SYSCALL(2) Linux Programmer's Manual SYSC

NAME top

syscall - indirect system call

SYNOPSIS top

```
#define _GNU_SOURCE         /* See feature_test_macros(7) */
#include <unistd.h>
#include <sys/syscall.h>   /* For SYS_xxx definitions */

long syscall(long number, ...);
```

DESCRIPTION top

syscall() is a small library function that invokes the system c
whose assembly language interface has the specified *number* with
specified arguments. Employing **syscall**() is useful, for exampl
when invoking a system call that has no wrapper function in the
library.

syscall() saves CPU registers before making the system call, re
the registers upon return from the system call, and stores any
code returned by the system call in errno(3) if an error occurs

Symbolic constants for system call numbers can be found in the
file *<sys/syscall.h>*.

RETURN VALUE top

The return value is defined by the system call being invoked.
general, a 0 return value indicates success. A -1 return value
indicates an error, and an error code is stored in *errno*.

NOTES top

syscall() first appeared in 4BSD.

Architecture-specific requirements

Each architecture ABI has its own requirements on how system ca

arguments are passed to the kernel. For system calls that have glibc wrapper (e.g., most system calls), glibc handles the deta copying arguments to the right registers in a manner suitable f architecture. However, when using **syscall**() to make a system c the caller might need to handle architecture-dependent details; requirement is most commonly encountered on certain 32-bit architectures.

For example, on the ARM architecture Embedded ABI (EABI), a 64- value (e.g., *long long*) must be aligned to an even register pai Thus, using **syscall**() instead of the wrapper provided by glibc, **readahead**() system call would be invoked as follows on the ARM architecture with the EABI:

```
syscall(SYS_readahead, fd, 0,
        (unsigned int) (offset >> 32),
        (unsigned int) (offset & 0xFFFFFFFF),
        count);
```

Since the offset argument is 64 bits, and the first argument (*f* passed in *r0*, the caller must manually split and align the 64-b value so that it is passed in the *r2*/*r3* register pair. That me inserting a dummy value into *r1* (the second argument of 0).

Similar issues can occur on MIPS with the O32 ABI, on PowerPC w the 32-bit ABI, and on Xtensa.

The affected system calls are fadvise64_64(2), ftruncate64(2), posix_fadvise(2), pread64(2), pwrite64(2), readahead(2), sync_file_range(2), and truncate64(2).

Architecture calling conventions

Every architecture has its own way of invoking and passing argu to the kernel. The details for various architectures are liste the two tables below.

The first table lists the instruction used to transition to ker mode, (which might not be the fastest or best way to transition the kernel, so you might have to refer to vdso(7)), the registe to indicate the system call number, and the register used to re the system call result.

arch/ABI	instruction	syscall #	retval	Notes
arm/OABI	swi NR	-	a1	NR is syscal
arm/EABI	swi 0x0	r7	r0	
blackfin	excpt 0x0	P0	R0	
i386	int $0x80	eax	eax	
ia64	break 0x100000	r15	r10/r8	bool error/ errno value
parisc	ble 0x100(%sr2, %r0)	r20	r28	
s390	svc 0	r1	r2	See below
s390x	svc 0	r1	r2	See below
sparc/32	t 0x10	g1	o0	

```
sparc/64   t 0x6d              g1          o0
x86_64     syscall             rax         rax

For s390 and s390x, NR (the system call number) may be passed
directly with "svc NR" if it is less than 256.

The second table shows the registers used to pass the system ca
arguments.

arch/ABI   arg1   arg2   arg3   arg4   arg5   arg6   arg7
──────────────────────────────────────────────────────────
arm/OABI   a1     a2     a3     a4     v1     v2     v3
arm/EABI   r0     r1     r2     r3     r4     r5     r6
blackfin   R0     R1     R2     R3     R4     R5     -
i386       ebx    ecx    edx    esi    edi    ebp    -
ia64       out0   out1   out2   out3   out4   out5   -
parisc     r26    r25    r24    r23    r22    r21    -
s390       r2     r3     r4     r5     r6     r7     -
s390x      r2     r3     r4     r5     r6     r7     -
sparc/32   o0     o1     o2     o3     o4     o5     -
sparc/64   o0     o1     o2     o3     o4     o5     -
x86_64     rdi    rsi    rdx    r10    r8     r9     -

Note that these tables don't cover the entire calling conventio
architectures may indiscriminately clobber other registers not
here.
```

EXAMPLE top

```
#define _GNU_SOURCE
#include <unistd.h>
#include <sys/syscall.h>
#include <sys/types.h>
#include <signal.h>

int
main(int argc, char *argv[])
{
    pid_t tid;

    tid = syscall(SYS_gettid);
    tid = syscall(SYS_tgkill, getpid(), tid, SIGHUP);
}
```

SEE ALSO top

_syscall(2), intro(2), syscalls(2), vdso(7)

COLOPHON top

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Linux **2014-05-10** **SYSC**
