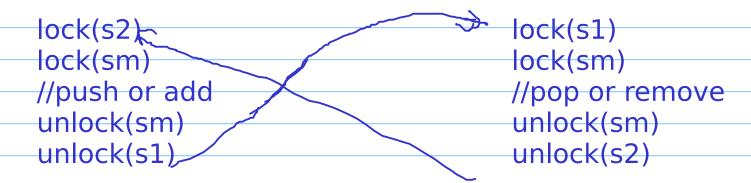
```
IPC:-
Mutex
                             - shared mem
Unnamed semaphores
                             - shared mem
Named semaphores
                             - any process (file path)
#define GNU SOURCE
#include<pthread.h>
(or)
gcc -D GNU SOURCE demo.c
gcc - DPI = 22.0/7.0 demo.c
https://visualgo.net/en
data structure visualization ==> cs.usfca.edu
Stack, Queue, LinkedList
Complexity - O(1), O(n), O(logn), O(nlogn), O(n^2)
```

Semaphore sm; sm.val=1 //Rule-3 Sempahore s1; s1.val=0; //Rule-1 Semaphore s2; s2.val=n; //Rule-2

Prod Cons



sm can be mutex instead of semaphore

Semaphores for processes (diff addr space)

- \* Named seamphores
- \* Sys V semaphores

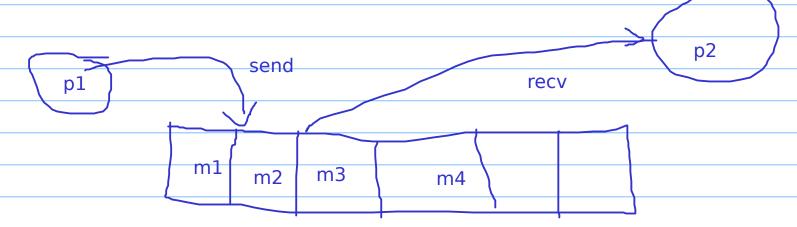
## Message Queues:-

Sender --> store messages

Receiver --> receive messages

- \* internal sync support
  - \* if there are no message (Q ie smpty) receiver will block
  - \* if Q is full, sender will block
- \* message descreter m1:10, m2:20 , one recv : m1

\* Typically FIFO based, but can be priority based



POSIX Message Queues
Sys V Message Queues ( later)

```
POSIX Message Queue:-
                       mq_notify
mq_open
mq send
                           mq setattr
mq_recv
mq close
mq_unlink
mq getattr
mqd_t mqid;
mqid = mq_open("sample", O_WRONLY);
                                                  //O_RDONLY
(or)
mqid=mq open("sample", O WRONLY|O CREAT, 0666, NULL);
(or)
struct mq_attr attr;
attr.mq maxmsg = 10;
attr.mq msgsize = 128;
mqid=mq open("sample", O RDONLY|O CREAT, 0666, &attr);
```

```
sending:-
     char msg[]= "Hello Linux";
                                       //strcut Student
     int len = strlen(msg);
                                        //sizeof(struct Student)
     int prio = 1;
     mq send(mqid, msg, len, prio);
                                       //mg send(mgid, &s1, len, prio)
Receive:-
     char buf[64];
     int maxlen=64;
     int prio;
     mq recv(mqid, buf, maxlen, &prio);
Close:-
     mq close(mqid);
-Irt (or) -Ipthread
Is /dev/mqueue
                    ==>
cat /proc/sys/fs/mqueue/msg max
cat /proc/sys/fs/mqueue/msgsize default
mq unlink
msgget, msgsnd, msgrcv, msgctl ==> Sys V Message Queues
```

```
File System:-
```

In UNIX/Linux (\*nix) everything is a file

UART device/port ==> /dev/ttyS0

What is a file system?

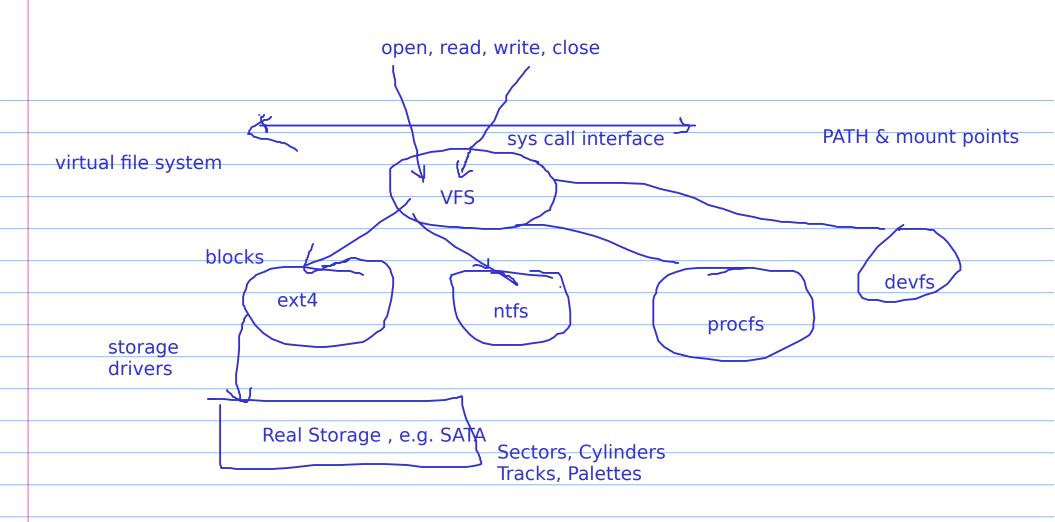
File System -- logical, Storage Device/Disk -- physical SATA/USB/Flash/Optical/IDE

