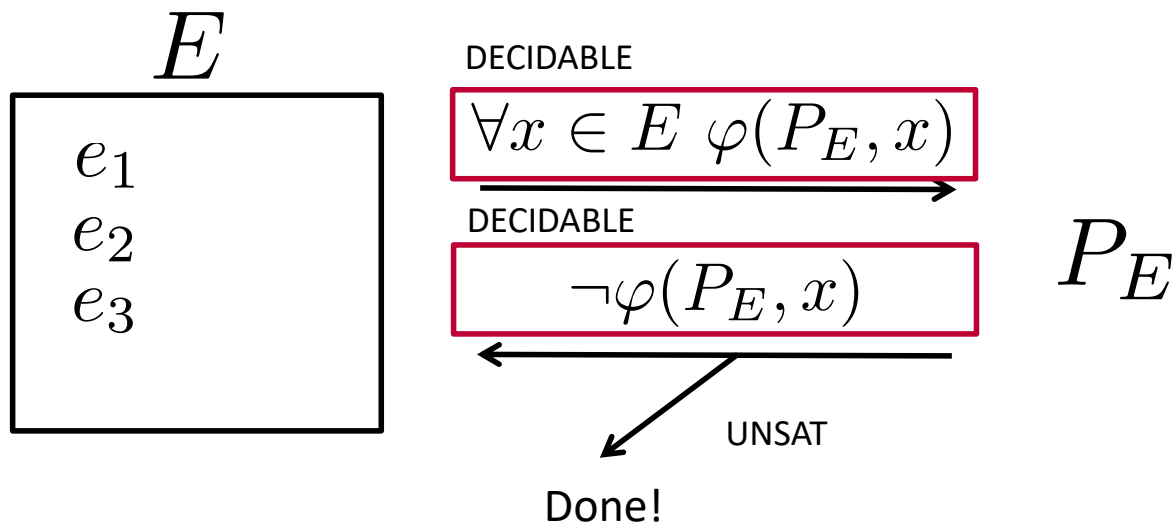
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Inductive synthesis

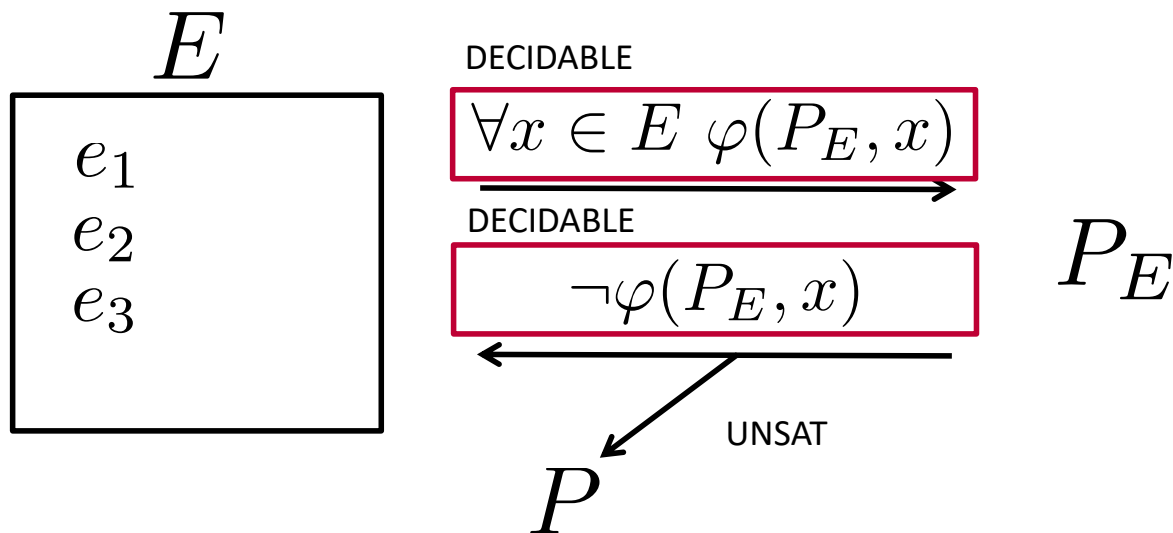
$$\exists P \in \mathcal{S} \ \forall x \ \varphi(P, x)$$



Does this algorithm always terminate?

Inductive synthesis

$$\exists P \in \mathcal{S} \ \forall x \ \varphi(P, x)$$



Does this algorithm always terminate  
if a solution to the synthesis problem exists?

In general NO

$$\exists P \in \mathcal{S} \ \forall x \ \varphi(P, x)$$

$$\mathcal{S} = \{P(x)=c \mid c \geq 0\}$$

$$\exists P \in \mathcal{S} \ \forall x \ P(x) > x$$

No matter what the example set is, there is a program in the search space consistent with it

$$E = \{1, 4, 7, 8\}$$

$$P_E(x) = 8$$



# Synthesis: Dreams $\Rightarrow$ Programs

ZOHAR MANNA AND RICHARD WALDINGER

Will we keep dreaming forever?

Will my synthesizer ever terminate?

How long will my synthesizer run for?

# This talk

Problem: How long will my synthesizer run for?

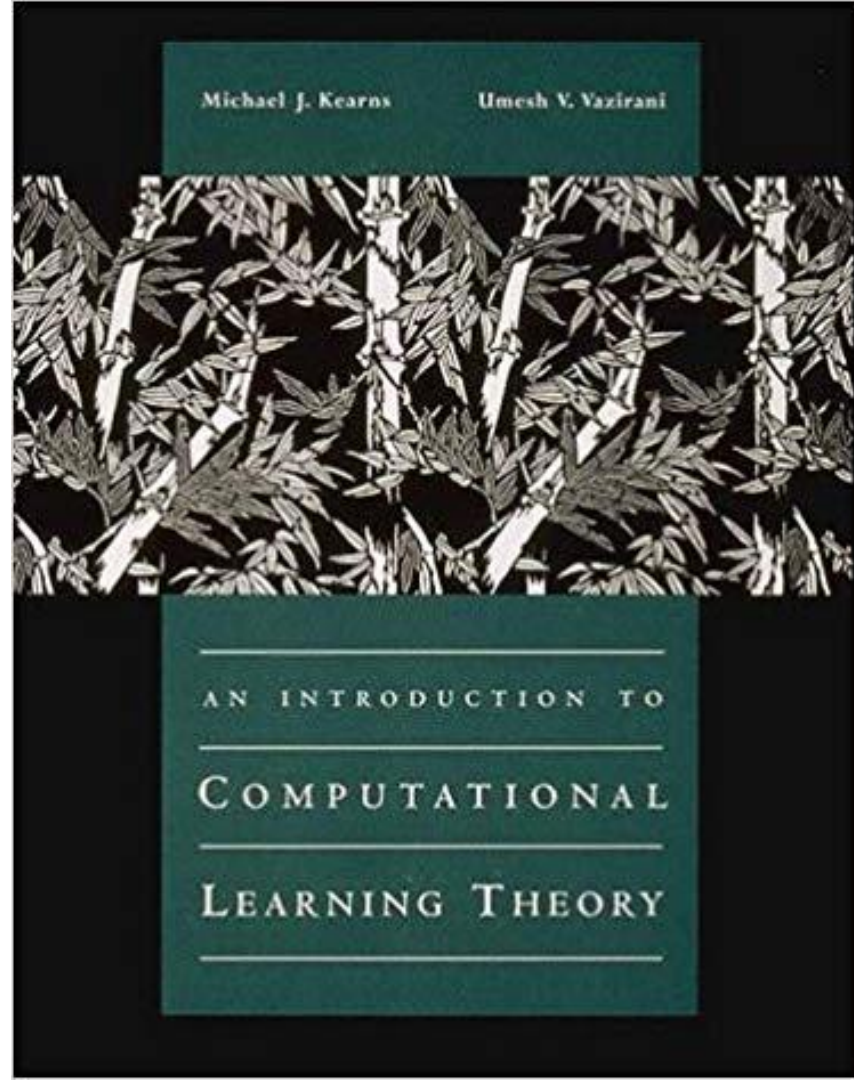
## **A connection to inductive learning**

DIGITS: Distribution Guided Inductive Synthesis

IS THIS REALLY A NEW QUESTION?

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When can functions be learned?



# Computational learning theory in a nutshell

## Probably Approximately Correct (PAC) learning

Certain types of functions

Can be learned approximately

Using polynomially many IO examples

# What is so different about program synthesis?

## Inductive learning

Target function I'm trying to learn

$$f(x) = 2x+2$$

Labeled input-output examples

(1,4), (2,6)

## Inductive synthesis

Target relation

$$f(x) > x$$

Input examples but no outputs!

(1,??), (2,??)

# What is so different about program synthesis?

## Inductive learning

Certain types of functions

Can be learned approximately

Using polynomially many IO examples

## Inductive synthesis

Search space is a mess

We can't be sort of right

We don't have labeled examples

## How long will my synthesizer run for?

To answer this question, we need to address these problems



Can we develop a *computational synthesis theory*?

# This talk

Problem: How long will my synthesizer run for?

A connection to inductive learning

**DIGITS: Distribution Guided Inductive Synthesis**



“The machine learning algorithm  
wants to know if we’d like a  
dozen wireless mice to feed the  
Python book we just bought.”

# PROGRAM REPAIR UNDER UNCERTAINTY

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REPAIRING PROGRAMS WITH PROBABILISTIC INPUTS

WITH: SAMUEL DREWS, AWS ALBARGHOUI [OOPSLA17, CAV17, CAV19]

