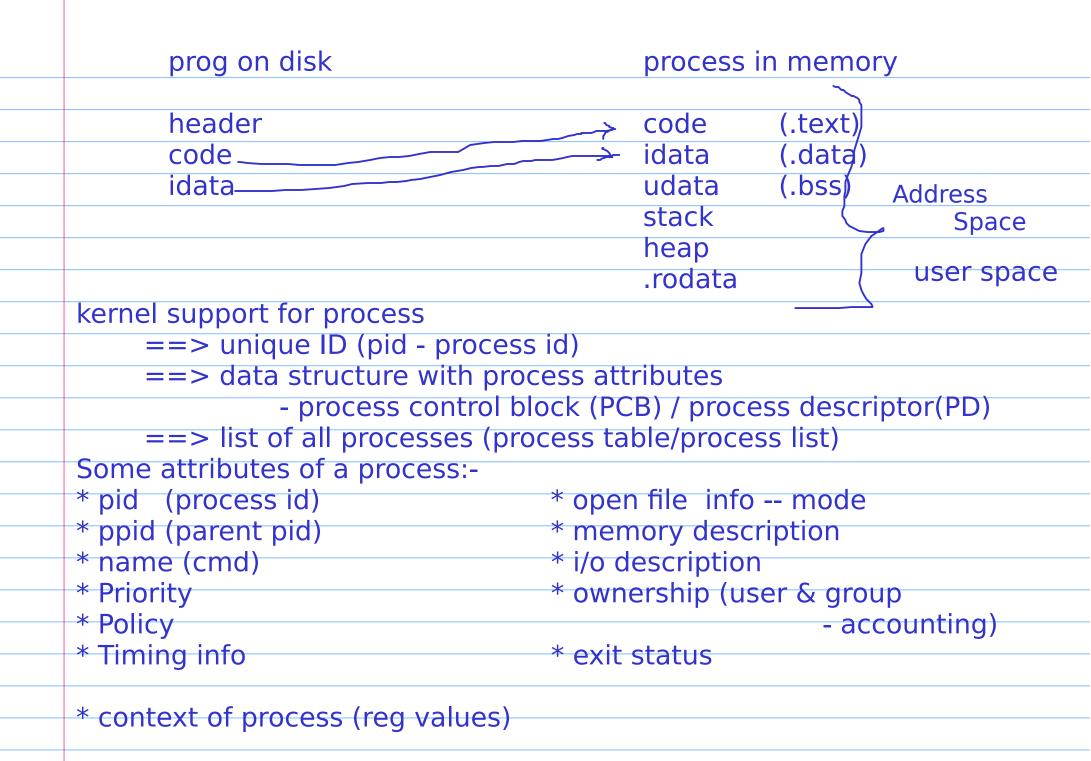
```
Examples of
 Monolithic
              -- Traditional Unix
 Micro
              -- RTOS like QNX
Hybrid Approach -- Modular Approach
Modules -- static vs dynamic modules
              /lib/modules
Ismod,
Simple system calls:-
fd=open("sample.txt",O_WRONLY|O_CREAT, 0666);
(or)
fd=open("sample.txt",O WRONLY);
(or)
fd=open("simple.txt",O RDONLY);
//system call wrapper in user space
char msg[] = "Hello Unix!";
nbytes= write(fd, msg, len);
```

```
system calls return -ve value on failures
0 or +ve means successful
TODO:-
wrsample.c
rdsample.c
alpha.txt/simple.txt ==> 36 bytes
int len=10;
nbytes=read(fd,buf,len); //A...J , nbytes=10
nbytes=read(fd,buf,len); //K...T , nbytes=10
nbytes=read(fd,buf,len); //U..0123 , nbytes=10
nbytes=read(fd,buf,len); //456789 , nbytes=6
nbytes=read(fd,buf,len); //nbytes will be zero, end of file
printf/fprintf vs write
                                                            POSIX std
lib APIs vs system calls:-
* interoperability (across OS)
* ease of use
* efficient
```

| strace <command/> |
|---|
| ctraco ca |
| strace cal strace ./a.out |
| |
| Process Management:- |
| Concept of a Process:- |
| |
| what is a process? |
| program loaded in memory for execution user oriented (user space) |
| program on disk> process in memory (loading) |
| flow of control function calls |
| |
| |
| |
| |
| |