FAT32, NTFS, ext4 xfs, zfs and many more

Logical FS maps to physical storage with the help of drivers (storage drivers)

uspace:- open, read, write, close

Making a new file system (like format) mkfs.\*



Pseudo file system , e.g. devfs, procfs, sysfs, tmpfs pipefs, debugfs

No data on real disk, only file entries

UART device/port ==> /dev/ttyS0 /proc/cpuinfo Memory based file system - ramdisk

## one partition of a disk - one file system

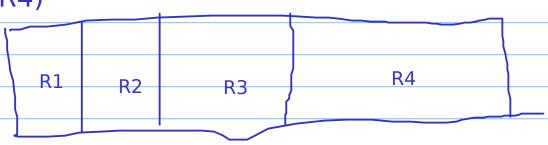
#### FS Layout:-

Boot block (R1)

Super block (R2) - meta data about filesystem

i-node blocks (R3) - meta data abt files (attributes)

data blocks (R4)



## Multistage bootloader

primary bootloader (bootstrap) stored in ROM, minimal hardware init & look for storage devices e.g. BIOS

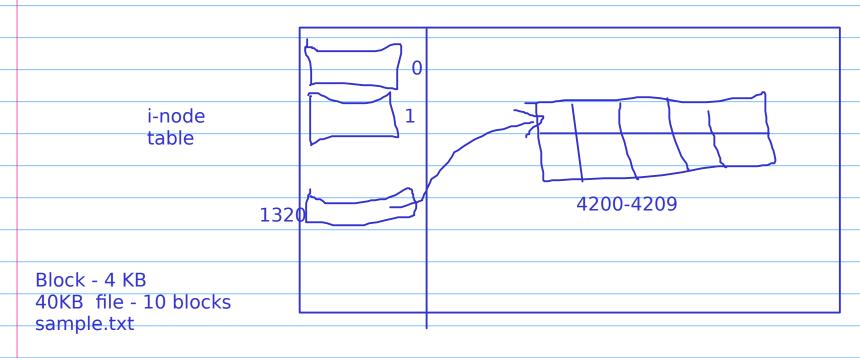
second stage bootloader stored first few blocks of selected storage e.g. GRUB, u-boot

last stage boot loader, locate kernel loads into memory (uncompress) kernel will do full init of system

at end load initial process in user space, i.e. init

(based on init scripts in /etc dir and run levels)

# Beginning of Disk - Master Boot Record (MBR) Linux Boot Sequence



Typical file attributes stored in i-node:-

- \* data block information
- \* mode : type + permissions
- \* inode ref/link count
- \* credentials user, group
- \* size of the file (actaul, no.of blocks)
- \* timestamps mtime, atime, ctime

```
ls -i
ls -li
Is -lid code
stat <file name>
inode ref count
cp f1 f2
rm command: unlink file name with inode
                 if no more links, then release i-node, data blocks
Hard links:-
  In f1 f2
file name vs inode no.s are maintained by directory
soft links / symolic links:-
```

```
Mounting:-
partitions --> primary, extended (logical)
                /dev/sda1
1st
        ==>
        ==> /dev/sda2 or /dev/sda5
2nd
        ==> /dev/sda6
3rd
                 /dev/sda7
                                  (Linux)
4th
        ==>
USB:- /dev/sdb1, /dev/sdb2, /dev/sdb3 etc
SD Card:- /dev/mmcblk0p1, /dev/mmcblk0p2
    mkdir /mnt/c # mount point
    # /mnt/c is blank before
    mount /dev/sda1 /mnt/c
    # can see sda1 (C:) content under /mnt/c
    umount /mnt/c
    #again blank
    C:\samples\hello.c
                         ==> /mnt/c/samples/hello.c
```

```
dev/sda, dev/sdb, dev/mmcblk0 ==> for unpartitioned disks
sudo fdisk -l
sudo fdisk -l /dev/sda
mount
mount /dev/sdb1 /mnt/usb
umount /mnt/usb
df -kh
attach/mapping
single origin for entire FS - /
Activity:-
* Go through FS notes, tutorials on Pipes/Fifos
* Any tutorials on relevant IPC concepts
* Assignments
* Please try first-steps, under kprog-drivers. username:- root
                                                 no pasword
```