Let's focus on a concrete problem

# Everyone can now use machine learning

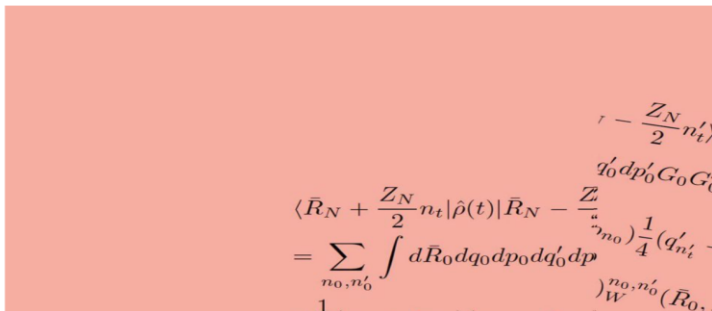


Machine learning  
library



Decision tree  
SVM  
DNN

# Who do you blame when an algorithm gets you fired?



## TheUpshot

HIDDEN BIAS

## When Algorithms Discriminate



Claire Cain Miller @clairecm JULY 9, 2015

SundayReview | OPINION

## Artificial Intelligence's White Guy Problem

By KATE CRAWFORD JUNE 25, 2016



Bianca Bagnarelli

# Algorithmic fairness for programs (example)

$$\{v \sim \mathcal{M}\}$$



**Population model:**  
probability distribution of inputs

$$h \leftarrow \mathcal{D}(v)$$



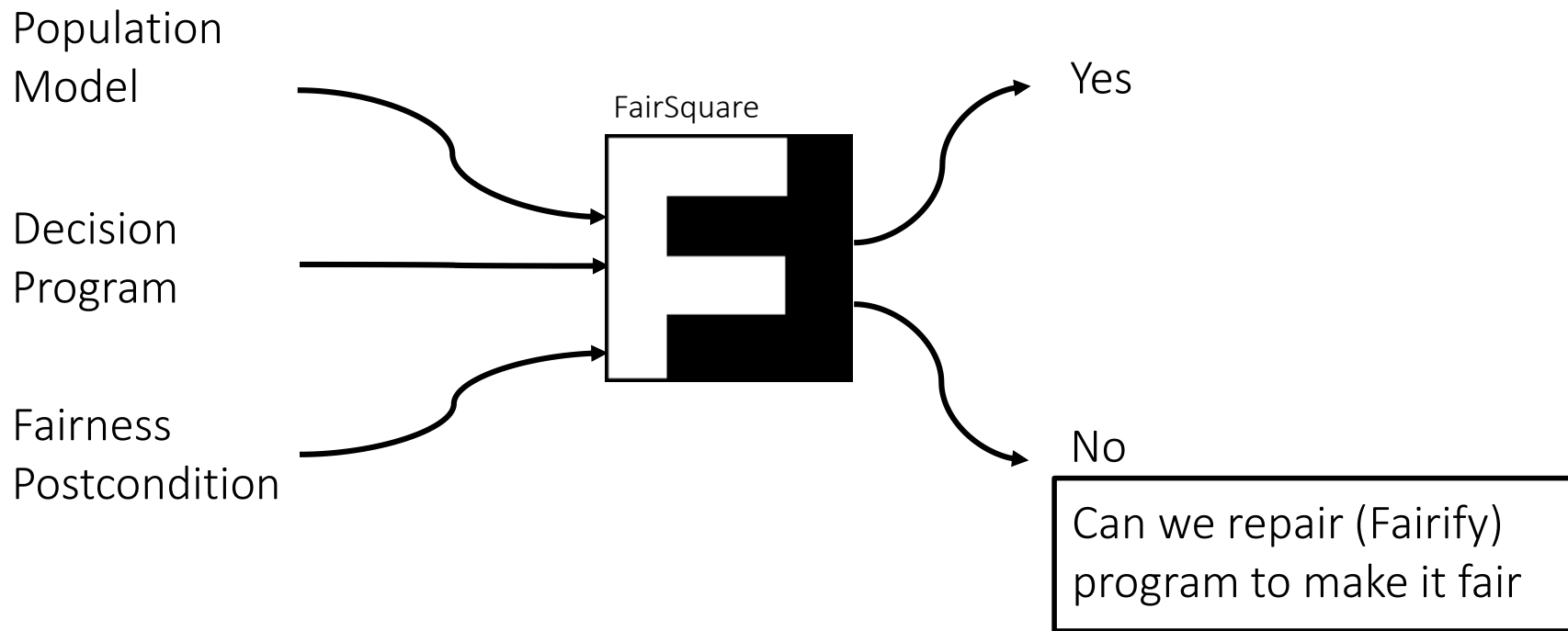
**Decision making program:**  
returns true or false

$$\left\{ \frac{\Pr[h \mid v_s]}{\Pr[h \mid \neg v_s]} > 1 - \epsilon \right\}$$



**Probabilistic postcondition:**  
Probability of hiring given that candidate is woman should be close to probability of hiring given that is man

# Verifying algorithmic fairness [OOPSLA17]





# Probabilistic program repair: definition

Input

$$\{v \sim \mathcal{M}\}$$

$$h \leftarrow \mathcal{D}(v)$$

$$\{post\}$$

The postcondition **does not hold**

Output

Program  $\mathcal{D}'$  in repair model  $RM(\mathcal{D})$

$$\{v \sim \mathcal{M}\}$$

$$h \leftarrow \mathcal{D}'(v)$$

$$\{post\}$$

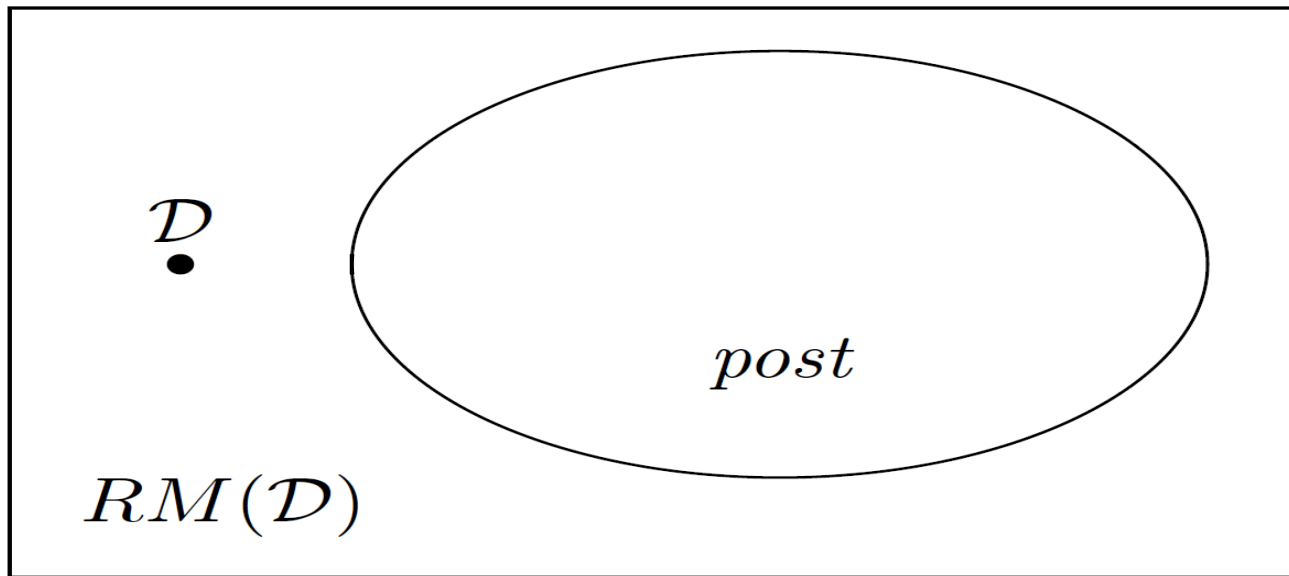
The postcondition **holds**

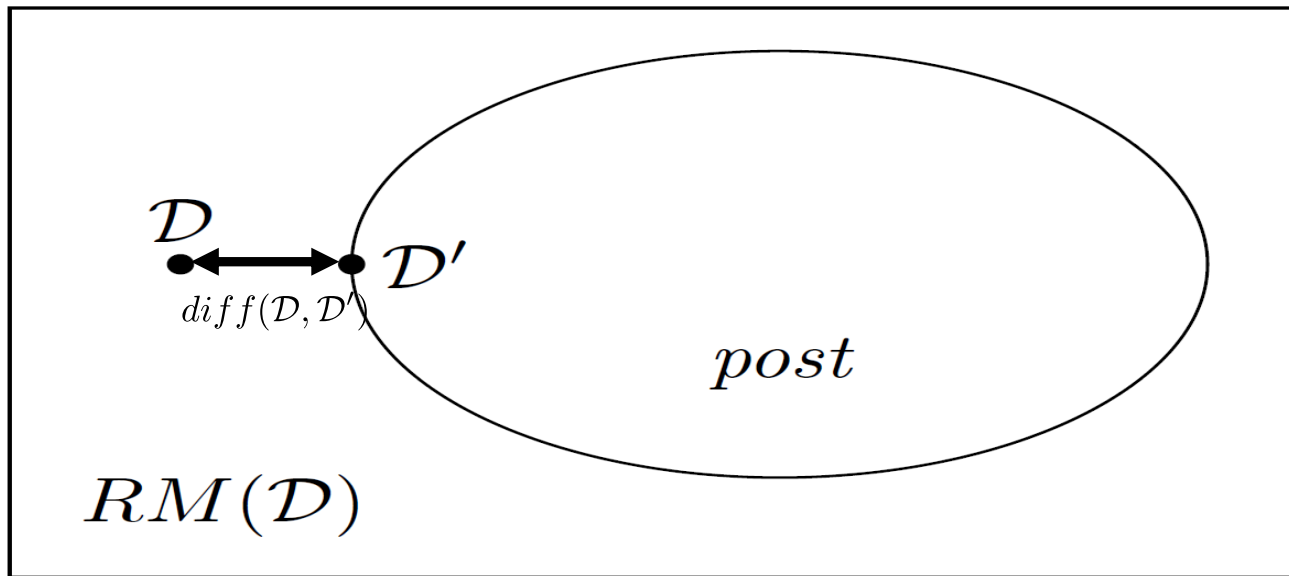
Difference between  $\mathcal{D}$  and  $\mathcal{D}'$  is minimal

$$diff(\mathcal{D}, \mathcal{D}') = Pr(\mathcal{D}(v) \neq \mathcal{D}'(v) \mid v \sim \mathcal{M})$$

$\mathcal{D}$   
•

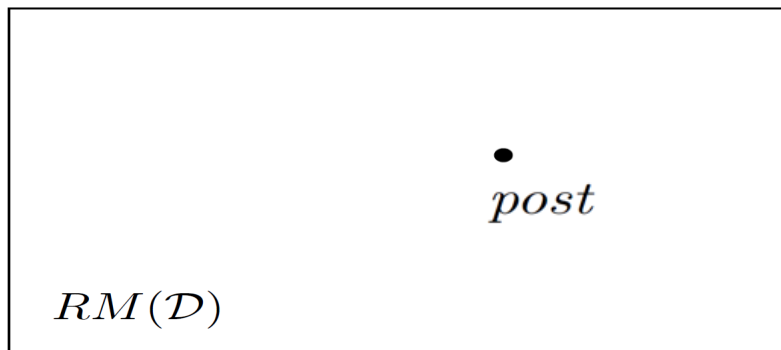
$RM(\mathcal{D})$



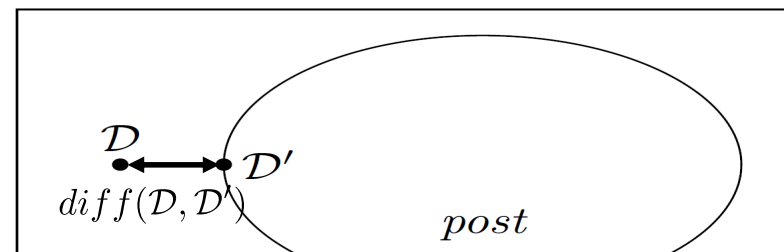


# What is so different about program synthesis? [reprised]

## Learning



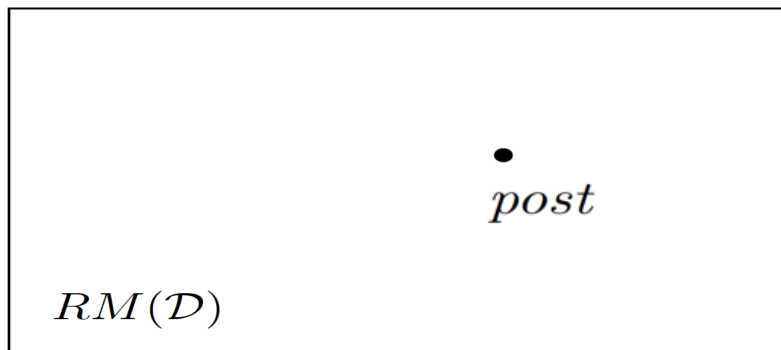
## Synthesis



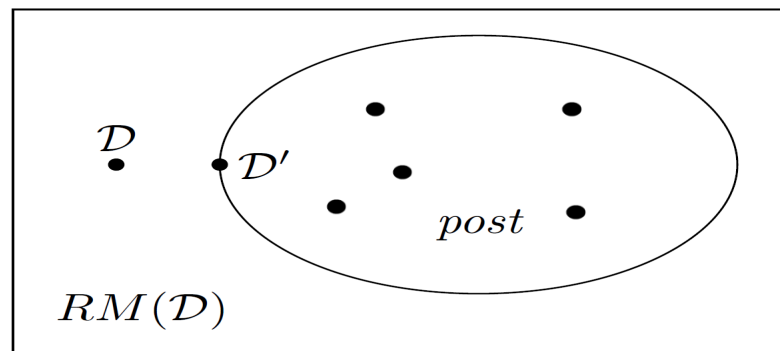
**Challenge:** no ground truth.  
What should the output of  $\mathcal{D}'$   
be on each input?

