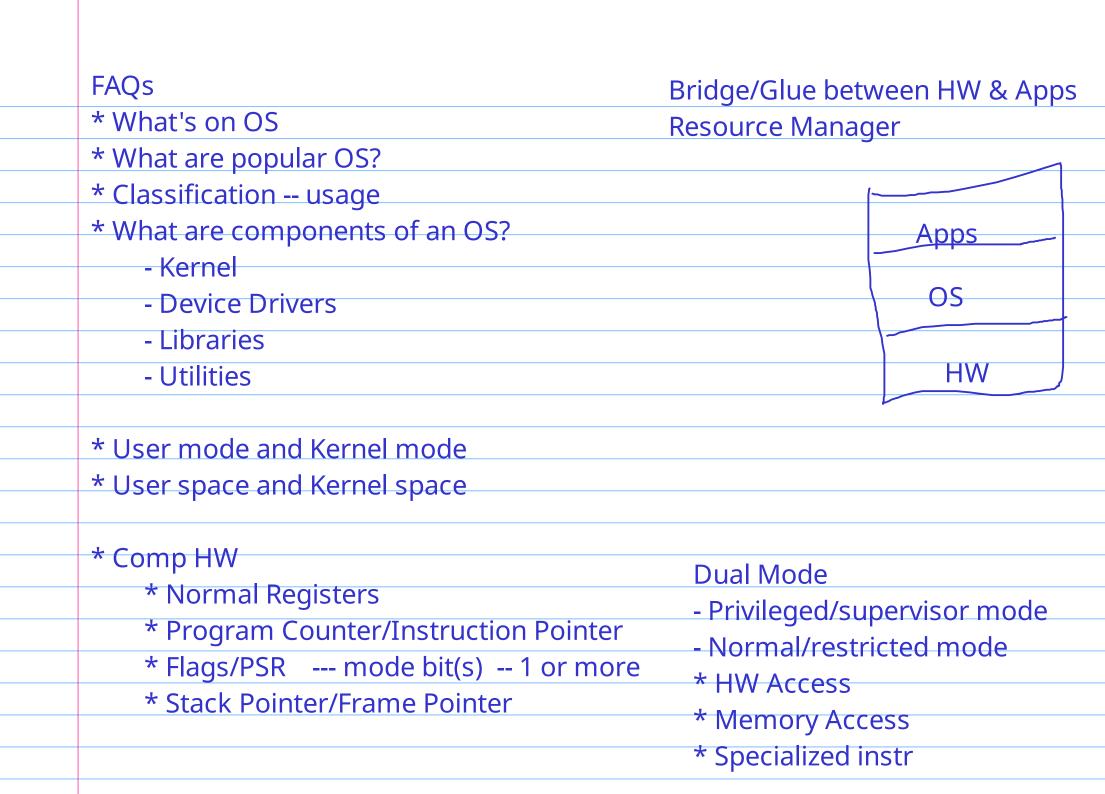
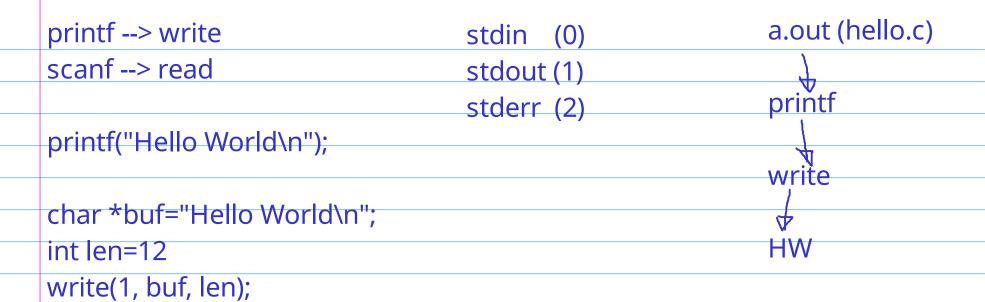
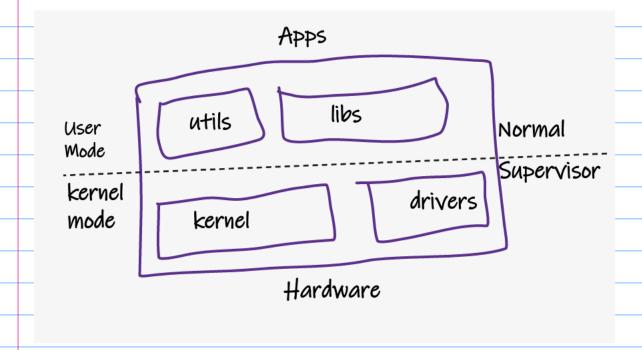
OS Topics:-* Intro & Arch (Kernel, System Calls) * Process Management * Threads * Scheduling * Signals * IPC * Memory Management * File System Need for learning OS Practical Examples / Assignments Commands - Analysis PBL / Micro Proj / Prob Statement https://github.com/caia-techblr/linux-os-sys-prog/







