Machine Problem 4 - File System

CSIE3310 - Operating Systems National Taiwan University

Total Points: 100 Release Date: May 16

Due Date: May 29, 23:59:00

TA e-mail: ntuos@googlegroups.com

TA hours: Thu. 14:00-15:00 before the due date, CSIE Building B04

Contents

1	Overview	2
2	Environment Setup	2
3	Problem 1: Large Files (30 points) 3.1 Description 3.2 Guidelines and Hints 3.3 How to Test	2 2 2 3
4	Problem 2: Symbolic Links to Files (30 points) 4.1 Description	3 3 4
5	Problem 3: Symbolic Links to Directories (30 points) 5.1 Description	
6	Report (10 points) 6.1 Description	5
7	Submissions 7.1 xv6 code	5 5
8	Grading Policy	6

1 Overview

In this MP, you will learn the fundamental knowledge of the file system by adding two features to xv6: large files and symbolic links. We strongly recommend you read Chapter 8 (file system) in xv6 hand book while you trace code. This gives you a quick overview of how xv6 implements its file system.

2 Environment Setup

We provide the whole xv6 code you need for this MP. All the necessary files are given, and you only need to modify the missing parts. The following instructions will help you complete the setup.

1. Downloadthe MP4.zip from NTUCOOL, unzip it, and enter the folder.

```
$ unzip MP4.zip
$ cd mp4
```

2. Pull the docker image.

```
$ docker pull ntuos/mp4
```

3. Run the following command to start a container and mount the directory.

```
$ docker run -it --name mp4 -v $(pwd)/xv6:/home/xv6 ntuos/mp4
```

4. After getting into the container, run the command to start an xv6.

```
(container)$ make qemu
```

If everything is fine, then you are able to start working on the MP.

3 Problem 1: Large Files (30 points)

3.1 Description

In this problem, you have to increase the maximum size of an xv6 file. Currently xv6 files are limited to 268 blocks, or 268*BSIZE bytes (BSIZE is 1024 bits in xv6). This limitation comes from the fact that an xv6 inode contains 12 "direct" block numbers and one "singly-indirect" block number, which refers to a block that holds up to 256 more block numbers, for a total of 12 + 256 = 268 blocks.

Your task is to change the xv6 file system code to support large files. You need to implement "doubly-indirect" blocks, containing 256 addresses of singly-indirect blocks, each of which can contain up to 256 addresses of data blocks. By modifying one direct block into a "doubly-indirect" block, a file will be able to consist of up to 65803 blocks (11 from direct blocks, 256 from singly-indirect blocks and 256 × 256 from doubly-indirect blocks). However, it is still not sufficient to accomplish our goal. You will need to implement extra doubly-indirect blocks to achieve up to 66666 blocks.

3.2 Guidelines and Hints

- 1. You can pass problem 1 with modifying only: fs.h and file.h, fs.c under xv6/kernel/.
- 2. Checkout TODO in the skeleton code.
- 3. kernel/fs.h describes the structure of an on-disk inode. The address of the data block is stored in addrs. Note that the length of addrs is always 13.
- 4. If you change the definition of NDIRECT, you probably have to change the declaration of addrs[] in struct inode in kernel/file.h. Make sure that struct inode and struct dinode have the same number of elements in their addrs[] arrays.

- 5. If you change the definition of NDIRECT, make sure to run make clean to delete fs.img. Then run make qemu to create a new fs.img, since mkfs uses NDIRECT to build the file system.
- 6. Make sure you understand bmap(). Write out a diagram of the relationships among ip->addrs[], indirect blocks, doubly-indirect blocks, and data blocks. Make sure you understand why replacing a direct block with a doubly-indirect block increases the maximum file size by 256 × 256 1 blocks.
- 7. You should allocate indirect blocks and doubly-indirect blocks only as needed, like the original bmap().
- 8. Don't forget to brelse() each block that you bread().
- 9. Make sure itrunc() frees all blocks of a file, including doubly-indirect blocks.

3.3 How to Test

There are 2 public testcases (6% each) and 3 private testcases (6% each). You can run make bigfile in the container to check your grade for public testcases. The public testcases are implemented in bigfile. You can change variable target to test your bigfile. As to private testcases, they will be added to bigfile_private.c while grading. Note that the public testcases do not reach 66666 blocks, so make sure your code can handle such large file to pass private testcases.

```
# make bigfile
......qemu log......
Testing large files: (38.2s)
  Large files: public testcase 1 (6 points): OK
  Large files: public testcase 2 (6 points): OK
Score: 12/12
```

4 Problem 2: Symbolic Links to Files (30 points)

4.1 Description

In this problem, you have to add symbolic links to xv6. Symbolic links (or soft links) refer to a linked file by pathname; when a symbolic link is opened, the kernel follows the link to the referred file. Symbolic links resemble hard links, but hard links are restricted to pointing to files on the same disk, while symbolic links can cross disk devices. Although xv6 doesn't support multiple devices, implementing this system call is a good exercise to understand how pathname lookup works. You will implement the symlink(char *target, char *path) system call, which creates a new symbolic link at path that refers to a file named target. In addition, you also need to handle open when encountering symbolic links. If the target is also a symbolic link, you must recursively follow it until a non-link file is reached. If the links form a cycle, you must return an error code. You may approximate this by returning an error code if the depth of links reaches some threshold (e.g., 20). However, when a process specifies O_NOFOLLOW flags, open should open symbolic links (not targets).