# How do we get our input/output examples?

Learning



Synthesis



Can we approximately solve the synthesis problem using boundedly many examples?

# DISTRIBUTION GUIDED INDUCTIVE SYNTHESIS

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# Distribution-Guided InducTive Synthesis (DIGITS)

I'm going to present a stupid algorithm

Show that it actually works

Inspired by the following

It works -> It ain't stupid

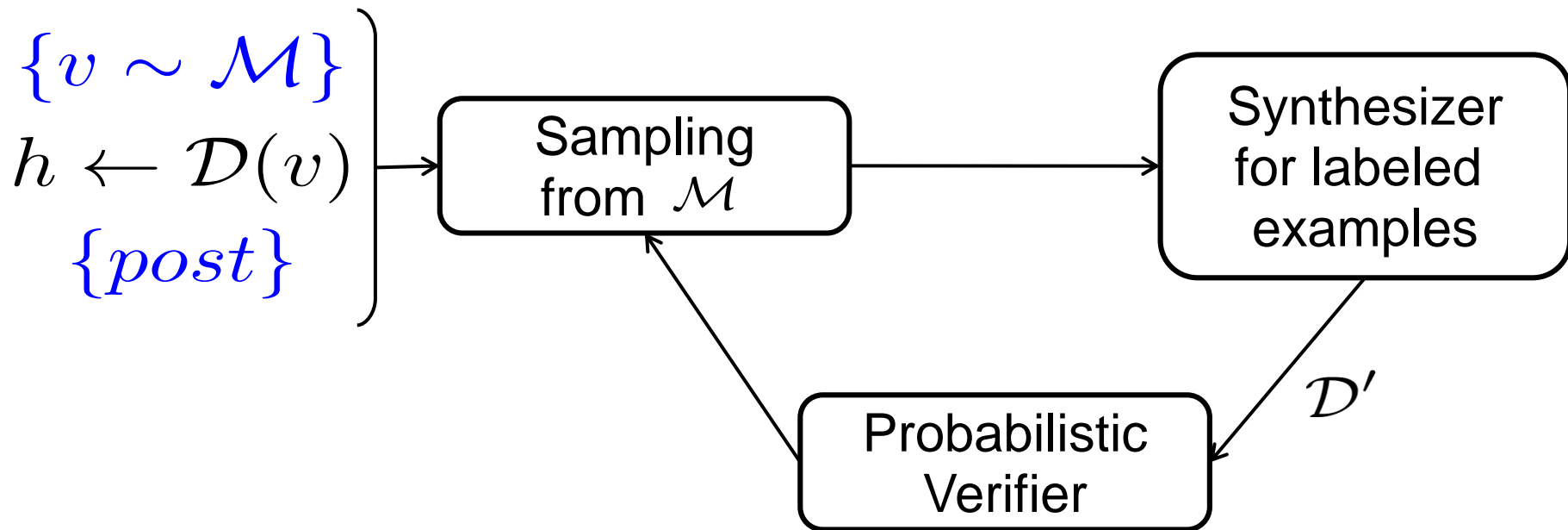
I prove the algorithm isn't that stupid and terminates using few examples [CAV17]

**Tuesday:** Sam will show you that, not only it's not that stupid, it's actually pretty smart!

Great complexity bound [Synthesis session at CAV19]



# DIGITS Components



# DIGITS Distribution directed inductive synthesis

Sample  $n$  inputs from  $\mathcal{M}$       Synthesize one program per labeling  
Check the postcondition

inp <sub>1</sub>	inp <sub>2</sub>	inp <sub>3</sub>	...	inp <sub>n</sub>
T	T	T	...	T
T	T	T	...	F
T	T	F	...	T
T	T	F	...	F
T	F	T	...	T
..	..	..	...	...

$\{v \sim \mathcal{M}\} \mathcal{D}_1 \{post\}$

~~$\{v \sim \mathcal{M}\} \mathcal{D}_2 \{post\}$~~

UNSAT

$\{v \sim \mathcal{M}\} \mathcal{D}_4 \{post\}$

~~$\{v \sim \mathcal{M}\} \mathcal{D}_5 \{post\}$~~

Output  $\mathcal{D}_i$  that  
with minimal  
 $diff(\mathcal{D}, \mathcal{D}_i)$

Synthesis using sketching  
Verification and difference using FairSquare

*$\mathcal{D}$*

```
fun hire(min,urank)
    dec = 1 <= urank <= 10
return dec
```

$\mathcal{D}$

```
fun hire(min,urank)
    dec = 1 <= urank <= 10
return dec
```

$RM(\mathcal{D})$

```
fun hireRep(min,urank)
    dec =  $\bullet_1$  <= urank <=  $\bullet_2$ 
return dec
```

```
fun hireRep(min,urank)
    dec = ●1 <= urank <= ●2
    return dec
```

$$\varphi = ((H_1 \leq urank \leq H_2) == dec)$$

$$((H_1 \leq 12 \leq H_2) == \textit{true})$$

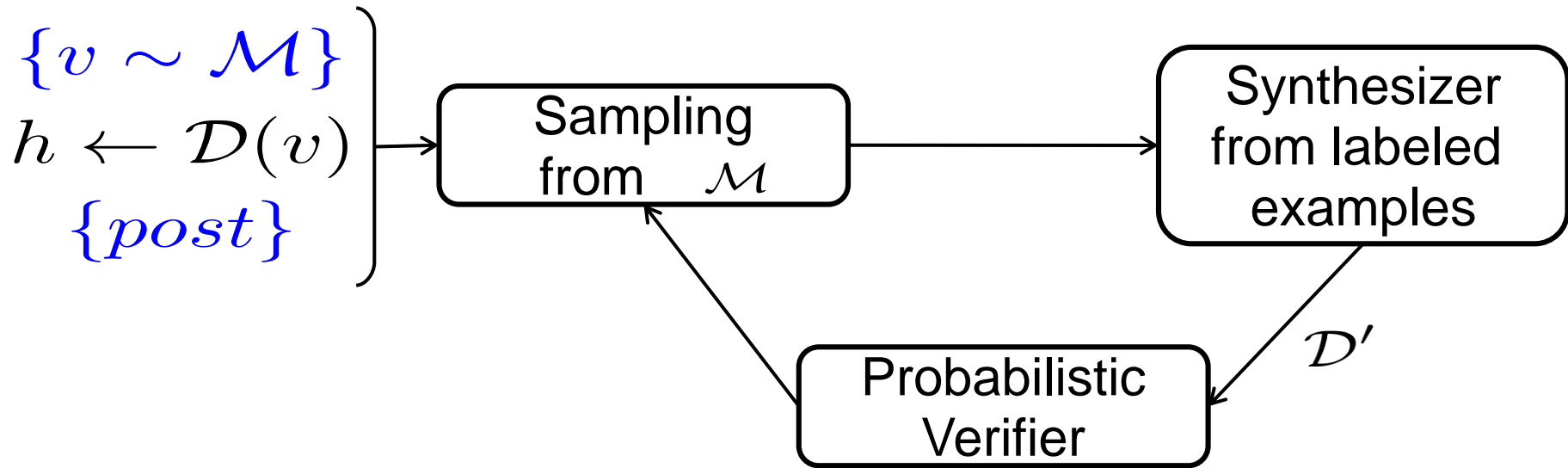
$$((H_1 \leq 17 \leq H_2) == \textit{false})$$

$((H_1 \leq 12 \leq H_2) == \text{true})$

$((H_1 \leq 17 \leq H_2) == \text{false})$

$H_1 = 10, H_2 = 15$

```
fun hireRep1(min,urank)
    dec = 10 <= urank <= 15
    return dec
```



How long will my synthesizer run for?

How many samples to find a solution?



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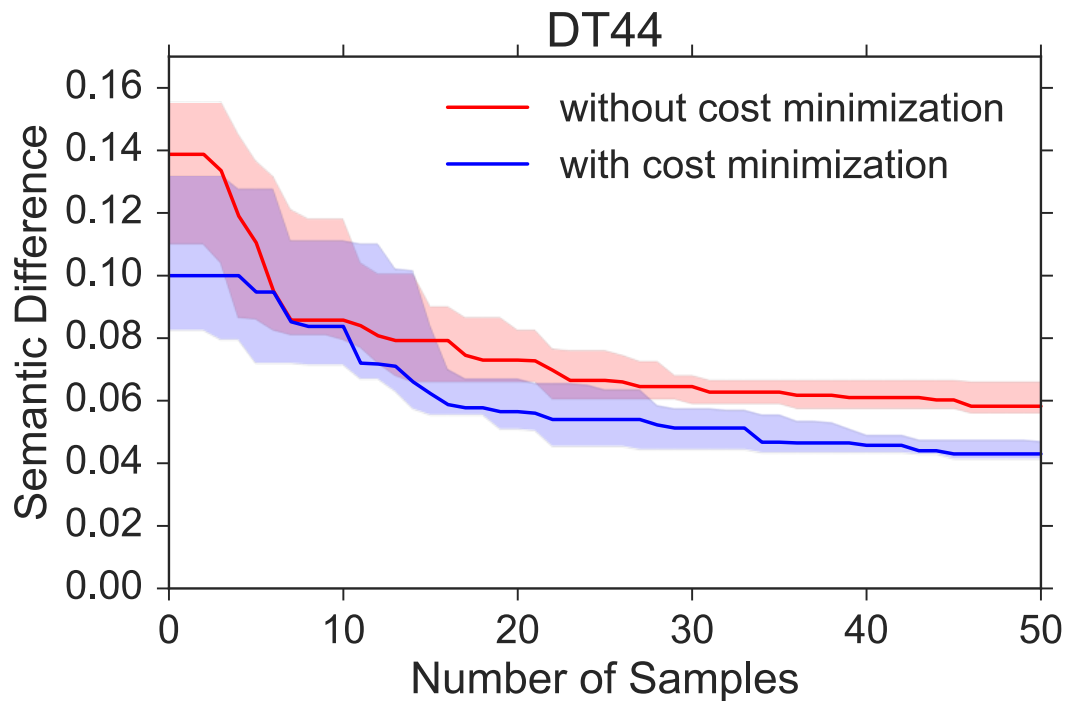
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Great complexity bound [Synthesis session at CAV19]