Trap Instruction

int 80h sysenter

ARM:swi/svc

```
examples on file handing
- open,read,write,close,lseek
syscall(SYS_write, fd, buf, len);
write(fd, buf, len);
Assignment Ideas
- Copy files (cp file1 file2)
- Display file contents (cat file1)
- How to read large file
    file contents: - "ABCEDFGHIJKLMNOPQRSTUVWXYZ0123456789"
    char buf[11];
    int buflen=10;
    while(1)
         nbytes = read(fd, buf, len);
         if(nbytes==0) break;
         //print buf
```

```
sudo apt update
sudo apt install build-essential
gcc wrsample.c -o wrsample
./wrsample
gcc rdsample.c -o rdsample
./rdsample
int main() {
int i;
for (i = 1; i <= 5; i++)
  printf("hello:%d\n", i);
 return 0;
                                       strace echo "Hello Linux"
gcc hello.c -o hello
strace./hello
```

Process	PCB
- program loaded in memory for execution	
- program on disk (passive), process in memory(active)	pid
- memory sections a process	ppid
- stack	name
- code	context (??)
- idata (.data)	scheduling info
- udata (.bss)	memory info
- heap	file info
- process table / process list	cred (uid, gid)
- process control block / process descriptor	exit status
- process states, process life cycle	
- context switching - context saving + context loading	
process hierarchy : parent & child	

Boot Sequence

boot sequence
ps
ps -el
ps aux
ps -e -o pid,ppid,stat,cmd
pstree -np
ret = waitpid(-1, &status, 0);
if(WIFEXITED(status))
printf("normal,child exit status = %d\n", WEXITSTATUS(status));
else
printf("abnormal termination\n");



```
#include <pthread.h>
                                                   Threads
#include <stdio.h>
                                                   - Concurrency
                                                   - Resource Sharing
void *task_body(void *pv) {
 char *plabel = pv;
 int i;
 printf("%s--welcome\n",plabel);
 for (i = 1; i <= 100; i++)
  printf("%s--%d\n", plabel, i);
 // pthread_exit(NULL);
int main() {
 pthread_t pt1, pt2, pt3; // thread handle
 pthread_create(&pt1, NULL, task_body, "A1");
 pthread_create(&pt2, NULL, task_body, "B2");
 pthread_create(&pt3, NULL, task_body, "C3");
 pthread_join(pt1,NULL);
 pthread_join(pt2,NULL);
 pthread_join(pt3,NULL);
 printf("main--thank you\n");
```

IPC	Semaphore
- Race conditions	- integral value/counter
- Critical Section	- Q to hold processes/threads
- Mutual Exclusion	
- Semaphores, Mutex	wait/lock/block
- Deadlock	1A. if value>0, value, go ahead
	1B. if value==0, block and add to Q
- Synchronization	
	release/unlock/unblock
	2A. if Q is not empty, remove one Q and resume
	2B. if Q is empty, value++
Semaphores, Mutex	
•	