4.2 Guidelines and Hints

- 1. You can pass problem 2 with modifying only: sysfile.c under xv6/kernel/.
- 2. Checkout TODO in the skeleton code.
- 3. Checkout kernel/sysfile.c. There is an unimplemented function sys_symlink(). Note that the system call symlink() is already added in xv6, so you don't need to worry about that.
- 4. Checkout kernel/stat.h. There is a new file type T_SYMLINK, which represents a symbolic link.
- 5. Checkout kernel/fcntl.h. There is a new flag O_NOFOLLOW that can be used with the open system call.
- 6. The target does not need to exist for the system call to succeed.
- 7. You will need to store the target path in a symbolic link file, for example, in inode data blocks.
- 8. symlink should return an integer representing O(success) or -1(failure) similar to link and unlink.

- 9. Modify the open system call to handle paths with symbolic links. If the file does not exist, open must fail.
- 10. Don't worry about other system calls (e.g., link and unlink). They must not follow symbolic links; these system calls operate on the symbolic link itself.
- 11. You do not have to handle symbolic links to directories in this part.

4.3 How to Test

There are 2 public testcases (9% each) and 3 private testcases (4% each). You can run make symlinkfile in the container to check your grade for public testcases. The public testcases are implemented in symlinkfile. Feel free to check it out. As to private testcases, they will be added to symlinkfile_private.c while grading.

```
# make symlinkfile
......qemu log......
Testing symbolic links to files (public): (2.1s)
   Symbolic links to files: public testcase 1 (9 points): OK
   Symbolic links to files: public testcase 2 (9 points): OK
Score: 18/18
```

5 Problem 3: Symbolic Links to Directories (30 points)

5.1 Description

Instead of just implementing symbolic links to files, now you should also consider symbolic links to directories. We expect that a symbolic link to a directory should have these properties:

- 1. It can be part of a path, and will redirect to what it links to.
- 2. You can cd a symbolic link if it links to a directory.

For example, symlink("/y/", "/x/a") creates a symbolic link /x/a links to /y/. The actual path of /x/a/b should be /y/b. Thus, if you write to /x/a/b, you actually write to /y/b. Also, if you cd into /x/a, your working directory should become /y/.

5.2 Guidelines and Hints

- 1. You can pass problem 3 with modifying only: sysfile.c and fs.c under xv6/kernel/.
- 2. Checkout TODO in the skeleton code.
- 3. You can leave sys_symlink function unchanged, since symbolic links store paths as strings. There is no difference between a file path and a directory path.
- 4. You have to handle paths that consist of symbolic links. Check namex function in fs.c.
- 5. You have to handle symbolic links in sys_chdir function. Like problem 2, you need to avoid infinite loops.

5.3 How to Test

There are 2 public testcases (9% each) and 3 private testcases (4% each). You can run make symlinkdir in the container to check your grade for public testcases. The testcases are implemented in symlinkdir. Feel free to check it out. As to private testcases, they will be added to symlinkdir_private.c while grading.

```
# make symlinkdir
.....qemu log.....
Testing symbolic links to directories (public): (2.1s)
  Symbolic links to directories: public testcase 1 (9 points): OK
  Symbolic links to directories: public testcase 2 (9 points): OK
  Score: 18/18
```

6 Report (10 points)

6.1 Description

Your report should include two parts:

- 1. Briefly answer each problem. If your solutions are reasonable and correct, you can get points. Note that you can get points even if you fail to solve the problems. (10 points)
 - (a) Problem 1: Large Files (4 points)
 Please draw your inode block structure and explain how to access 66666th block in your implementation.
 - (b) Problem 2: Symbolic Links to Files (3 points)
 Function symlink("/b", "/a") creates a symbolic link /a links to file /b, explain how to open symbolic link /a.
 - (c) Problem 3: Symbolic Link to Directories (3 points)

 Function symlink("/y/", "/x/a/") creates a symbolic link /x/a links to directory /y/, explain how to write to /x/a/b.
- 2. We encouraged you to help other students. Please describe how you helped other students here. You should make the descriptions as short as possible, but you should also make them as concrete as possible (e.g., you can screenshot how you answered other students' questions on NTU COOL). Please note that you will not get any penalty if you leave it empty here. Please also note that this bonus is not for you to do optimization, so we will not release the grading criteria and the grades. Regarding the final letter grades, it is very likely that this does not help you will get promoted to the next level only if you are near the boundary of levels and you have significant contributions. (0 points)

7 Submissions

$7.1 \quad xv6 \text{ code}$

Please compress your xv6 source code as <whatever>.zip and upload it to NTUCOOL. The filename does not matter since NTUCOOL will rename your submissions. Never compress files we do not request, such as .o, .d, .asm files. You can run make clean in the container before you compress. Make sure your xv6 can be compiled by make qemu with the Makefile provided by TA.

Note that the English letters in the <student_id> must be lowercase. E.g., it should be b12123a23 instead of B12123A23.

7.2 Report

Submit your report to Machine Problem 4 on Gradescope in one PDF file.

8 Grading Policy

- You will get 0 points if we cannot compile your submission.
- You will be deducted 10 points if the folder structure is wrong. Using uppercase in the <student_id> is also a type of wrong folder structure.
- If your submission is late for n days, your score will be $max(raw_score 20 \times \lceil n \rceil, 0)$. Note that you will not get any points if $\lceil n \rceil \geq 5$.
- You can submit your work as many times as you want, but only the last submission will be graded. Previous submissions will be ignored.