# Convergence on repairing a decision tree



# Distribution-Guided InducTive Synthesis (DIGITS)

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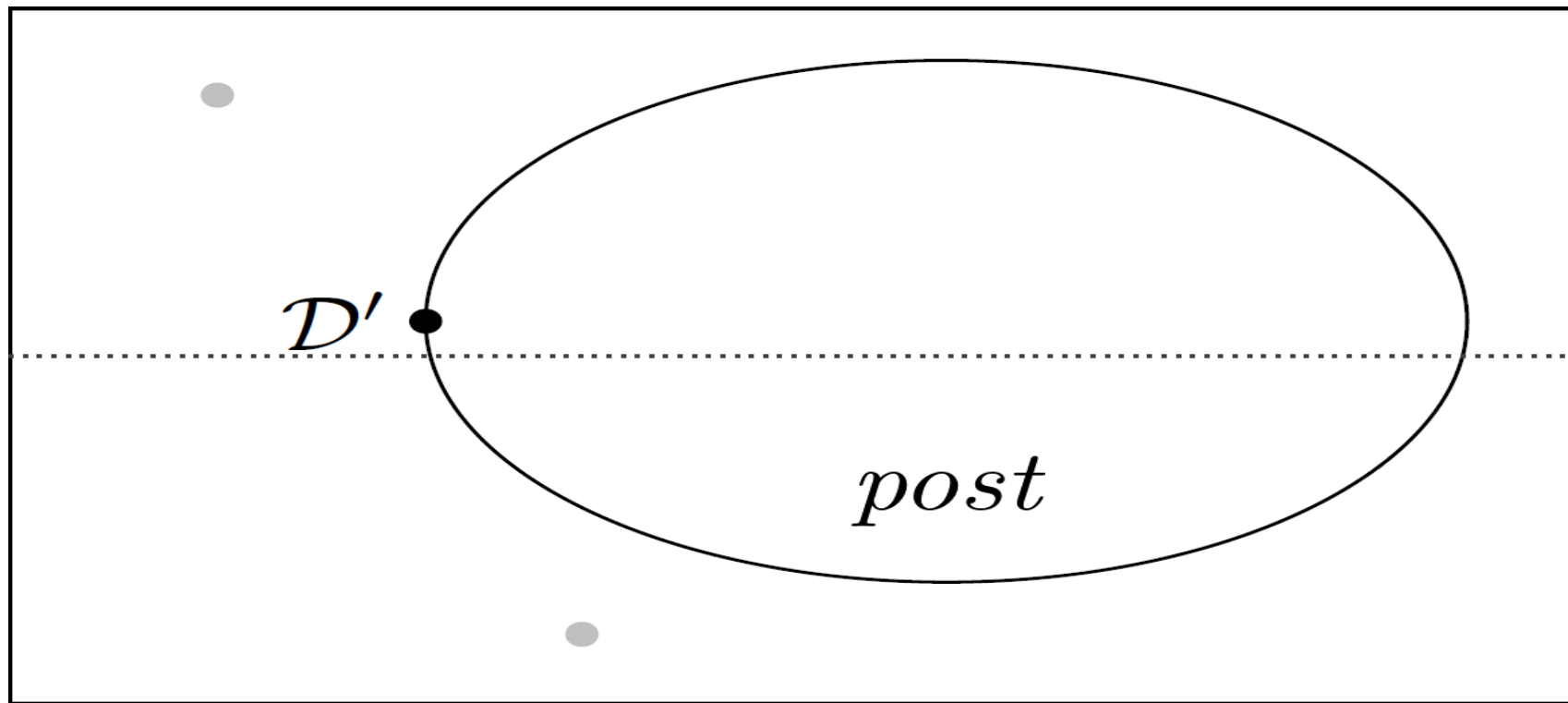
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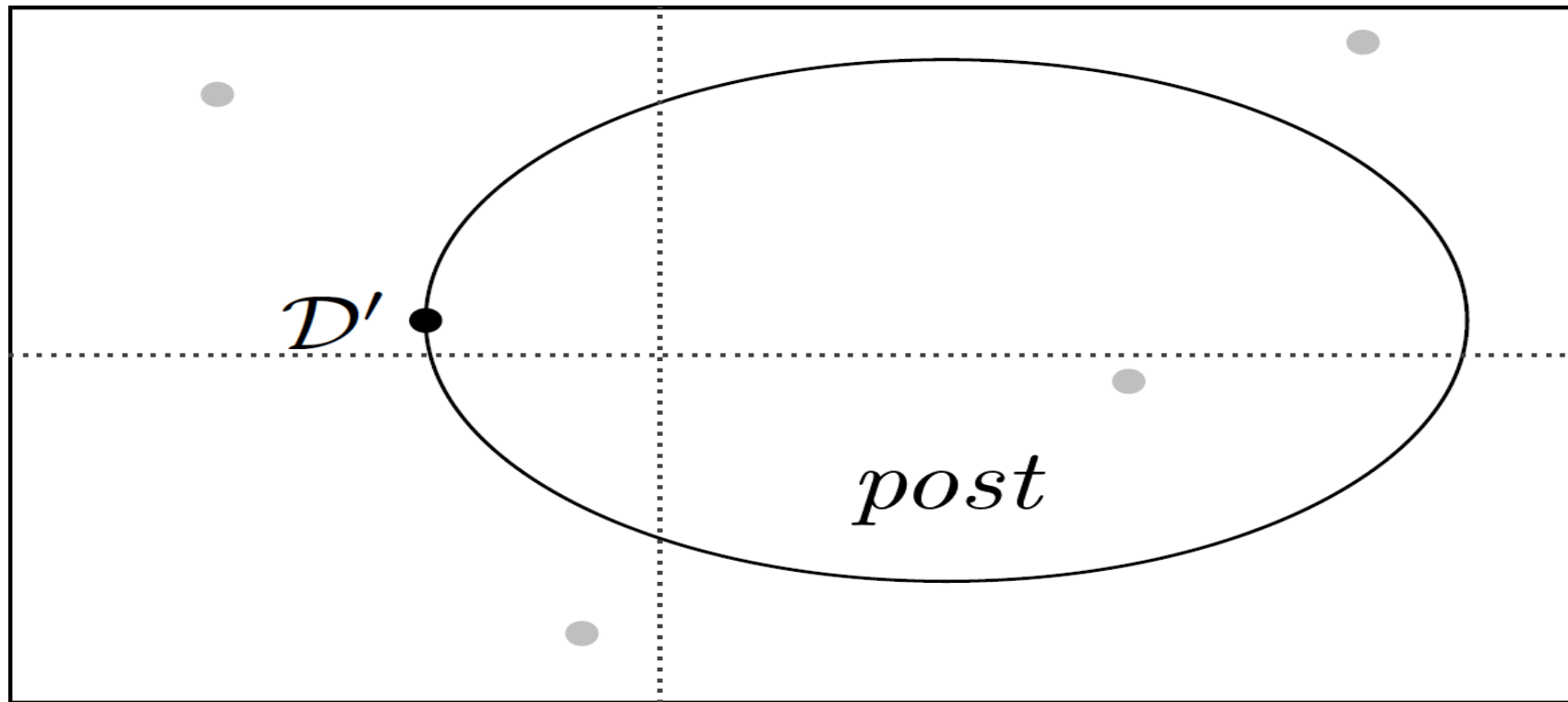
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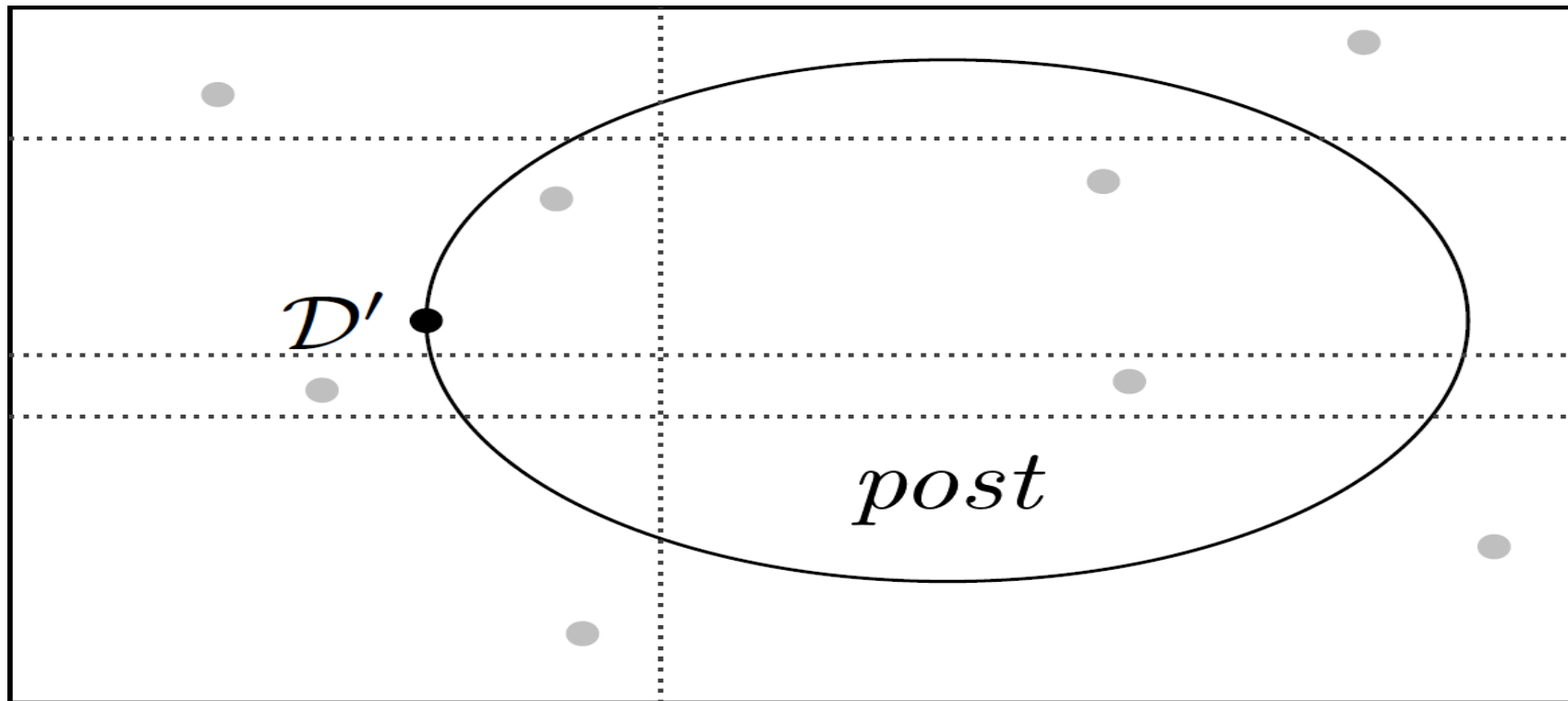
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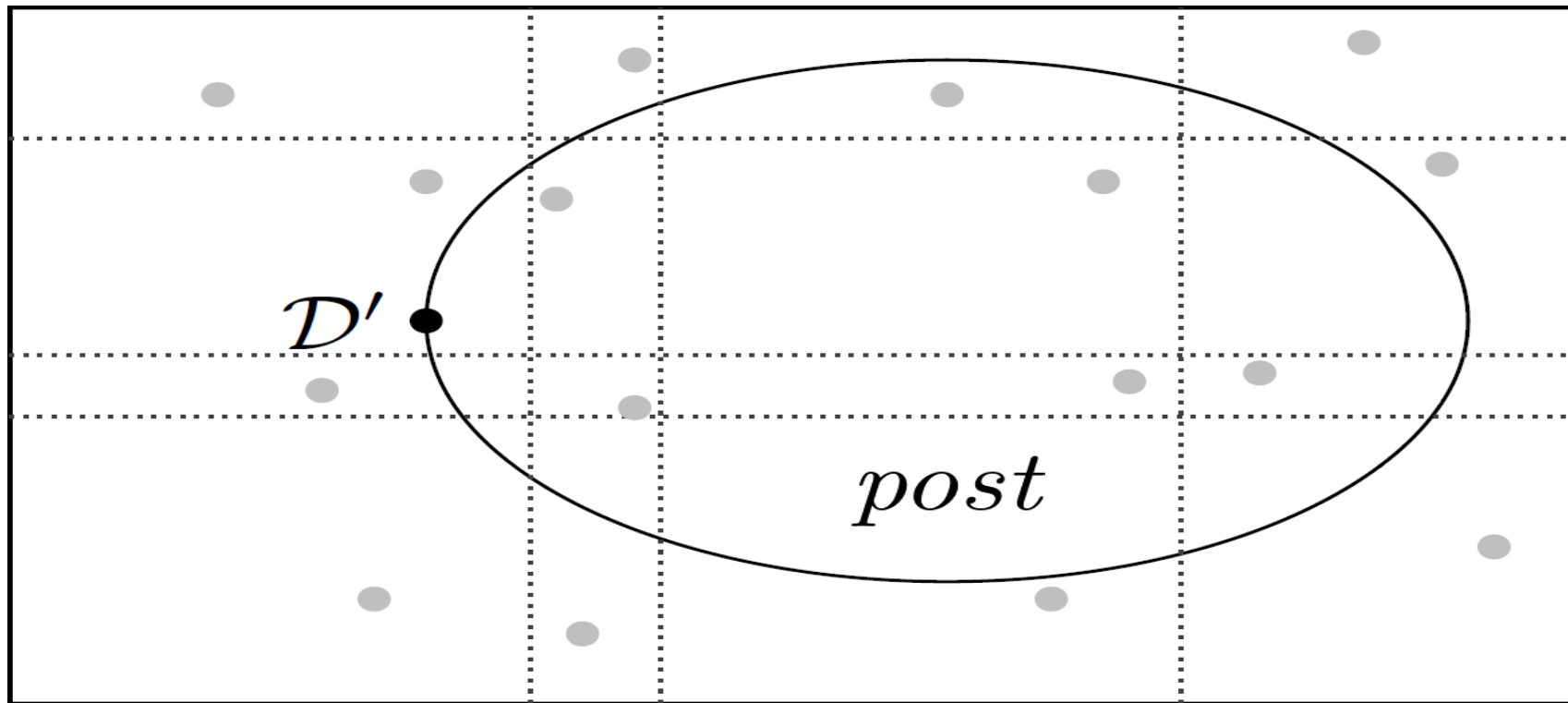
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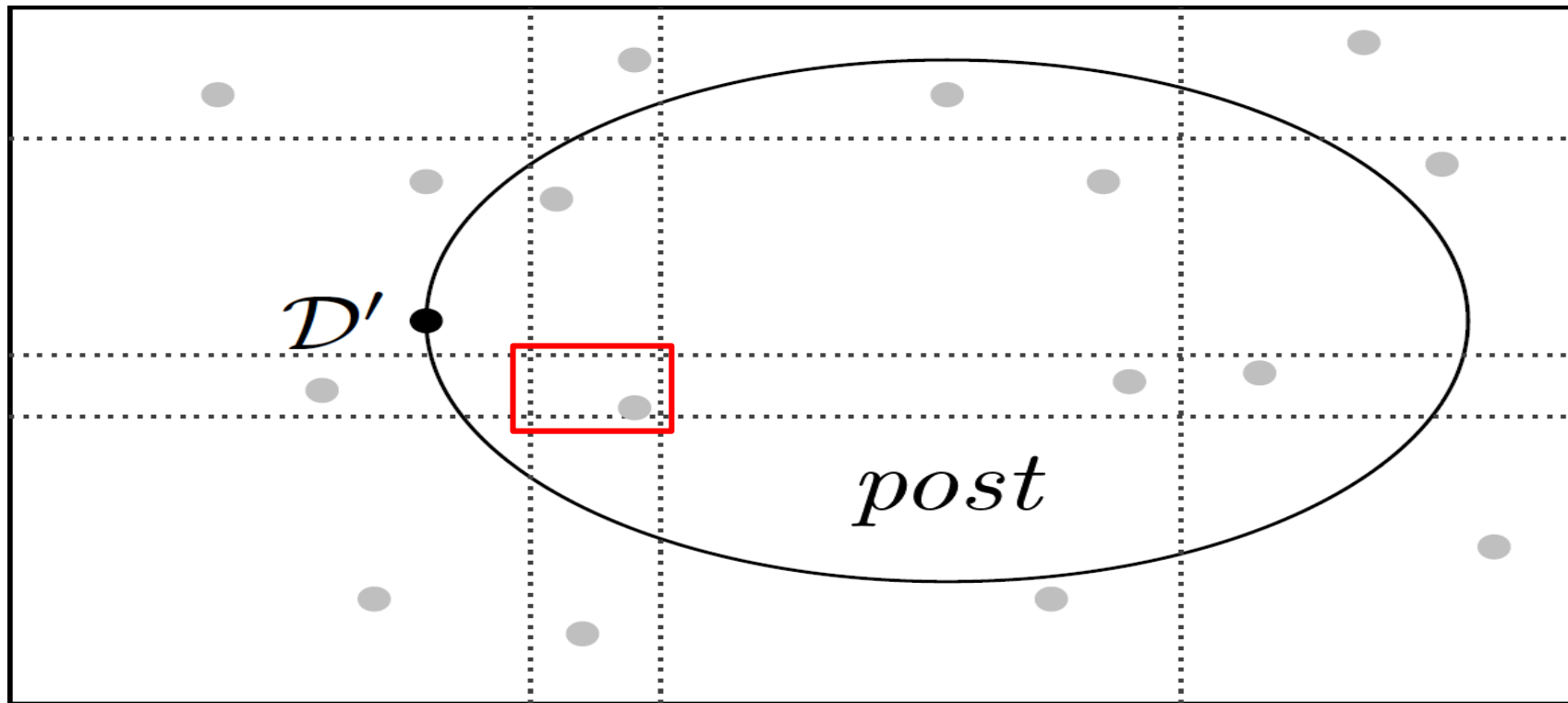
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# VC dimension of a set of programs