Mutex API pthread_mutex_t m1; //global decl pthread_mutex_init(&m1, NULL); //before pthread_create pthread_mutex_destroy(&m1); //after join pthread_mutex_lock(&m1); // before val++ or val-pthread_mutex_unlock(&m1); // after val++ or val--

```
sem_t sm; //global / shared
                sem init(&sm,pshared,ivalue);//pshared-0, ivalue-1
           Thread-1
                                            Thread-2
while(1) {
                                 while(1) {
   r*sem wait(sm)
                                    r*sem wait(sm)
    //critical section
                                   //critical section
   sem post(sm)
                                    sem post(sm)
                  sem_destroy(&sm)
Semaphores for synchronization (Dependency/sequencing)
```

```
//create a semaphore, initial val = 0
Prod
                      Cons
```

```
sem_wait(&s1);
//add
                    //remove
sem_post(&s1);
```

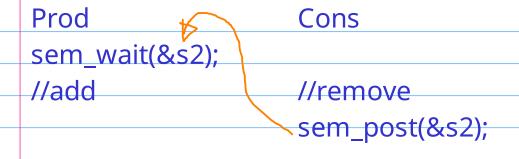
Mutex is not just Binary Semaphore

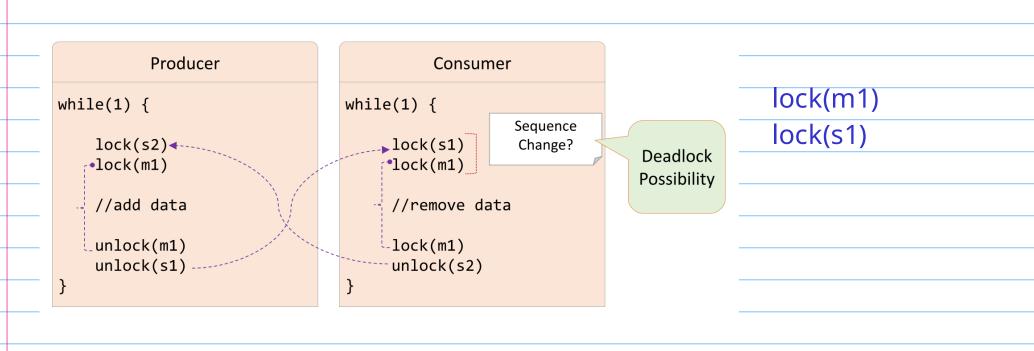
Mutex Features

- Strictly for mutual exclusion
- Ownership applies, whichever thread locks mutex, same can unlock
- Unlocking before locking/ unlocking more than once allowed
- In RTOS, Mutex provides solutions to Priority Inversion problem (e.g. priority inheritance/ceiling solutions)

Producer Consumer Prob

- There is a common data source, accessible by multiple threads
- Threads known as producers add data
- Thread known as consumers remove data
- R1 If buffer is empty, consumer should block, only producer is allowed
- R2 If buffer is full, producer should block, only consumer is allowed
- R3 If both are allowed (partially filled), any one can access at a time





Mutex pm;	//Semaphore pm, ival=1; Print	ter
Mutex sm;	//Semaphore sm, ival=1; Scar	nner
A1	A2	Prevention At a ratio la alaire at / a a autima
		- Atomic locking / acquire
lock(sm)	lock(pm)	all resources in one go
lock(pm)	lock(sm)	- Try to lock resources
		in particular order
//copy	//copy	- Apply mutual exclusion
		after resolving
unlock(pm)	unlock(pm)	dependencies
unlock(sm)	unlock(sm)	

0) Examples, commands, techniques
1)Design your own shell (mini shell / tiny shell)
==> Display some string as prompt, e.g. "myshell>"
==> Take command name as user input
==> Create a child process using fork
==> Launch requested requested using execl/execlp inside child
==> Parent(shell) wait for child (command exec) using waitpid and print exit status
==> Repeat above steps, until user input is "quit"
2) Write a multithreaded application to perform parallel sum of large array
3) Modify/enhance TCP server code, to communicate to multiple clients
using threads (concurrent server)
4) Producer-Consumer problem using threads practically
5) Implement own commands like cp, cat (open,read,write,close)