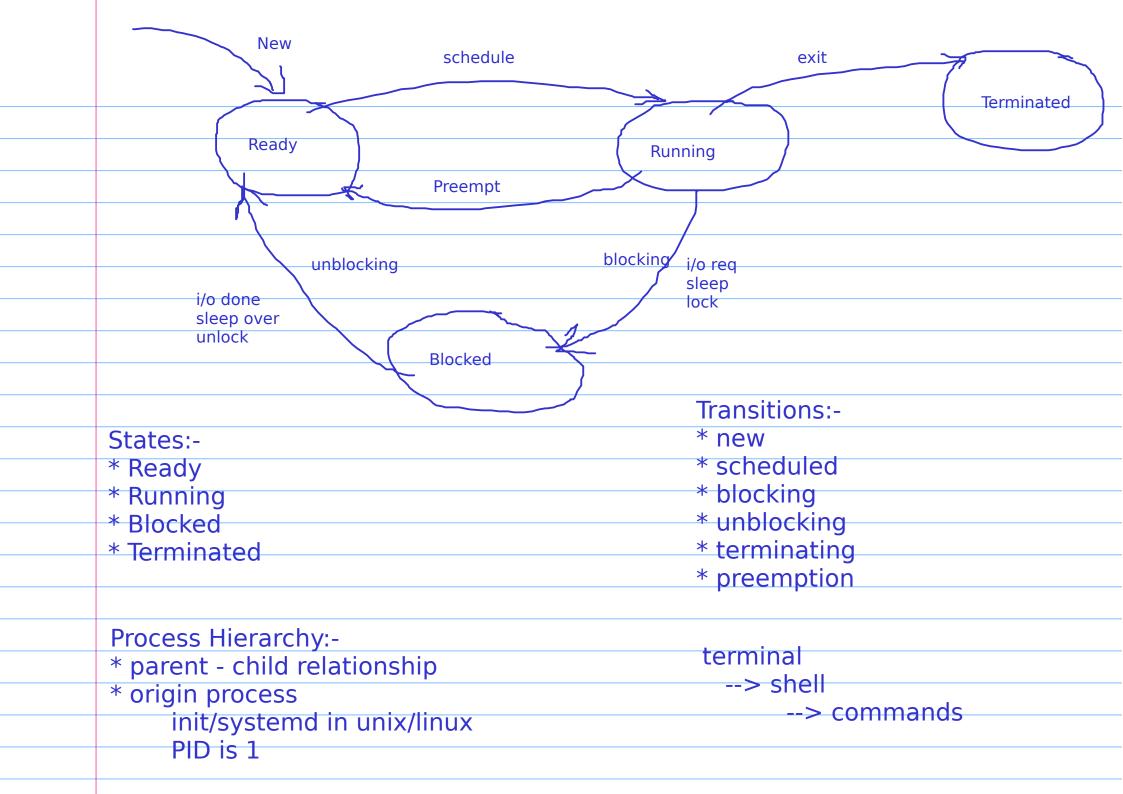
context area -- backup(snapshot) of registers

... typically in memory

... on top of process stack in X86

cotext saving -- copying CPU reg values in context area context loading -- filling CPU regs as per context area (saved reg values) context switching -- context saving + context loading

every process will have independent stack (stack frames) stack frames -- local vars, parameters of functions



```
webminal.org
repl.it
katacoda.com [ Ubuntu Playground ]
CoCalc Linux
Commands:-
ps
ps aux
ps -el
ps -e -o pid, ppid, stat, cmd
pstree
pstree -np
                   # q - quit
top
```

```
System Calls & APIs
getpid
                       sleep
getppid
                       exit
fork
waitpid
execl, execlp
exit
fork - create a new process, known as child
int ret;
ret=fork();
if(ret<0)
 perror("fork");
if(ret==0) {
 //some code for child
 exit(0);
else { //ret > 0
 //some code for parent
```

```
Example1
                    - simple fork
Example2a - concurrency, lengthy loop, no delay
Example2b - concurrency, small loop, with delay(sleep)
If parent terminates before child, init/systemd becomes parent
of running child (reparenting/adopting)
waitpid:-
At end of parent:-
    waitpid(-1, &status, 0);
    //print WEXITSTATUS(status);
waitpid 1st param:
    -1 means, waiting for any one child
     +ve val means, specific child
```

```
execl:-
                                                                    usr
if(ret==0) {
                        //child
  execl("/usr/bin/cal","cal",NULL);
  printf("Thank you"); //not reachable if execl succeeds
                                                                    which cal
                                                                    which gcc
else {
                                                                    which Is
     //waitpid
    //print status
//execl("/usr/bin/cal","cal","10", "2015", NULL);
//execlp("cal", "cal", "10", "2015", NULL);
                                                           //refer PATH variable
execl, execlp
execv, execvp
execle, execvpe
TODO:- kernel, system calls & process management (concepts ,hands-on)
Further:-
* Signal Handling
* Threads
* Scheduling
```

| Beginning Linux Programming - Richard Stones/Neil Mathew (11-15/16) |
|---|
| The Linux Programming Interface(TLPI) by michael kerrsick |
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