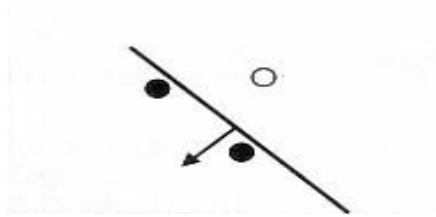


image: V. Kecman, 2001

# VC dimension of a set of programs

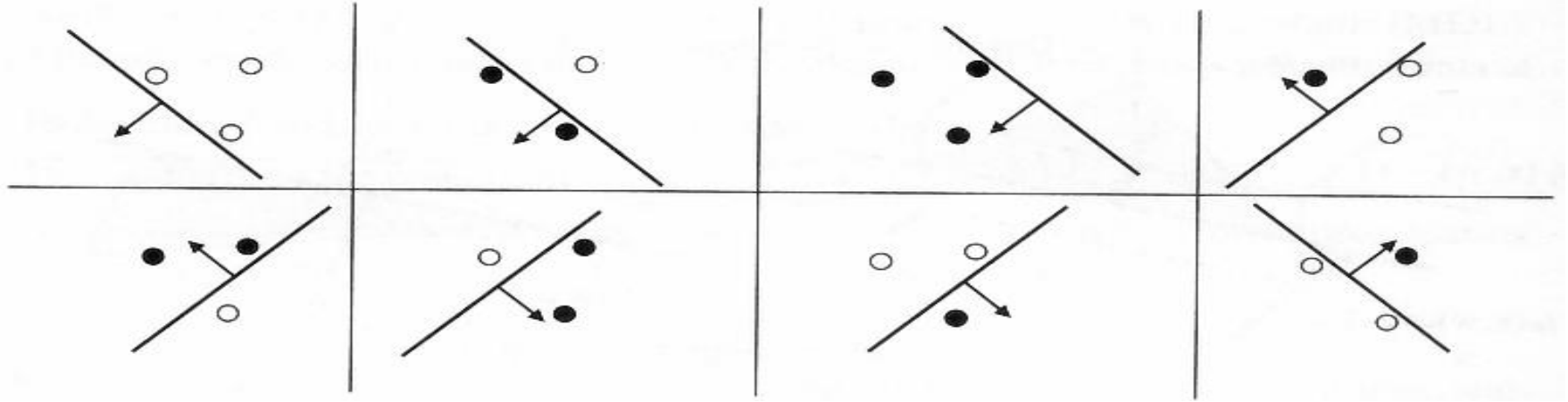


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# VC dimension of a set of programs

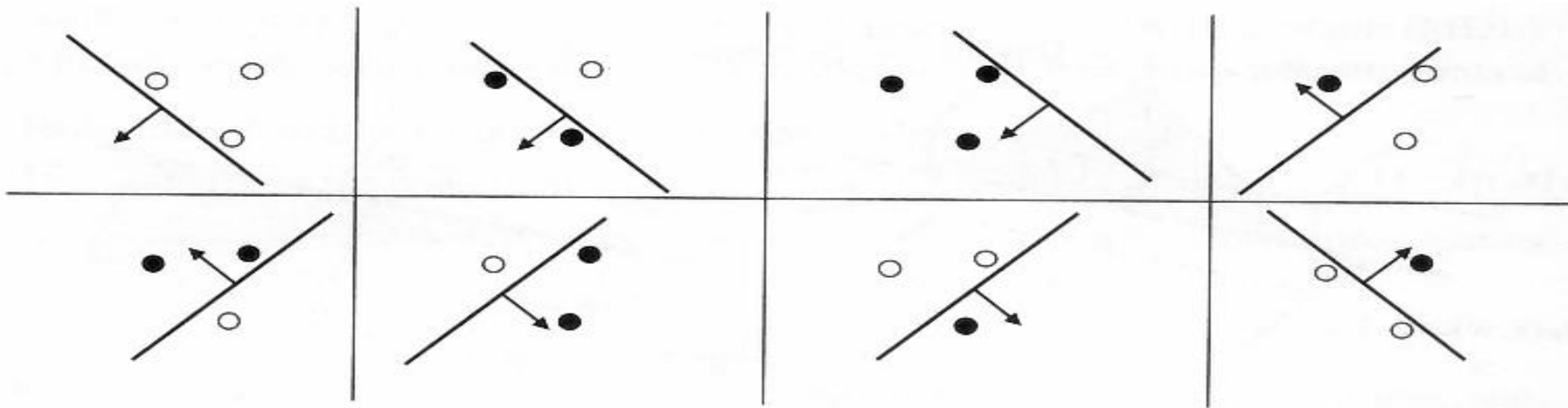
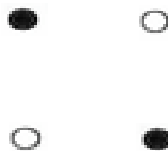


image: V. Kecman, 2001



No line can separate these four points

```
fun hireRep(min,urank)
    dec = ●1 <= urank <= ●2
    return dec
```

$$\varphi = ((H_1 \leq urank \leq H_2) == dec)$$

VC Dimension is 2

urank = 2	true
urank = 10	false
urank = 20	true

# Remember PAC learning?

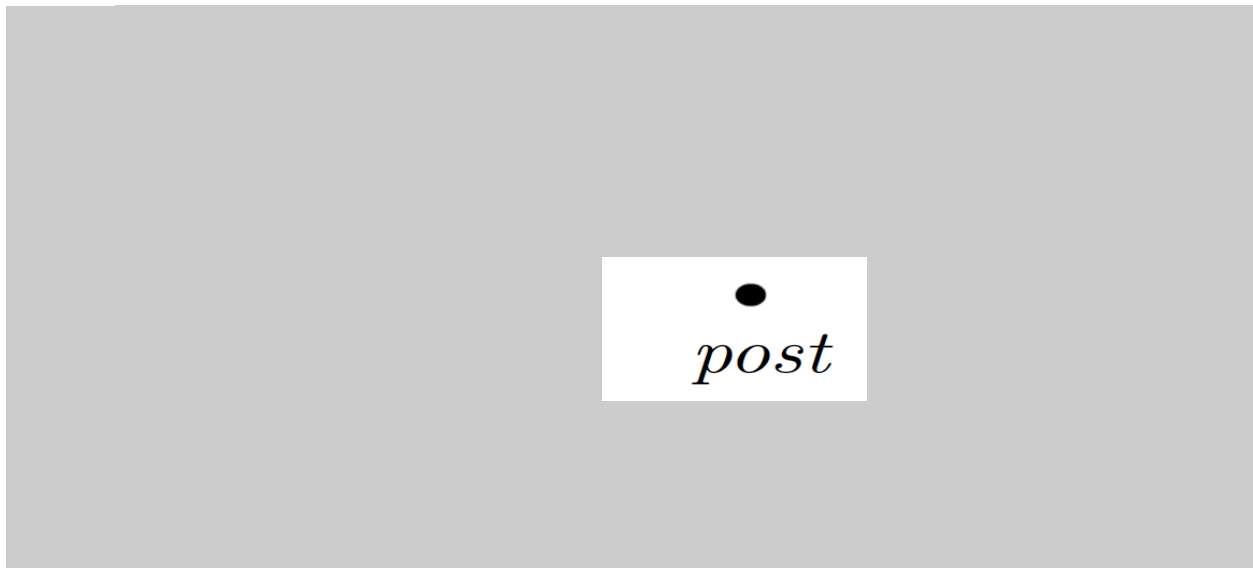
Finite VC dimension



$\bullet$   
*post*

# Remember PAC learning?

Finite VC dimension



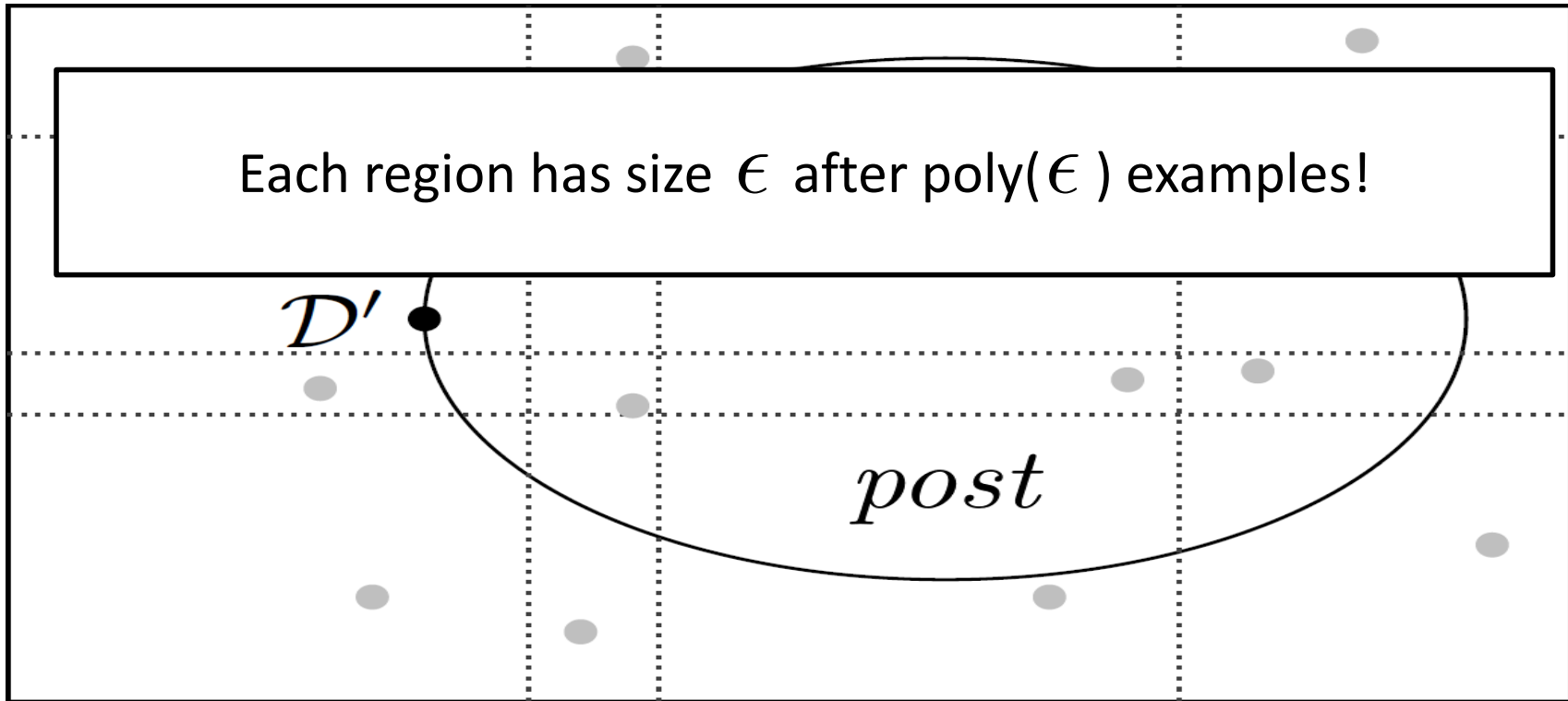
The white square has size  $\epsilon$  after  $\text{poly}(\epsilon)$  examples

Finite VC dimension

Each region has size  $\epsilon$  after  $\text{poly}(\epsilon)$  examples!

$\mathcal{D}'$

*post*



# What is so different about program synthesis?

## Inductive learning

Certain types of functions

Can be learned approximately

Using polynomially many IO examples

## Inductive synthesis

~~Search space is a mess~~

~~We can't be sort of right~~

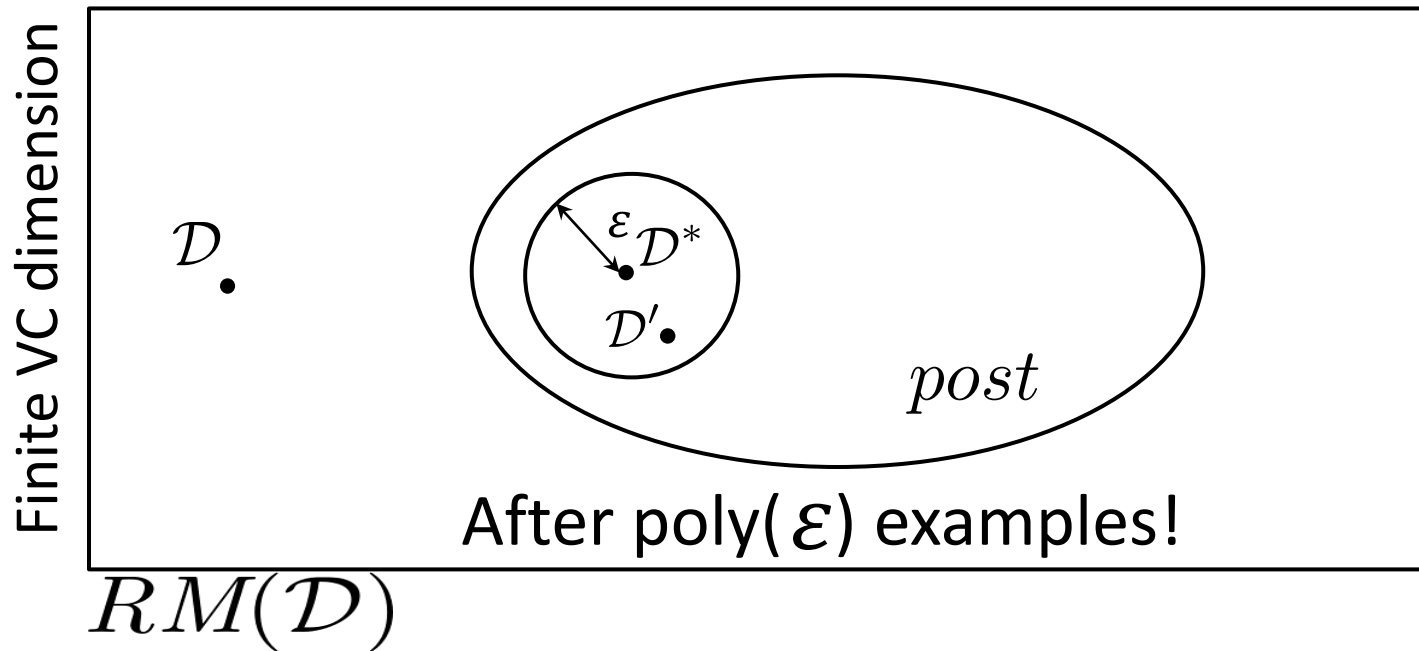
~~We don't have labeled examples~~

## How long will my synthesizer run for?

To answer this question, we need to address these problems



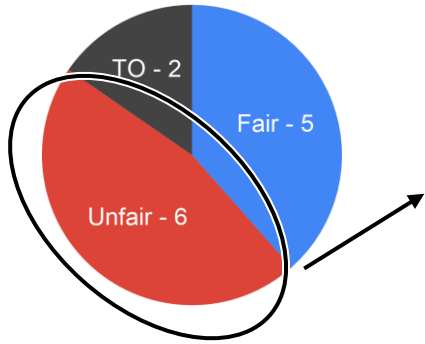
# Theorem: Convergence of Digits



DOES IT WORK?

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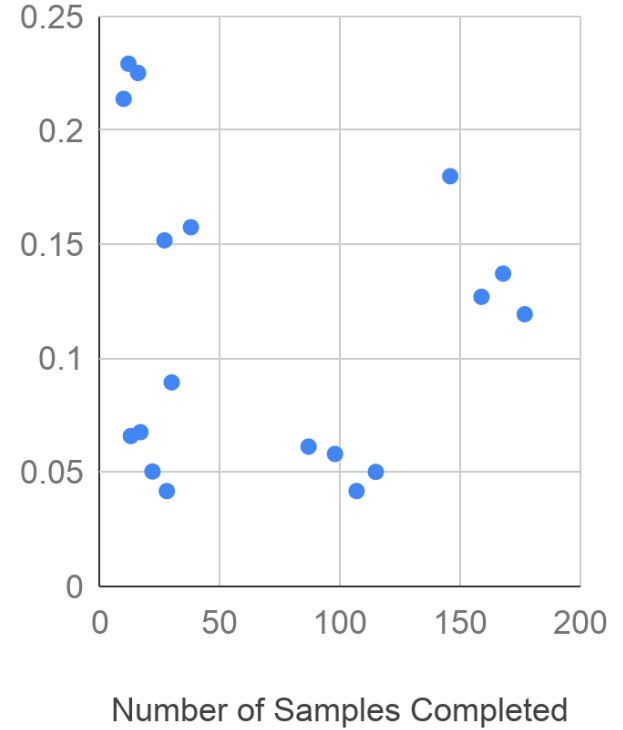
# DIGITS Case Study



Manually replace some  
constants with "holes":  
18 Repair Problems

**All Repaired!**

Best Solution (minimal  $\Pr[P(x) \neq P'(x)]$ )



# Distribution-Guided InducTive Synthesis (DIGITS)

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Show that it actually works

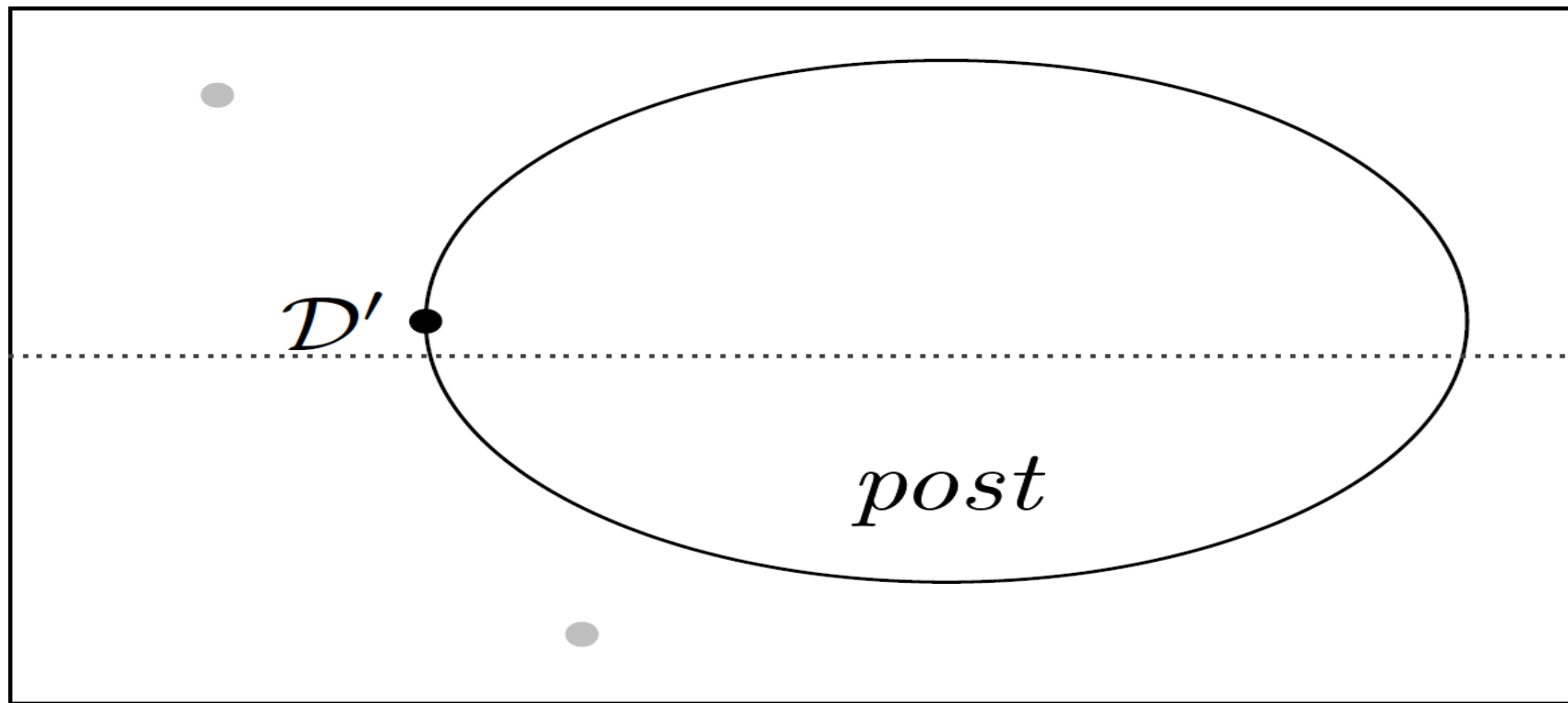
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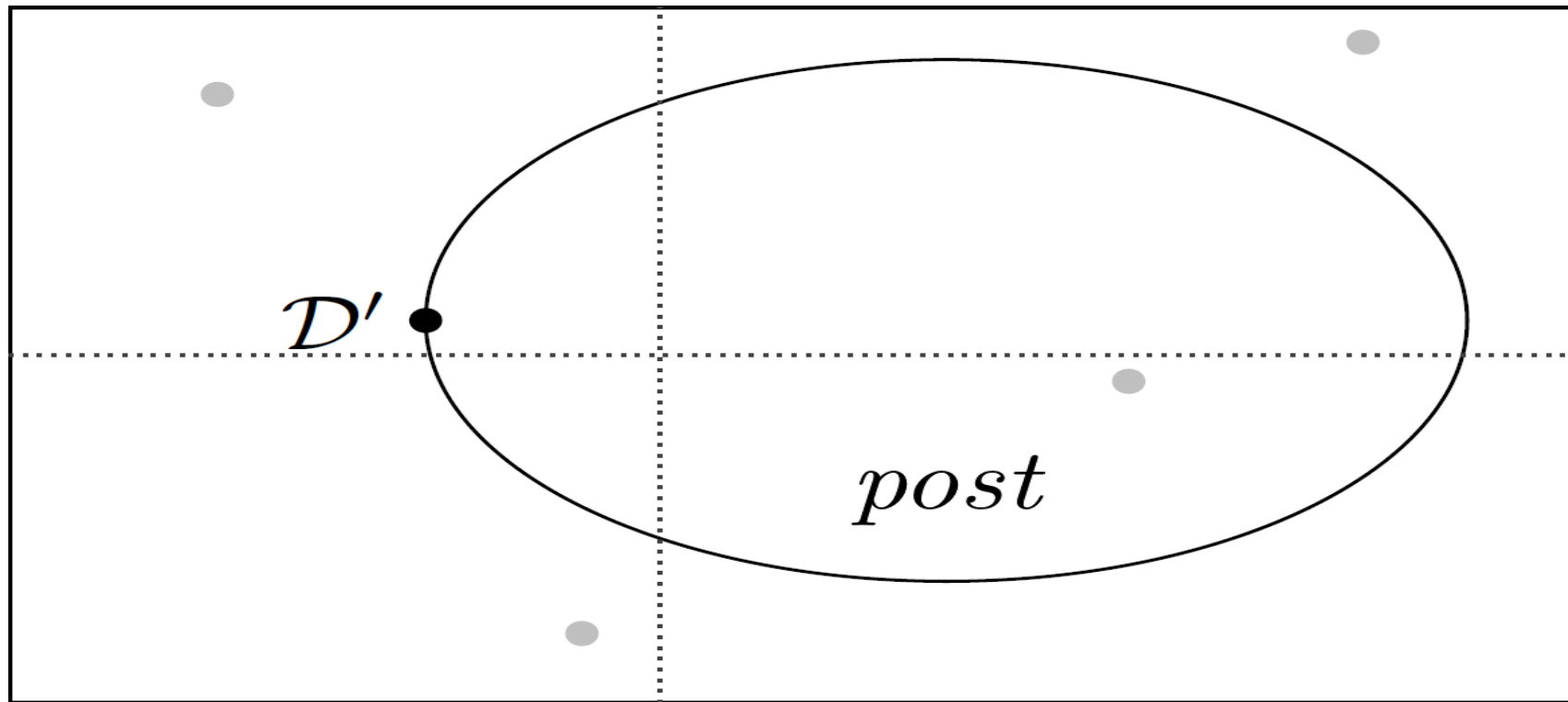
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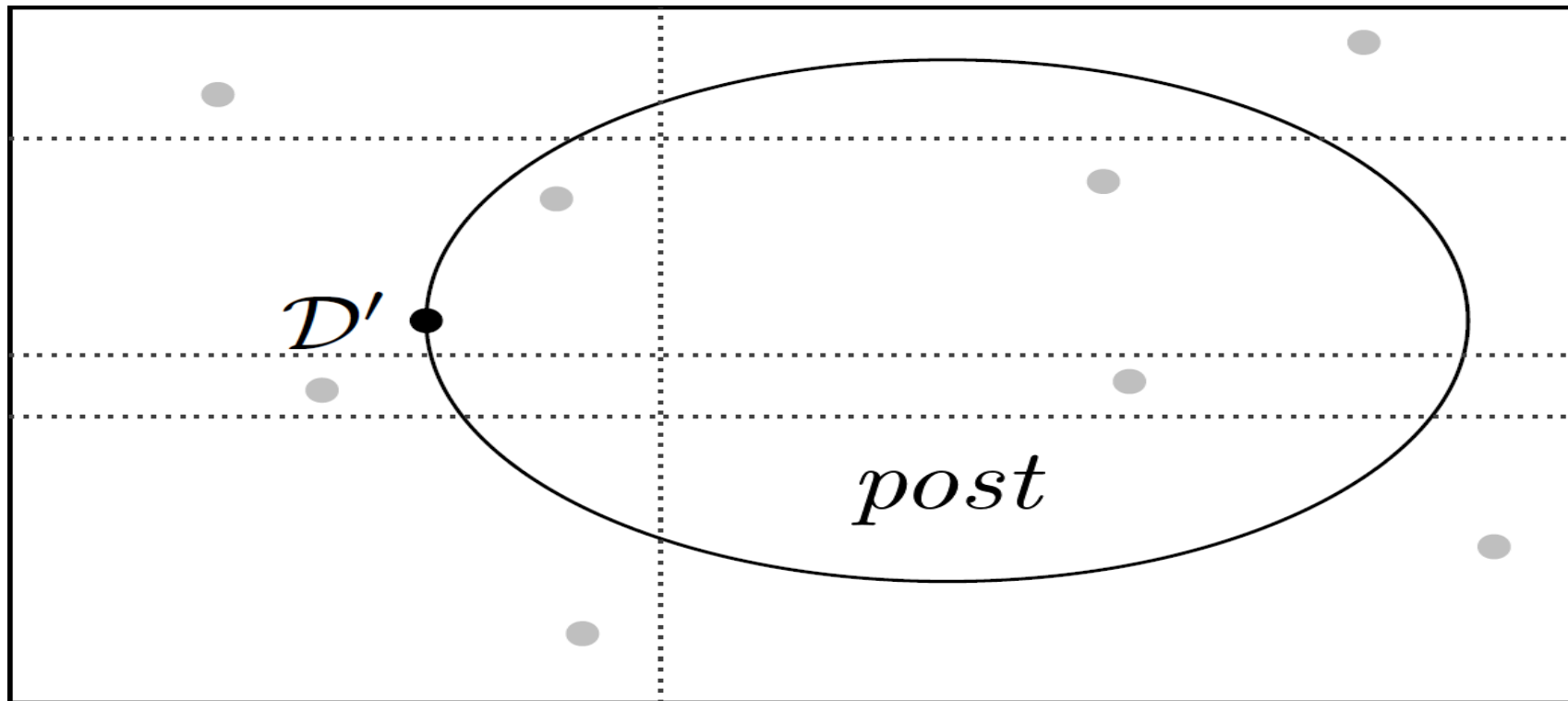
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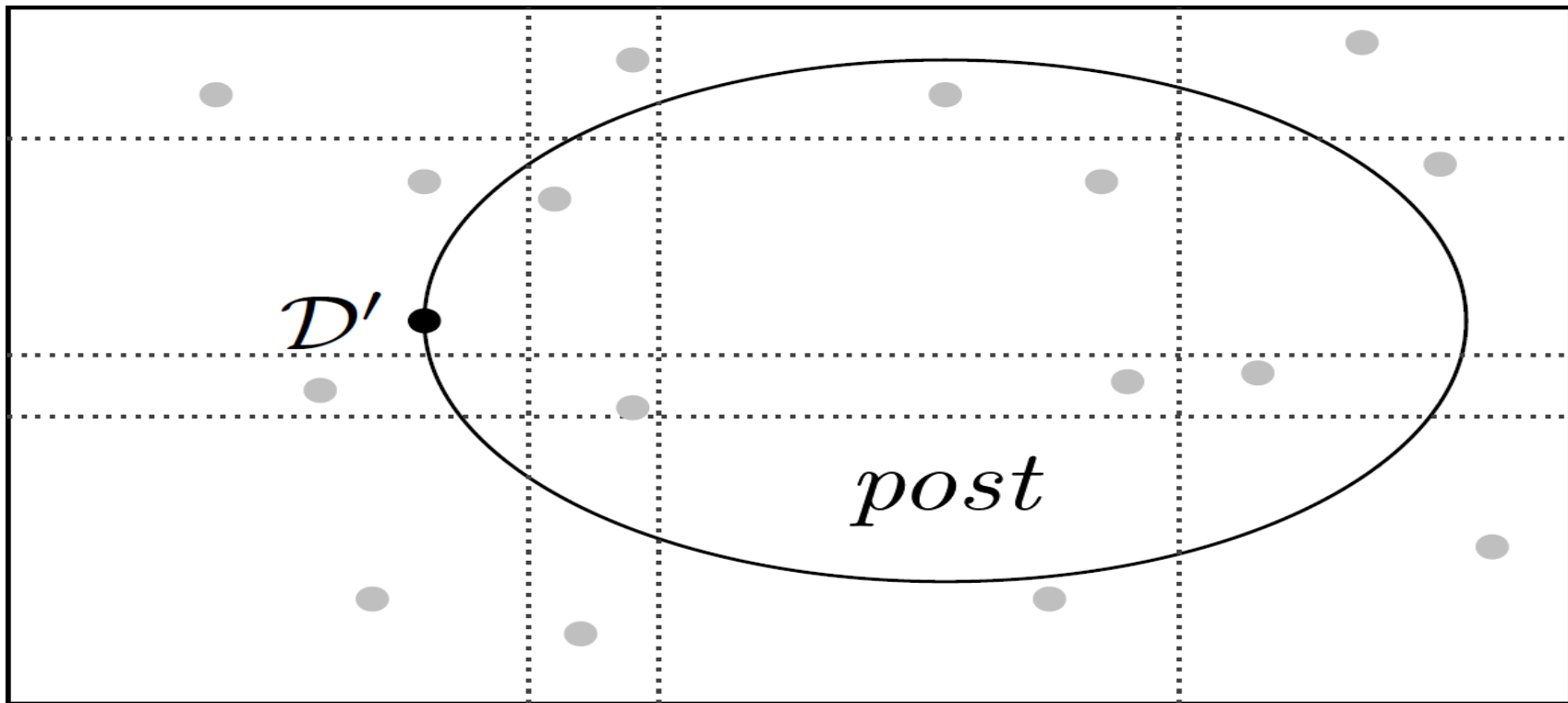
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Great complexity bound [Synthesis session at CAV19]



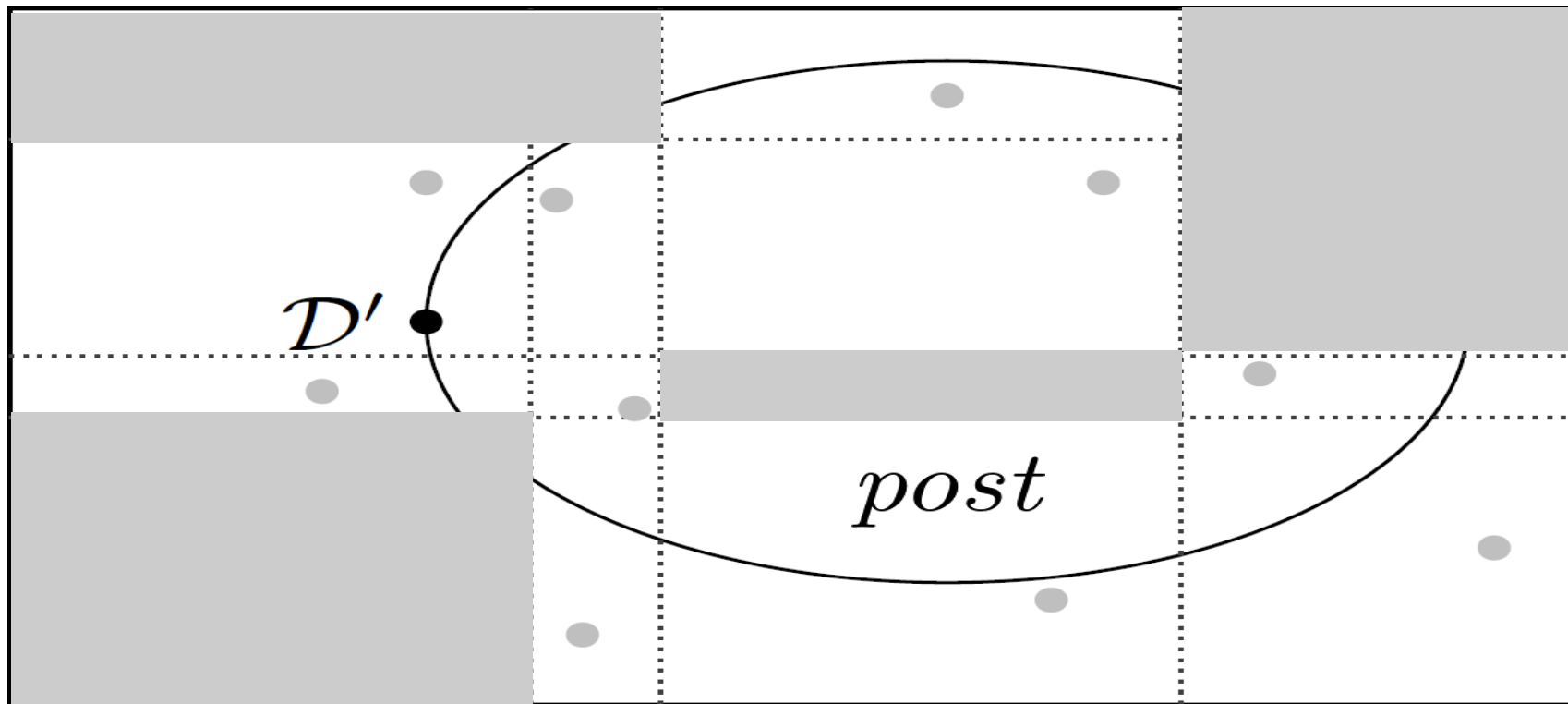






Number of dots in this plot is exponential in the number of examples





But maybe not all these regions admit a solution...

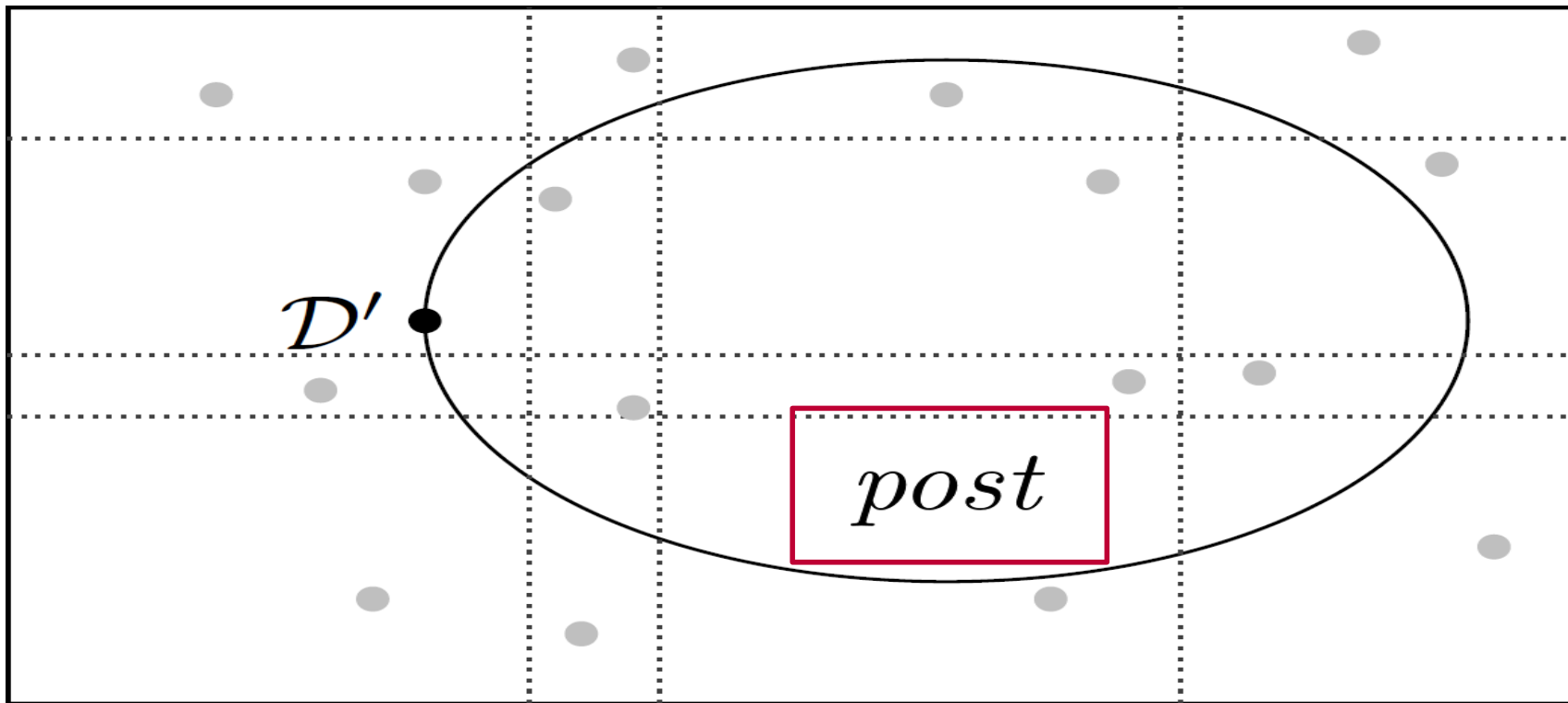
**Tuesday's talk:** only  $O(\text{poly}(\varepsilon))$  regions have a solution!

$\mathcal{D}'$

*post*

But maybe not all these regions admit a solution...

Can we speed up the search? [Sam on Tuesday]



Can we handle more expressive programs?

$$\{v \sim \mathcal{M}\}$$

$$h \leftarrow \mathcal{D}(v) \longrightarrow \begin{array}{l} \text{Decision making program:} \\ \text{returns true or false} \end{array}$$

$$\left\{ \frac{\Pr[h \mid v_s]}{\Pr[h \mid \neg v_s]} > 1 - \epsilon \right\}$$

What about programs that compute values?

## CONCLUSION

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